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A Time Complexity Question

What is the time complexity of following function fun()? Assume that log(x) returns log value in base 2.

```
void fun()
{
    int i, j;
    for (i=1; i<=n; i++)
        for (j=1; j<=log(i); j++)
            printf("GeeksforGeeks");
}
```

Time Complexity of the above function can be written as $\theta(\log 1) + \theta(\log 2) + \theta(\log 3) + \dots + \theta(\log n)$ which is $\theta(\log n!)$

Order of growth of $\log n!$ and $n \log n$ is same for large values of n , i.e., $\theta(\log n!) = \theta(n \log n)$. So time complexity of `fun()` is $\theta(n \log n)$.

The expression $\theta(\log n!) = \theta(n \log n)$ can be easily derived from following [Stirling's approximation \(or Stirling's formula\)](#).

$$\log n! = n \log n - n + O(\log(n))$$

Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

Sources:

http://en.wikipedia.org/wiki/Stirling%27s_approximation

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Rashmi • 9 months ago

where can i find the complete tutorial on time complexities right from the start?

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Saurabh Gupta → Rashmi • 8 months ago

@Rashmi you can follow this stanford tutorial on algorithms, the time complexity concept has been explained in less than first 10 videos.

<https://class.coursera.org/alg...>

Also to juice out the maximum of this tutorial, you can use the book CLRS, also known as the bible of algorithms

as the size of algorithm.

<http://www.amazon.in/Introduct...>

Just read the book for the part of complexity while you are watching the videos, that will help you the best way.

-Keep coding.

^ | v • Reply • Share ›



gaurav • 10 months ago

what is the time complexity of rsa algorithm for encryption and decryption

^ | v • Reply • Share ›



Devesh_D • a year ago

One more reason is there for use of quick sort more than merge sort because merge sort complexity is fixed in all the cases (best , avg, worst) that means it is not depend on the input type(like sorted or unsorted etc).And in case of quick sort if we are using randomized quick sort then we are sure that there is no worst case situation will occur. and the time taken by the quick sort is always less than or equal to merge sort.

^ | v • Reply • Share ›



Prince • 2 years ago

Yes harsh, correct..n! 's upper bound is n^n so will beome $n \log n$.

^ | v • Reply • Share ›



Har?sh • 2 years ago

$n! = 1 \times 2 \times 3 \dots n$

$n! \leq n \times n \times \dots n$

$n! \leq n^n$

$\log n! \leq \log n^n$

$\log n! \leq n \log n$

$T(n) = O(n \log n)$ (Asymptotic tight upperbound).

8 ^ | v • Reply • Share ›



Guest → Har?sh • 2 years ago

But here they are using Theta notation, but you have used Big-oh notation?

^ | v • Reply • Share ›



Mihir • 2 years ago

I might be wrong, but we do not know anything about the complexity of the function that calculates the log value here. How can we comment about the overall complexity without knowing this? (Unless it is a library function of known complexity. Then it is fine.)

2 ^ | v • Reply • Share ›



Nidhi • 2 years ago



Shouldn't we consider the first loop "i" while calculating complexity ?

```
/* Paste your code here (You may delete these lines if not writing code) */
```

^ | v • Reply • Share ›



neha2210 → Nidhi • 2 years ago

We are indeed considering the first loop also that is why we are adding up $\theta(\log 1) + \theta(\log 2) \dots$

As you see if both loops are repeated for n times then the complexity would be $n+n+n \dots n$ times i.e n^2 . We do it here in a similar manner.

1 ^ | v • Reply • Share ›



vivek • 4 years ago

Does it mean that the given code is as efficient as following code?

```
void fun()
{
    int i, j;
    for (i=1; i<=n; i++)
        for (j=1; j<=log(n); j++)
            printf("GeeksforGeeks");
}
```

1 ^ | v • Reply • Share ›



Sandeep → vivek • 4 years ago

@vivek: Both codes are of same time complexity as both are asymptotically same, but code given by you will take more time in general. As an example, both $2n$ and $10000n$ are of same time complexity $O(n)$, but a code that takes $10000n$ time will definitely take more time. What we conclude from this is Asymptotic analysis is not perfect, but it is the only available way to compare algorithms irrespective of different computing powers, architectures.. etc. As another example, we can consider Merge Sort and Quick Sort. Although, worst case complexity of Quick Sort is more and average case is same as Merge Sort. Quick Sort is preferred because of the hidden greater constants in Merge Sort.

1 ^ | v • Reply • Share ›



Doom → Sandeep • 4 years ago

@Sandeep: Could you please tell me more about the hidden greater constants in Merge Sort - the reason why we prefer quick sort over mergesort?

^ | v • Reply • Share ›



Bhagat Vishal → Doom • a year ago

what i think , in case of deciding priority between the merge sort and

quick sort is the extra ammount of space taken by merge sort .but this could be avoided by implementing it through linked list without any extra space.

^ | v • Reply • Share ›



Amit → Doom • 4 years ago

@Doom: These constants depends upon cost of executing instructions on the specific machine....

Like some machine incurs more cost in assignment than other machine....

^ | v • Reply • Share ›



viresh → Amit • 4 years ago

mergesort is not inPlace.. so each time u ask for a extra space in the merge function there might be a page fault and also mergesort involves too many assignments that makes mergesort costlier.. in contrast quicksort is inplace and quicksort deals with changing the index rather than swapping the elements in each inner loop(avg case swaps 2s in d inner loop)..

1 ^ | v • Reply • Share ›



Doom → Amit • 4 years ago

@Amit: yes, you are right. But discussing about the architecture is a different thing. I am more concerned about the specific instructions which make make mergesort an expensive operation than quick sort.

^ | v • Reply • Share ›



Sandeep → Doom • 4 years ago

@Doom: These constants are not fixed and vary from machine to machine. All we can say is that the constants for MergeSort are greater than QuickSort in general.

^ | v • Reply • Share ›



Sandeep → Sandeep • 4 years ago

@Doom: These constants are greater if number of CPU cycles involved in loop or recursion are more. The CPU cycles might be used for things like data comparisons, data movement etc.

For example, consider the following two loops. Time complexity of both of them is $O(n)$, but the constants involved in Loop 2 are more than Loop 1 and these constants vary from machine to machine.

```
// Loop 1
for(i = 0; i < n; i++)
```

```
for (i = 0; i < n; i++)  
{  
    // Loop 2  
    for(i = 0; i < n; i++)  
    {  
        printf("GeeksforGeeks");  
    }  
}
```

Hope I made things clear this time.

^ | v • Reply • Share ›



Doom ↗ Sandeep • 4 years ago

@Sandeep: But what are these constants related to? I mean what exactly are u trying to point out? is it like the no. of comparisons involved? or is it like the data movement in the array? plz give some examples about the constants.

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