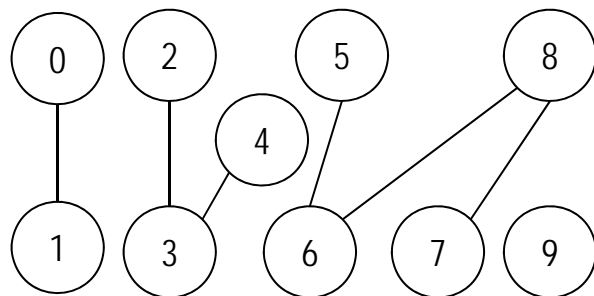


서로소집합과 union-find

서로소집합(disjoint set) 표현: 개념



5과 8이 연결되어 있는가?

$\{0,1\}, \{2,3,4\}, \{5,6,7,8\}, \{9\}$

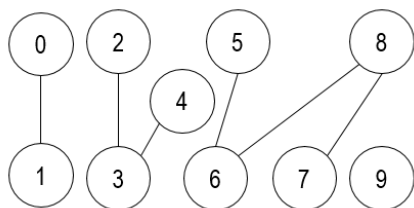
	0	1	2	3	4	5	6	7	8	9
parent	1	1	4	4	4	8	8	8	8	9

$\text{find}(5) == \text{find}(8) ?$

	0	1	2	3	4	5	6	7	8	9
parent	1	1	3	4	4	6	8	8	8	9

$\text{find}(5) == \text{find}(8) ?$

서로소집합(disjoint set) 표현: 개념



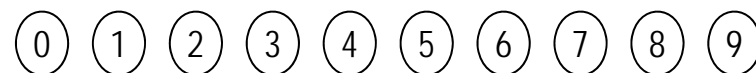
$union(a, b) \rightarrow a$ 가 속한 집합과 b 가 속한 집합을 합집합

$\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\}$

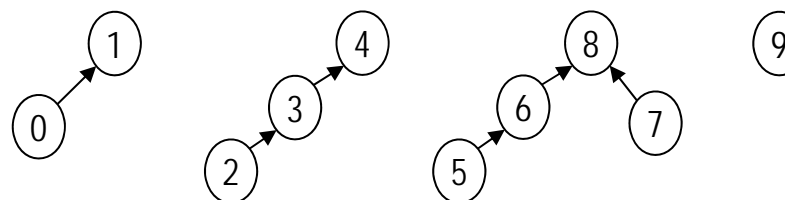
union(0,1)
union(2,3), union(3,4)
union(5,6), union(7,8), union(6,8)

$\{0,1\}, \{2,3,4\}, \{5,6,7,8\}, \{9\}$

0	1	2	3	4	5	6	7	8	9
0	1	2	3	4	5	6	7	8	9



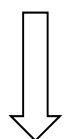
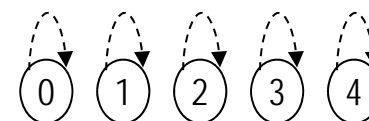
0	1	2	3	4	5	6	7	8	9
1	1	3	4	4	6	8	8	8	9



서로소집합(disjoint set) 개념: union

{0}, {1}, {2}, {3}, {4}

	0	1	2	3	4
parent	0	1	2	3	4



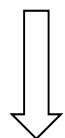
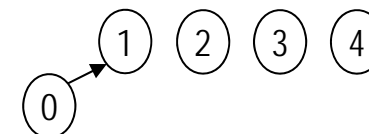
union(0,1)

0이 속한 집합과 1이 속한 집합을 합집합

노드 0의 parent는 노드 1이다.

