

## 1 Question 1

### 1.1 Part a

$$P = (C \boxplus \bar{K}_1) \oplus K_0$$

where  $\bar{K}_1$  is the complement of  $K_1$ .

### 1.2 Part b

Let our plaintext messages be  $P_a, P_b$  and their corresponding ciphertexts be  $C_a, C_b$ . We know that

$$C_a = (P_a \oplus K_0) \boxplus K_1$$

$$C_b = (P_b \oplus K_0) \boxplus K_1$$

We can define  $K_0$  in terms of  $C_a$ ,  $P_a$ , and  $K_1$  by rearranging the variables.

$$(P_a \oplus K_0) \boxplus K_1 = C_a$$

$$P_a \oplus K_0 = C_a \boxplus \bar{K}_1$$

$$K_0 = (C_a \boxplus \bar{K}_1) \oplus P_a$$

Likewise,

$$K_0 = (C_b \boxplus \bar{K}_1) \oplus P_b$$

Which means

$$(C_a \boxplus \bar{K}_1) \oplus P_a = (C_b \boxplus \bar{K}_1) \oplus P_b$$

$$(C_a \boxplus \bar{K}_1) \oplus P_a \oplus P_b = C_b \boxplus \bar{K}_1$$

$$((C_a \boxplus \bar{K}_1) \oplus P_a \oplus P_b) \boxplus C_b = \bar{K}_1$$

We can then use  $\bar{K}_1$  to find  $K_0$  via

$$K_0 = (C_a \boxplus \bar{K}_1) \oplus P_a$$

## 2 Question 2

I will solve this problem for the general case. Let me denote the encryption of message  $m$  as  $E(m)$ , since the key is not relevant here. The defined encryption scheme being linear means

$$E(m_a) \oplus E(m_b) = E(m_a \oplus m_b)$$

This also implies that a message made of all zeroes will be encrypted to a message of all zeroes.

Now, let each input  $m_i$  and corresponding output  $E(m_i)$  be  $l$  bits long. Choose  $l$  ciphertexts  $E(m_i)$  such that

$$E(m_1) = 1000\dots$$

$$E(m_2) = 0100\dots$$

$$E(m_3) = 0010 \dots$$

...

$$E(m_{l-1}) = \dots 0010$$

$$E(m_l) = \dots 0001$$

I denote this set  $\mathcal{E}$ . If I know the corresponding plaintexts  $\{m_1, m_2, \dots, m_l\}$ , I can decipher any message using the linear property of  $E$ . Let  $E(m_k)$  be an intercepted message.  $E(m_k)$  can be described by XORing a unique subset of  $\mathcal{E}$ . Now because of the linear property of  $E$ , I can retrieve  $m_k$  by XORing the  $m$ 's that correspond to the aforementioned unique subset.

For example, let

$$E(m_k) = E(m_1) \oplus E(m_{17}) \oplus E(m_{42})$$

We can retrieve  $m_k$  via

$$m_k = m_1 \oplus m_{17} \oplus m_{42}$$