Team Notebook

Ari Blondal (SFU)

March 11, 2021

Contents		2.1 Dijkstra	2	3.1 Longest Increasing Subsequence	2
1 0 - Template	2	2.2 Set-Union	2	4 Mathematics	2
2 Basics	$2 \mid_3$	Dynamic Programming	2	4.1 FFT and Multiplication	2

SFU

1 0 - Template

```
#include <bits/stdc++.h>
using namespace std:
// incomplete
#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef long long 11;
typedef pair<int, int> pii;
typedef vector<int> vi;
int main() {
ios_base::sync_with_stdio(false);
cin.tie(NULL):
int t;
cin >> t:
while(t--){
}
```

2 Basics

2.1 Dijkstra

```
/* Taken from cp-algorithms.com */
/* O(m log m) */
const int INF = 1e9:
vi adi[MAXN]:
void dij(int s, int n, vi &d, vi &p){
d.assign(n, INF);
p.assign(n, -1);
d[s] = 0;
priority_queue<pii, vector<pii>, greater<pii>> q;
q.push({0, s});
   while (!q.empty()) {
      int v = q.top().second;
      int d_v = q.top().first;
      q.pop();
      if (d v != d[v])
          continue;
       for (auto edge : adj[v]) {
          int to = edge.first;
```

```
int len = edge.second;

if (d[v] + len < d[to]) {
    d[to] = d[v] + len;
    p[to] = v;
    q.push({d[to], to});
    }
}</pre>
```

2.2 Set-Union

```
/* Written by Ari */
/* If joining by size, -UF[root] is size of set */
const int INF = 1e7;
int UF[MAXN] = {0};

int find(int x){
  return UF[x] < 0 ? x : UF[x] = find(UF[x]);
}

void init(int n){
  rep(i,0,n) UF[i] = -1;
}

bool join(int a, int b){
  a = find(a), b = find(b);
  if (a==b) return false;
  if (UF[b] < UF[a]) swap(a,b);
  UF[a] += UF[b];
  UF[b] = a;
  return true;
}</pre>
```

3 Dynamic Programming

3.1 Longest Increasing Subsequence

```
// From doggy-sweat-cheatsheet by Erfaniaa
void reconstruct_print(int end, int a[], int p[]) {
  int x = end;
  stack<int> s;
  for (; p[x] >= 0; x = p[x]) s.push(a[x]);
  printf("[%d", a[x]);
  for (; !s.empty(); s.pop()) printf(", %d", s.top());
```

```
printf("]\n"):
int main() {
 int n = 11, A[] = \{-7, 10, 9, 2, 3, 8, 8, 1, 2, 3, 4\};
 int L[MAX N], L id[MAX N], P[MAX N];
 int lis = 0, lis_end = 0;
 for (int i = 0; i < n; ++i) {</pre>
   int pos = lower_bound(L, L + lis, A[i]) - L;
   L[pos] = A[i];
   L id[pos] = i:
   P[i] = pos ? L_id[pos - 1] : -1;
   if (pos + 1 > lis) {
    lis = pos + 1;
     lis_end = i;
   printf("LIS ending at A[%d] is of length %d: ", i, pos +
   reconstruct_print(i, A, P);
   printf("\n");
 printf("Final LIS is of length %d: ", lis);
 reconstruct_print(lis_end, A, P);
 return 0;
```

4 Mathematics

4.1 FFT and Multiplication

```
//From "You Know Izad?" team cheatsheet
#define base complex<double>
void fft (vector<base> & a, bool invert){
   if (L(a) == 1) return:
   int n = L(a);
   vector \langle base \rangle a0(n / 2), a1(n / 2);
   for (int i = 0, j = 0; i < n; i += 2, ++j){
       a0[i] = a[i]:
       a1[j] = a[i + 1];
   }
   fft (a0, invert);
   fft (a1. invert):
   double ang = 2 * PI / n * (invert ? -1 : 1);
   base w(1), wn(cos(ang), sin(ang));
   fore(i, 0, n / 2) {
       a[i] = a0[i] + w * a1[i];
```

SFU

```
3
```

```
3
    a[i + n / 2] = a0[i] - w * a1[i];
    if (invert)
        a[i] /= 2, a[i + n / 2] /= 2;
    w *= wn;
}
```