{0,1}, {2}, {3}, {4}

	0	1	2	3	4
parent	1	1	2	3	4

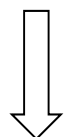
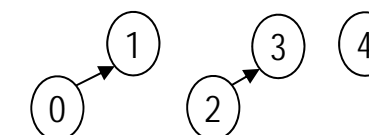


union(2,3)

2가 속한 집합과 3이 속한 집합을 합집합

{0,1}, {2,3}, {4}

	0	1	2	3	4
parent	1	1	3	3	4

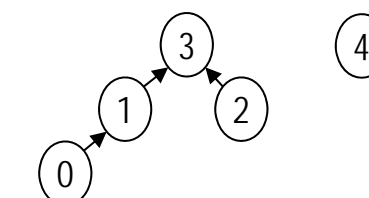


union(1,3)

1이 속한 집합과 3이 속한 집합을 합집합

{0,1,2,3}, {4}

	0	1	2	3	4
parent	1	3	3	3	4

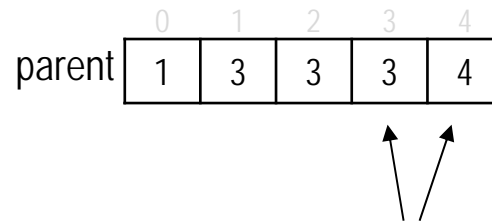


서로소집합(disjoint set) 연산: find

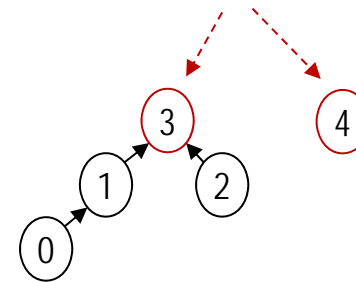
find(i)

- 원소 i가 속한 집합의 대표 원소 id를 반환

{0,1,2,3}, {4}

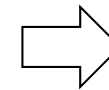


집합의 대표 원소는 집합 표현
트리의 루트 노드



$i == \text{parent}[i]$ 을 만족하는 i가 집합의 대표 원소 id

```
public int find(int i) {
    while( i != parent[i] ) i=parent[i];
    return i;
}
```

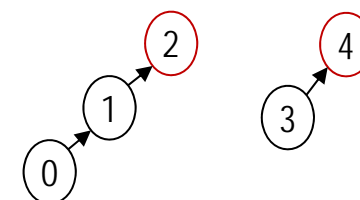
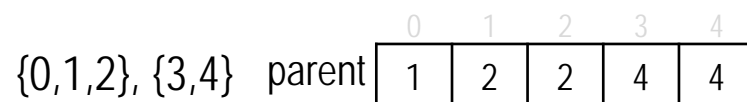


find(0) → 3
 find(1) → 3
 find(2) → 3
 find(3) → 3
 find(4) → 4

서로소집합(disjoint set) 연산: union

union(i, j)

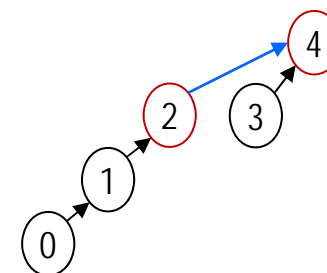
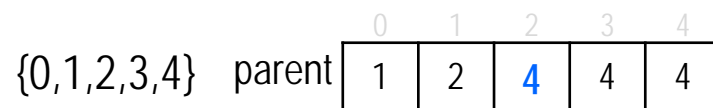
- 원소 i가 속한 집합과 원소 j가 속한 집합을 합집합



```
public int find(int i) {
    while(i != parent[i]) i = parent[i];
    return i;
}
```

```
public void union(int i, int j) {
    i = find(i);
    j = find(j);
    if(i == j) return;
    parent[i] = j;
}
```

union(0,3)



서로소집합(disjoint set) 연산: find by path compression



경로 압축 기반 find 연산

```
public int find(int i) {
    if(i != parent[i]) parent[i] = find(parent[i]);
    return parent[i];
}
```

find(0)
0 != parent[0]
parent[0] = find(parent[0])
return parent[0]

find(1)
1 != parent[1]
parent[1] = find(parent[1])
return parent[1]

