5/21/22, 8:42 PM Subset Sum Equal To K



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C++ (g++ 5.4)

Subset Sum Equal To K

Difficulty: MEDIUM

Problem

Avg. time to solve Success Rate **30 min 65%**

Problem Statement Suggest Edit

Submissions

You are given an array/list 'ARR' of 'N' positive integers and an integer 'K'. Your task is to check if there exists a subset in 'ARR' with a sum equal to 'K'.

Note: Return true if there exists a subset with sum equal to 'K'. Otherwise, return false.

For Example:

If 'ARR' is $\{1,2,3,4\}$ and 'K' = 4, then there exists 2 subsets with sum = 4. These are $\{1,3\}$ and $\{4\}$. Hence, return true.

Input Format:

The first line contains a single integer T representing the number of test cases.

The first line of each test case contains two spaceseparated integers 'N' and 'K' representing the size of the input 'ARR' and the required sum as discussed above.

The next line of each test case contains 'N' single space-separated integers that represent the elements of the 'ARR'.

Output Format:

For each test case, return true or false as discussed above.

Output for each test case will be printed in a separate line.

Note:

You don't need to print anything, it has already been taken care of. Just implement the given function.

Constraints:

```
1 <= T <= 5

1 <= N <= 10^3

0 <= ARR[i] <= 10^9

0 <= K <= 10^9

Time Limit: 1 sec
```

Sample Input 1:

```
2
4 5
4 3 2 1
5 4
2 5 1 6 7
```

Sample Output 1:

true false

Explanation For Sample Input 1:

```
In example 1, 'ARR' is \{4,3,2,1\} and 'K' = 5. There exist 2 subsets with sum = 5. These are \{4,1\} and \{3,2\}. Hence, return true.
```

In example 2, 'ARR' is $\{2,5,1,6,7\}$ and 'K' = 4. There are no subsets with sum = 4. Hence, return false.

Sample Input 2:

```
2
4 4
6 1 2 1
5 6
1 7 2 9 10
```

Sample Output 2:

.

```
exist
ence,
e are
```

```
1 • /*
   In both the approach differ in space complexity
2
3
        Approach 1: Space Complexity = O(n * sum of all elements)
4
5
        Approach 2: Space COmplexity = O(n * k)
6
    */
7
8
            //Approach 1: Vary or keep the count of Addition of
                         //Elements in subset (dp array size = [n][sum of all elements])
9
10
11
12 bool solve(int n, int sumTillNow, int k, vector<vector<int>> &dp, vector<int> &arr){
13
        if(sumTillNow == k)
14
15
            return dp[n][sumTillNow] = true;
16
        if(n == 0)
17
            return dp[n][sumTillNow] = false;
18
        if(dp[n][sumTillNow] != -1)
19
20
            return dp[n][sumTillNow];
21
        bool include = solve(n - 1, sumTillNow + arr[n - 1], k, dp, arr);
22
        bool nInclude = solve(n - 1, sumTillNow, k, dp, arr);
23
24
        return dp[n][sumTillNow] = (include or nInclude);
25
26 ▼ bool subsetSumToK(int n, int k, vector<int> &arr) {
27
28
        // Write your code here.
29
        int sum = 0;
        for(int i = 0; i < n; i++)</pre>
30
31
            sum += arr[i];
32
33
34
35
        vector<vector<int>> dp(n + 1, vector<int>(sum + 1, - 1));
36
        bool ans = solve(n, 0, k, dp, arr);
37
        return ans;
38
39
40
41
            //Approach 2: Vary or keep the count of Subtraction from Target/K
42
43 • /*
44
45
    bool solve(int n, int k, vector<vector<int>> &dp, vector<int> &arr){
46
        if(k < 0)
47
            return false;
        if(k == 0)
48
49
            return dp[n][k] = true;
        if(n == 0)
50
51
            return dp[n][k] = false;
52
53
        if(dp[n][k] != -1)
54
            return dp[n][k];
55
        bool include = solve(n - 1, k - arr[n - 1], dp, arr);
        bool nInclude = solve(n - 1, k, dp, arr);
56
57
58
        return dp[n][k] = (include or nInclude);
59
    bool subsetSumToK(int n, int k, vector<int> &arr) {
60
61
62
        // Write your code here.
63
64
65
        vector < vector < int >> dp(n + 1, vector < int > (k + 1, - 1));
66
67
        bool ans = solve(n, k, dp, arr);
68
        return ans;
69
70
71
72 */
```

Console ^

Previous