

# The Merkle-Hellman Cryptosystem

ALGORITMOS E ESTRUTURAS DE  
DADOS

Prof. Tomás Silva

André Butuc - 103530 – 60%

Gonçalo Silva - 103668 - 40%

# Índice

---

- Introdução	2
- Material Adaptado	3
- Abordagem ao problema	4
- Iterative Brute Force	5
- Gráficos Iterative Brute Force	6
- Recursive Brute Force	7
- Gráficos Recursive Brute Force	8
- Clever Brute Force	9
- Gráficos Clever Brute Force	10
- Horowitz-Sahni	11
- Gráficos Horowitz-Sahni	14
- Schroeppe-Shamir	16
- Gráficos Schroeppe-Shamir	19
- Comentários/Interpretações	21
- Gráficos Finais	23
- Soluções Obtidas	27
- Código Desenvolvido	53

# Introdução

---

## The Merkle-Hellman Cryptosystem

Este trabalho prático desenvolvido no âmbito da cadeira Algoritmos e Estruturas de Dados tem como objetivo a implementação e o estudo de inúmeras técnicas de resolução do sistema de criptografia de Merkle e Hellman.

O criptosistema de Merkle-Hellman é aquilo que se chama um sistema de chave pública, ou seja, são usadas duas chaves, uma delas sendo pública usada para a criptografia e uma privada usada na descryptografia.

Este sistema é baseado no problema de somas de um subconjunto, sendo este um caso do mais complexo knapsack problem.

Seja dado um conjunto de números inteiros e um número  $s$ , existe um subconjunto em que a soma dos seus elementos dê  $s$ ? Este é o nosso problema geral: verificar se existe este subconjunto, identificá-lo e traduzi-lo numa sequência binária que denota o peso que cada número inteiro do conjunto inicial tem na soma.

Tendo como dados iniciais,  $n$  (o número de elementos do conjunto inicial),  $p$  (o vetor que contém os números inteiros do conjunto inicial) e  $desired\_sum$  (a soma que desejamos encontrar) iremos percorrer os dois ficheiros com os problemas: "103530.h" e "103668.h".

Os ficheiros ".h" estão construídos de forma a terem um intervalo de valores de " $n=10$ " a " $n=57$ ", onde em cada " $n$ " há um vetor  $p$  com " $n$ " números inteiros e em cada " $n$ " temos "20  $desired\_sum$ 's" às quais teremos que encontrar a sequência binária solução.

## Material adaptado

- Código relacionado com o algoritmo de ordenação "Quicksort" (swap, partition, quicksort) foi adaptado do site "GeeksforGeeks", presente na seguinte hiperligação:  
<https://www.geeksforgeeks.org/quick-sort/>
- Código relacionado com o algoritmo de ordenação "Merge Sort" (insertion\_sort e merge\_sort) foi adaptado dos slides teóricos de AED do professor Tomás Silva.

# Abordagem ao Problema

---

Na tentativa de resolução deste problema foram testadas diversas técnicas computacionais com o intuito de desenvolver algoritmos capazes de resolver o problema em causa, avaliando a eficiência dos mesmos.

Foram desenvolvidos 7 algoritmos:

- Iterative Brute Force ("iterative\_brute\_force")
- Recursive Brute Force ("recursive\_brute\_force")
- Clever Recursive Brute Force ("clever\_recursive\_brute\_force")
- Horowitz and Sahni ("horowitz\_sahni")
  - QuickSort version
  - MergeSort version
- Schroeppeel and Shamir ("shroeppel\_shamir")
  - QuickSort version
  - MergeSort version

# Iterative Brute Force

---

Complexidade Computacional:  $O(n2^n)$

A implementação “Iterative Brute Force” consiste em percorrer iterativamente todas as possibilidades de combinação dos “n” elementos do conjunto “p”, logo a complexidade computacional desta abordagem reduz-se a

$$O(2^n).$$

Em cada iteração é calculada a “test\_soma”, sendo esta incrementada por “p[bit]” dentro de um loop que percorre o valor de “bit” de 0 a n-1, que corresponde ao índice de cada elemento de “p”. No entanto, esta incrementação só ocorre se o valor de “bit” pertence à combinação binária de “comb” (podendo esta propriedade ser observada se a posição em que o valor bit a “um” se encontra em “bit” corresponde à posição em que um dos bits a “um” se encontra em “comb”).

As operações bitwise em C, facilitam bastante a extensão de operandos que não têm a mesma dimensão, permitindo fazer cálculos em base binária que se enquadram com a natureza do problema.

Esta última abordagem contribui com um acréscimo de complexidade de  $O(n)$  em cada iteração de  $2^n$  possibilidades de combinações binárias.

Totalizando as complexidades computacionais das abordagens internas ao algoritmo teremos de facto como complexidade final

$$O(n2^n)$$

## Gráficos Iterative Brute Force

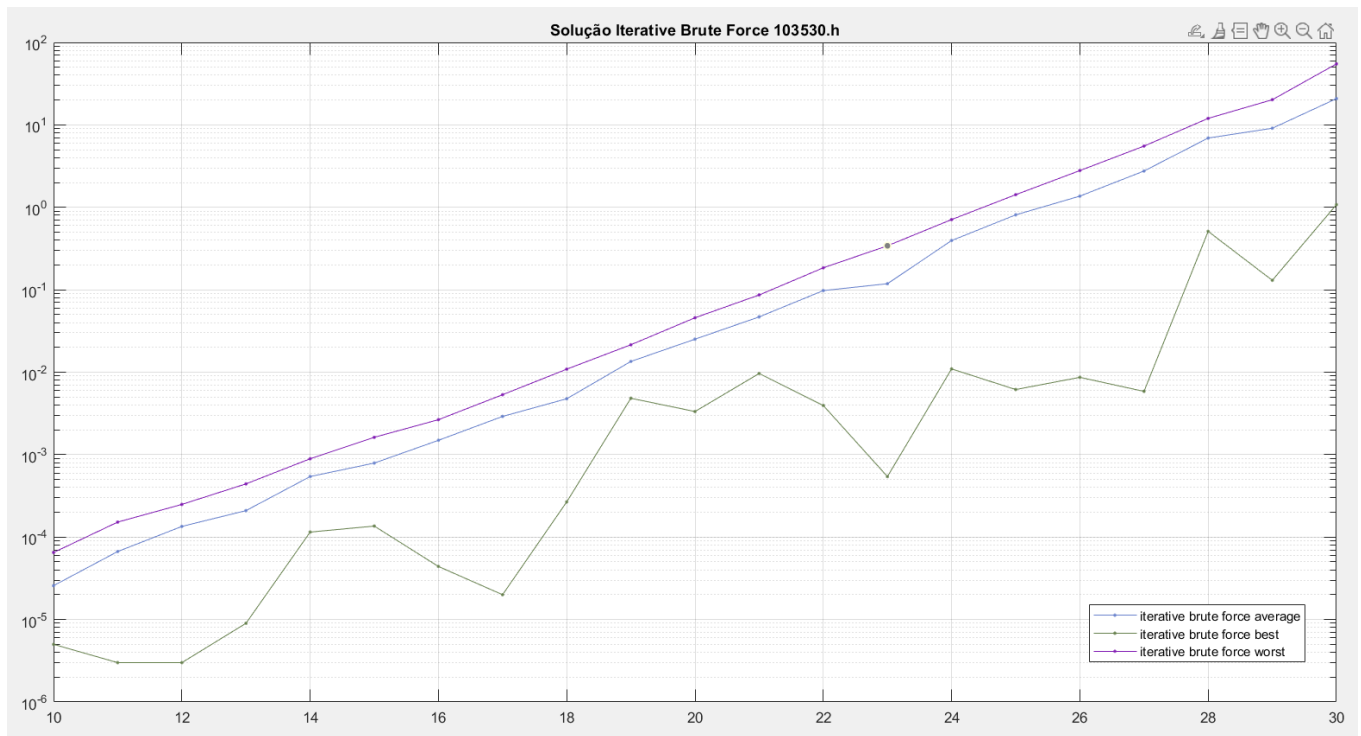


Figura 1 - Iterative Brute Force 103530.h

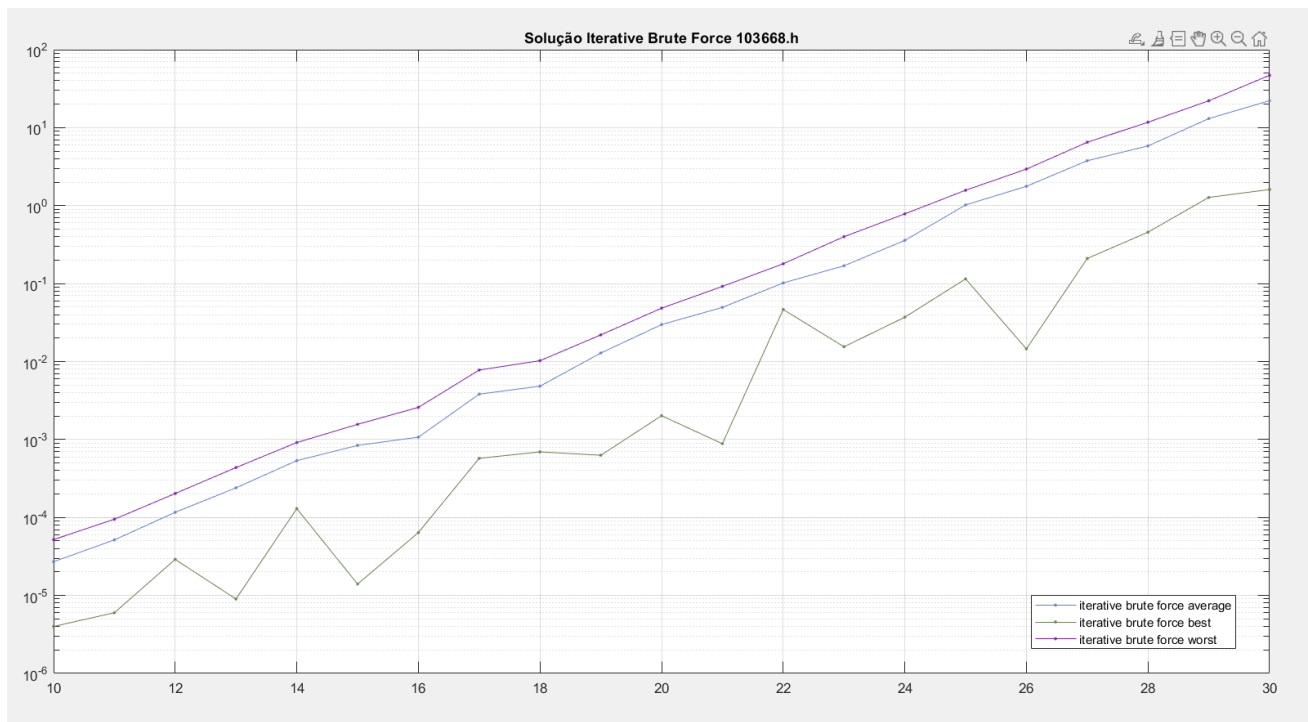


Figura 2 - Iterative Brute Force 103668.h

# Recursive Brute Force

---

Complexidade computacional:  $O(2^n)$

O algoritmo `recursive_brute_force`, como o nome indica, encontra a solução recursivamente. Em cada valor de "p" há um "branching": o bit do valor de "p" fica a 1 ou fica a 0. Em cada ramo, a soma parcial ("partial\_soma") é incrementada (no ramo do bit igual a 1) ou mantém-se igual (no ramo do bit igual a 0). De seguida é incrementado o índice do array "p" em cada ramo e é chamada de novo a função "recursive\_brute\_force" em cada ramo. Esta abordagem permite criar uma "árvore" com todas as possibilidades de combinações binárias que mapeiam a "presença" ou "não presença" do elemento a ser avaliado de "p" na soma parcial. Devido ao facto da função ser chamada duas vezes de forma recursiva e pelo facto de no pior caso esta o ser chamada "n" vezes, a complexidade computacional do algoritmo é totalizado a  $O(2^n)$ .

As condições de paragem da função recursiva são:

- `if(partial_soma == desired_soma)`, que significa que foi encontrada a solução. O retorno do valor 1 irá desencadear o desenvolvimento das sucessivas chamadas recursivas prévias da função. Sendo que o bit da "folha" que desencadeou a terminação da função com sucesso irá ser armazenado no array "b" e assim sucessivamente até à raiz;
- `if(current_index == n)`, que significa, caso verdadeiro, que a recursividade chegou ao fim do array "p" e como já não há mais elementos ou combinações a somar e como ainda não foi encontrada a solução, retorna o valor 0 que por sua vez indica que não foi encontrada a solução-soma.



