HPSC - Lecture 19

Outline:

- Review MPI, reduce and bcast
- · MPI send and receive
- Master–Worker paradigm

References:

- \$/codes/mpi
- MPI Standard
- OpenMPI

MPI — Simple example

```
program test1
    use mpi
    implicit none
    integer :: ierr, numprocs, proc num,
    call mpi init(ierr)
    call mpi comm size (MPI COMM WORLD, numprocs, ierr)
    call mpi_comm_rank(MPI_COMM_WORLD, proc num, ierr)
    print *, 'Hello from Process ', proc num, &
    ' of ', numprocs, ' processes'
    call mpi finalize(ierr)
end program test1
Always need to: use mpi,
Start with mpi init,
End with mpi finalize.
```

Compiling and running MPI code (Fortran)

Try this test:

```
$ cd $UWHPSC/codes/mpi
$ mpif90 test1.f90
$ mpiexec -n 4 a.out
```

You should see output like:

```
Hello from Process number 1 of 4 processes
Hello from Process number 3 of 4 processes
Hello from Process number 0 of 4 processes
Hello from Process number 2 of 4 processes
```

Note: Number of processors is specified with mpiexec.

MPI Communicators

All communication takes place in groups of processes.

Communication takes place in some context.

A group and a context are combined in a communicator.

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MPI_COMM_WORLD is a communicator provided by default that includes all processors.

```
MPI_COMM_SIZE(comm, numprocs, ierr) returns the number of processors in communicator comm.
```

MPI_COMM_RANK(comm, proc_num, ierr) returns the rank of this processor in communicator comm.

mpi module

The mpi module includes:

```
Subroutines such as mpi_init, mpi_comm_size,
   mpi_comm_rank, ...
```

Global variables such as

```
MPI_COMM_WORLD: a communicator,
MPI_INTEGER: used to specify the type of data being sent
MPI_SUM: used to specify a type of reduction
```

Remember: Fortran is case insensitive: mpi init is the same as MPI INIT.

MPI functions

There are 125 MPI functions.

Can write many program with these 8:

- MPI_INIT(ierr) Initialize
- MPI_FINALIZE(ierr) Finalize
- MPI_COMM_SIZE(...) Number of processors
- MPI_COMM_RANK(...) Rank of this processor
- MPI_SEND(...) Send a message
- MPI_RCV(...) Receive a message
- MPI BCAST(...) Broadcast to other processors
- MPI REDUCE (...) Reduction operation

MPI Reduce

Examples: Compute $|x| = \max_i |x_i|$ for a distributed vector: (each process has some subset of x elements)

```
xnorm proc = 0.d0
! set istart and iend for each process
do i=istart, iend
    xnorm proc = max(xnorm proc, abs(x(i)))
    enddo
call MPI REDUCE (xnorm proc, xnorm, 1, &
         MPI DOUBLE PRECISION, MPI MAX, 0, &
         MPI COMM WORLD, ierr)
if (proc num == 0) print "norm of x = ", xnorm
```

Processors do not exit from MPI_REDUCE until all have called the subroutine.

Normalize the vector x: Replace x by $x/|x|_{\infty}$

```
! compute xnorm proc on each process as before.
call MPI REDUCE (xnorm proc, xnorm, 1, &
            MPI DOUBLE PRECISION, MPI MAX, 0, &
            MPI COMM WORLD, ierr)
! only Process O has the value of xnorm
call MPI BCAST (xnorm, 1, &
            MPI DOUBLE PRECISION, 0, &
            MPI COMM WORLD, ierr)
! now every process has the value of xnorm
do i=istart, iend
    x(i) = x(i) / xnorm
    enddo
```

MPI AllReduce

To make a reduction available to *all* processes:

One-step alternative: simpler and perhaps more efficient...

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We would have to:

- Send parts of x to other processes,
- Compute xnorm_proc on each process,
- Use MPI_ALLREDUCE to combine into xnorm and broadcast to all processes,
- Normalize part of x on each process,
- Send each part of normalized x back to Process 0.

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Communication cost will probably make this much slower than just normalizing all of *x* on Process 0!

Might be worthwhile if much more work is required for each element of *x*.

Suppose all of vector *x* is stored on memory of Process 0,

Want to solve an expensive differential equation with different initial conditions given by elements of x,

and then collect all results on Process 0.

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Master–Worker paradigm:

- Process 0 sends different chunks of x to Process 1, 2, . . .
- Each process grinds away to solve differential equations
- Each process sends results back to Process 0.

