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Divide and Conquer | Set 4 (Karatsuba algorithm for fast multiplication)

Given two binary strings that represent value of two integers, find the product of two strings. For example, if the first bit string is “1100” and second bit string is “1010”, output should be 120.

For simplicity, let the length of two strings be same and be n .

A **Naive Approach** is to follow the process we study in school. One by one take all bits of second number and multiply it with all bits of first number. Finally add all multiplications. This algorithm takes $O(n^2)$ time.

```

X = 101001 = 41
Y = 101010 = 42
-----
      1010010
      101001
+   101001
-----
11010111010 = 1722

```

Using **Divide and Conquer**, we can multiply two integers in less time complexity. We divide the given numbers in two halves. Let the given numbers be X and Y.

For simplicity let us assume that n is even

$X = X_l * 2^{n/2} + X_r$ [X_l and X_r contain leftmost and rightmost $n/2$ bits of X]
 $Y = Y_l * 2^{n/2} + Y_r$ [Y_l and Y_r contain leftmost and rightmost $n/2$ bits of Y]

The product XY can be written as following.

$$\begin{aligned}
 XY &= (X_l * 2^{n/2} + X_r)(Y_l * 2^{n/2} + Y_r) \\
 &= 2^n X_l Y_l + 2^{n/2}(X_l Y_r + X_r Y_l) + X_r Y_r
 \end{aligned}$$

If we take a look at the above formula, there are four multiplications of size $n/2$, so we basically divided the problem of size n into for sub-problems of size $n/2$. But that doesn't help because solution of recurrence $T(n) = 4T(n/2) + O(n)$ is $O(n^2)$. The tricky part of this algorithm is to change the middle two terms to some other form so that only one extra multiplication would be sufficient. The following is tricky expression for middle two terms.

$$X_l Y_r + X_r Y_l = (X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r$$

So the final value of XY becomes

$$XY = 2^n X_l Y_l + 2^{n/2} * [(X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r] + X_r Y_r$$

With above trick, the recurrence becomes $T(n) = 3T(n/2) + O(n)$ and solution of this recurrence is $O(n^{1.59})$.

What if the lengths of input strings are different and are not even? To handle the different length case, we append 0's in the beginning. To handle odd length, we put $\text{floor}(n/2)$ bits in left half and $\text{ceil}(n/2)$ bits in right half. So the expression for XY changes to following.

$$XY = 2^{2\text{ceil}(n/2)} X_l Y_l + 2^{\text{ceil}(n/2)} * [(X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r] + X_r Y_r$$

The above algorithm is called Karatsuba algorithm and it can be used for any base.

Following is C++ implementation of above algorithm.

```

// C++ implementation of Karatsuba algorithm for bit string multiplication.
#include<iostream>
#include<stdio.h>

using namespace std;

```

```

// FOLLOWING TWO FUNCTIONS ARE COPIED FROM http://goo.gl/q00hZ
// Helper method: given two unequal sized bit strings, converts them to
// same length by adding leading 0s in the smaller string. Returns the
// the new length
int makeEqualLength(string &str1, string &str2)
{
    int len1 = str1.size();
    int len2 = str2.size();
    if (len1 < len2)
    {
        for (int i = 0 ; i < len2 - len1 ; i++)
            str1 = '0' + str1;
        return len2;
    }
    else if (len1 > len2)
    {
        for (int i = 0 ; i < len1 - len2 ; i++)
            str2 = '0' + str2;
    }
    return len1; // If len1 >= len2
}

// The main function that adds two bit sequences and returns the addition
string addBitStrings( string first, string second )
{
    string result; // To store the sum bits

    // make the lengths same before adding
    int length = makeEqualLength(first, second);
    int carry = 0; // Initialize carry

    // Add all bits one by one
    for (int i = length-1 ; i >= 0 ; i--)
    {
        int firstBit = first.at(i) - '0';
        int secondBit = second.at(i) - '0';

        // boolean expression for sum of 3 bits
        int sum = (firstBit ^ secondBit ^ carry)+'0';

        result = (char)sum + result;

        // boolean expression for 3-bit addition
        carry = (firstBit&secondBit) | (secondBit&carry) | (firstBit&carry);
    }

    // if overflow, then add a leading 1
    if (carry) result = '1' + result;

    return result;
}

