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Program for nth Catalan Number

Catalan numbers are a sequence of natural numbers that occurs in many interesting counting problems like following.

- 1) Count the number of expressions containing n pairs of parentheses which are correctly matched. For $n = 3$, possible expressions are $((()))$, $()(())$, 000 , $(0)0$, (00) .
- 2) Count the number of possible Binary Search Trees with n keys (See [this](#))
- 3) Count the number of full binary trees (A rooted binary tree is full if every vertex has either two children or no children) with $n+1$ leaves.

See [this](#) for more applications.

The first few Catalan numbers for $n = 0, 1, 2, 3, \dots$ are **1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, ...**

Recursive Solution

Catalan numbers satisfy the following recursive formula.

$$C_0 = 1 \quad \text{and} \quad C_{n+1} = \sum_{i=0}^n C_i C_{n-i} \quad \text{for } n \geq 0;$$

Following is C++ implementation of above recursive formula.

```
#include<iostream>
using namespace std;

// A recursive function to find nth catalan number
unsigned long int catalan(unsigned int n)
{
    // Base case
    if (n <= 1) return 1;

    // catalan(n) is sum of catalan(i)*catalan(n-i-1)
    unsigned long int res = 0;
    for (int i=0; i<n; i++)
        res += catalan(i)*catalan(n-i-1);

    return res;
}

// Driver program to test above function
int main()
{
    for (int i=0; i<10; i++)
        cout << catalan(i) << " ";
    return 0;
}
```

Output :

1 1 2 5 14 42 132 429 1430 4862

Time complexity of above implementation is equivalent to nth catalan number.

$$T(n) = \sum_{i=0}^{n-1} \{T(i)*T(n-i)\} \quad \text{for } n \geq 0;$$

The value of nth catalan number is exponential that makes the time complexity exponential.

Dynamic Programming Solution

We can observe that the above recursive implementation does a lot of repeated work (we can see the same by drawing recursion tree). Since there are overlapping subproblems, we can use dynamic programming for this. Following is a Dynamic programming based implementation in C++.

```
#include<iostream>
using namespace std;
```

```
// A dynamic programming based function to find nth
// Catalan number
unsigned long int catalanDP(unsigned int n)
{
    // Table to store results of subproblems
    unsigned long int catalan[n+1];

    // Initialize first two values in table
    catalan[0] = catalan[1] = 1;

    // Fill entries in catalan[] using recursive formula
    for (int i=2; i<=n; i++)
    {
        catalan[i] = 0;
        for (int j=0; j<i; j++)
            catalan[i] += catalan[j] * catalan[i-j-1];
    }

    // Return last entry
    return catalan[n];
}

// Driver program to test above function
int main()
{
    for (int i = 0; i < 10; i++)
        cout << catalanDP(i) << " ";
    return 0;
}
```

Output:

```
1 1 2 5 14 42 132 429 1430 4862
```

Time Complexity: Time complexity of above implementation is $O(n^2)$

Using Binomial Coefficient

We can also use the below formula to find nth catalan number in $O(n)$ time.

[Tex] $C_n = \frac{1}{n+1} \binom{2n}{n}$ [/Tex]

We have discussed a [O\(n\) approach to find binomial coefficient nCr](#).

```
#include<iostream>
using namespace std;

// Returns value of Binomial Coefficient C(n, k)
unsigned long int binomialCoeff(unsigned int n, unsigned int k)
{
    unsigned long int res = 1;
```

```

// Since C(n, k) = C(n, n-k)
if (k > n - k)
    k = n - k;

// Calculate value of [n*(n-1)*---*(n-k+1)] / [k*(k-1)*---*1]
for (int i = 0; i < k; ++i)
{
    res *= (n - i);
    res /= (i + 1);
}

return res;
}

// A Binomial coefficient based function to find nth catalan
// number in O(n) time
unsigned long int catalan(unsigned int n)
{
    // Calculate value of 2nCn
    unsigned long int c = binomialCoeff(2*n, n);

    // return 2nCn/(n+1)
    return c/(n+1);
}

// Driver program to test above functions
int main()
{
    for (int i = 0; i < 10; i++)
        cout << catalan(i) << " ";
    return 0;
}

```

Output:

1 1 2 5 14 42 132 429 1430 4862

Time Complexity: Time complexity of above implementation is $O(n)$.

We can also use below formula to find nth catalan number in $O(n)$ time.

$$C_n = \frac{(2n)!}{(n+1)!n!} = \prod_{k=2}^n \frac{n+k}{k} \quad \text{for } n \geq 0$$

References:

http://en.wikipedia.org/wiki/Catalan_number

Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above

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Ted Wilson • a month ago

Here is the dynamic solution in Haskell:

```
catalan n = head $ catem n [1,1]
where catem n lis = if (length lis > n) then lis
else catem n ((sum $ zipWith (*) lis (reverse lis)):lis)
```

^ | v • Reply • Share ›

aa1992 • 7 months ago

excellent solution with o(n).