## Gráficos Recursive Brute Force

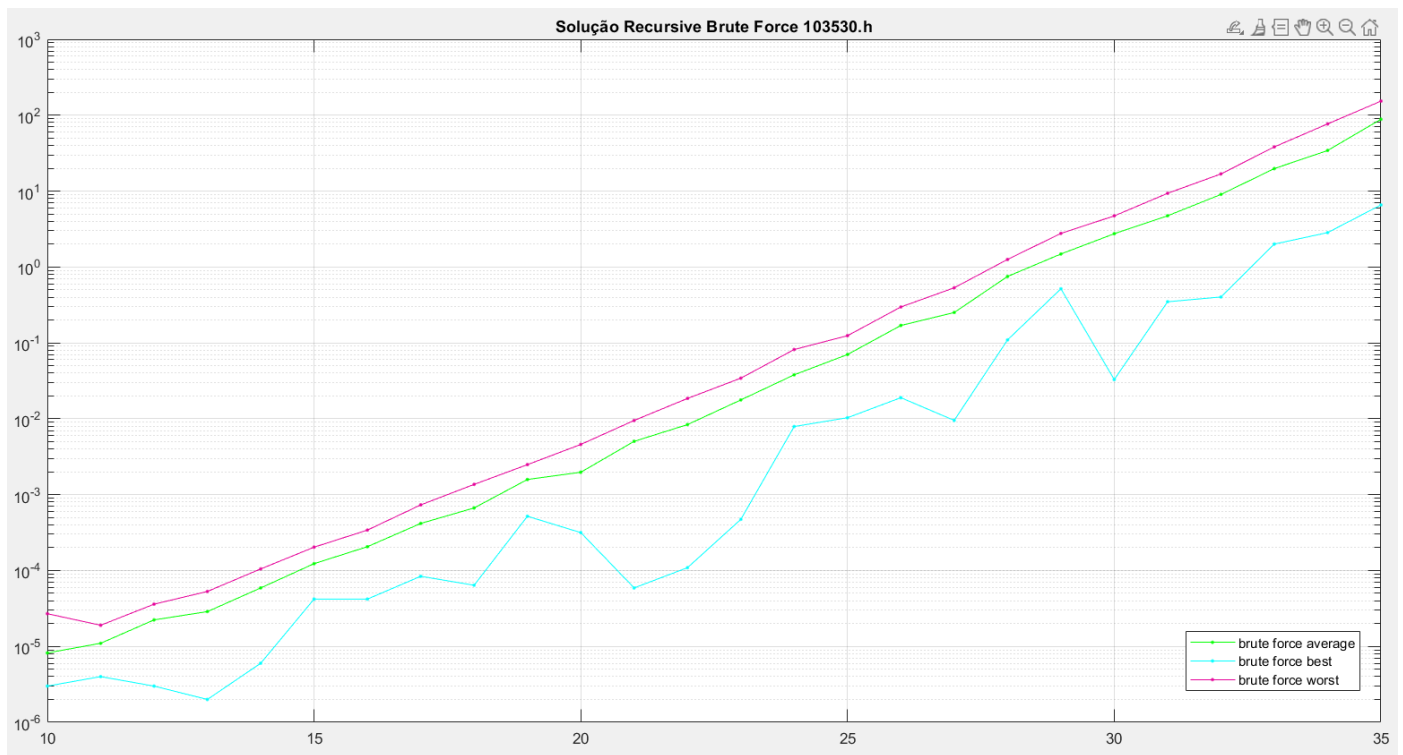


Figura 3 - Recursive Brute Force 103530.h

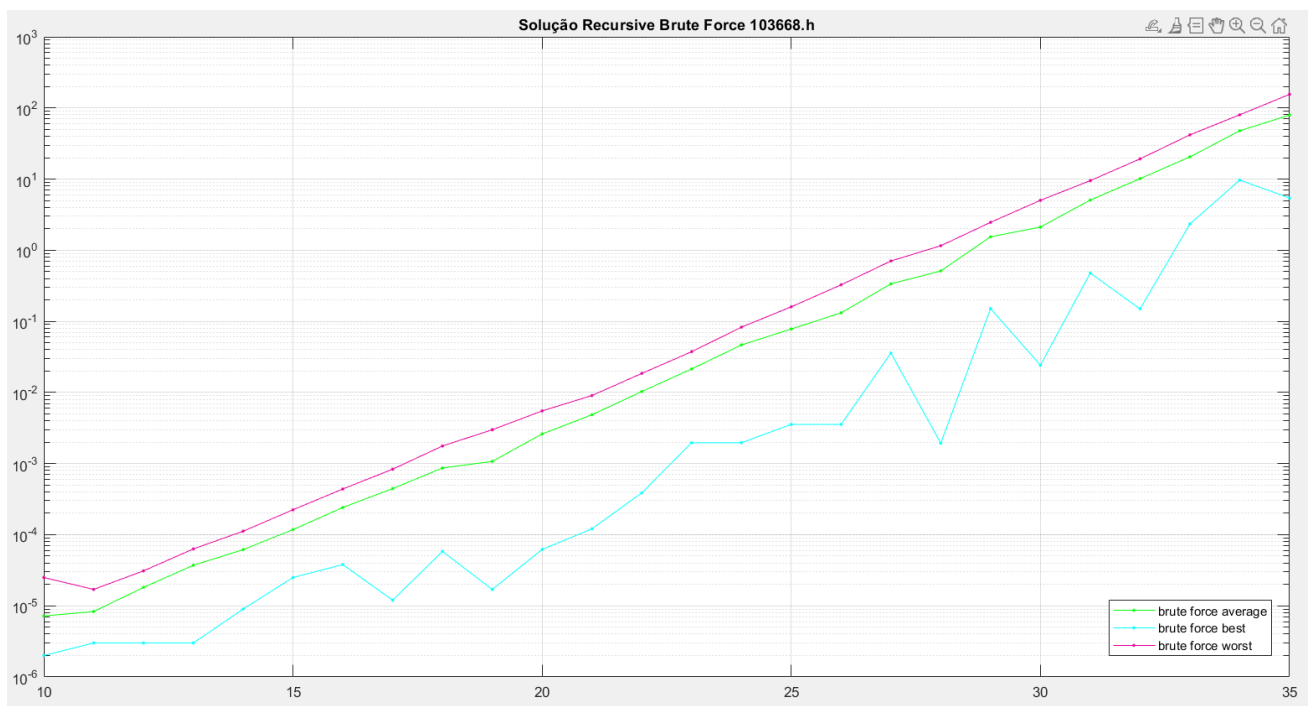


Figura 4 - Recursive Brute Force 103668.h

# Clever Recursive Brute Force

---

Complexidade computacional:  $O(2^n)$

O algoritmo `clever_recursive_brute_force`, é semelhante ao algoritmo `recursive_brute_force`, no entanto, apresenta mais dois parâmetros de entrada que servem para calcular uma soma parcial a partir dos últimos índices para os primeiros. E para além de apresentar as condições de paragem do algoritmo `recursive_brute_force`, também apresenta mais duas condições de paragem.

Este algoritmo não difere em complexidade computacional do anterior mencionado, no entanto, difere na capacidade de conseguir parar a sua execução quando, após observar provas concretas, interpreta que não será possível encontrar uma solução ao problema.

Os dois parâmetros novos serão utilizados para interpretar se é de facto possível chegar à solução com os valores dispostos ou a soma será sempre inferior à soma desejada. O parâmetro `current_last_index` será decrementado em cada chamada da função recursiva e o parâmetro `bigger_soma` irá guardar a soma dos valores sucessivos de `p[current_last_index]`. Quando o `current_index` é igualado a `current_last_index + 1` significa que já passámos a metade de `p`, logo se a soma de `bigger_soma` com `partial_soma` continuar a ser inferior a `desired_soma` então não é possível alcançar a soma desejada com os números presentes no vetor `p`.

A condição de paragem restante é `if(partial_soma > desired_soma)`. Pelo facto do array `p` se encontrar ordenado e pelo facto de apresentar valores inteiros positivos, sabendo que também a `partial_soma` é sempre incrementada e não decrementada, então a partir do momento em que `partial_soma` é superior a `desired_soma`, não há como "voltar atrás" e assim é possível afirmar que não é encontrada a solução.

## Gráficos Clever Recursive Brute Force

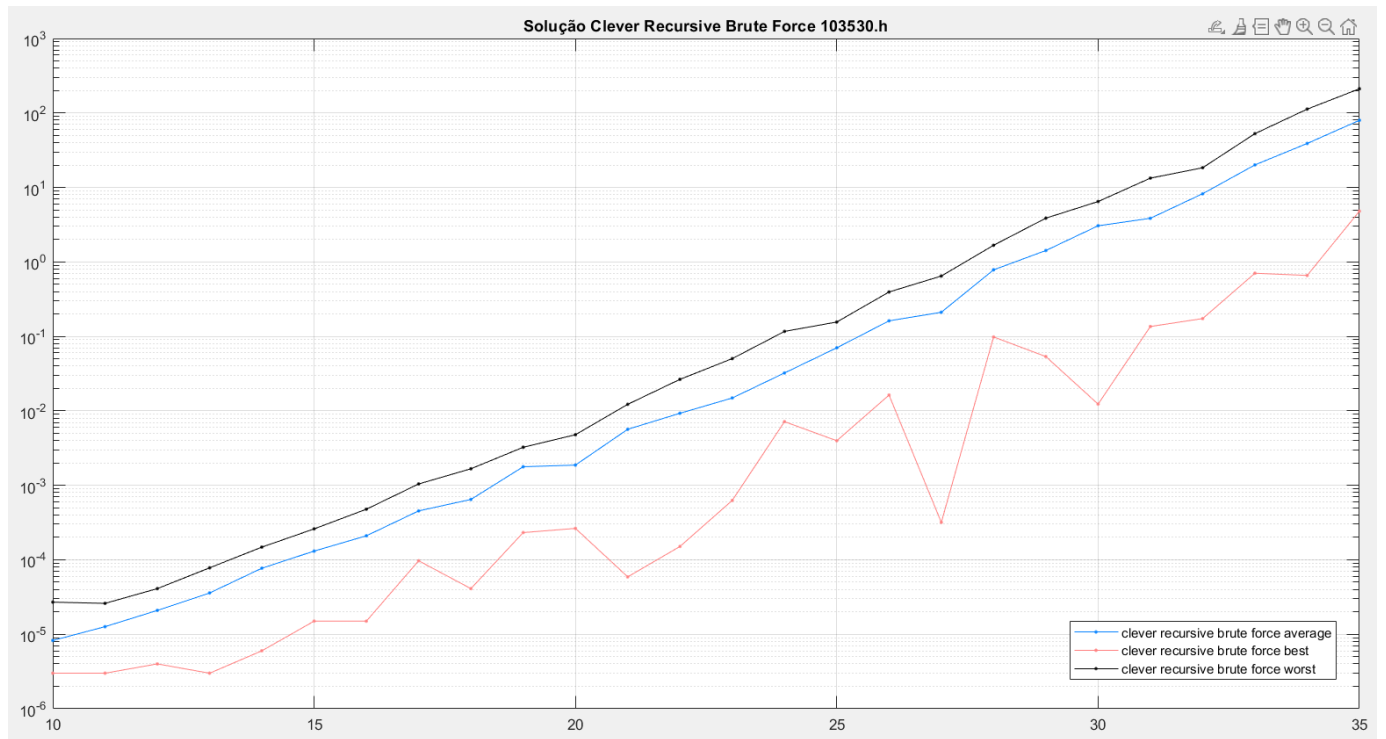


Figura 5 - Clever Recursive Brute Force 103530.h

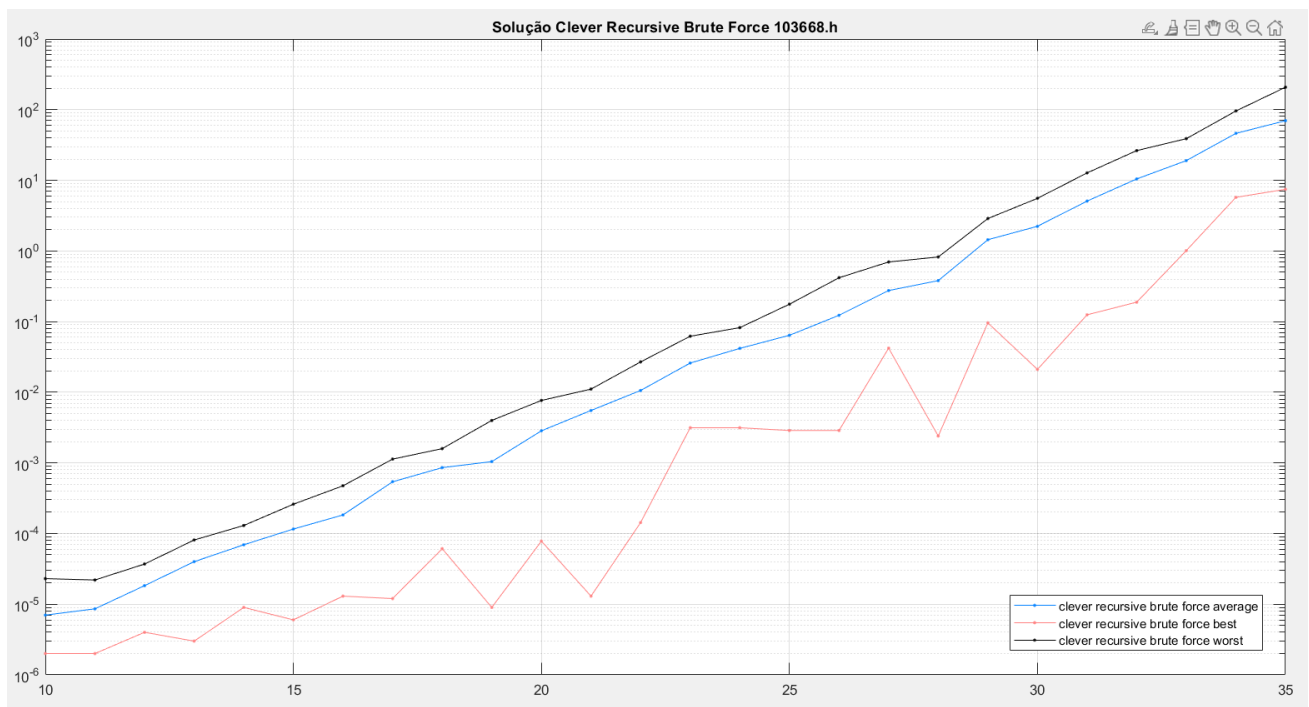


Figura 6 - Clever Recursive Brute Force 103668.h

# Horowitz-Sahni

Complexidade computacional:  $O(2^{\frac{n}{2}} \frac{n}{2})$

O algoritmo de Horowitz-Sahni levou-nos a criar uma struct ("mascara\_struct\_t") que guardasse e mapeasse de melhor forma as somas parciais e as máscaras binárias associadas a elas.

Esta abordagem utiliza a técnica meet-in-the-middle.

Inicialmente o array "p" é dividido de forma aproximadamente igual em dois arrays, "p1" e "p2". Pelo facto do array "p" se encontrar ordenado, "p1" terá os valores menores de "p" e "p2" os valores maiores de "p". De seguida são calculadas todas as somas internas possíveis de "p1" e de "p2", estas somas são calculadas através da função "soma\_all\_subsets", que consiste na adaptação das funções recursivas anteriormente referidas, para que estas estejam em concordância com a struct "mascara\_struct\_t", não existindo no entanto uma condição de paragem à exceção de "if(current\_index == n)" (pois queremos todas as combinações possíveis e não a combinação-solução como anteriormente).

Em cada chamada recursiva é feita uma operação bitwise OR ("|") do valor de "mascara" com o valor de 1 bit deslocado para a esquerda de "current\_index" bits, caso o bit do elemento de "p" seja "um".

Por exemplo, vamos supor que temos o valor inicial da "mascara" 0 e que o elemento que estamos a avaliar tem índice de 5 no vetor "p" e vamos colocar este elemento com bit 1, ou seja, participa na soma parcial, então "mascara | (1<<5)" = "mascara | 100000" = "000000 | 100000" = "100000" (o valor da mascara é extendido com bits "zeros" mais significativos até ser igual à dimensão com a qual estamos a fazer a operação bitwise OR, o mesmo acontece inversamente), esta operação e propriedade lógica permite-nos "armazenar" o valor do bit do elemento (que pode ser 0 ou 1) na máscara.

No entanto, é ainda de realçar, dentro da função "soma\_all\_subsets" que a forma de armazenamento dos dados nas estruturas no array "somas", array que contém todas as struct "mascara\_struct\_t" inicializadas, teve que ser feita da seguinte forma:

Para não haver conflito quando quisermos guardar elementos no array "somas", quando há o branching do ramo que iguala o bit do elemento a ser avaliado a 0 (não participa na soma parcial) é incrementado o índice dos índices ímpares e o ramo que iguala o bit do elemento a ser avaliado a 1 (participa na soma parcial) é incrementado o índice dos índices pares. Implementámos esta abordagem pois observámos que se incrementássemos de forma sucessiva o valor de "index", por exemplo com "index++" iria existir "over-write" de dados já presentes no array "somas".