```

```
// A utility function to multiply single bits of strings a and b
int multiplySingleBit(string a, string b)
{ return (a[0] - '0')*(b[0] - '0'); }

// The main function that multiplies two bit strings X and Y and returns
// result as long integer
long int multiply(string X, string Y)
{
    // Find the maximum of lengths of x and Y and make length
    // of smaller string same as that of larger string
    int n = makeEqualLength(X, Y);

    // Base cases
    if (n == 0) return 0;
    if (n == 1) return multiplySingleBit(X, Y);

    int fh = n/2; // First half of string, floor(n/2)
    int sh = (n-fh); // Second half of string, ceil(n/2)

    // Find the first half and second half of first string.
    // Refer http://goo.gl/1Lmgn for substr method
    string Xl = X.substr(0, fh);
    string Xr = X.substr(fh, sh);

    // Find the first half and second half of second string
    string Yl = Y.substr(0, fh);
    string Yr = Y.substr(fh, sh);

    // Recursively calculate the three products of inputs of size n/2
    long int P1 = multiply(Xl, Yl);
    long int P2 = multiply(Xr, Yr);
    long int P3 = multiply(addBitStrings(Xl, Xr), addBitStrings(Yl, Yr));

    // Combine the three products to get the final result.
    return P1*(1<<(2*sh)) + (P3 - P1 - P2)*(1<<sh) + P2;
}
```

```
// Driver program to test above functions
int main()
{
    printf ("%ld\n", multiply("1100", "1010"));
    printf ("%ld\n", multiply("110", "1010"));
    printf ("%ld\n", multiply("11", "1010"));
    printf ("%ld\n", multiply("1", "1010"));
    printf ("%ld\n", multiply("0", "1010"));
    printf ("%ld\n", multiply("111", "111"));
    printf ("%ld\n", multiply("11", "11"));
}
```

Output:

120
60

30
10
0
49
9

Time Complexity: Time complexity of the above solution is $O(n^{1.59})$.

Time complexity of multiplication can be further improved using another Divide and Conquer algorithm, fast Fourier transform. We will soon be discussing fast Fourier transform as a separate post.

Exercise

The above program returns a long int value and will not work for big strings. Extend the above program to return a string instead of a long int value.

References:

[Wikipedia page for Karatsuba algorithm](#)

[Algorithms 1st Edition by Sanjoy Dasgupta, Christos Papadimitriou and Umesh Vazirani](#)

<http://courses.csail.mit.edu/6.006/spring11/exams/notes3-karatsuba>

<http://www.cc.gatech.edu/~ninamf/Algos11/lectures/lect0131.pdf>

Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

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prashant saxena · 3 months ago

This is a solution for large Numbers. Try out any big number.

```
#include <string>
#include <math.h>
#include <cassert>
#include <sstream>

int adjustLength(std::string& s1, std::string& s2)

{

int len1 = s1.size();

int len2 = s2.size();

std::string s_adjusted;

std::string s_orig;

if (len1 != len2)
```

[see more](#)[^](#) | [v](#) • [Reply](#) • [Share](#) ›**prashant saxena** → prashant saxena • 3 months ago

Ok, the format sucks.

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**Diptesh Patel** • 5 months ago

When n is odd can't we just add one more bit ahead of strings, that does not changes the values and reduces the difficulty.

Your approach of floor and n-floor is also simpler. but just curious to know will my solution work?

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**Arjun K** • 8 months ago

Check this for analysis of algorithm which one is best. <http://www.msccomputerscience....>

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**n@733d** • 9 months ago

PLEASE explain how we got this $X = Xl \cdot 2^{n/2} + Xr$

[^](#) | [v](#) • [Reply](#) • [Share](#) ›**rohit** → n@733d • 4 months ago

shift left by n/2 bits the first half bits of X (i.e. Xl) and then add other half.. so 1101 (11 = xl , 01 = xr) $(2^{n/2} \cdot xl = 1100 + 0001 (xr) = x(1101)$

[^](#) | [v](#) • [Reply](#) • [Share](#) ›

   **blaze** · a year ago

how did the recurrence expression comes

$$t(n) = 4t(n/2) + o(n)$$

  · Reply · Share ›**RK** → blaze · a year ago

The problem of size(n) is reduced to 4 problems of size(n/2) and it would take extra $O(n)$ time to multiply with 2^n , $2^{(n/2)}$

  · Reply · Share ›**Amit** · a year agoFor normal high school multiplication, the complexity is $O(n^2)$ right?1   · Reply · Share ›**RK** → Amit · a year ago

yes

  · Reply · Share ›**Spanky** · a year agoMistake : $O(n^{159}) \rightarrow O(n^{158}) \approx \log_2(3)$   · Reply · Share ›**anvesh** · a year ago

can any one write code for complex numbers multiplication using FFT

  · Reply · Share ›**frenzydude** · a year agoThis can be done even faster in $O(n \log n)$ using FFT.  · Reply · Share ›**sudhanshu** · a year ago

can any1 pls tell me wht are floor n ceil values ??

