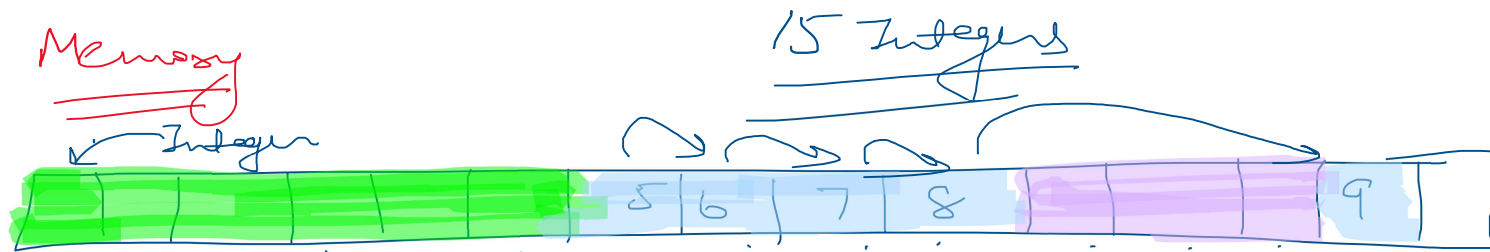
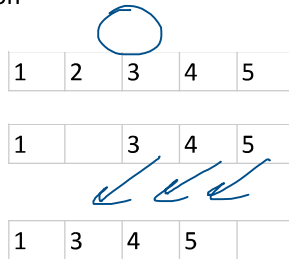


Disadvantages of Array

1. Similar data types
2. Fixed Size
3. Costly Deletion $O(n)$
4. Costly Insertion (When array is full, or at any particular index)
5. Contiguous Memory Allocation



4 Bytes
↓
32 bits

A → 6 integers

B → 4 integers

C → 3 integers

5, 6, 7, 8, 9

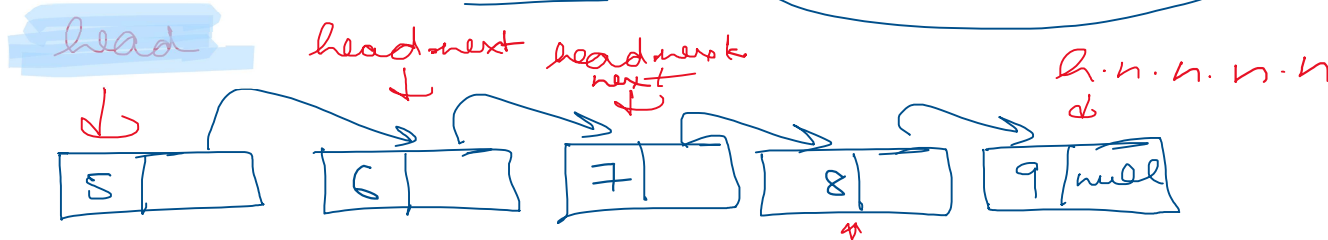
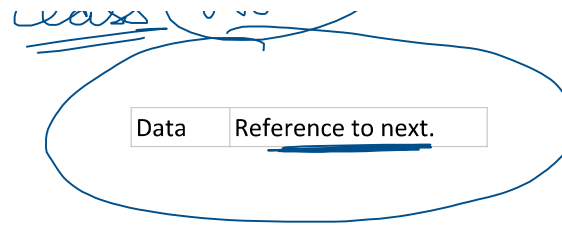
← D → 5 integers

... (node)

Solution offered:

Store two things:

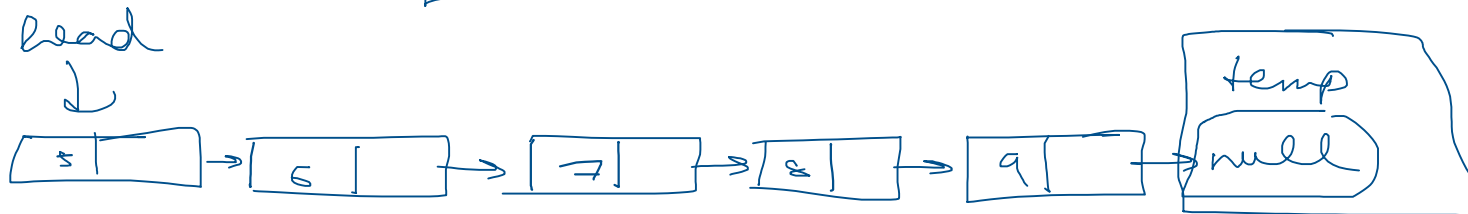
1. Own Data → data
2. Link to the next place → next