Após termos os arrays "a" e "c" com as "mascara\_struct\_t" atualizadas em concordância com todas as somas possíveis de cada array "p1" e "p2", sendo que "a" tem as somas de "p1" e "c" as somas de "p2", é necessário ordenar os arrays por ordem crescente. Para este efeito elaboramos duas versões, uma em que a ordenação era feita através do algoritmo de ordenação "QuickSort" e uma em que a ordenação era feita através do algoritmo de ordenação "MergeSort".

O último passo é efetivamente o algoritmo de Horowitz-Sahni, onde percorremos num while loop os arrays "a" e "c", "a" a começar do primeiro elemento e "c" a começar do último até um dos índices de "a" ou de "c" "transbordar" ou até encontrarmos a soma desejada.

Em cada iteração é avaliado se a soma parcial é superior à soma desejada, caso o seja, então o índice de "c" é decrementado, caso contrário, se a soma parcial for inferior à soma desejada então é incrementado o índice de "a".

Quando for atingida a soma desejada, a solução consistirá na "junção" das máscaras das somas parciais, a máscara da soma parcial de "c" terá o peso dos bits mais significativos e a máscara da soma parcial de "a" terá o peso dos bits menos significativos.

Relativamente à complexidade computacional:

- “soma\_all\_subsets”, por si só tem uma complexidade computacional de  $O(2^n)$ , no entanto, no enquadramento do problema, irá contribuir com

$$O(2^{\frac{n}{2}+1})$$

- o algoritmo de ordenação “QuickSort”, no caso médio tem complexidade de  $O(n \log(n))$ , mas enquadrando-o ao nosso problema terá  $O(2^{\frac{n}{2}+1} \log_2(2^{\frac{n}{2}+1}))$ , que simplificado ficará

$$O(2^{\frac{n}{2}+1} (\frac{n}{2} + 1))$$

- o algoritmo de ordenação “MergeSort”, no pior caso tem complexidade de  $O(n \log(n))$ , mas enquadrando-o ao nosso problema terá  $O(2^{\frac{n}{2}+1} \log_2(2^{\frac{n}{2}+1}))$ , que simplificado ficará  $O(2^{\frac{n}{2}+1} (\frac{n}{2} + 1))$
- o while-loop nos piores dos casos itera “length\_a”+“length\_c” vezes, logo tem como complexidade computacional  $O(2^{\frac{n}{2}} + 2^{\frac{n}{2}})$   
=

$$O(2^{\frac{n}{2}+1})$$

A complexidade do algoritmo, tanto na versão de QuickSort como na versão MergeSort, assume a componente de complexidade de maior peso, logo o algoritmo terá

$$O(2^{\frac{n}{2}} \frac{n}{2}).$$

## Gráficos Horowitz-Sahni (103530.h)

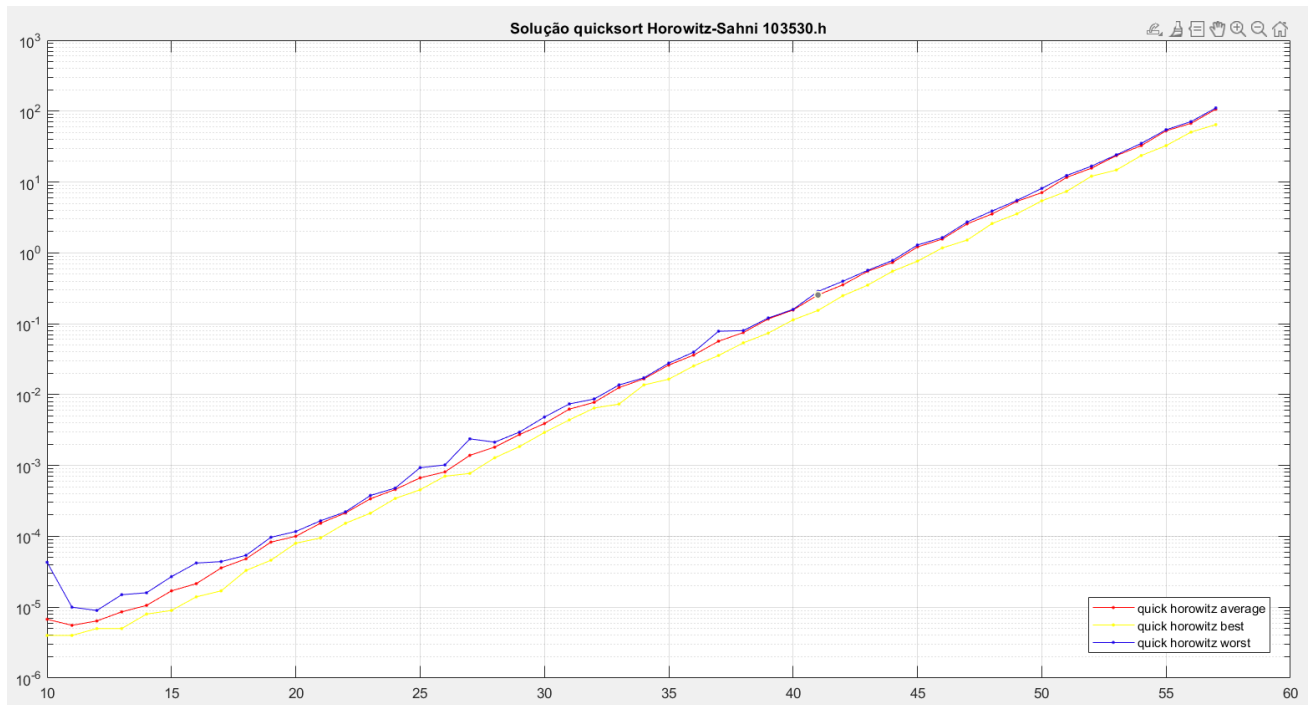


Figura 7 - Horowitz-Sahni QuickSort 103530.h

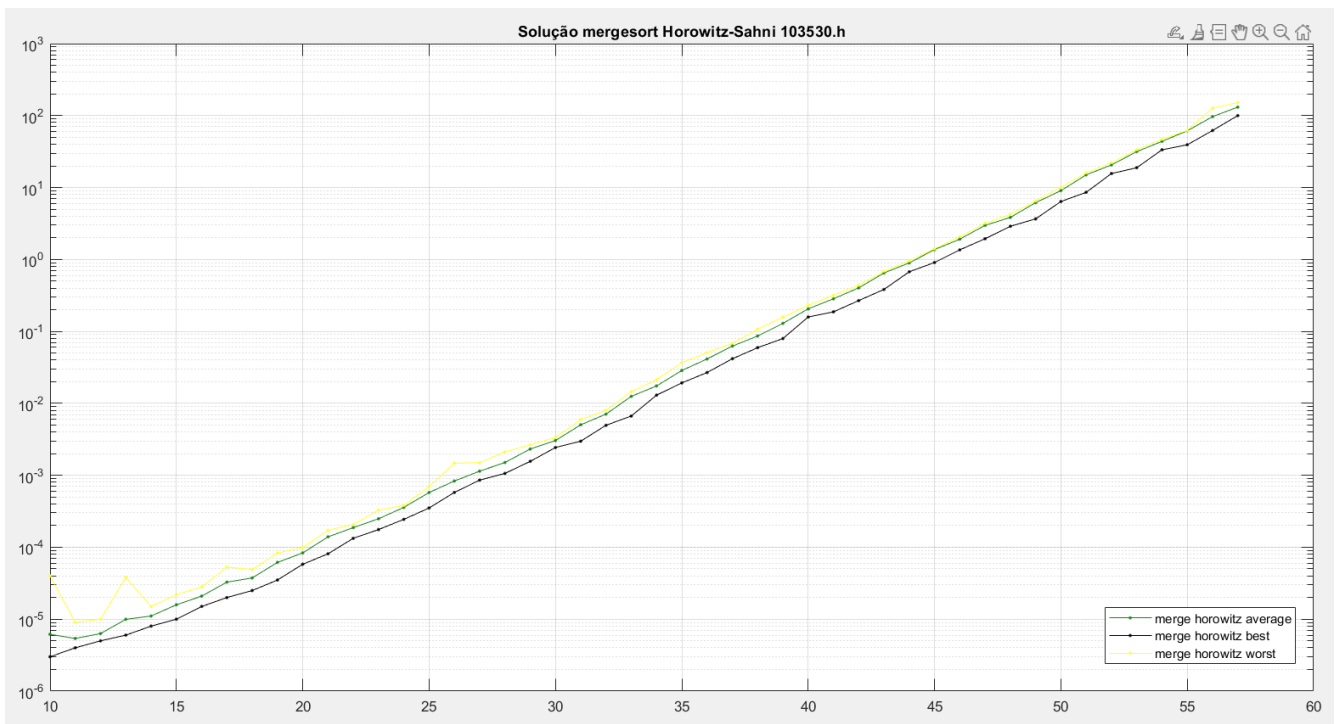


Figura 8 - Horowitz-Sahni MergeSort 103530.h

## Gráficos Horowitz-Sahni (103668.h)

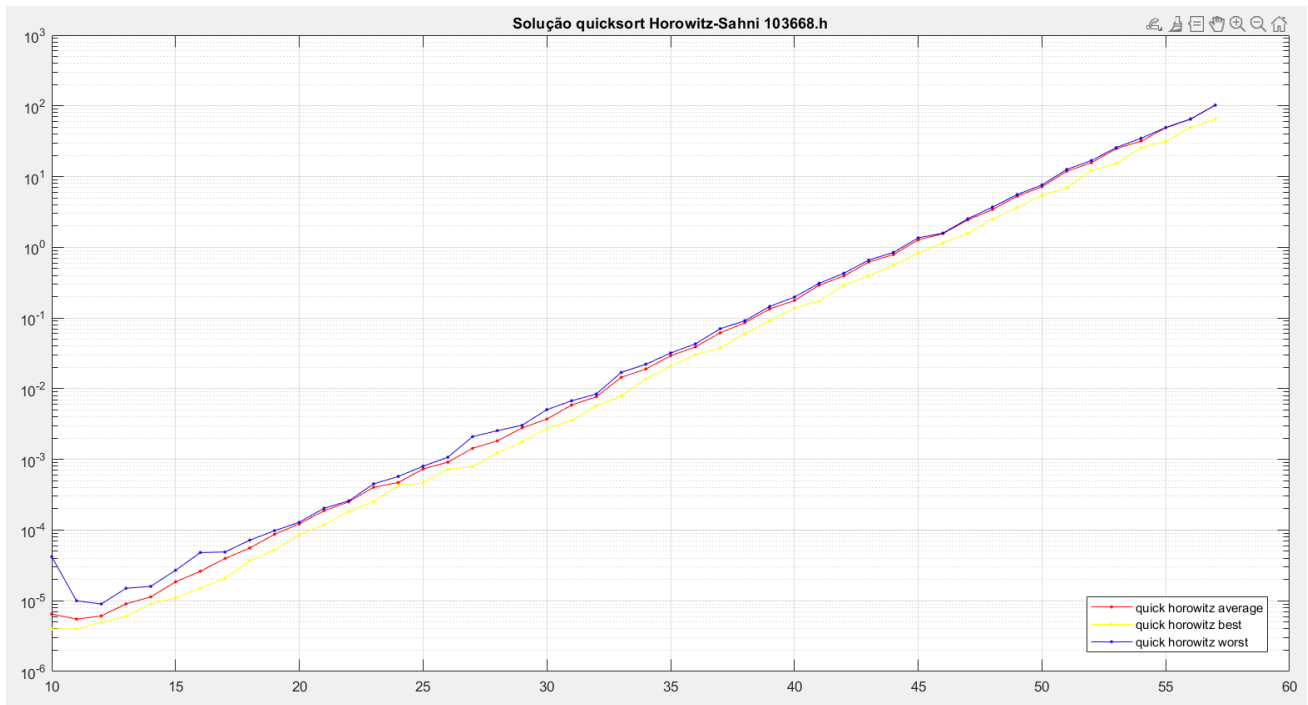


Figura 9 - Horowitz-Sahni QuickSort 103668.h

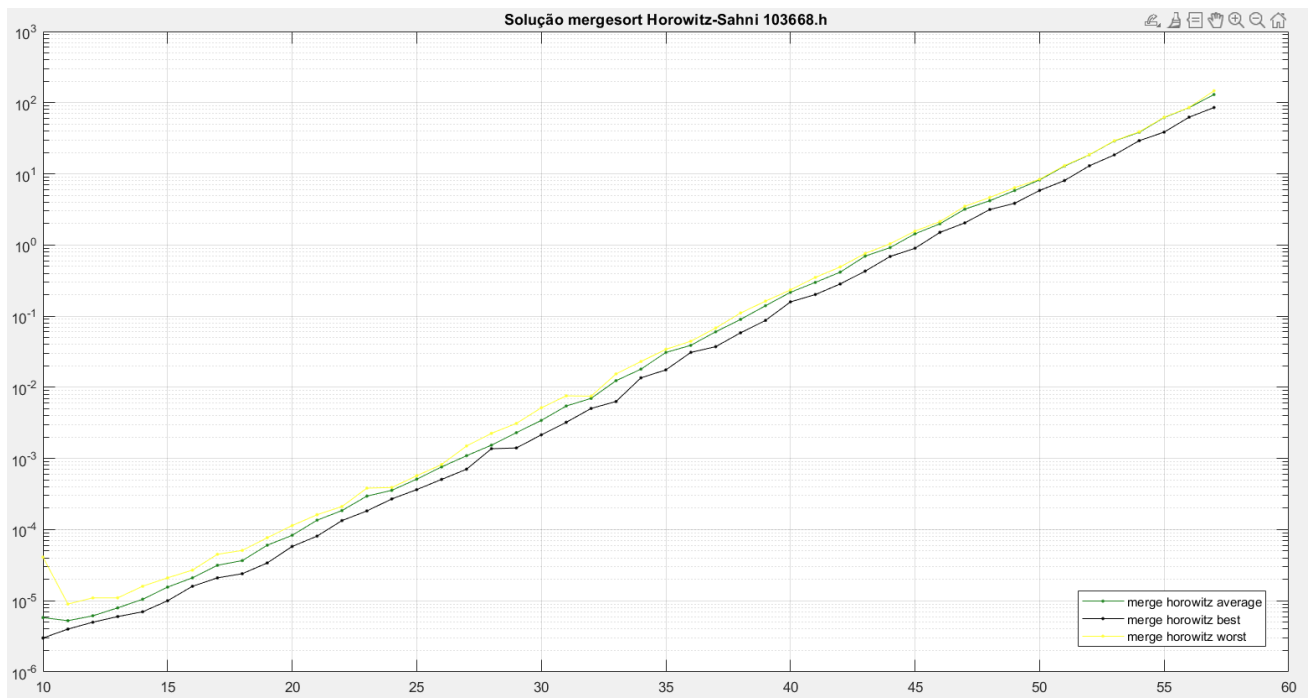


Figura 10 - Horowitz-Sahni MergeSort 103668.h



# Schroeppe-Shamir

---

Complexidade computacional:  $O(2^{\frac{n}{2}} \frac{n}{2})$

O algoritmo de Schroeppe-Shamir é um algoritmo bastante semelhante ao algoritmo de Horowitz-Sahni. Aplica também a técnica meet-in-the-middle, no entanto é uma abordagem que requer menos memória, pois a soma parcial é calculada através da utilização de uma min-heap e de uma max-heap.

