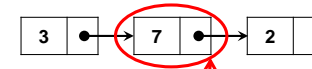


Lists of Nodes

- Linked lists use pointers to go to the next element



- each block is called a **node**

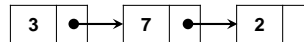
Let's implement it:

- a node consists of
 - a data element — an *int* here
 - a pointer to the next node

```
struct list_node {
    int data;
    struct list_node* next;
};
```

- The whole list is a pointer to its first node

Linked Lists



Lists of Nodes

```
struct list_node {
    int data;
    struct list_node* next;
};
```

- Linked lists are a **recursive type**
 - a *struct list_node* is defined in terms of itself

- What if we don't have this pointer?

a node that contains an *int* and
a node that contains an *int* and
a node that contains an *int* and
...

- It would take an *infinite amount of memory!*
- The C0 compiler disallows this
 - recursion can only occur behind a pointer (or an array)



Lists of Nodes

```
struct list_node {
    int data;
    struct list_node* next;
};
```

- Let's make it more readable

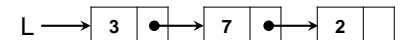
```
typedef struct list_node list; // ADDED

struct list_node {
    int data;
    list* next;
}; // MODIFIED
```

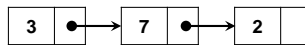
This can go before
or after the struct

- Implementing this linked list

```
list* L = alloc(list);
L->data = 3;
L->next = alloc(list);
L->next->data = 7;
L->next->next = alloc(list);
L->next->next->data = 2;
```



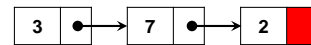
Lists of Nodes



- Does this help us implement queues?
 - Linked lists can be arbitrarily large or small
 - use just the nodes we need
 - size is not fixed like arrays
 - It's easy to insert an element at the beginning
 - allocate a new node and point its next field to the list
 - In fact, it's easy to insert an element between any two nodes
 - allocate a new node and move pointers around
- What about inserting an element at the end?
 - How do we indicate the end of a linked list?

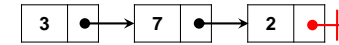
So far we just drew an empty box ...

The End of a List



We need to make the pointer in the last node **special**

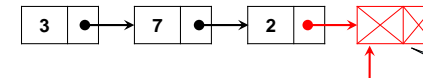
- Use the NULL pointer



This is a great idea if we don't need direct access to the end of the list

➢ This is a **NULL-terminated list**

- Point it to a special node we keep track of somewhere

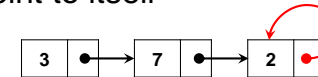


This is a great idea if we **do** need direct access to the end of the list

➢ We know we reached the end of the list if its next field is equal to the address of the dummy node

This node is called the **dummy node** or the **sentinel**

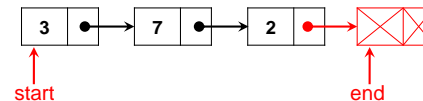
- Have it point to itself



This works too, but nobody does that

Lists with a Dummy Node

- We need to keep track of *two* pointers



- **start**: where the first node is
- **end**: the address in the next field of the last node
 - the address of the dummy node

- What's in the dummy node?

- some values that are not important to us
 - some number and some pointer
- we say its fields are *unspecified*
 - no way to test for "unspecified"

These values are not special in any way:

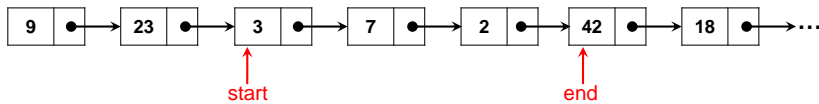
- data could be any element
- next may or may not be NULL

- A dummy value is a value we **don't care** what it is

List Segments

List Segments

- There may be more nodes before and after



- The pair of pointers **start** and **end** identify our list exactly

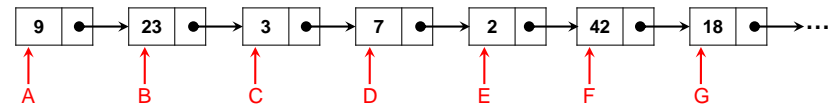
- **start** is **inclusive** (the first node of the list)
- **end** is **exclusive** (one past the last node of the list)

points to the dummy node

- They identify the **list segment** [start, end)

- here it contain values 3, 7 and 2
- similar to array segments A[lo, hi)

- There are many list segments in a list



- The list segment [C, F) contains elements 3, 7, 2
 - its dummy node contains 42 and the pointer G
- The list segment [A, G) contains 9, 23, 3, 7, 2, 42
 - its dummy node contains 18 and the some pointer
- The list segment [B, D) contains 23, 3
 - its dummy node contains 7 and the pointer E
- The list segment [C, C) contains no elements
 - its dummy node contains 3 and the pointer D
 - this is the **empty segment**
 - any segment where **start** is the same as **end**
 - [A, A), [B, B), ...

