

CS 331: Theory of Computing

Problem Set 7

Iowa State University
Computer Science Department
March 21, 2013

Due by the midnight of March 29, 2013

Problem 1 (10 points)

Let $\mathcal{K} = \{\langle M \rangle \mid \langle M \rangle \notin \mathcal{L}(M)\}$. Prove that \mathcal{K} is **not** Turing recognizable.

Hint: Russell's Paradox.



Problem 2 (20 points)

Prove the following statements.

- ① $\mathcal{L}_1 \leq_m \mathcal{L}_2$ and $\mathcal{L}_2 \leq_m \mathcal{L}_3$ imply $\mathcal{L}_1 \leq_m \mathcal{L}_3$. (10 points)
- ② $\mathcal{L}_1 \leq_T \mathcal{L}_2$ implies that $\overline{\mathcal{L}_1} \leq_T \overline{\mathcal{L}_2}$. (10 points)

Note: \leq_T is Turing reduction (introduced in Chapter 5.1) and \leq_m is many-one reduction (introduced in Chapter 5.3).



Problem 3 (10 points)

Prove that $A_{TM} \not\leq_m \overline{A_{TM}}$.

Hint: Proof By Contradiction and Problem 2.



Problem 4 (20 points)

Which of the following PCP problems has a solution? Justify your answer.

$$(1) \left\{ \begin{array}{|c|} \hline ab \\ \hline a \\ \hline \end{array} \begin{array}{|c|} \hline bb \\ \hline ab \\ \hline \end{array} \begin{array}{|c|} \hline aa \\ \hline ba \\ \hline \end{array} \begin{array}{|c|} \hline cc \\ \hline bc \\ \hline \end{array} \begin{array}{|c|} \hline aa \\ \hline ca \\ \hline \end{array} \begin{array}{|c|} \hline d \\ \hline cd \\ \hline \end{array} \right\}$$
$$(2) \left\{ \begin{array}{|c|} \hline ab \\ \hline a \\ \hline \end{array} \begin{array}{|c|} \hline bb \\ \hline ab \\ \hline \end{array} \begin{array}{|c|} \hline aa \\ \hline ba \\ \hline \end{array} \begin{array}{|c|} \hline c \\ \hline bc \\ \hline \end{array} \begin{array}{|c|} \hline aa \\ \hline ca \\ \hline \end{array} \begin{array}{|c|} \hline d \\ \hline cd \\ \hline \end{array} \right\}$$



Problem 5 (20 points)

Does the following PCP problem P have a solution?

$$\left\{ \begin{array}{|c|c|c|c|c|c|} \hline a & b & c & c & dddd & ddde \\ \hline ab & ccc & b & d & & e \\ \hline \end{array} \right\}$$

Hint: Prove that P has a solution if and only if $\exists n (3^n \bmod 4) = 3$.



Problem 6 (20 points)

Let $\Sigma = \{0, 1, -\}$ (where $-$ denotes whitespace) be the tape alphabet for all TMs in this problem. Define the busy beaver function $BB : \mathbb{N} \rightarrow \mathbb{N}$ as follows. For each value of k , consider all k -state TMs that halt when started with a blank tape. Let $BB(k)$ be the maximum number of steps that can be performed by a k -state Turing machine. Show that BB is not a computable function.

Hint: Halting Problem.