^ | v • Reply • Share ›

kaushik Lele • 9 months ago

@GeeksforGeeks

A full binary tree with 3 nodes can be formed in 6 ways.
But catalan number for n=3 is 5. Which does not match.
Is my understanding incorrect ?

For a binary tree with 3 nodes viz. (a,b,c) I could see below 6 combinations.

a -> b

-> c

a -> c

-> b

b -> a

-> c

b -> c

-> a

c -> a

-> b

c -> b

-> a

^ | v • Reply • Share ›

kaushik Lele ➔ kaushik Lele • 9 months ago

Article says that

"A rooted binary tree is full if every vertex has either two children or no children"

But the example given on <http://www.findstat.org/Binary...> page (which is related to Catalan number wikipedia)

gives examples of Binary tree with 3 nodes. But those are not full binary trees.

So

1) is requirement of full trees given in this article incorrect ?

2) We are just thinking on different arrangements of nodes where every node is treated same ?

^ | v • Reply • Share ›

Anurag Singh ➔ kaushik Lele • 8 months ago

It needs to be reworded.

On <http://en.wikipedia.org/wiki/C...>

"Successive applications of a binary operator can be represented in terms of a full binary tree.[3] (A rooted binary tree is full if every vertex has either two children or no children.) It follows that C_n is the number of full binary trees with $n + 1$ leaves:"

So here $n = 3$ is three operators. So 3 operators and 4 (i.e. $n + 1$) operands can be represented in 5 ways.

^ | v • Reply • Share ›

kaushik Lele ➔ Anurag Singh • 8 months ago

It still does not answer my question ->

What are those 5 ways of creating full binary tree?

The diagram there shows 5 Binary tree. But I see that there is only 1 tree which satisfies the condition of two child or no child

WHICH SATISFIES THE CONDITION OF TWO-CHILD-OR-NO-CHILD.

^ | v • Reply • Share ›

Anurag Singh ➔ kaushik Lele • 8 months ago

Statement is: "It follows that C_n is the number of full binary trees with $n + 1$ leaves:"

i.e. total $2n + 1$ nodes are involved while making the tree (n internal nodes as operator and $n+1$ leaf nodes as operands).

For $n = 3$ operator and $n+1 = 4$ operands

In the diagram, there are 3 operators (internal nodes) and 4 operands (the leaves). Total 7 nodes are involved.

1 ^ | v • Reply • Share ›

kaushik Lele ➔ Anurag Singh • 8 months ago

Ohh .. its about "binary operator" !! I completely ignored that word and I was thinking in terms of normal Binary tree concept. Thanks for explaining patiently till I understood. **@GeeksforGeeks** please modify the sentence in this article to include word "operator" and add above explanation by **@Anurag Singh** This will help other learners.

^ | v • Reply • Share ›

kaushik Lele ➔ kaushik Lele • 8 months ago

Description in this article also mentions " $n+1$ leaves" but I by mistake took it as " $n+1$ total nodes" and caused all this confusion.

However it is better to explain it in terms of nodes than leaves; as depending on number of nodes; we can draw different arrangements. Leaves will automatically fall in place.

So we can say -> full binary tree with n -nodes (& hence $n+1$ leaves) can be formed with C_n ways :)

^ | v • Reply • Share ›

Jun • 10 months ago

Method 2

<http://ideone.com/YbNst4>

^ | v • Reply • Share ›

Jun • 10 months ago

Method 1

<http://ideone.com/7q8wtV>

^ | v • Reply • Share ›

valar morphulis · 10 months ago

dp soln can be optimized:-

```
for (int j=0; j
```

3 ^ | v · Reply · Share ›

Manish M Berwani · a year ago

A better approach to this problem would be

$$C[n+1] = (2 * (2 * n + 1) * C[n]) / (n + 2)$$

3 ^ | v · Reply · Share ›

<HoldOnLife!#> → **Manish M Berwani** · 10 months ago

It would also require O(n) only I guess?

^ | v · Reply · Share ›

arjomanD · a year ago

Good !

Does this site have contests ?

^ | v · Reply · Share ›

GOPI GOPINATH → **arjomanD** · a year ago

No. Just Concepts

5 ^ | v · Reply · Share ›

Karshit Jaiswal → **GOPI GOPINATH** · a year ago

which site are you talking about?

^ | v · Reply · Share ›

GOPI GOPINATH → **Karshit Jaiswal** · a year ago

[Geeksforgeeks.org](http://www.geeksforgeeks.org)

^ | v · Reply · Share ›



kabeer · a year ago

nice!

^ | v · Reply · Share ›

Vivek VV · a year ago

Please correct the spelling mistake in 2nd usage example :P

2 ^ | v · Reply · Share ›

GeeksforGeeks Mod → **Vivek VV** · a year ago

Thanks for pointing this out. We have corrected the typo.

^ | v · Reply · Share ›



Guest → GeeksforGeeks • 6 months ago

Here $T(n - i)$ should be changed to $T(n - i - 1)$

^ | v • Reply • Share ›



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