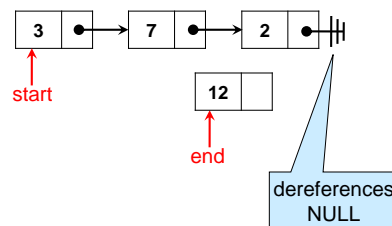
Checking for List Segments

```
typedef struct list_node list;
struct list_node {
    int data;
    list* next;
};
```

- We want to write a specification function that checks that two pointers **start** and **end** form a list segment

- Follow the next pointer from **start** until we reach **end**

```
bool is_segment(list* start, list* end) {
    list* l = start;
    while (l != end) {
        l = l->next;
    }
    return true;
}
```



- Does this work?

- the dereference l->next may not be safe
 - we need NULL-checks!
- we never return false

✗

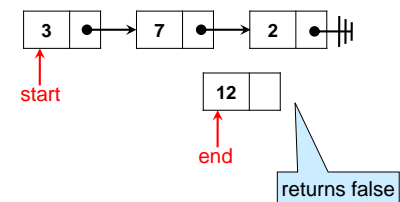
Checking for List Segments

```
typedef struct list_node list;
struct list_node {
    int data;
    list* next;
};
```

- We want to write a specification function that checks that two pointers **start** and **end** form a list segment

- Follow the next pointer from **start** until we reach **end**

```
bool is_segment(list* start, list* end) {
    list* l = start;
    while (l != NULL) {
        if (l == end) return true; // MODIFIED
        l = l->next;
    }
    return false; // MODIFIED
}
```



- Does this work?

- if there is a list segment from **start** to **end**, it will return true
- if it returns false, there is no list segment from **start** to **end**

- It works then ...

returns false

Checking for List Segments

```
typedef struct list_node list;
struct list_node {
    int data;
    list* next;
};
```

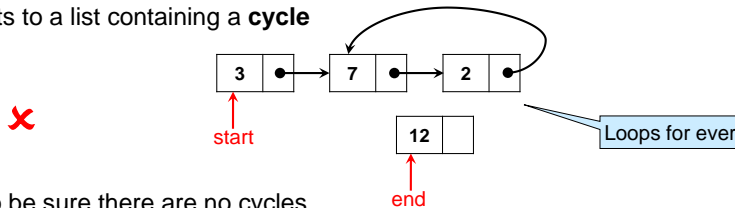
- A function that checks that **start** and **end** form a list segment

```
bool is_segment(list* start, list* end) {
    list* l = start;
    while (l != NULL) {
        if (l == end) return true;
        l = l->next;
    }
    return false;
}
```

- if there is a list segment from **start** to **end**, it will return true
- if it returns false, there is no list segment from **start** to **end**

- Can there be no list segment but it does not return false

- if start points to a list containing a **cycle**



- We need to be sure there are no cycles

Checking for List Segments

```
typedef struct list_node list;
struct list_node {
    int data;
    list* next;
};
```

- A function that checks that **start** and **end** form a list segment

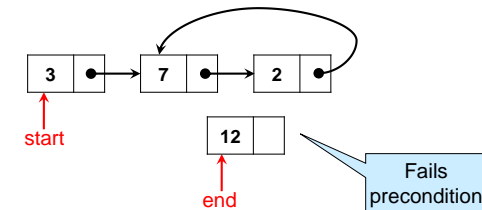
- We need to be sure there are no cycles

```
bool is_segment(list* start, list* end)
//@requires is_acyclic(start); // ADDED
{
    list* l = start;
    while (l != NULL) {
        if (l == end) return true;
        l = l->next;
    }
    return false;
}
```

We will implement it later

- Does this work?

- Yes!



Checking for List Segments

```
typedef struct list_node list;
struct list_node {
    int data;
    list* next;
};
```

- A function that checks that **start** and **end** form a list segment

```
bool is_segment(list* start, list* end)
//@requires is_acyclic(start);
{
    list* l = start;
    while (l != NULL) {
        if (l == end) return true;
        l = l->next;
    }
    return false;
}
```

- Notes:

- returns false if start == NULL
- or if end == NULL
 - ❑ NULL is not a pointer to a list node
 - ❑ subsumes NULL-check for both start and end

Checking for List Segments

```
typedef struct list_node list;
struct list_node {
    int data;
    list* next;
};
```

- We can also write it more succinctly

- using a for loop

```
bool is_segment(list* start, list* end)
//@requires is_acyclic(start);
{
    for (list* l = start; l != NULL; l = l->next) {
        if (l == end) return true;
    }
    return false;
}
```

All 3 versions are equivalent

- recursively

```
bool is_segment(list* start, list* end)
//@requires is_acyclic(start);
{
    if (start == NULL) return false;
    return start == end
        || is_segment(start->next, end);
}
```

Detecting Cycles

- How to check if a list is cyclic?