  · Reply · Share ›**Himanshu Dagar** → sudhanshu · a year ago

floor means greatest integer less than or equal to that no
and ceil is least integer greater than or equal to that no

e.g floor(2.6)=2;

ceil(2.6)=3;

defined in math.h

  · Reply · Share ›



Sanjeet · a year ago

/*****Sanjeet from NIT allahabad *****/ you can simply adjust the value of i & j to get the multiplication of two strings .

```
#include<stdio.h>
#include<math.h>
int convertbin(char a[]);
int main()
{
char string1[4]="0111";
char string2[4]="0111";
int DecimalNum1=0;
int DecimalNum2=0;
int Product=0;
DecimalNum1=convertbin(string1);
DecimalNum2=convertbin(string2);
Product=DecimalNum1*DecimalNum2;
printf("Product of two string =%d\n",Product);
return 0;
}
```

[see more](#)

^ | v · [Reply](#) · [Share](#) ›



Shabaz Ahmed → [Sanjeet](#) · a year ago

The point here is to do the multiplication in less than $O(n^2)$..!

1 ^ | v · [Reply](#) · [Share](#) ›



Raunak Lakhwani · a year ago

```
public class Main {

static String first = "1", second = "00110";

/**

* @param args

*/

public static void main(String[] args) {

//int n = makeEqualString(first, second);

System.out.println(multiplication(first, second));

}
```



```
public static int makeEqualString(StringBuffer a, StringBuffer b) {
```

```
if (a.length() > b.length()) {
```

[see more](#)

^ | v • Reply • Share ›



Trilok • 2 years ago

Typo mistake ..

$$XY = 2n \cdot X_l Y_l + 2 \cdot \text{ceil}(n/2) \cdot [(X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r] + X_r Y_r$$

^ | v • Reply • Share ›



komal • 2 years ago

```
#include
```

```
#include
```

```
#include
```

```
using namespace std;
```

```
int pow(int x,int n)
```

```
{
```

```
if(x==0) return 0;
```

```
if(n==0) return 1;
```

```
if(n==1) return (x);
```

```
return (x*pow(x,n-1));
```

```
}
```

```
void print(string s)
```

```
{
```

```
for(int i=0;i<=s.size();i++)
```

```
cout<<s[i]<<" ";
```

```
cout<<endl;
```

[see more](#)

^ | v • Reply • Share ›



akshat gupta → komal • 2 years ago

- 1.you are entering a number in binary representation(serves no purpose here),
- 2.converting it to decimal,
- 3.printing the multiplication...

PROBLEM:

question asks you to find an algorithm for "fast multiplication of 2 numbers"..

^ | v • Reply • Share ›



akshat gupta • 2 years ago



```
int karatsuba(int n1,int n2)
{
    int l1,l2,l;
    int i,mask1=1;
    int a,b,c,d;
    int res1,res2,res3;

    l1=dig(n1);//returns number of digits
    l2=dig(n2);

    if(n1==0 || n2==0)
        return(0);

    l=min(l1,l2);

    if(l>1)
    {
        for(i=1;i<=l/2;i++)
            mask1=mask1*10;
```

[see more](#)

^ | v • Reply • Share ›



akshat gupta • 2 years ago

C IMPLEMENTATION:

```
int karatsuba(int n1,int n2)
{
    int l1,l2,l;
    int i,mask1=1;
    int a,b,c,d;
    int res1,res2,res3;
    l1=dig(n1);//returns number of digits
    l2=dig(n2);
    if(n1==0 || n2==0)
        return(0);

    l=min(l1,l2);

    if(l>1)
    {
        for(i=1;i<=l/2;i++)
            mask1=mask1*10;
```

[see more](#)

^ | v • Reply • Share ›

**Sandeep Jain** · 2 years ago

Thanks for pointing this out. We have corrected it.

^ | v · Reply · Share ›

**Anwit Roy** · 2 years agoThere is a typo: it should be $XIYr + XrYl = (Xl + Xr)(Yl + Yr) - XIYl - XrYr$.

1 ^ | v · Reply · Share ›



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Thanks. Very interesting lectures.

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forgot to see that part ;)

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thanks

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