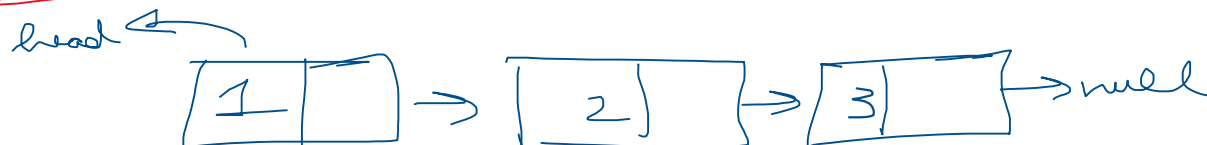
Data → int data
Node next;

1. n. n. n. n

```
class Node
{
    int data;
    Node next;
}
```



Insertion

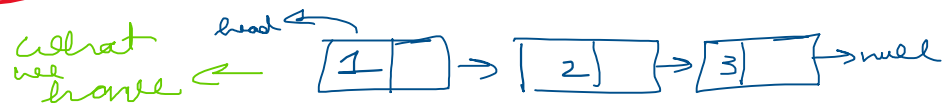


1. Insertion at Head

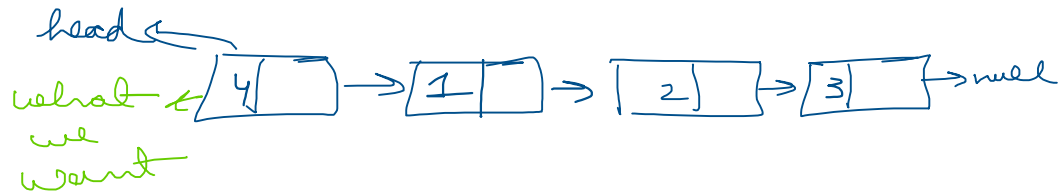
2 Insertion at End

3 Insertion at n^{th} position.

1 Insertion at Head

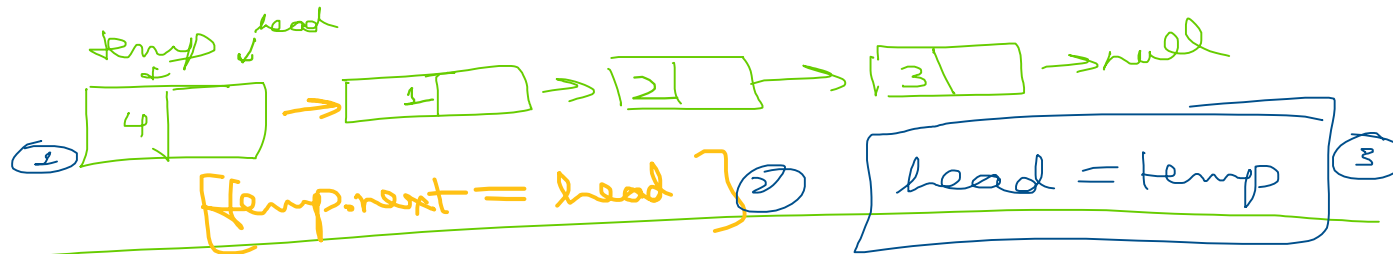


Task InsertAtHead(4)



Brahmastra:

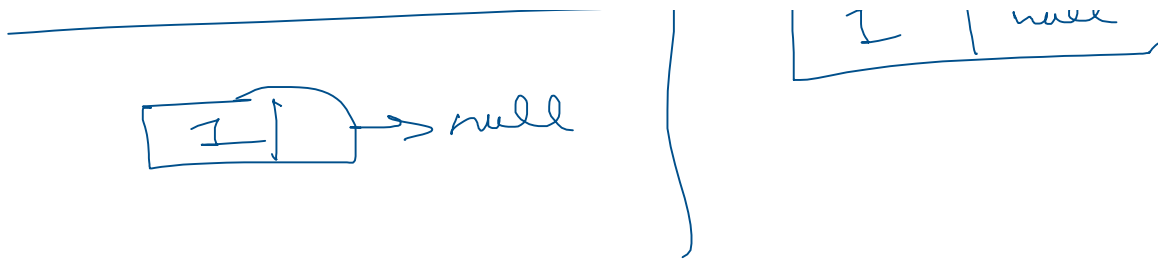
1. Draw what you have.
2. Draw what you want.
3. Draw the intermediate steps to reach what you want one by one.