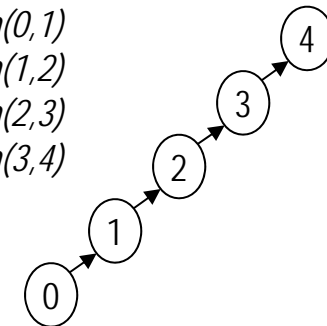
find(2)
2 != parent[2]
parent[2] = find(parent[2])
return parent[2]

find(3)
3 != parent[3]
parent[3] = find(parent[3])
return parent[3]

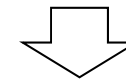
find(4)
4 == parent[4]
return parent[4]

	0	1	2	3	4
parent	1	2	3	4	4

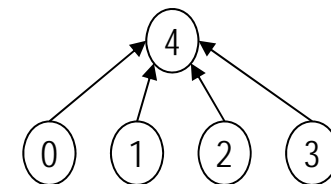
union(0, 1)
union(1, 2)
union(2, 3)
union(3, 4)



find(0)



	0	1	2	3	4
parent	4	4	4	4	4

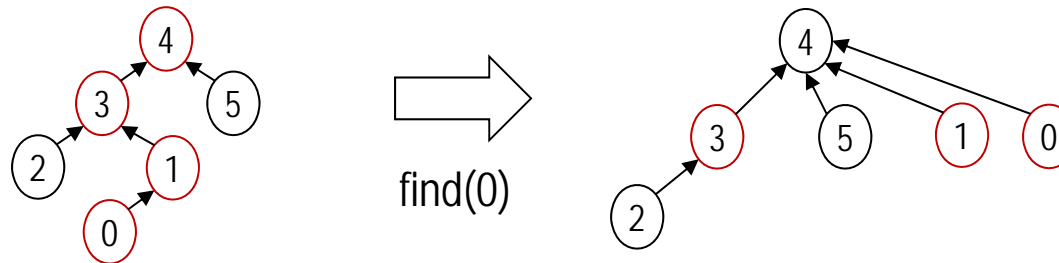


서로소집합(disjoint set) 연산: find by path compression



경로 압축 기반 find 연산

```
public int find(int i) {  
    if(i != parent[i]) parent[i] = find(parent[i]);  
    return parent[i];  
}
```



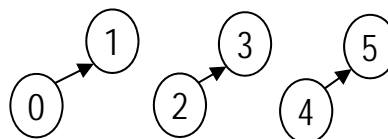
서로소집합(disjoint set) 연산: union by rank

rank → 특정 집합에 대응하는 트리의 높이와 유사

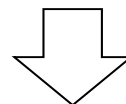
rank 기반 union

```
public void union(int i, int j) {
    i=find(i);
    j=find(j);
    if(i==j) return;
    if(rank[i]<rank[j]) parent[i]=j;
    else if(rank[i]>rank[j]) parent[j]=i;
    else{
        parent[i]=j;
        rank[j]++;
    }
}
```

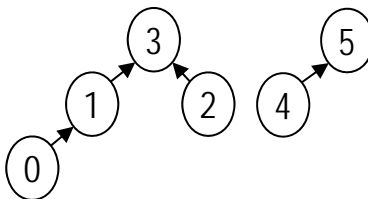
	0	1	2	3	4	5
rank	0	1	0	1	0	1



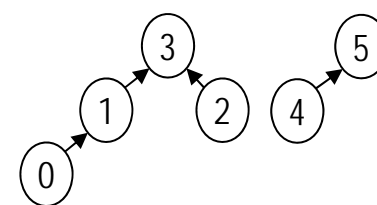
union(1,3)



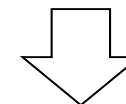
	0	1	2	3	4	5
rank	0	1	0	2	0	1



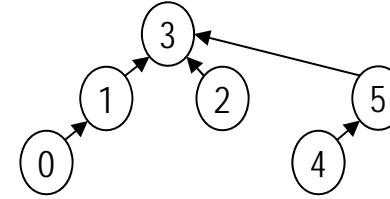
	0	1	2	3	4	5
rank	0	1	0	2	0	1



union(3,5)