Nesta abordagem o array "p" é dividido de forma aproximadamente igual em quatro arrays, "p1", "p2", "p3" e "p4". De seguida são calculadas todas as somas possíveis internas de cada array e ordenadas através do QuickSort (ou do MergeSort).

Foi necessário construir uma struct de elementos que vamos guardar nas heaps, sendo que este elemento tem uma "mascara" (combinação binária), "soma", "indice1" e "indice2", estes últimos atributos irão mapear os elementos dos arrays "a1" e "a2" (para a min-heap) e "b1" e "b2" (para a max-heap), respetivamente.

A max-heap e min-heap são inicializados com os valores de "b1" e "a1", respetivamente, somados com o último valor de "b2" para a max-heap e com o primeiro valor de "a2" para a min-heap. O mesmo acontece para as máscaras com o procedimento mencionado no algoritmo de "Horowitz-Sahni".

Para manipular as heaps, foram utilizados os métodos "push" e "pop". A função "push" insere um elemento na heap e reordena-a consoante o seu tipo, se o tipo for "1", então corresponde à min-heap, se o seu tipo for "2", então corresponde à max-heap. A função "pop" remove um elemento da heap e retorna-o, re-ordenando de seguida a heap consoante o seu tipo como descrito na função "push".

No algoritmo de "Schroeppe-Shamir", ao contrário do que acontece no algoritmo de "Horowitz-Sahni", em cada iteração é calculada a soma dos elementos que se encontram no topo das heaps. Caso a soma seja inferior à desejada, é retirado da min-heap o elemento que

se encontra no seu topo, é incrementado o atributo "indice2", caso o valor de "indice2" não seja superior ao tamanho de "a2", então é criada uma nova "heapStruct\_t" com os valores atualizados do elemento retirado e é voltada a ser colocada na min-heap. A condição de ser inferior ao tamanho de "a2" deve-se ao facto de, caso esta seja verdadeira, já foram "testadas" todas as combinações do valor "indice1" do vetor "a1" com os valores de "a2", logo não deve permanecer na min-heap.

O mesmo acontece na max-heap, mas agora com os arrays "b1" e "b2" e o decremento do valor de "indice2".

Quando é encontrada a soma-solução, então a solução é guardada tal como foi guardada no algoritmo de "Horowitz-Sahni".

Relativamente à complexidade computacional:

- "soma\_all\_subsets", por si só tem uma complexidade computacional de  $O(2^n)$ , no entanto, no enquadramento do problema, irá contribuir com

$$O(2^{\frac{n}{4}+2})$$

- o algoritmo de ordenação "QuickSort", no caso médio tem complexidade de  $O(n \log(n))$ , mas enquadrando-o ao nosso problema terá  $O(2^{\frac{n}{4}+2} \log_2(2^{\frac{n}{4}+2}))$ , que simplificado ficará

$$O(2^{\frac{n}{4}+2} (\frac{n}{4} + 2))$$

- o algoritmo de ordenação "MergeSort", no pior caso tem complexidade de  $O(n \log(n))$ , mas enquadrando-o ao nosso problema terá  $O(2^{\frac{n}{4}+2} \log_2(2^{\frac{n}{4}+2}))$ , que simplificado ficará

$$O(2^{\frac{n}{4}+2} (\frac{n}{4} + 2))$$

- o while-loop nos piores dos casos itera até ser falsa a sua condição de while loop, o que significa que uma das heaps ficou vazia, o que nos dá uma complexidade computacional de

$$O(2^{\frac{n}{2}+2} \frac{n}{2})$$

A complexidade do algoritmo, tanto na versão de QuickSort como na versão MergeSort, assume a componente de complexidade de maior peso, logo o algoritmo terá  $O(2^{\frac{n}{2}+2} \frac{n}{2})$ .

