

# Analysis of Algorithms

## Homework 4

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December 21, 2020

### 1 Dijkstra's Algorithm with negative weights

#### Part A

Figure 1 shows a graph with negative weights such that if we apply Dijkstra's algorithm to find the shortest path between vertices  $S$  and  $D$ , it will return the wrong path.

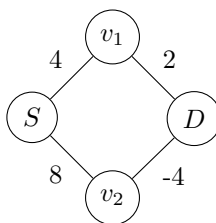


Figure 1: Graph for which Dijkstra doesn't work

Starting off, we assign the unvisited vertices  $v_1$  and  $v_2$  with the distances 4 and 8 respectively, marking  $S$  as visited. Then we take the unvisited vertex with the smallest distance,  $v_1$ , and check its neighbours, namely  $D$ . We assign it the distance 6 and mark  $v_1$  as visited. Now that our destination vertex is the unvisited vertex with the shortest distance, the algorithm would claim that it has finished, with the shortest path going through  $v_1$  with a distance of 6.

However, in reality the shortest path goes through  $v_2$  and has a total distance of  $8 - 4 = 4$ . This path was not considered by the algorithm because the path to the intermediate vertex  $v_2$  has a larger distance than the path it found first.

#### Part B

If we take the example graph in figure 1 and modify it so that the edges are directed (away from  $S$  or towards  $D$ ), then that would form a directed acyclic

graph for which Dijkstra's algorithm would not work for a similar reason to that of part A.

## 2 Floyd-Warshall with Negative Cycles

We are to add pseudocode to the Floyd-Warshall algorithm which checks for negative cycles. First, let us take a look at the algorithm.

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function FLOYDWARSHALLNEGATIVECYCLES( $V, E$ )
  for  $(u, v) \in V \times V$  do                                 $\triangleright$  Initialise all distances to infinity
     $D_{u,v} := \infty$ 
  for  $(u, v) \in E$  do                                       $\triangleright$  Apply distances of each edge
     $D_{u,v} := \text{WEIGHT}(u, v)$ 
  for  $v \in V$  do                                           $\triangleright$  Set distance to itself to zero
     $D_{v,v} := 0$ 
  for  $k \in V$  do     $\triangleright k$  is a possible intermediate vertex between all  $(u, v)$ 
    for  $u \in V$  do
      for  $v \in V$  do
         $D_{u,v} := \text{MIN}(D_{u,k} + D_{k,v}, D_{u,v})$      $\triangleright$  Seek shorter path via  $k$ 

```

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