Dynamic Programming | Set 37 (Boolean Parenthesization Problem)

Given a boolean expression with following symbols.

```
Symbols
    'T' ---> true
    'F' ---> false
```

And following operators filled between symbols

```
Operators
    &   ---> boolean AND
    |   ---> boolean OR
    ^   ---> boolean XOR
```

Count the number of ways we can parenthesize the expression so that the value of expression evaluates to true.

Let the input be in form of two arrays one contains the symbols (T and F) in order and other contains operators (&, | and ^}

Examples:

```
Input: symbol[]    = {T, F, T}
       operator[]  = {^, &}
Output: 2
The given expression is "T ^ F & T", it evaluates true
in two ways "((T ^ F) & T)" and "(T ^ (F & T))"

Input: symbol[]    = {T, F, F}
       operator[]  = {^, |}
Output: 2
The given expression is "T ^ F | F", it evaluates true
in two ways "( (T ^ F) | F )" and "( T ^ (F | F) )".

Input: symbol[]    = {T, T, F, T}
       operator[]  = {|, &, ^}
Output: 4
The given expression is "T | T & F ^ T", it evaluates true
in 4 ways ((T|T)&(F^T)), (T|(T&(F^T))), (((T|T)&F)^T)
and (T|((T&F)^T)).
```

Solution:
Let **T(i, j)** represents the number of ways to parenthesize the symbols between i and j (both inclusive) such that the subexpression between i and j evaluates to true.

$$T(i,j) = \sum_{k=i}^{j-1} \begin{cases} T(i,k) * T(k+1,j) & \text{If operator[k] is '\&'} \\ Total(i,k) * Total(k+1,j) - F(i,k) * F(k+1,j) & \text{If operator[k] is '|'} \\ T(i,k) * F(k+1,j) + F(i,k) * T(k+1) & \text{If operator[k] is '\^'} \end{cases}$$

$$\text{Total(i, j)} = \text{T(i, j)} + \text{F(i, j)}$$

Let **F(i, j)** represents the number of ways to parenthesize the symbols between i and j (both inclusive) such that the subexpression between i and j evaluates to false.

$$F(i,j) = \sum_{k=i}^{j-1} \begin{cases} Total(i,k) * Total(k+1,j) - T(i,k) * T(k+1,j) & \text{If operator[k] is '\&'} \\ F(i,k) * F(k+1,j) & \text{If operator[k] is '|'} \\ T(i,k) * T(k+1,j) + F(i,k) * F(k+1) & \text{If operator[k] is '\^'} \end{cases}$$

$$\text{Total(i, j)} = \text{T(i, j)} + \text{F(i, j)}$$

Base Cases:

```
T(i, i) = 1 if symbol[i] = 'T'
T(i, i) = 0 if symbol[i] = 'F'

F(i, i) = 1 if symbol[i] = 'F'
F(i, i) = 0 if symbol[i] = 'T'
```

If we draw recursion tree of above recursive solution, we can observe that it many overlapping subproblems. Like other dynamic programming problems, it can be solved by filling a table in bottom up

manner. Following is C++ implementation of dynamic programming solution.

```
#include<iostream>
#include<cstring>
using namespace std;

// Returns count of all possible parenthesizations that lead to
// result true for a boolean expression with symbols like true
// and false and operators like &, | and ^ filled between symbols
int countParenth(char symb[], char oper[], int n)
{
    int F[n][n], T[n][n];

    // Fill diaginal entries first
    // All diagonal entries in T[i][i] are 1 if symbol[i]
    // is T (true).  Similarly, all F[i][i] entries are 1 if
    // symbol[i] is F (False)
    for (int i = 0; i < n; i++)
    {
        F[i][i] = (symb[i] == 'F')? 1: 0;
        T[i][i] = (symb[i] == 'T')? 1: 0;
    }

    // Now fill T[i][i+1], T[i][i+2], T[i][i+3]... in order
    // And F[i][i+1], F[i][i+2], F[i][i+3]... in order
    for (int gap=1; gap<n; ++gap)
    {
        for (int i=0, j=gap; j<n; ++i, ++j)
        {
            T[i][j] = F[i][j] = 0;
            for (int g=0; g<gap; g++)
            {
                // Find place of parenthesization using current value
                // of gap
                int k = i + g;

                // Store Total[i][k] and Total[k+1][j]
                int tik = T[i][k] + F[i][k];
                int tkj = T[k+1][j] + F[k+1][j];

                // Follow the recursive formulas according to the current
                // operator
                if (oper[k] == '&')
                {
                    T[i][j] += T[i][k]*T[k+1][j];
                    F[i][j] += (tik*tkj - T[i][k]*T[k+1][j]);
                }
                if (oper[k] == '|')
                {
                    F[i][j] += F[i][k]*F[k+1][j];
                    T[i][j] += (tik*tkj - F[i][k]*F[k+1][j]);
                }
                if (oper[k] == '^')
```

```
                {
                    T[i][j] += F[i][k]*T[k+1][j] + T[i][k]*F[k+1][j];
                    F[i][j] += T[i][k]*T[k+1][j] + F[i][k]*F[k+1][j];
                }
            }
        }
    }
    return T[0][n-1];
}

// Driver program to test above function
int main()
{
    char symbols[] = "TTFT";
    char operators[] = "|&^";
    int n = strlen(symbols);

    // There are 4 ways
    // ((T|T)&(F^T)), (T|(T&(F^T))), (((T|T)&F)^T) and (T|((T&F)^T))
    cout << countParenth(symbols, operators, n);
    return 0;
}
```

Output:

```
4
```

Time Complexity: $O(n^3)$
Auxiliary Space: $O(n^2)$

References:
http://people.cs.clemson.edu/~bcdean/dp_practice/dp_9.swf

Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above

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T(i,j) of length 2
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helper • 6 months ago
a very nice code.
a small suggestion:
int tik = T[i][k] + F[i][k];
int tkj = T[k+1][j] + F[k+1][j];
we can put these two assignments in that these two assignments in second if. another suggestion. the later two if's can be made else if.
all this will save some time.
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excellent program.
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what is the application on boolean parenthesization??
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I think a dry run example and proper explanation for the formula is required for this topic.
Although the video is fine to understand. its really a good problem.
Thanks guys.
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ok i understand with the help of video
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Bhagwat Singh • a year ago

kindly provide the logic behind the mathematical formula.

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Rainer Hoffmann • a year ago

Hello,
running the program with this input from top of article:

Input: symbol[] = {T, F, F}
operator[] = {^, |}
gives Output = 2 instead of above mentioned
Output: 1
Could you please check?
Thank you!

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GeeksforGeeks Mod ➚ Rainer Hoffmann • a year ago

Thanks for pointing this out. The code seems to be giving the correct output. The output mentioned in example was incorrect. We have updated the example now.

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GOPI GOPINATH • a year ago

Wow the explanation given in the reference is pretty OSM..Nice one

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