MPI Send and Receive

MPI_BCAST sends from one process to all processes.

Often want to send selectively from Process i to Process j.

Use MPI_SEND and MPI_RECV.

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Need a way to tag messages so they can be identified.

The parameter tag is an integer that can be matched to identify a message.

Tag can also be used to provide information about what is being sent, for example if a Master process sends rows of a matrix to other processes, the tag might be the row number.

MPI Send

Send value(s) from this Process to Process dest.

General form:

where:

- start: starting address (variable, array element)
- count: number of elements to send
- datatype: type of each element
- dest: destination process
- tag: identifier tag (integer between 0 and 32767)
- comm: communicator

MPI Receive

Receive value(s) from Process source with label tag.

General form:

where:

- source: source process
- tag: identifier tag (integer between 0 and 32767)
- comm: communicator
- status: integer array of length MPI_STATUS_SIZE.

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where:

- source: source process
- tag: identifier tag (integer between 0 and 32767)
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- status: integer array of length MPI STATUS SIZE.

source could be ${\tt MPI_ANY_SOURCE}$ to match any source.

tag could be MPI ANY TAG to match any tag.

Processor 3 will print j = 55

```
if (proc num == 4) then
    i = 55
    call MPI SEND(i, 1, MPI INTEGER, 3, 21, &
                  MPI COMM WORLD, ierr)
    endif
if (proc num == 3) then
    call MPI RECV(j, 1, MPI INTEGER, 4, 21, &
                  MPI COMM WORLD, status, ierr)
    print *, "j = ", j
    endif
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The tag is 21. (Arbitrary integer between 0 and 32767)

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Blocking Receive: Processor 3 won't return from MPI_RECV until message is received.

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Run-time error if num_procs <= 4 (Procs are 0,1,2,3)

Send/Receive example

Pass value of i from Processor 0 to 1 to 2 ... to num_procs-1

```
if (proc num == 0) then
   i = 55
   call MPI SEND(i, 1, MPI INTEGER, 1, 21, &
             MPI COMM WORLD, ierr)
   endif
else if (proc num < num procs - 1) then
  else if (proc num == num procs - 1) then
   call MPI RECV(i, 1, MPI INTEGER, proc num-1, 21, &
  print *, "i = ", i COMM_WORLD, status, ierr)
   endif
```

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MPI Receive — status argument

Elements of the status array give additional useful information about the message received.

In particular,

```
status(MPI_SOURCE) is the source of the message,
   May be needed if source = MPI_ANY_SOURCE.
status(MPI_TAG) is the tag of the message received,
   May be needed if tag = MPI_ANY_TAG.
```

Another Send/Receive example

Master (Processor 0) sends jth column to Worker Processor j, gets back 1-norm to store in anorm (j), j = 1, ..., ncols

Note: Master may receive back in any order!

MPI_ANY_SOURCE will match first to arrive.

The tag is used to tell which column's norm has arrived (j j).

Send and Receive example — worker code

```
Master (Processor 0) sends jth column to Worker Processor j, gets back 1-norm to store in anorm (j), j = 1, ..., ncols
```

```
! code for Workers (Processors 1, 2, ...):
if (proc num /= 0) then
  call MPI RECV(colvect, nrows, MPI DOUBLE PRECISION, &
                    O, MPI ANY TAG, &
                   MPI COMM WORLD, status, ierr)
                          ! this is the column number
  j = status(MPI TAG)
                          ! (should agree with proc num)
  colnorm = 0.d0
  do i=1, nrows
      colnorm = colnorm + abs(colvect(i))
      enddo
 call MPI_SEND(colnorm, 1, MPI DOUBLE PRECISION, & 0, j, MPI COMM WORLD, Terr)
  endif
```

Note: Sends back to Process 0 with tag *j*.

Send may be blocking

Both processors *might* get stuck in MPI_SEND!

May depend on size of data and send buffer.

Blocking send: MPI_SSEND. See documentation

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There are also non-blocking sends and receives:

```
MPI ISEND, MPI IRECV
```

Non-blocking receive

```
call MPI_IRECV(start, count, datatype, &
    source, tag, comm, request, ierror)
```

Additional argument: request.

Program continues after initiating receive,

Can later check if it has finished with

```
call MPI_TEST(request, flag, status, ierror)
flag is logical output variable.
```

Or can later wait for it to finish with

```
call MPI_WAIT(request, status, ierror)
```