## Gráficos Schroepel-Shamir (103530.h)

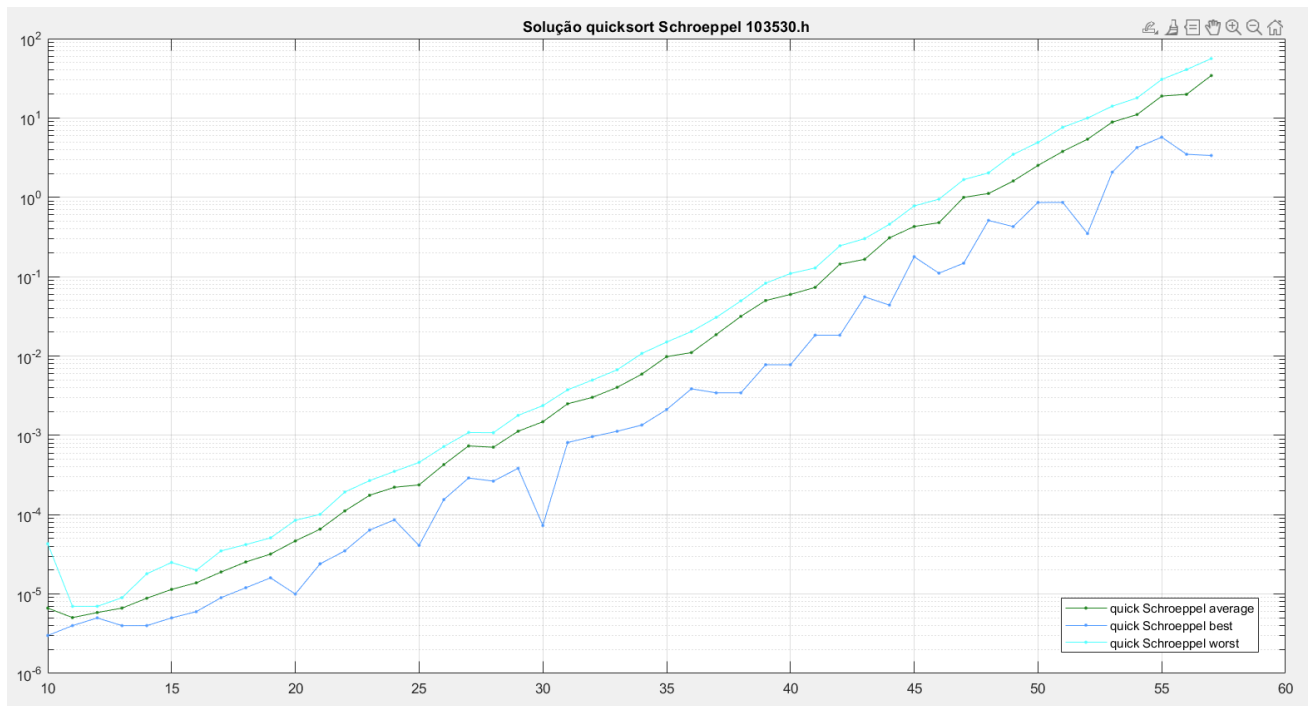


Figura 11 - Schroepel-Shamir QuickSort 103530.h

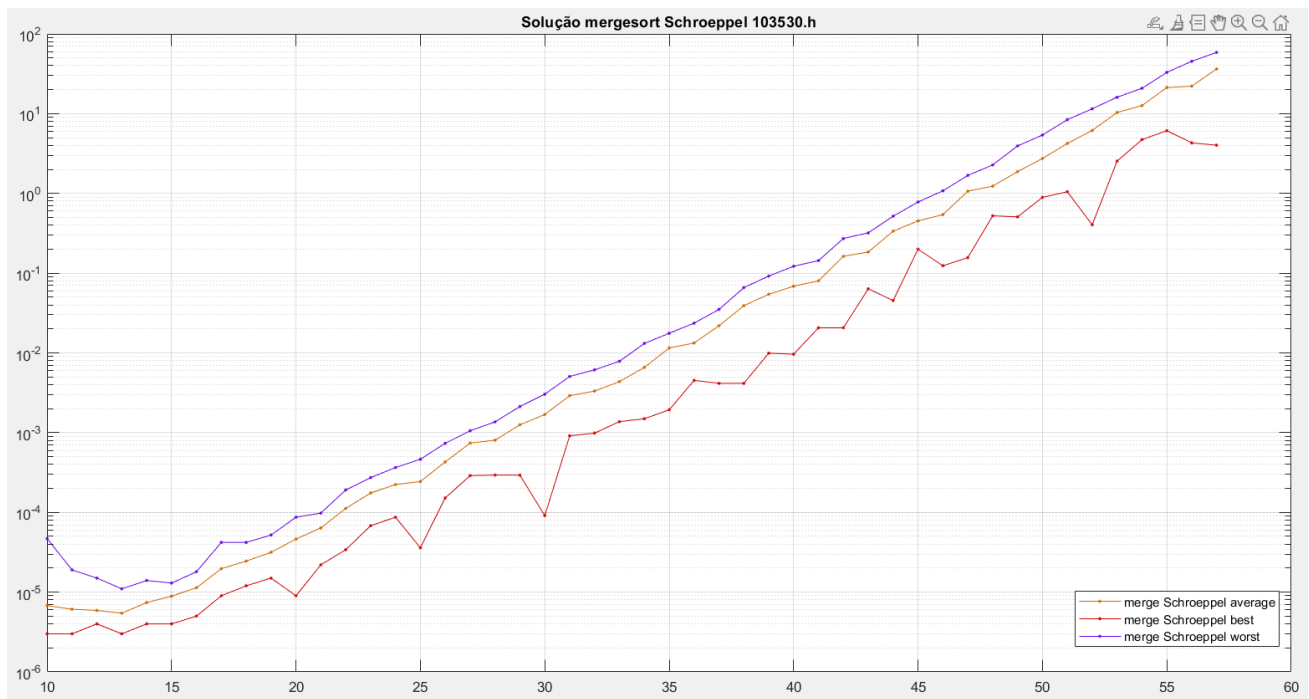


Figura 12 - Schroepel-Shamir MergeSort 103530.h

## Gráficos Schroeppel-Shamir (103668.h)

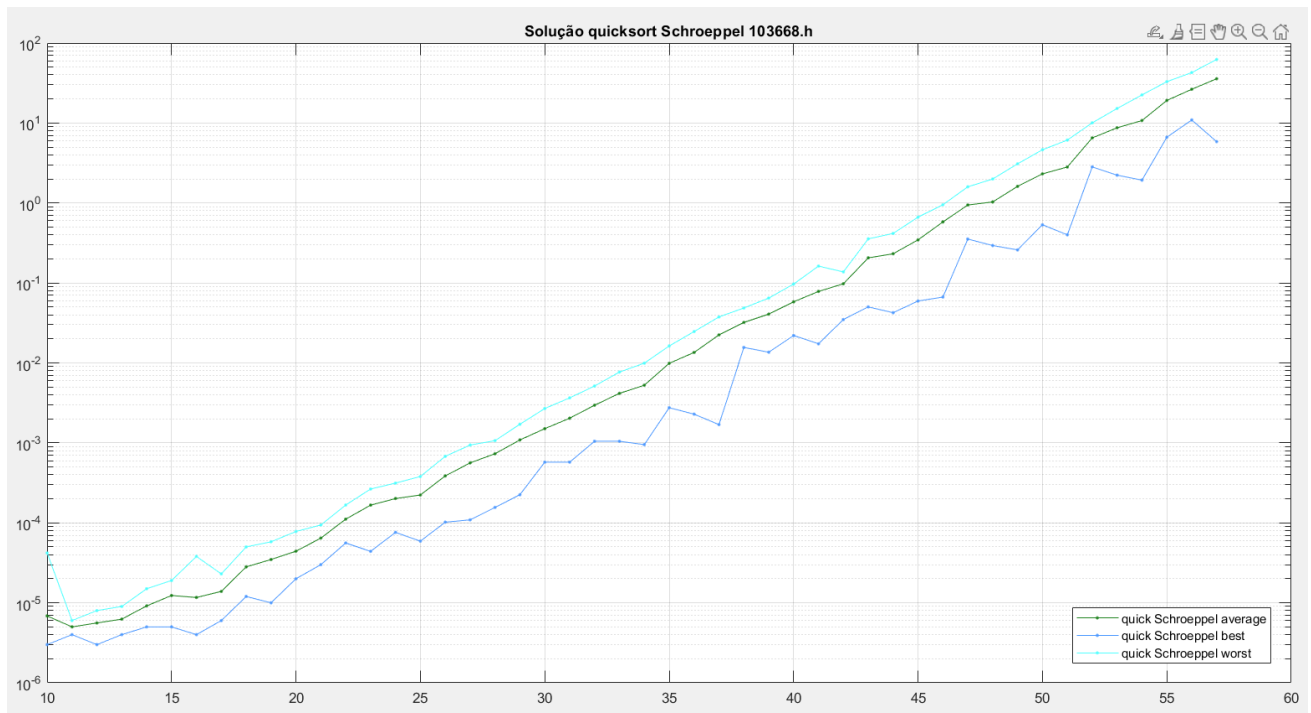


Figura 13 - Schroeppel-Shamir QuickSort 103668.h

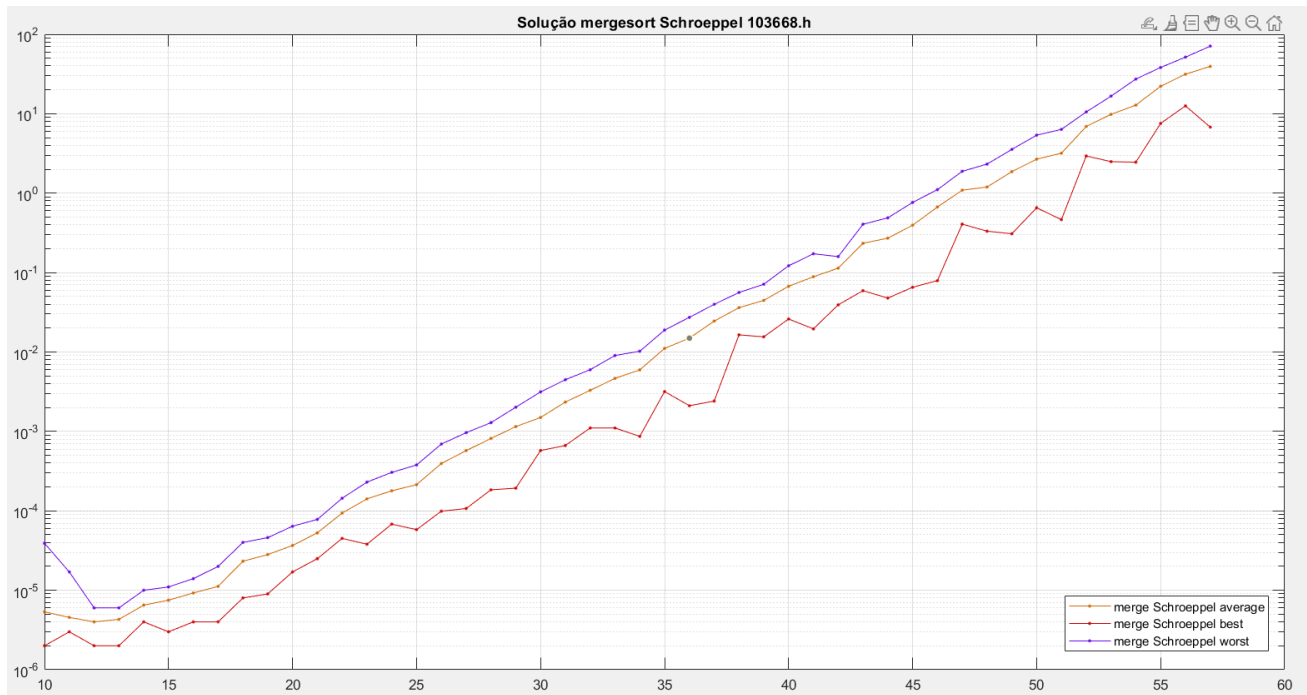


Figura 14 - Schroeppel-Shamir MergeSort 103668.h

# Comentários/Interpretações

---

O primeiro algoritmo desenvolvido foi o “Iterative Brute Force”, apesar da sua complexidade computacional ser a pior das soluções propostas, este foi extremamente importante na resolução do problema, pois serviu de algoritmo “piloto” que desencadeou a resolução do problema de outras maneiras. Para além disso, foi interessante identificar que em alguns momentos, esta abordagem, apesar de trivial, conseguia encontrar a solução em tempos extremamente inferiores aos restantes algoritmos.

Relativamente às abordagens recursivas, “Recursive Brute Force” e “Clever Recursive Brute Force”, apesar de serem teoricamente diferentes, apresentaram resultados extremamente semelhantes. Isto deve-se ao facto, da otimização que o “Clever Recursive Brute Force” apresenta reduz a quantidade de cálculos quando não existe uma solução ao problema, o que nunca aconteceu neste trabalho, pois todos os problemas presentes nos ficheiros “103530.h” e “103668.h” tinham solução. Observámos, com maior detalhe, que o próprio “Clever Recursive Brute Force” era ligeiramente “pior” (pode ser observado nas figuras 15 e 16), devido aos dois parâmetros extra na chamada recursiva.

As abordagens mais eficientes são claramente “Horowitz-Sahni” e “Schroeppel-Shamir” tendo chegado ao final de ambos os ficheiros-problema (até  $n=57$ ). Nestas abordagens tivemos também a curiosidade de estudar o melhor algoritmo de ordenação para o efeito, tendo utilizado o algoritmo de “QuickSort” e de “MergeSort”, ambos têm complexidades computacionais semelhantes, diferindo somente no pior caso, onde o “QuickSort” apresenta uma complexidade de  $n^2$ , enquanto o “MergeSort” apresenta  $n \log(n)$ , no entanto, o pior caso do “QuickSort”, no nosso contexto, nunca acontece. Dito isto, foi observado tanto na comparação numérica como gráfica (figuras 17, 18, 19 e 20) que o algoritmo de ordenação “QuickSort” aparentava ser mais eficiente do que o algoritmo de ordenação “MergeSort”, portanto nos gráficos de solução final colocámos só a versão do “QuickSort” das abordagens (figuras 21 e 22).

A abordagem que se destacou foi a de "Schroeppel-Shamir", tendo alcançado  $n=57$  de forma mais eficaz em comparação à abordagem "Horowitz-Sahni", no entanto ambas apresentam a mesma complexidade computacional.

É através da abordagem "divide-and-conquer" que este problema diminui de complexidade computacional. Se compararmos as abordagens dos algoritmos recursivos com os últimos dois algoritmos descritos, a complexidade do cálculo de todas as possibilidades das somas parciais (que corresponde a um passo nos algoritmos que requer muitos recursos) vai diminuindo à medida que dividimos cada vez mais o array "p" (para as recursivas,  $O(2^n)$ , para "Horowitz-Sahni",  $O(2^{\frac{n}{2}+1})$  e para "Schroeppel-Shamir",  $O(2^{\frac{n}{4}+2})$ ), notando-se um padrão de  $O(2^{\frac{n}{x}+\frac{x}{2}})$  de complexidade, sendo "x" o número de partes aproximadamente iguais em que o array "p" é dividido.

Era expectável que a complexidade computacional do algoritmo de "Schroeppel-Shamir" fosse menor e não igual à do "Horowitz-Sahni". Na verdade, o algoritmo de "Schroeppel-Shamir" apresenta mais eficiência (figura 21 e 22), mas isto deve-se ao facto dos recursos de memória não serem tão requeridos pelo algoritmo em si, o que poderá por sua vez influenciar a performance da CPU.