Condition check → If LL is empty

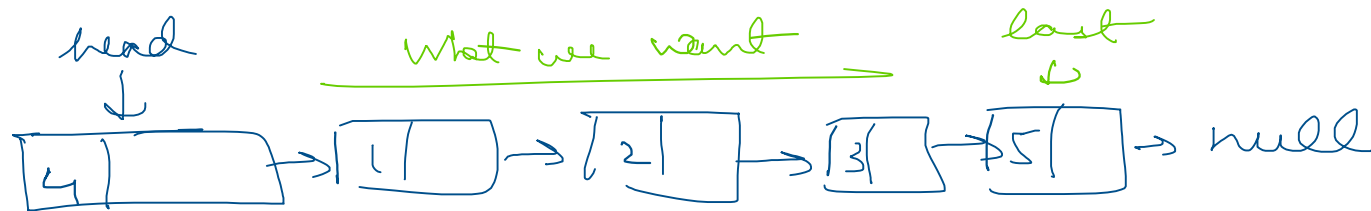
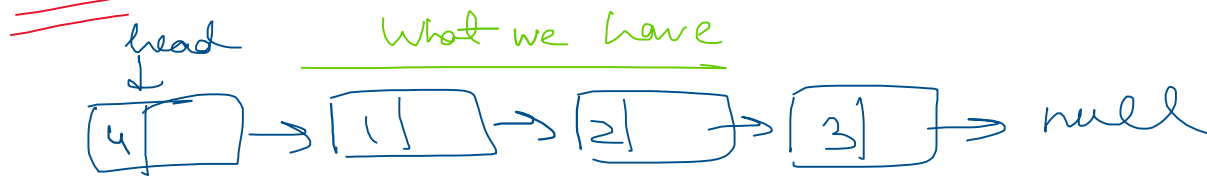
$\text{head} == \text{null}$





This case has been automatically handled by our code.

2 Insertion at End

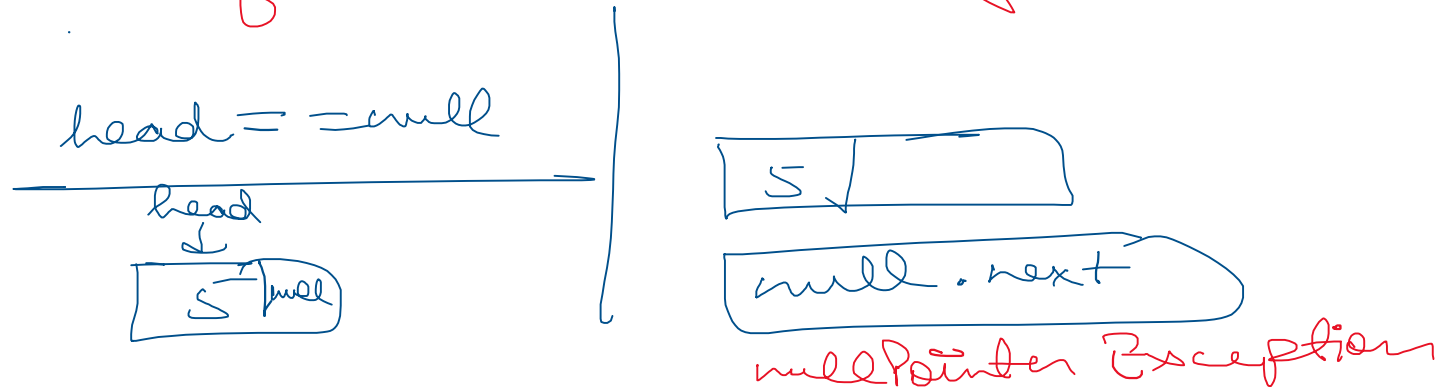


1 → Create a node ✓

2 → Find Last Node

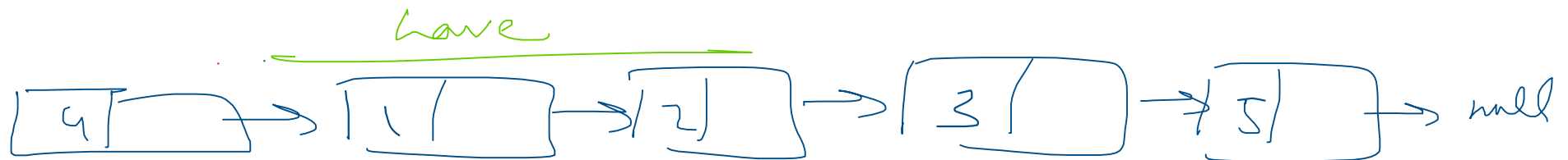
3 → last.next = temp ✓

what if Linked List is Empty.



In this case we will have to handle this explicitly

3 Insertion at nth position. H.W



insert at Position (2) element (6).

want



$[4] \rightarrow [6] \rightarrow [1] \rightarrow [2] \rightarrow [3] \rightarrow [5] \rightarrow \text{null}$