- Use a counter and look for overflows

➤ very inefficient!

➤ also, C0 pointers are 64 bits but **ints** are 32 bits

In C0, there are more pointers than integers!

- Keep track of visited nodes somewhere

➤ in an array?

how big to make it?

array indices are 32 bits

➤ in another list?

how do we check it has no cycles?

- Add a "visited" field to the nodes (a boolean)

➤ we need to know the list is acyclic to initialize it to false!

- What then?

Detecting Cycles

- The tortoise and hare algorithm

by this dude



Robert W. Floyd

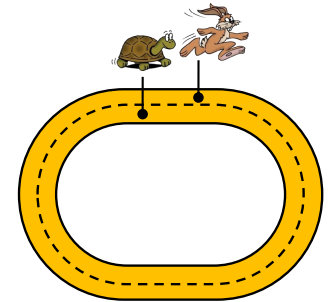
- Traverse the list using two pointers

➤ the tortoise starts at the beginning and moves by 1 step

➤ the hare starts just ahead of the tortoise and moves by 2 steps

- If the hare ever overtakes the tortoise, there is a cycle

```
bool is_acyclic(list* start) {
    if (start == NULL) return true;
    list* t = start;           // tortoise
    list* h = start->next;     // hare
    while (h != t) {
        if (h == NULL || h->next == NULL) return true;
        //@assert t != NULL; // hare hits NULL quicker
        t = t->next;           // tortoise moves by 1 step
        h = h->next->next;      // hare moves by 2 steps
    }
    //@assert h == t;           // hare has overtaken tortoise
    return false;
}
```



Detecting Cycles

- The tortoise and hare algorithm

```
bool is_acyclic(list* start) {
    if (start == NULL) return true;
    list* t = start;           // tortoise
    list* h = start->next;     // hare
    while (h != t) {
        if (h == NULL || h->next == NULL) return true;
        //@assert t != NULL; // hare hits NULL quicker
        t = t->next;           // tortoise moves by 1 step
        h = h->next->next;      // hare moves by 2 steps
    }
    //@assert h == t;           // hare has overtaken tortoise
    return false;
}
```

- Returns

➤ true if there is no cycle

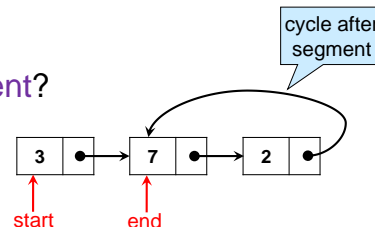
➤ false if there is a cycle

- Does it fix our problem with **is_segment**?

- Too aggressive

- Exercise: fix it!

Hint: you need to account for end



Manipulating List Segments

Deleting an Element

- How do we remove the node **at the beginning** of a non-empty list segment **[start, end)**?

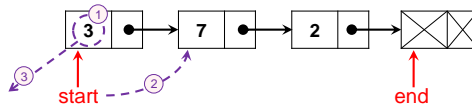
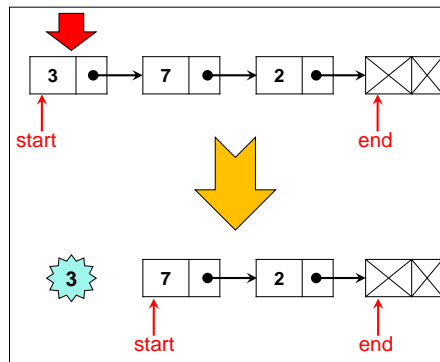
➤ and return the value in there

- grab the value in the **start** node
- move **start** to point to the next node
- return the value

```
1 int x = start->data;
2 start = start->next;
3 return x;
```

○ **Complexity:** $O(1)$

Note: we are not "deleting" the node, just making the segment shorter



Deleting an Element

- How do we remove the **last** node of a non-empty list segment **[start, end)**?

➤ and return the value in there

- we must go from **start**

❑ **end** is one node too far

- follow next until just before **end**
- move **end** to that node
- return its value

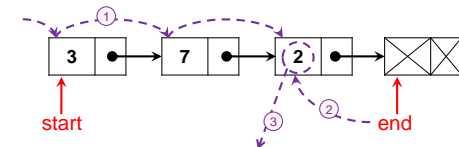
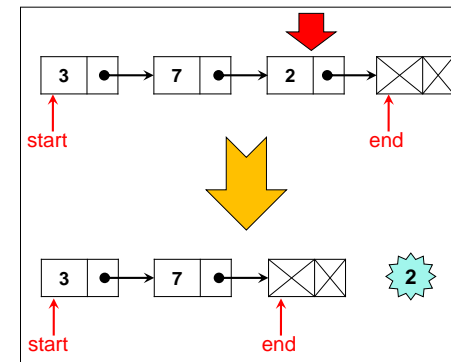
```
1 list* l = start;
2 while (l->next != end)
3   l = l->next;
4 end = l;
5 return l->data;
```

○ **Complexity:** $O(n)$

Expensive!