## Gráficos Finais

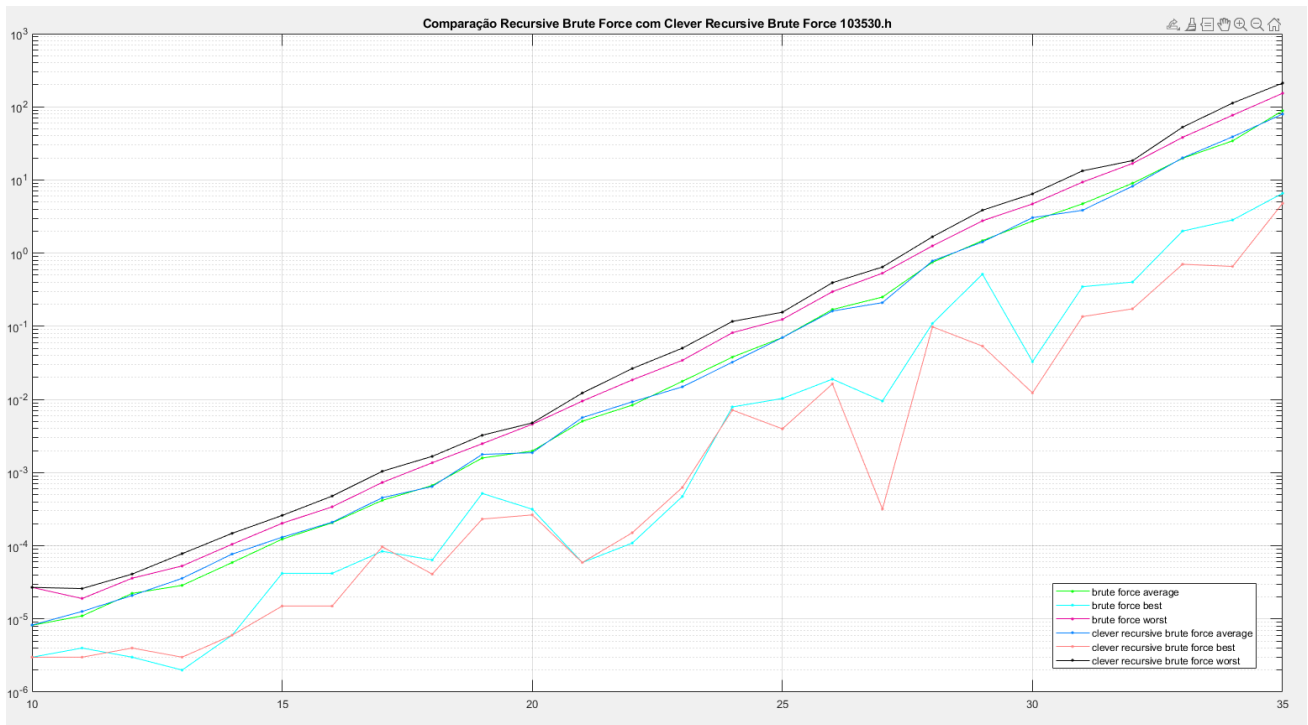


Figura 15 - Estudo Recursive Brute Force e Clever Recursive Brute Force 103530.h

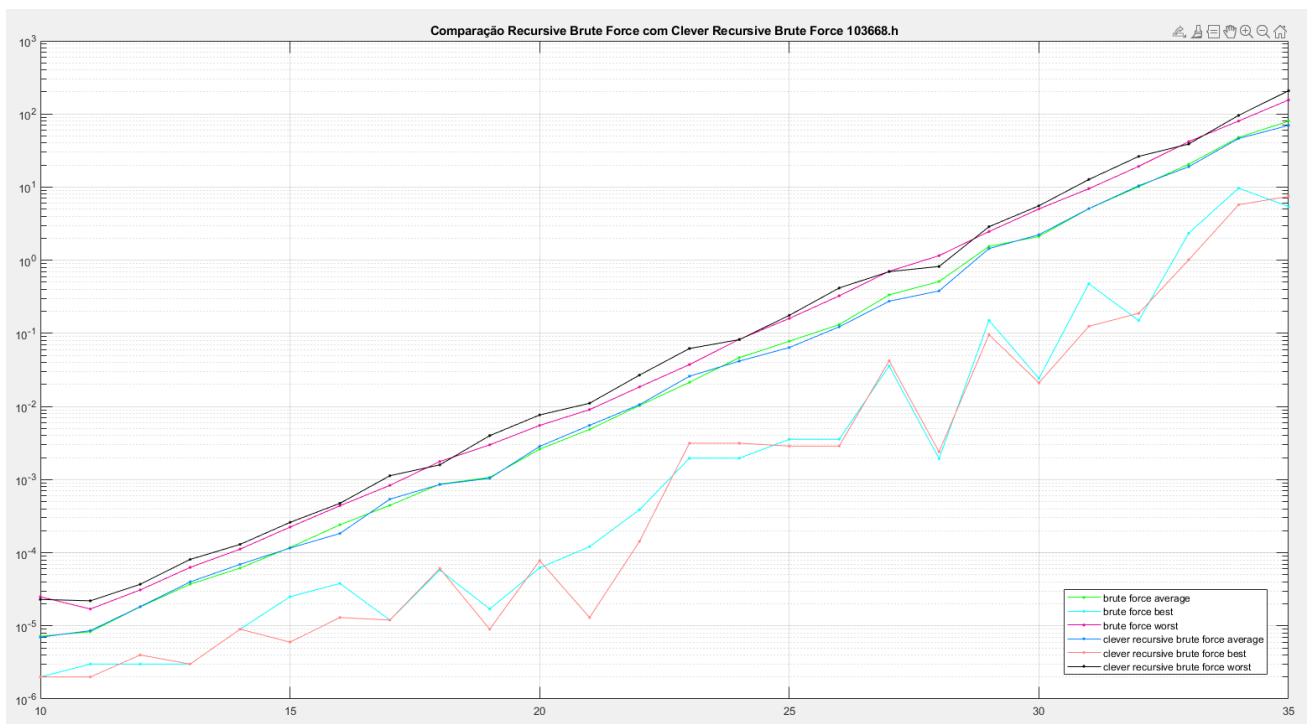


Figura 16 - Estudo Recursive Brute Force e Clever Recursive Brute Force 103668.h



## Gráficos Finais

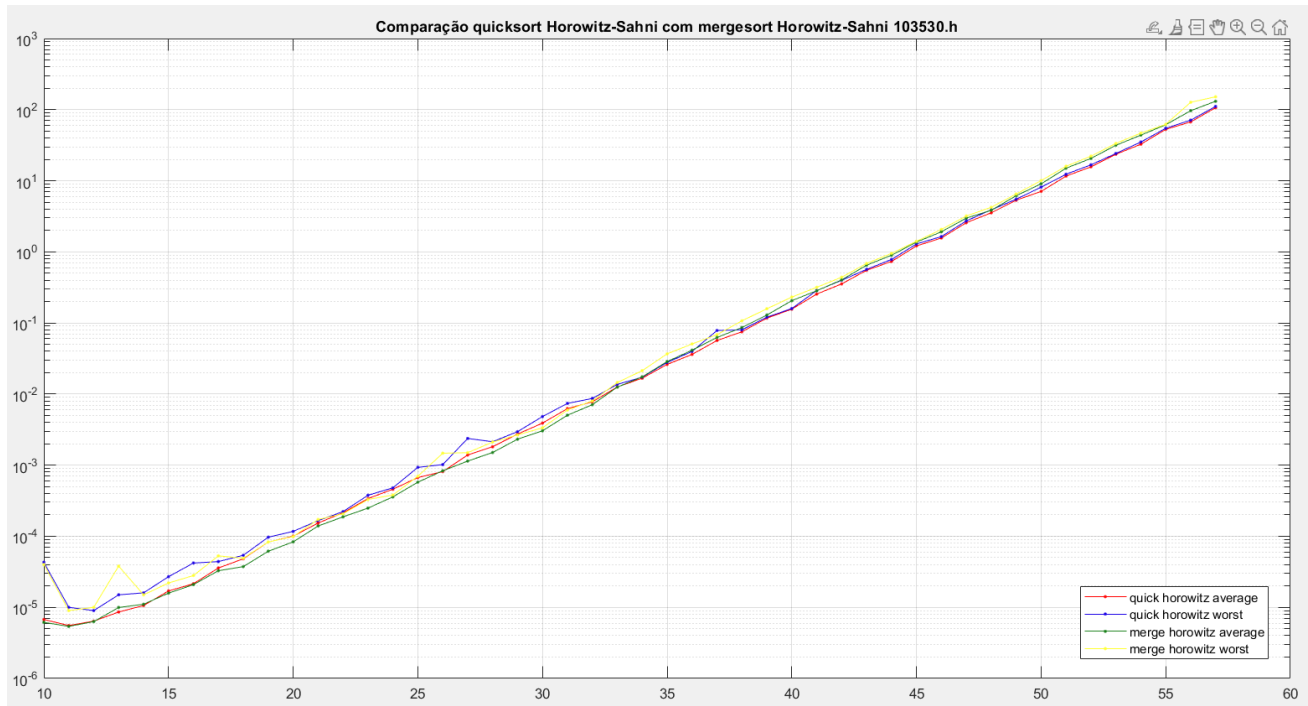


Figura 17 - Estudo MergeSort/QuickSort em Horowitz-Sahni 103530.h

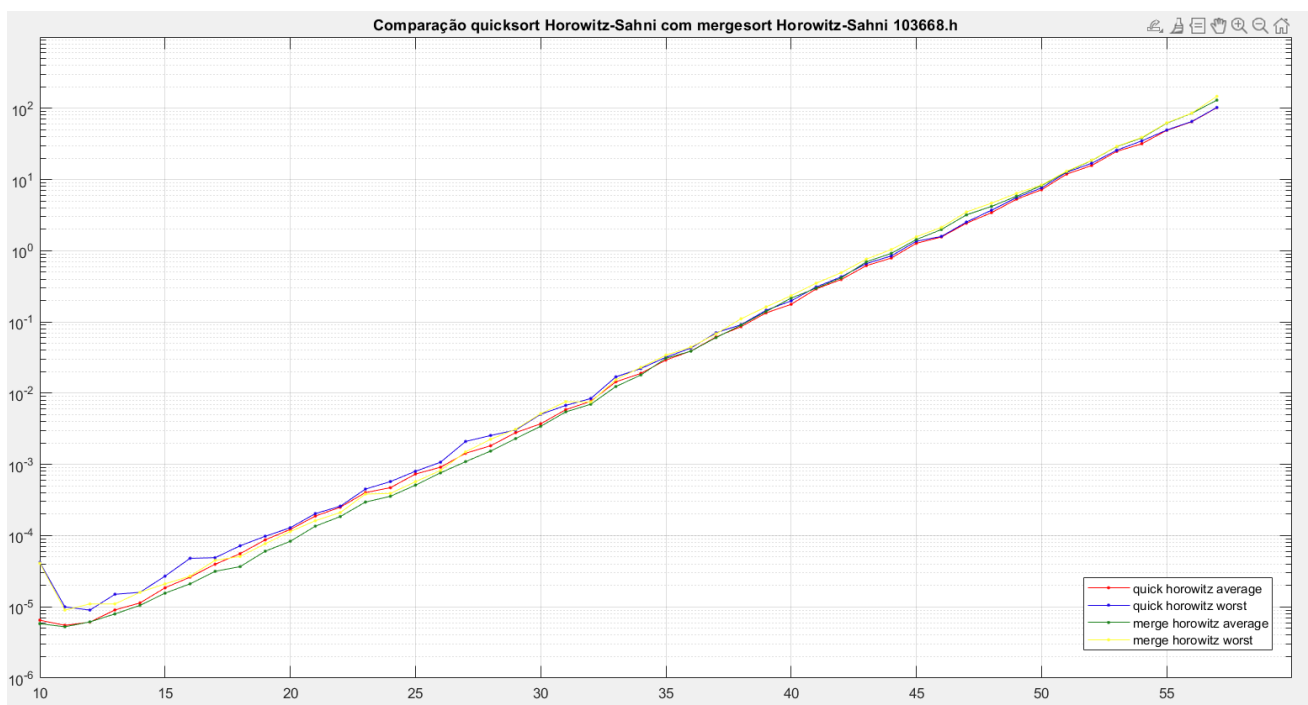


Figura 18 - Estudo MergeSort/QuickSort em Horowitz-Sahni 103668.h

## Gráficos Finais

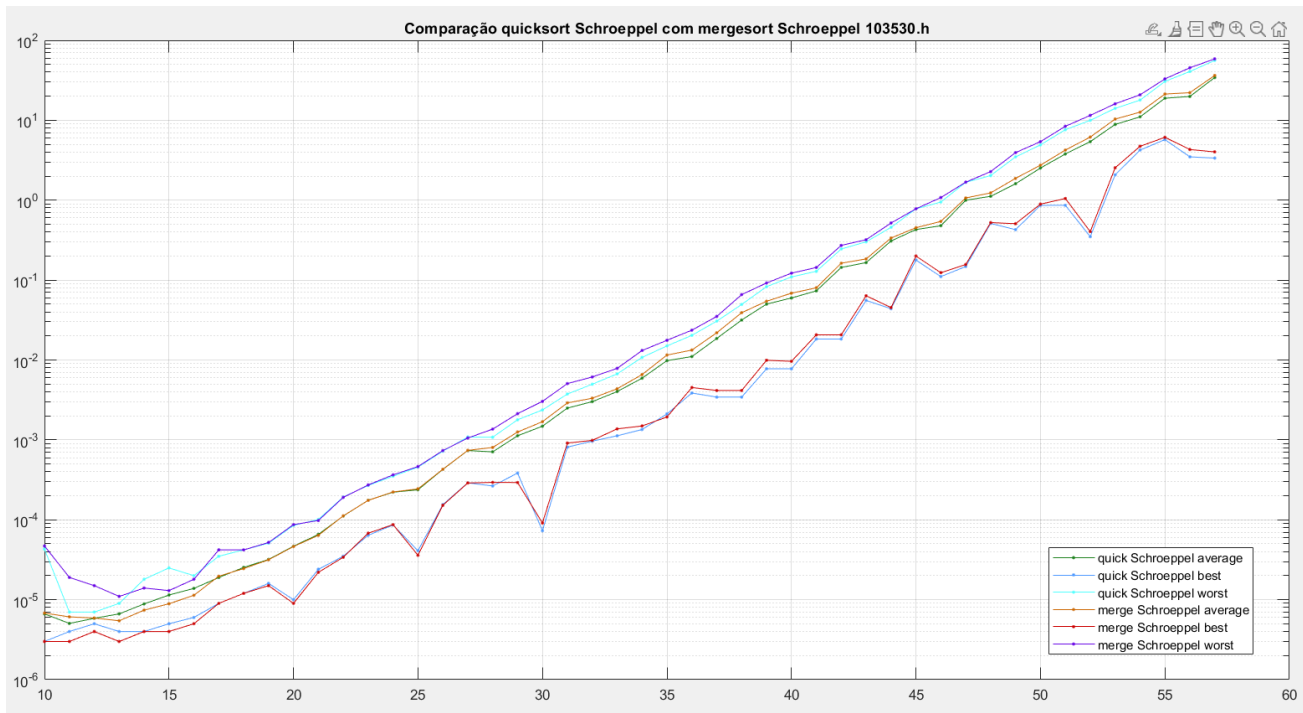


Figura 19 - Estudo MergeSort/QuickSort em Schroeppe-Shamir 103530.h

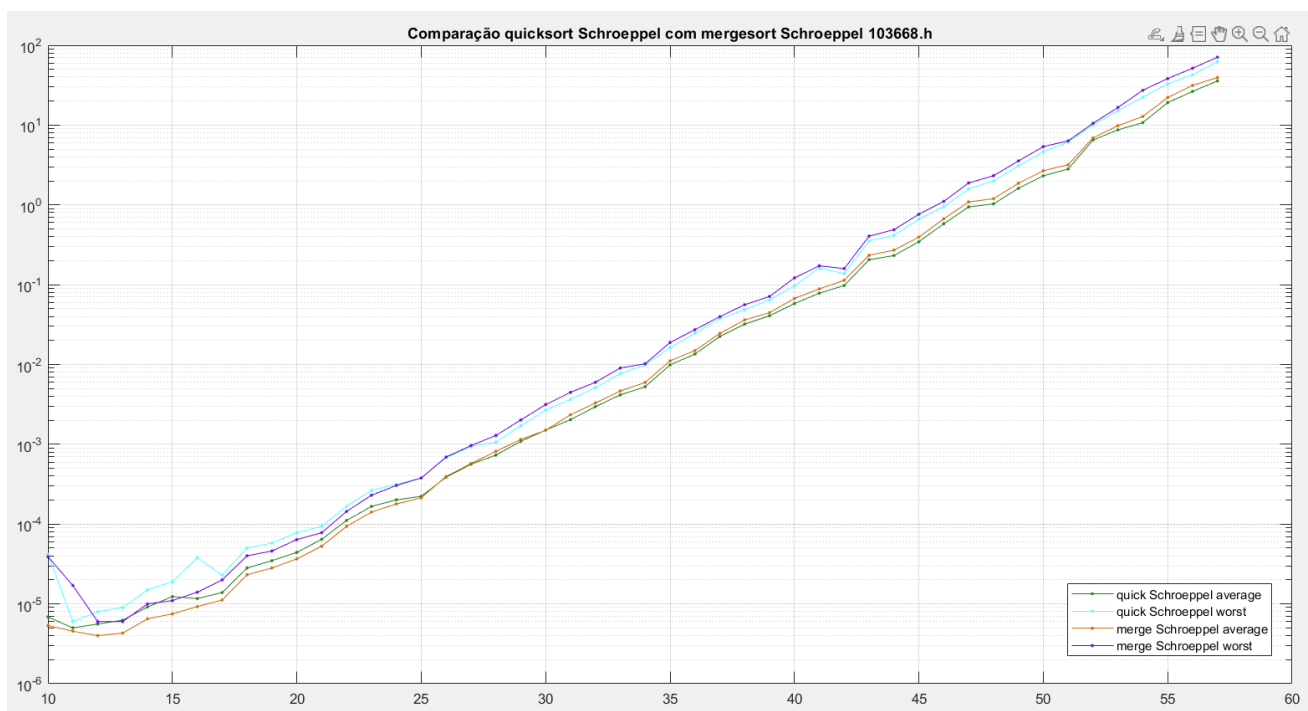


Figura 20 - Estudo MergeSort/QuickSort em Schroeppe-Shamir 103668.h

## Gráficos Finais

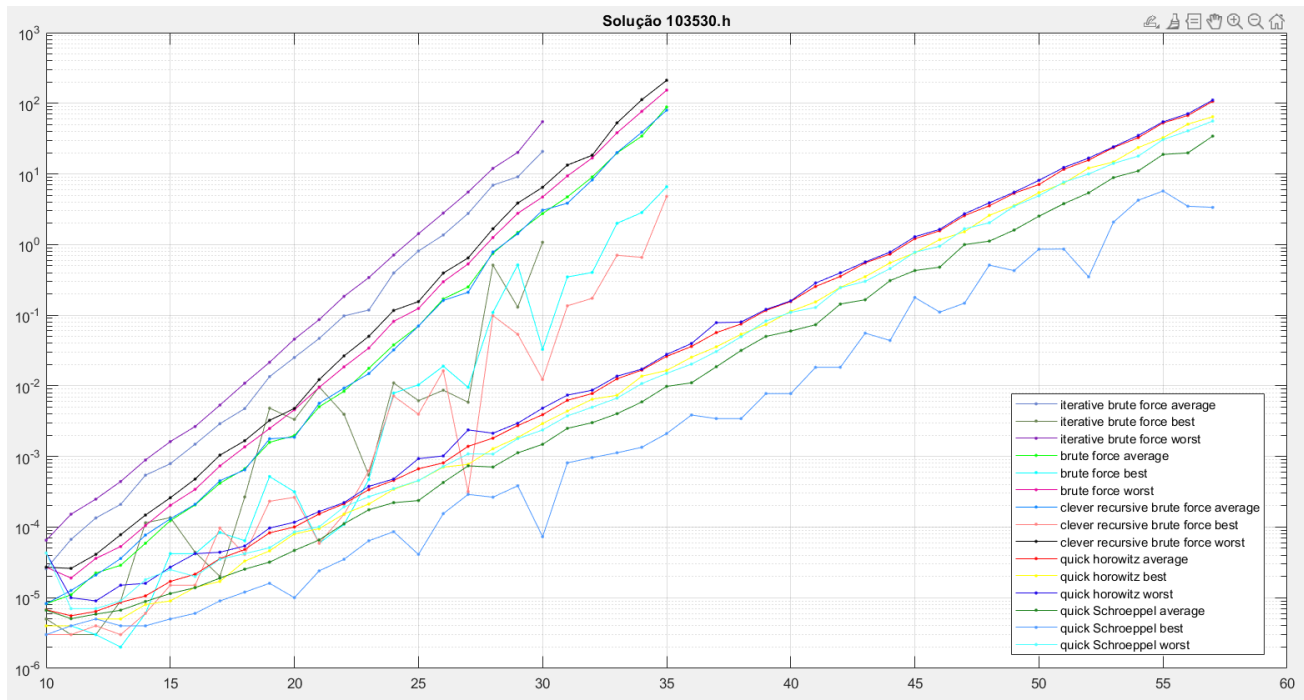


Figura 21 - Solução 103530.h

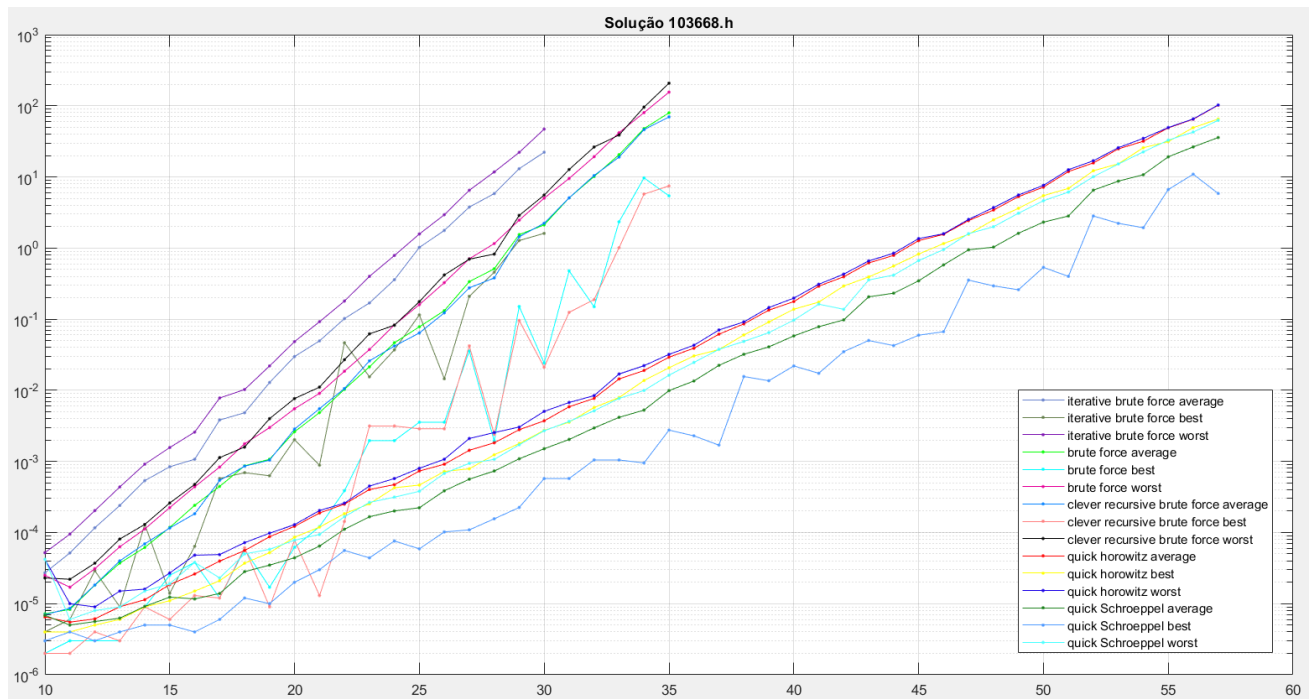


Figura 22 - Solução 103668.h

# Soluções obtidas

---

## Solução 103530.h

n=10

0110100001  
1100010000  
1110110100  
0100000110  
1110110000  
1010000010  
0101100001  
1101110100  
0110001101  
1010001010  
1011100000  
0010010111  
0110111110  
1010111100  
0001100010  
1110011011  
1000001110  
0110101000  
1101010101  
1000000011

n=11

01000100101  
11111111011  
00010000011  
00101101111  
11011000110  
00011110100  
10000100000  
11101100100  
10001111011  
00101011001  
11110111101  
11101011100  
11001100110  
01111101010  
11000001110  
01000011110  
00111111111  
00100000111  
11000001111  
10001110110

n=12

110100010001  
101110000110  
010110001011  
011110011101  
001010111101  
101100011101  
111000100000  
110101000001  
011101000100