	0	1	2	3	4	5
rank	0	1	0	2	0	1



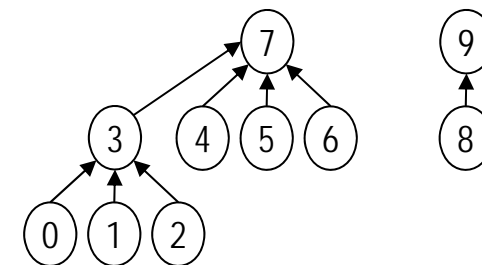
Union-find 자료구조

```
public class UF {
    int parent[], rank[], count;
    public UF(int N) {
        parent=new int[N];
        rank=new int[N];
        for (int i = 0; i < parent.length; i++){
            parent[i]=i;
        }
    }
    public void union(int i, int j) { // union by rank
        i=find(i);
        j=find(j);
        if(i==j) return;
        if(rank[i]<rank[j]) parent[i]=j;
        else if(rank[i]>rank[j]) parent[j]=i;
        else{
            parent[i]=j;
            rank[j]++;
        }
    }
    public int find(int i) { // find by path compression
        if(i!=parent[i]) parent[i]=find(parent[i]);
        return parent[i];
    }
    @Override
    public String toString() { return Arrays.toString(parent); }
}
```

실습: 그래프 내
연결요소의 개수가
count에 저장되도록
클래스 UF를
수정하시오

실습: find 함수를
비재귀적으로
구현하시오

```
public class Test {
    public static void main(String[] args) {
        int N=10;
        UF uf=new UF(N);
        System.out.println(uf);
        uf.union(0,1);
        uf.union(2,3);
        uf.union(4,5);
        uf.union(6,7);
        uf.union(8,9);
        uf.union(0,2);
        uf.union(4,6);
        uf.union(0,4);
        System.out.println(uf);
        System.out.println(uf.find(1)==uf.find(6));
        //System.out.println("연결요소 개수 = "+uf.count);
    }
}
```



Reference: <https://algs4.cs.princeton.edu/15uf/UF.java.html>, GPLv3

Reference: https://en.wikipedia.org/wiki/Disjoint-set_data_structure, CC-BY-SA

서로소집합과 union-find

Union-find 시간복잡도

- Find

- ◆ $O(\alpha(n))$

- Union

- ◆ $O(\alpha(n))$

$\alpha(n)$

- inverse Ackermann function
- 대부분의 경우 $\alpha(n) < 5$
- Ackermann function $A(4,4) \approx 2^{2^{65536}}$

Union-find 응용

- 동적으로 변하는 비방향 그래프에서 연결요소들 표현
- 그래프 내 임의의 두 노드 간 연결 경로 존재 여부 판단
- Kruskal의 최소신장트리 알고리즘에서 활용

Reference: https://en.wikipedia.org/wiki/Disjoint-set_data_structure, CC-BY-SA

Reference: https://en.wikipedia.org/wiki/Ackermann_function, CC-BY-SA

References

- ✚ C로 쓴 자료구조론 (Fundamentals of Data Structures in C, Horowitz et al.). 이석호 역. 사이텍미디어. 1993.
- ✚ 쉽게 배우는 알고리즘: 관계 중심의 사고법. 문병로. 한빛아카데미. 2013.
- ✚ C언어로 쉽게 풀어 쓴 자료구조. 천인국 외 2인. 생능출판사. 2017.
- ✚ 프로그래밍 콘테스트 챌린징, Akiba 등 공저, 로드북, 2011.
- ✚ <https://introcs.cs.princeton.edu/>
- ✚ Introduction to Algorithms, Cormen et al., 3rd Edition (The MIT Press)