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## 10. VC1

**Program Name:** VC1.java

**Input File:** vc1.dat

In distributed computing, there is a concept known as a “vector clock”. Assume you have a distributed system consisting of  $k$  processes, and each process runs a series of events, and communicates with the other processes by sending messages. These events can either be local events, a send event, or a receive event. The vector clock of an event  $e$  contains, for each process, the entire history of events that “happen before”  $e$ . Instead of a long list of events though, you can show that there is one last event on each process that happens before  $e$ , so you can represent the entire history by storing the event number of the last event of each process that happens before  $e$  in a vector, creating a “vector clock”. With this, we can build a relation between vector clocks called “happens-before”, or  $\rightarrow$ . The syntax we will use is: given  $p$  processes, a vector clock  $VC$  for an event  $e$  is an array of integers of length  $p$ , where  $VC[i]$  is the last event number on process  $i$  that “**happened-before**”  $e$ . For example, for event  $e$  = event 3 on process 1 (0-indexed), with 3 total processes, the vector clock could be  $[2, 3, 1]$ , or  $[0, 3, 0]$ , or  $[4, 1, 7]$ , etc....

Vector clocks have the following properties with respect to the “happens-before” relation:

If two events are on the same process, then the earlier event  $\rightarrow$  the later event.

If event  $a$  is the send event for a message  $M$  and event  $b$  is the receive event for a message  $M$ , then  $a \rightarrow b$ .

If  $a \rightarrow b$  and  $b \rightarrow c$ , then  $a \rightarrow c$  (transitivity).

Given two events  $a$  and  $b$ , if neither  $a \rightarrow b$  nor  $b \rightarrow a$ , then they are said to be “concurrent”, or  $a \parallel b$ .

Since vector clocks contain the last event that “happens-before”  $e$  on each process, given two vector clocks  $A$  and  $B$  for events  $a$  and  $b$  respectively,  $a \rightarrow b$  when for all processes  $i$ ,  $A[i] \leq B[i]$ . In other words, for every entry  $i$  in the vector clock, the last event number  $A$  has for process  $i$  must be less than or equal to the last event number  $B$  has for process  $i$ .

Also, if no event on process  $i$  has  $\rightarrow e$ , we say  $E[i] = 0$ .

Given the vector clocks of two events  $a$  and  $b$ , your job is to figure out whether  $a \rightarrow b$ ,  $b \rightarrow a$ , or  $a \parallel b$  (they are 'concurrent' -- neither  $a \rightarrow b$  nor  $b \rightarrow a$ ).

### Input

The first line will contain the number of test cases  $T$ .

Each test case contains three lines. The first line is the number of processes  $P$ . The second line is the vector clock  $A$  represented by  $P$  space-separated integers. The third line is the vector clock  $B$ , in the same format as  $A$ .

### Output

Output “ $A \rightarrow B$ ” if  $A \rightarrow B$ , “ $B \rightarrow A$ ” if  $B \rightarrow A$ , and “ $A \parallel B$ ” if they are concurrent. It is guaranteed that  $A$  and  $B$  will never be equal, so exactly one of these is always true.

### Constraints

$1 \leq T \leq 10$

$0 \leq P \leq 20$

$1 \leq \text{each event number} \leq 10^9$

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**Example Input File**

```
3
2
1 1
2 1
3
3 3 3
1 2 3
2
1 2
2 1
```

**Example Output to Screen**

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A -> B
B -> A
A || B
```

**Explanation of the example**

In the first test case, all numbers in A are  $\leq$  all numbers in B. In the second case, all numbers in B are  $\leq$  all numbers in A. In the third case,  $A[0] < B[0]$ , but  $A[1] > B[1]$ , so they are concurrent.