## The Merkle-Hellman Cryptosystem

```
110111110010
011001011100
001011101011
000011000001
011101101111
100111001111
110010000000
011001001001
111100110010
111100000111
110111101110
```

```
n=13
1000110110000
1010101001000
1000111000000
0000100001000
1100110011111
1011011100100
0110111011111
1101000001100
0111100000010
0101100101000
0011011110000
1010101100110
1000001100010
0100000010100
1100101010100
0001110000100
0101111010100
0001000001000
1111010001000
1100010111100
```

```
n=14
11001110110001
01001010001100
1011111111001
01110000111001
11010000010010
10111111100110
00111011100101
01000110011010
01101010001111
11010100001010
00110011111000
00111111110100
10100001101101
10101100001101
10000001000111
10110001011000
00001010000110
01010001001001
11010000000101
11110100011111
```

```
n=15
010010101001111
110011111011110
110110101100111
100000111011000
011100010101111
100111011011010
010011000101010
101100000001111
011010111010100
010011100010110
110000001101000
010110010101010
111100001010010
110011010010101
00101111001100
110100101011010
11100001110101
100110010101010
100001101011101
001100000000101
```

## The Merkle-Hellman Cryptosystem

n=16

```
1101011011010011
0110000100001101
0100110101011101
0110110010100001
0100001011001000
1110000100011001
0010001110100101
1001001010100101
1001001001111100
0011100110111101
010111111100110
101110111110101
1011100000011000
101000101111110
1010010110011010
1100011101010001
1101000011101011
1001110101100000
0011110100100000
1100110100111010
```

n=17

```
11100110110100010
11111110010000101
01010110011100110
10001001100100010
11010100001100010
00100111001000110
11100011100000000
00101111101011010
10000100000001010
11111110001010010
11111111100000000
01101010000000101
11010001100110100
10011111100000010
00011100101101100
01110010011100000
00110101111101010
00111101010000100
10110011011110010
01101100010001010
```

n=18

```
111011010101010000
100111001110100000
110001110111001000
011000110110010100
010100001001101100
000010001011101000
111010010100100001
010000110001000001
100100001100101000
011010100000110010
000110011100001000
100100100011001000
011110100010000000
101110100100101100
000011100011011010
011110110000110110
010000100011110010
111011101111000000
100110101010011111
011011100101010011
```

n=19

```
0111100111011100001
1100010101100110110
1110010001111000001
1010011110100101001
1111011110001011110
0100000011101010010
1100011010110010001
1001110000110000111
1001111000011100010
```

## The Merkle-Hellman Cryptosystem

```
1100010111001000001
1010011001101011011
0101000000100000011
1110001101011101001
1100010010111010101
1001111000011110111
1100101010010001010
0100100001010101100
1000111100001011111
0101011100101110111
1101010111101011011
```

n=20

```
10000000000111100111
10010000100110000011
00100111010100011011
10010110100101110110
0111111001101010010
00010001000111100010
10100010000001111111
01001001101000011010
01011000010010101000
11110011111011001000
00011001101100011110
01010001001010011001
0100111111100001101
10010000001100100101
01000010111000111110
01010111000010011000
10001110110100011010
00010111111001101011
00110010111100101111
01001001110101100111
```

n=21

```
000110100001000101100
001101110000000101110
011110011101110011010
111110110110100011000
10001011111011110010
11000111011111010100
011110110000010001110
010101100101101110000
011100011001000000110
001110111000101001100
111110000011100010110
010010111101111011110
111101110001110011100
001001010010001000110
110000011011010110000
101101110011001010010
101000100110111101010
01011010111111111100
000000011000111011000
111101010000111010100
```

n=22

```
0001001000011110000010
1011011110011100011111
1101000000010110100000
0110010001000011101110
1000000011100110111111
1110111001101101111010
0000101101111011001001
0001000011000001101010
1001101110101101000000
0001100000000101011000
1010101000000000100101
1101101101101011100001
0111010110100011001001
0110110001111101000100
0110111111011110110010
0010011011100100001000
0000000101110101111100
0100001011100001011000
0111011011100100110010
0101111001011001110001
```

## The Merkle-Hellman Cryptosystem

n=23

```
00010001110000011011000
11011001010010101110100
00110000111101110100101
10001110010101011101010
01100011011100000100100
00101011111101011100110
10000001010000101000010
10010111000010010001010
01111101111010110000010
10000010111001000000000
1010000111111101001101
00101100001000100101000
00000010101111111111001
10001010011101011110111
10001101100101000100000
11010010101011010001100
11101100110010111111001
01001110110010100111100
01011001100100000110110
10001010011010100100001
```

n=24

```
001010001110000000110101
001101000100110001011011
000110011111001101001011
110000010111001001000010
011101001000100100001101
000110010010001111001100
010110001110010000010110
100110110011010001000001
11110111110101011110111
110110100111000010010010
011001010010000010010110
001101001100001110101001
100000100011010011010001
110011000111101001011011
001101011110110000100101
011110010100111010011000
010111110010011000001101
011100010000110010000101
010001000110101111000000
111111001110011011010111
```

n=25

```
1010100100011100100010000
0010000010001010110110110
0111110001011011100010000
1010000010010011001101110
000011111101100111110000
100011010111101110001100
0101010000100110001111100
0011111100110111001010110
0100001010101010110010100
0110011000100100011011000
1011100011101000101100010
0111010010000011110001110
0111001010000101100100110
1010011100010100010000000
0001011011100000101000100
0111101001000000011110010
1011011011010111111011010
0010111111001110011110110
010111111110101110100000
1001000011100001101010000
```

n=26

```
11101001111111000010000100
11001010011011110100000000
10100010001110110000000010
11110001000000010101010010
00111111010111110011100000
01111000010101010001100101
00110000001011001001100110
10101000011001000001100000
01101101011110100101010000
```



## The Merkle-Hellman Cryptosystem

```
10101010111000110110001010
01101101101011011100100100
00010000001110001010001101
01110011101000101010010100
10010100100110100101100101
101101101010010100010011
01011111100001101001000010
10010100000100010010010010
11111101010111010100000101
10111010000110110100001100
01011100100010110011010110
```

n=27

```
000111010110100011000010000
010001101011111110000000000
001110110111000111010011011
101110110011000001000101100
110100110010101101110110010
000100101011100011100100101
01011001011111111100010000
011011100101000110011011100
011100001100111111011100110
101001110101000001110000100
000100101000100000100000000
01111110101010101000111110100
11011010000000101000001010
100101111010101111011010110
000000111111000110001110011
001010110101010010100010001
100100100000111010110000001
0111011011111111001010011101
101111000101100010100010101
001011101000001100001111100
```

n=28

```
1001011101101000010011111010
1010000001110001000010101101
1011110001011000011010000111
1011010010100000111110000111
1010011100011101100101010000
1111101101110011010000001011
1011010101101010111000001111
0101111011000110111000011110
0001111000101100111010101000
1001101111001100000001011001
1111110010111010001111011100
1101000110110100010100011111
0100011010101001100100001001
0010110101110001010111110100
1101000000011101110110010110
1011111110011010001110001111
0111000011100000010011110010
1100110110001100111111001001
0101110010101011000010111010
0110011011100000110100100100
```

n=29

```
11000110000001001000100010000
01010010100111000011011111110
00111011000010110001011010000
00111110001100111011100111110
01111111000010010111110011110
11111110011111110101000001100
10101101011110000100000100000
10101010110001010110110000010
00110001011010110110100111000
00111111011111010000100011110
01000101010101110110110100000
0100111111100101010011111100
11111011010111010110110001110
01111100110010000010000010100
10000011001101110011000000000
10101101100010000011111001100
10000000101111110101010000000
11100101110010010010101001110
11111001100001000101100000010
11000000010100101111101000100
```

## The Merkle-Hellman Cryptosystem

n=30

```
111011011001001111011100101000
000101001000100011001000100000
110010000100000011011101010010
110101000111110100111100000011
010000100111110101011010101001
101101001100011110011110100100
101110100001110111001100011000
01111001111011100111100100011
100110001110110110111000010000
000000011010000011100000000110
000100001101110001011010001100
010110100000110101100010010001
111001001110011010010000001001
10011100111101010001010110100
110011000100110000011100001101
10010001111000010100110010100
101001111100001011111000110000
110010001110010111101100010101
110011101100011000111101001100
001010111010100000100100001010
```

n=31

```
0000110011001110011110011001110
010101011110110110111110011010
0110011001100011111110011100101
0000111011101111100111011001010
1010111010000100011000000100101
0001010110011100010100000111101
0110110011101111111101010010100
0010001010101110001010000001000
1101001100000101101000001100000
1110100011110010100011000001010
1001010010000101110000110001010
111110100001011111000001110000
0111000000000010000111011100000
1111100000111010010111010100111
0101000101000110101101000001111
101011000111100001110010110010
1000011000100100110000000110100
0000100100111011000101111000000
1001101111110101110001011001111
1100110101100110100110000101100
```

n=32

```
10000011010010111111010110101011
10111110011100011101111101000110
01000101001100111110000011000100
10111001101100110010011011100000
00101001000101111001000110101101
01011011001001110101101100111100
11010001101101110001100100101000
0111010011100000000111100010000
10110100011110111101110000001010
01110100001101100011100101111100
1100100111110000100101101010110
100001111101010101010111101111
11011111010100010000011100010000
00010110100011001101010011001111
01001000000111100010010110110110
01000001000101111101011100110110
00000101011000001011101001000100
01011110101100110101100101110
1001001010110111100101000010100
011001001101100000001101011000
```

n=33

```
111111000100001010010011011010010
00001110100111000011111001010010
110110010100101111001101011111010
001011011101110000101000000000000
110011000011110100010110000000000
100100101010010000101010010010010
100101110001000100011000010101010
101010001101000000111010110101110
101010000110101000110101001111000
```

## The Merkle-Hellman Cryptosystem

```
00110100110011101101010010000001
101011100110100100010000011111010
110100101010000100011000101101000
000011010111000100001010010010010
001110000011001000110000111101110
1001010100111110000101101111000
100101011100011011000011011111000
00110111101110000001010010111100
110111001101110010000011011010100
0011110010001100001011011010110
011000010010100000100011111101100
```

n=34

```
000100011000111111110011100101110
111111011111010101110101000000101
111001001111100000111111011111111
0011100010011111000001010000110110
1111110110000011111110000011000
1011101011111011010101000010010101
00001011011011111101010101100001
1000101000001011011011011011000001
0000100101110001000000001000100110
1001000010111111100101010011100010
10001110000101011111000001001001
1001110111001000110010111011000001
0011100101001110010011101101000101
0010110000011011100010111010000110
01001100100011101001110111010000
1111000110110010011000000100000010
100010101010001000011110001010101
0011100011011001000000000110000010
0010001000011100110001100010110000
0001011000000000101101110101000010
```

n=35

```
00110000010101111111001010110010111
10101100000110100110000111110100010
11000010110010001000000000100011010
00010110001010101001011000001010100
11000011001011000101010000010010101
10111010000001010100110010100011010
0111100100110011010000101011110001
00010110110100100000011110000010101
11000000101000101101010000110110111
11111001011010011100001100100100111
00111010111100110101110010101000010
00100101101011100010111110101011000
1101011000111100100111110111100011
111111101111010111010111010110110
011010111101100001111101001111011
10000111000100101010010001000100101
10010110000011011000101000010101001
0111000011111100111000110111010110
110010110100111110011101111000010
11111111100001001111110100110000001
```

n=36

```
001101110000101001101011001100101101
101111000111000110000100010101101001
110101001001011000110011011001110010
000011101011011000010100111111011001
01110011000111011110101110100001001
001001101011100110100111101110001101
00010001010011110010100101000011110
101101100101000101100110100100010111
001000010111001010011110100000100011
011000000111011100110110111111011111
000010010110001101111011101110000011
100111101101000001010011100010001010
100010111000110000100000111110001111
0000110010110101100101011011001101
111010111101010010000100111000110001
110000101111101000010010011101011111
100101101101000101011101011100111011
110001001001010100011000110101110001
110101010101111010001100111100011000
10100111011001101110001010101111100
```

## The Merkle-Hellman Cryptosystem

n=37

```
000101101000000111101010010011110000
1110111110100110001000100110010010010
000011100011111110110010101011110110
1001101111001001001111011011110111100
000111000110010110101110011000110010
0111110110110111101010010001000011110
0000101000101011001101000010001111100
001100100000011011011010001110011000
1001111111001001000010101001001100000
1011110000101111010010000100101010010
0010000101110011110111110101111111100
100001010101110101011101000010001110
0101001101111011011011010010011100110
000111011011100010000000111100101010
111001101111000001010011011101110010
11110010011100001110011010000100010
0101011011001011101100111010011100100
1111011001010110001100011010010001110
1000010000100101100101011011001100100
100100101110000010111111001001011000
```

n=38

```
01000000100100111110110100110100001101
0101000101101100111100001010010110001
00111011111011110010011010001000110100
1100001010001001111100011110101000100
10111011111101001011010111000011000
0010111000101011110000101001000001100
01101011000000101100100101010001101000
01110100101101000010001101000010100011
01111001010011110100100011101001010100
11100100011001110110101000011010000001
11010001111110000111011100111111000100
00001110110111000001101110010001001011
10011111000110110010101111000100110000
10111101110101010011110100000000110100
01100110100111100001110101011000010010
1111001011101111000001010001010110001
11011101010110101001011011101000100011
00101001110111010010111100001010100110
11001111101010101101111011100010011000
01101010000110011001000000101000000110
```

n=39

```
100000000111011011001011111001000100100
001111100110000001001100011001011110000
110010110000101111101110110001011000011
111110010110011010101001000011111001000
000011001111111011111001101000100101001
110000001101110110001100111001101110101
100000010000001001011011110110000110000
001000000001011100000010111001110100001
101110000010011111101111011101000100010
10000011101011110111110010100101011001
00101001001110111111111100100000000100
0000010010001000101001001000101101111
001100100100110010010100100010101101110
0011111011110010001010100110110001011
011101010100101111010110011100001000101
010111010101110001101010100011110010111
111110111110100001000000000111001101001
101101111010000000000011110001000010001
000101001110010000010001100010000010111
010001001110010001011011101111010011101
```

n=40

```
01010011000000010111011111001111110110
011000000000001011011111110010010111111
0000010101000100111100001000110010100010
1110010000111010101110100011001001110001
1110110000101100110110000010001101101010
001000110000100010110111011101101000011
01001010111101101010110111100001101100
11100111011101011001000100100100100001
1100001101010111001110110010000111101011
```