Notes:

- The old last node becomes the new dummy node
- We are not "deleting" anything, just making the segment shorter



Inserting an Element

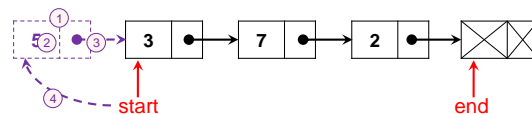
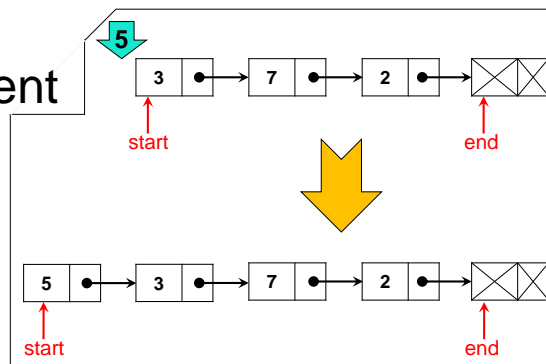
- How do we add a node **at the beginning** of a list segment **[start, end)**?

- create a new node
- set its data field to the value to add
- set its next field to **start**
- set **start** to it

```
1 list* l = alloc(list);
2 l->data = x;
3 l->next = start;
4 start = l;
```

○ **Complexity:** $O(1)$

Note: we are adding a brand new node



Inserting an Element

- How do we add a node **as the last** node of a list segment **[start, end)**?

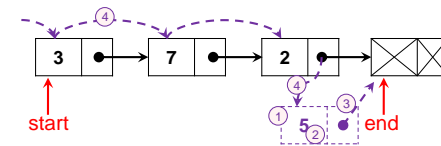
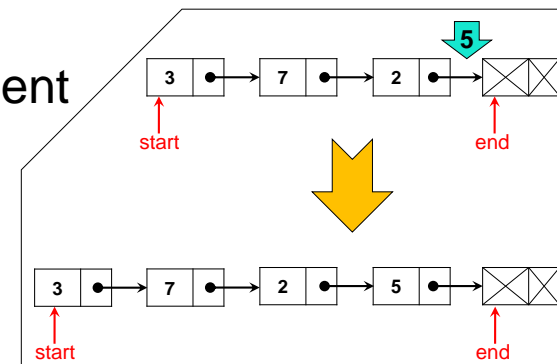
- create a new node
- set its data field to the value to add
- set its next field to **end**
- point the old last node to it

```
1 list* new_last = alloc(list);
2 new_last->data = x;
3 new_last->next = end;
4 list* l = start;
5 while (l->next != end)
6   l = l->next;
7 l->next = new_last;
```

○ **Complexity:** $O(n)$

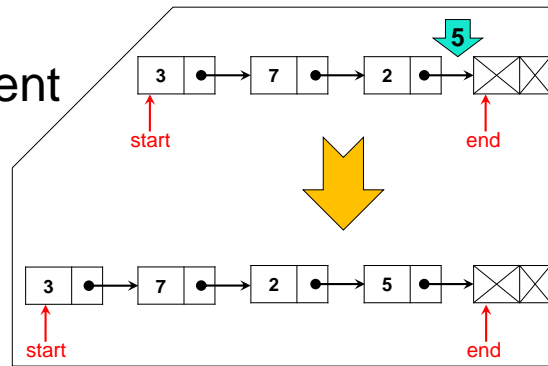
Expensive!

Note: we are adding a new last node, but we modify the next pointer of the old last node



Inserting an Element

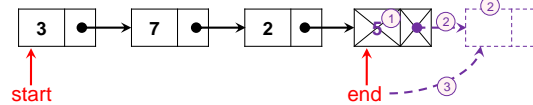
- How do we add a node as the last node of a list segment [start, end)?



Can we do better?

- set the data field of **end** to the value to add
- set its next field to a new dummy node
- set **end** to it

- end->data = x;
- end->next = alloc(list);
- end = end->next;



Complexity: $O(1)$

Much better!

Note: we are using the old dummy node as the new last node, and creating a new dummy

Summary

	at the beginning	at the end
Inserting	$O(1)$	$O(1)$
Deleting	$O(1)$	$O(n)$

Good

Bad

- We will use this as a guide when implementing queues (and stacks) to achieve their complexity goals