Distributional Sparsity of Kantorovich Solutions

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1 Writeup of results so far

In our project, we consider two sets of finite points X, Y where |X| = m and |Y| = n and $m \neq n$ necessarily. Further we look at the empirical measures

$$\mu = \frac{1}{m} \sum_{i=1}^{m} \delta_{x_i} \text{ and } \nu = \frac{1}{n} \sum_{j=1}^{n} \delta_{y_j}$$
 (1.1)

to motivate the problem of optimally transporting the mass from X to Y. The case when m=n is well studied, as Birkhoff's theorem implies that there is a transport plan bijection $\pi:X\to Y$ that permutes indices in Y such that π is optimal according to the Kantorovich problem. Recently, Steinerberger and Hosseini proved the following theorem:

Theorem 1. Let μ and ν be defined as in 1.1. Then there is a solution of the Kantorovich problem such that mass from each point in X is moved to at most $\frac{n}{\gcd(m,n)}$ different points in Y and that each point in Y receives mass from at most $\frac{m}{\gcd(m,n)}$ points in X.

In building off of this work and this theorem, we investigate the problem of knowing the distributions of the number of different points moved from X to Y, and vice-versa, as $\frac{n}{\gcd(m,n)}$ and $\frac{m}{\gcd(m,n)}$ may be large. As an example, if m=20 and n=30, then we know that each point in X is mapped to at most 3 points in Y, and each point in Y receives mass from at most 2 points in X. Yet, we seek to understand the frequency with which a point in X is mapped to at most 3 points in Y, for instance. On simple intuition, we know that



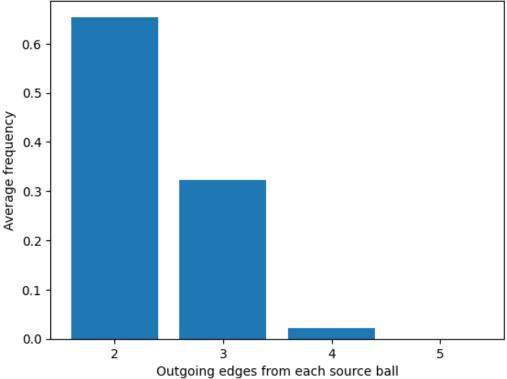


Figure 1

this exact scenario shouldn't happen which relatively high frequency, since points will optimally be matched to those points closest to them (on average). We investigate this question both numerically and theoretically. Specifically numerically, we look at specific interesting cases, such as when m, n are two (relatively) large, coprime numbers, and when m, n are again coprime but with large distance from each other. To our surprise, we found that the distribution of points tends to have small support, consistently grouping across a small number of points that are always close together.

Let's take a look at a couple more in-depth examples of what was motivated above. First, let's look at the distribution created by matching m=57 points in X to n=79 points in Y. We run a simulation of 200 trials and average the results. Note, clearly, m, n are coprime. The resulting distribution is shown in (1). Note that out of a maximum of 79 possible matchings, each

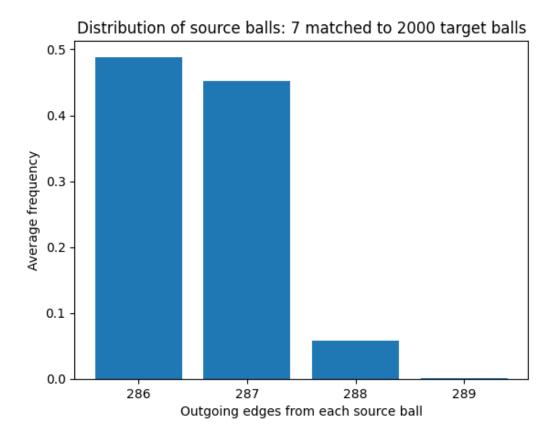


Figure 2

point in X only matches to 2, 3, 4, and 5 (with very low probability).

We now look at the case where m=7 and n=2000 to observe what happens when the distance between coprime integers becomes larger. We show the results in (2). Note that while the number of possible matchings must necessarily increase, the distribution of matchings from the source distribution to the target distribution is still sparse.

We seek to explain this behaviour theoretically. To see this, first we give a result on the transportation cost between n random and m points drawn from the same distribution but that are of different number.

Proposition 1. Given m points in X and n points in Y, with $m \neq n$, and μ and ν as defined in 1.1, then the transportation cost associated with

transporting mass μ to ν is of the order

$$\sqrt{m\log m} + \sqrt{n\log n}$$
.

Proof sketch. We use the triangle inequality property of Wasserstein distances to show this bound, along with the result from *On Optimal Matchings*. Let ζ_x be the flat measure on X and ζ_y be the flat measure on Y. It then follows that

$$W_p(\mu, \nu) \le W_p(\mu) \le W_p(\mu, \zeta_x) + W_p(\nu, \zeta_y)$$

$$\lesssim \sqrt{m \log m} + \sqrt{n \log n}.$$

From here, we are looking at relations to the distance of the nearest neighbors and other phenomena in forbidden subgraphs to see if there is such an explanation of the sparse distributions.

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