Cheat Sheet

Logical Equivalences

Identity: $p \wedge \mathbf{T} \equiv p$

 $p \vee \pmb{F} \equiv p$

Domination: $p \lor T \equiv T$

 $p \wedge \boldsymbol{F} \equiv \boldsymbol{F}$

Idempotent: $p \lor p \equiv p$

 $p \wedge p \equiv p$

Double negation: $\neg(\neg p) \equiv p$

Commutative: $p \lor q \equiv q \lor p$

 $p \wedge q \equiv q \wedge p$

Associative: $(p \lor q) \lor r \equiv p \lor (q \lor r)$

 $(p \land q) \land r \equiv p \land (q \land r)$

Distributive: $p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$

 $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

De Morgan's: $\neg (p \land q) \equiv \neg p \lor \neg q$

 $\neg(p \lor q) \equiv \neg p \land \neg q$

Absorption: $p \lor (p \land q) \equiv p$

 $p \land (p \lor q) \equiv p$

Negation: $p \lor \neg p \equiv T$

 $p \wedge \neg p \equiv \mathbf{F}$

Conditional Statements

$$p \to q \equiv \neg p \vee q$$

$$p \to q \equiv \neg q \to \neg p$$

$$p \vee q \equiv \neg p \to q$$

$$p \land q \equiv \neg (p \to \neg q)$$

$$(p \to q) \land (p \to r) \equiv p \to (q \land r)$$

$$(p \to r) \land (q \to r) \equiv (p \lor q) \to r$$

$$(p \to q) \lor (p \to r) \equiv p \to (q \lor r)$$

$$(p \to q) \lor (p \to r) = P \lor (q \lor r)$$

 $(p \to r) \lor (q \to r) \equiv (p \land q) \to r$

Biconditional Statements

$$\begin{array}{l} p \iff q \equiv (p \to q) \land (q \to p) \\ p \iff q \equiv \neg p \iff \neg q \\ p \iff q \equiv (p \land q) \lor (\neg p \land \neg q) \\ \neg (p \iff q) \equiv p \iff \neg q \end{array}$$

Rules of Inference

$\begin{array}{c} p \\ \underline{p \to q} \\ \vdots \\ q \end{array}$	Modus ponens
$ \begin{array}{c} \neg q \\ \underline{p \to q} \\ \vdots \neg p \end{array} $	Modus tollens
$\begin{array}{ c c }\hline p \rightarrow q \\ \hline q \rightarrow r \\ \hline \therefore p \rightarrow r \\ \hline \end{array}$	Hypothetical syllogism
$ \begin{array}{c} p \lor q \\ \hline \neg p \\ \hline \therefore q \end{array} $	Disjunctive syllogism
$\frac{p}{\therefore p \vee q}$	Addition
$\frac{p \wedge q}{\therefore p}$	Simplification
$\begin{array}{c} p \\ \underline{q} \\ \vdots p \wedge q \end{array}$	Conjunction
$ \begin{array}{c c} p \lor q \\ \hline \neg p \lor r \\ \hline \therefore q \lor r \end{array} $	Resolution

Logic Gates

