

Analysis of Algorithms

Homework 4

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1 Dijkstra's Algorithm with negative weights

Part A

Figure 1 shows a graph with negative weights such that if we apply Dijkstra's algorithm to find the shortest path between vertices S and D , it will return the wrong path.

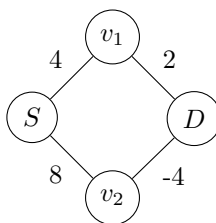


Figure 1: Graph for which Dijkstra doesn't work

Starting off, we assign the unvisited vertices v_1 and v_2 with the distances 4 and 8 respectively, marking S as visited. Then we take the unvisited vertex with the smallest distance, v_1 , and check its neighbours, namely D . We assign it the distance 6 and mark v_1 as visited. Now that our destination vertex is the unvisited vertex with the shortest distance, the algorithm would claim that it has finished, with the shortest path going through v_1 with a distance of 6.

However, in reality the shortest path goes through v_2 and has a total distance of $8 - 4 = 4$. This path was not considered by the algorithm because the path to the intermediate vertex v_2 has a larger distance than the path it found first.

Part B

If we take the example graph in figure 1 and modify it so that the edges are directed (away from S or towards D), then that would form a directed acyclic

graph for which Dijkstra's algorithm would not work for a similar reason to that of part A.

2 Floyd-Warshall with Negative Cycles

We are to add pseudocode to the Floyd-Warshall algorithm which checks for negative cycles. First, let us take a look at the algorithm.

```

function FLOYDWARSHALLNEGATIVECYCLES( $V, E$ )
  for  $(u, v) \in V \times V$  do                                ▷ Initialise all distances to infinity
     $D_{u,v} := \infty$ 
  for  $(u, v) \in E$  do                                       ▷ Apply distances of each edge
     $D_{u,v} := \text{WEIGHT}(u, v)$ 
  for  $v \in V$  do                                           ▷ Set distance to itself to zero
     $D_{v,v} := 0$ 
  for  $k \in V$  do      ▷  $k$  is a possible intermediate vertex between all  $(u, v)$ 
    for  $u \in V$  do
      for  $v \in V$  do
         $D_{u,v} := \text{MIN}(D_{u,k} + D_{k,v}, D_{u,v})$           ▷ Seek shorter path via  $k$ 

```