## The Merkle-Hellman Cryptosystem

```
0010100101011101100000111100010101010
011000111100000010111011101001111010
1101010111110011111010010001110000100
0001100011010101110000101010101100110
01011100010100100101011001001010111100
11010111000100101010100001011101101
01100101111000010101011101101001010010
0100101001111010001110001110000010110
000010111101010101000010001111000010001
100110011000110000000101001010001000001
101000100001011100011100101000011000111
```

n=41

```
1000110110111110100101010101010101000
011011110000010110001000010010010100101
11111001011100011001110101010100100010
1001101110000011000111111111001110000100
11011001010001000101000101000000000000
111101010100101000110110011000000100000
000000110011010100010101010011101100100
01000101010000001011111111101110110000
1000000110001001111000110001000000001000
10111000010100000001001010001101000000100
1101000100110010000010011110100000010001
000100010101010101101010101010000101101
10001000001100001101111000100101000110
0101101011101111010000001100111000110001
11010100010010100011000100001100101010001
100010011100000011110010000101010101000
010101110100110010111010101111000000110
011001111010100000010101000111101110110
0101010100011011000001010100010000001110
01111011000101000000100101000101101110110
```

n=42

```
011111110101001110011100111011011010000
10010100001011111010100100110101000111010
01001000110001011111101111100111110000000
011100100010111111001010101010000110100
01010011101111001101010001000000100001000
000101000111011101100001000000111000011001
1010100111101111000011011010111111000000
001010101001110111101110100100010001100010
00111100000011010101001110110011000101100
11011010001110100111110100101000000101000
00110111011100000101110101001101010101010
011100100111000110001101101001100100000111
110110110110011000001000011011101000110110
000111010010111010001000010011000010000010
01110110100001000000000110000000000101000
01100001010111100001111001000010101000110
101000110011000001010010011011110000100100
011010101110101100110000110110010010000111
0111011000001110101100010100101010001000
001101010001111101100100000010000111010010
```

n=43

```
011100001010101000100110110011101001010011
01100000000000011100000100111101100110000011
00100110101000111001110001010101010100001
11011010000000010101010010010100010011000
010100011010011010100001110000111011011101
1010011111001001011111010001110011011110000
0000000100011110110111000100011110000101011
100100111100011111110110100111001110111010
10010010111000010100110111001010101010001
1101011001100100101001011101111011010001110
1110111110000011011110101111001000011001011
0101101101001100001100010011011111011001111
0100110110001110010011010001010001100111001
010110010101001010100101110111111101101011
1001101111011101000101111011010011101010011
101000111000001111101000100010101011011000
100101010001110100111001111111110101010110
1000100100001100011101101111010100011010111
0100111010100100010000100000001011000100011
1111001101110001001011100100101010010010101
```

## The Merkle-Hellman Cryptosystem

n=44

```
00110010000110000101111010111100011101100010
0101001010110100100101010000100100111111011
0111101000110101011100010101001101110010111
01100101011010010000111001010011001110110101
101100010110100101011110111010111100111110
10100000001001010110111110000100111010101
10000010000110001110110010111000100100110
1100000100011010101011101101010001000010001
10011001000010010011000101001111001111011111
00010011001010100010010000111000010001100100
0110111111110111011110101001001010011100001
0111111111101010101010101010010111000110011
1101010100111111001010001110001110100000110
0111011010001110111000010111011001001101010
10110101000110011010101010010000110001000101
1001001111111110010110010001111001010100010
11001110111110111010100111010000001011010011
0110010011110101011111101010011100001001010
01010110110110001000100001000111001010000110
01110110111001000111100001000000110101000001
```

n=45

```
110101101111011100111000010000100111110010000
101000010011001110100111110000000100111100100
110110000111100100100110111010101110000010110
100001001000011100101011010100100010001001110
011111100100000111001000001110101101011001100
000111011010111000011110001110001000001101110
10001001011101011100010000001100111110010000
100110100110101000110111101011000111101111000
000010011000011001011001101011010100001011100
110000101010101001100000110100011100110101010
010110001011001001010111100000001000100111000
100000011100011001011111100110100100101110010
001011100100000100000001110011001011000000000
1001100011100011110111100000000101101111100
101010000101100011001100101001010010101110110
101010101100000100000110011111100000110100000
011000011000100101101010111011100011110000000
000101001011100010100101010111000001110000010
101100001000101110011100101101101101011110100
000010110001101000100100101110010111100101000
```

n=46

```
0110010010100101001100011000010010111100100001
1010101000111000011010111010011010001100001100
010101001111110110001001111011101100110001010
0110100001100101011100001000000011110001100000
100100010011100010000101100100100000000101011
0100110010100110000011011010001010000110010101
1111111101001010000010001101100011010001111100
1001100110011100101010000011001111101001010011
1100111110111000111100010100001111001010101010
1010010110000101010010110011110001101100000110
0001011001000110011001011000111011000100011100
000010010011101011110010110100010010001011000
0010110000011111100101000111001000010000100100
1101111011001100010101000011010011111001101100
111000100110100101110111100111100110010011000
01010000100010110010010101010111101000110010
100101100101001110000001011010000110011100010
100001001011001110111101110010000101110101000
11010001000010001110001110001100110100101100
1101001010000100010100000111110110000011010010
```

n=47

```
1001010100110000110011100100000111111111100000
000010011000110000001100011110001001001001010
10011000001000100000000001011100001110011100101
01101100000101101001001001000101101101001111110
11011011100111110011110011010010110000101001011
00000000110000001100010101000011011001000011100
11111110010101000100110110000000100000011001110
1000010111000011010111100110010111001101110110
10100100111010010101000010110001111001001101
```

## The Merkle-Hellman Cryptosystem

```
11111101001111110001000111000010001010010100101
10100110001001010010010101010100010100111000110
1001000101111001000111011101110001010101010000
0100001011101010110011100010001011111000000000
10011110101011110101000010111100110011110100000
000001011110001110100011101011110100001010001
1011110001101011001100010001011100010011010100
1000111000000101100100000000000101110000010110
001100001010101110010011100001111010101000110
01111010010011101100010010111100010100010011001
00100011010001111101001100010100001011101100100
```

n=48

```
010010010101010111000111001100010000000100000001
00011101010111100110101000010111011111101010001
0101001110000001001111010101010100110001011110
00000111000100001010000000010010100001101010101
110100100101000000010010000001110000101111101
110001010000010101101000001100110100111001001000
0111110100100111101011000010011001101000000000
010000110111101010001111001110111011110111000001
010100011001011000011111001010011110110011110010
11011010101111010110000010100011111010011001111
100111011000101110000001010100110000000000100101
011111011101010000011001111111011000011101000110
11100101010011000110111111100000101000011101100
101100111101100011010111101111011110110100001001
10010101010111100110100110111001010110010111011
00001001100011100010111110011111011100110010110
000100000000000001110111010011100010101110100011
000101001101001001100100001010010100101001000111
11111111111001111010100111101101110010111101
111101011001110101101111010000101100001101011111
```

n=49

```
0100000110110011010000011110010101111010100100110
1011011010011111101110110000000010010000101001100
0000100111011000100101010111010111011100100000000
1001010111101110100010101010101011100110101110100
0000101001101001010000011110011001011001001110000
000101010101010001111011100100001110011101010010
01100111111101011001101001010001101111111011010
1110001011111011111111111100110000111101001111110
0000000101100010100011100011111101100001001111100
0001000111100001111010101000001000111000011111110
1001100000111000011001001010000001001110010111100
1000010001100001000000101000011001110001110000010
101111010111011101110100000110000101000110111100
0000101001111000100010011101000000000100010100010
011100000100010100101010111110001110111101000100
011001101001101000100111001010011101111111101110
00111111000000001111000000000011101101101001000
1100111000101111010100000101001100110111010110100
000010000011000001011110000111010111010011111010
0111000110100011111110110000101100000111000100100
```

n=50

```
00101101101010011100100111010011101000010001101011
00000000100111001000101110110011101111000111010110
00110110011100101001100010011110011110010000110000
10111101110010111010101010111100000001000100101001
0011100000001110001000010111010000011010011101000
11111110010001111111001100101001110100100110100101
0100110010100000101010110010100010000000000001010
11100100100111011110000100011010111010000010001101
00000110011010101011111001001110100110001000010110
101010011111111101100100000010011111110000000100
101010101010110000011110100001110110011100000101
1000100001110101011000001111000010111011000110110
11011001010011100101000101000010000011000101010011
01000110010000101110001100110011101001010000100000
111001001100101010100000101001110011001110101000
100111100011100111110011110001100100000000101000001
011111011011001010001000000101001101000110010110
111111110011010010010010010011101110110011100100010
00101111010101010001010001100000100110101100000101
0100111110110010100110111111011101101110011000110
```

## The Merkle-Hellman Cryptosystem

n=51

```
011000101110011100000011101100101000110101100111010
101010010011100011000101011001110011100110101001010
10000000101001100100000110111101110010111110010011
0010111010001110011010010100000100100010001100110
100011001010101000011000101010010110001110001010010
010010011100100111001100110010011011011001111010000
00101011111111110001110101000011001001000010101100
00010011000001010101011111000001101110001110111
11100100000000000100111000001000011111011011000110
01010011101001111010100010011111010011001011011011
10110001110111001100001001011111011001101111000101
0011110000101000101101111001010101100000101101100
00011011011111101001000010000111110000100100001001
01011011000111111101000110100010001100110101011110
111001110001001110011100010000011101111001110000000
01001010101011101100100000011010101111101101111
100111100001001100000111101110111110110110000101011
011110101010010000001000100110001001100001110001000
1000101100111110101000011011011100100101011101011
01111100101100001010101001100010001101001100000010
```

n=52

```
1100010111110000101010001011101110101101010010010111
1011111001110101101101110100101110001000001100010
0000001001111000010111001111001001111000010010100010
010110100011010111000100110001111000110110001110111
1011100011011001011001101011000010110110010101110010
0010000110100010001000011001111011100100010110100010
1111110010111101010110010010001010100001000001110110
0100010101101001110010111100000110000001011100111
10000100100100001010001010100110111011011111011001
01101011010000000011101010101000011101101101100101
01011000100010001011110100101000000010010101111010
101011000101001110011001001111010100000101000010100
0110010001101001001110011001111010000110011011101110
1011101101010110101001101000011000001000001000011000
1110001111010000110110010111010010100101000100000110
1001001010001010011101001011000011110001000010000011
001011011110010101001111010111001001001010111001011
0001101100111110011001010100011100111100001100011001
1101100010011101010111100011010101011110101101001100
000000111101001100011100011010001111001111110110011
```

n=53

```
01000011010111101111010010001000010100000111110001010
11011001111001101100000101001101001110000010001000110
0100110111001100110000101011010000000011011001010110
0101001100111001010010000110000010101010000011111000
10111011000000101101100101101000000011101001101010000
001111100101010011111001111111000000000000000101100
010001101101010001011100011001011101000111101110100
011001111011010111110111100110110101111011111111000
100110001111001100011011101110000101010110001001001
010010110000101000100111111011001011010000011000100
01000000100111011010101111101101110101011100111010000
0010000010011101110111010100100001000010010000001000
11011101110100000110111001110101100011100110111001100
10111100110100011010010011000010111110011000110100110
01000010101001010010011010111010000011000110100101010
1001101110000010010111110101111000000000011101011100
10001101000100111001101000101001001110101110111100100
10100001010100101000110101110101010011100100011100
1000101000100011110010101110001010011001001110011000
1001011100111001010001111001110010101100110010100000
```

n=54

```
010000111101000100000111010011110001100011110101000001
011000111100011101000101100000111000011010001000001011
001000010110110100010101000101101011100101111100010
010111101101010101100101101111010101011010000001101
00011111110011011101101010000010010110110101010001000
001100010111000101111001100000101100111100010011100100
11011001111101011001001101000111111110010110000001010
0111011001101000010011100101101100010001000110011000
00011111011110100000011111111101011100000100000001011
```



## The Merkle-Hellman Cryptosystem

```
1011011000000011000011011010100111110000001100011011
101100110001111001101011001111011100010011000001101001
100111101101110100110000101110101100110111100110010011
110001110111000100010100100100101100010010000010101010
000011110100100111010110010000111110100010100100110001
0110010101100000001100010100001100111011110000010110
0000111010111110010101101100011110110010111100010100
101010001000001110100010100010001010010001101010001010
1100011000011001110011011010110010101011001000110100
11111111101011111101110001110111010100010101100001
001100001001100001001100101001110011010001011101001010
```

n=55

```
001000101010110001100110100010101000100010110101001
0111000011011000110110010101000011101000010011000010111
00100100000000000001000100110100010101110010001001000010
1000000101101101101101011111011111111010101001000100
111001101100111110110001101010101011110100010000010
11010000000100100111010110011010110110100000000110000
111001000010011011110001001011001000111100101101110101
000000001110001011100010111001111001011110100000110110
000000101111010000000001000001001001111010010101101010
100001001111111001001000110001010011111001001111010000
0010101100001001000010010001000110111001000101110100110
01111100000011110010100010000111110010000100011111110
0011010011010100101101110110100100010001110011011110001
0001111110101001111001101011000110010100010011100001111
111110011111010010100111001100101010010000000101000000
1010010101000101011011000000101110001100100010101100011
101100101111110011000111001011000001001110000001011010
001011010000000011111101010101011111001100110001111010
0000010001101110100010101000111010111010110110001011
110000010100101101110100110100110110010011110100
```

n=56

```
101010010001101101101110110110000001110110000101110000
0010011010011011010100100001000011111010001101110101111
010100100101010100000010001101000101101110111001010111
11001110001110010001010100010001111001100100111010110100
01100110010011110001100110111011001000100010101110100
01001011101001111000101010110111101010110000010110100011
0000001000110100001000010100010011110110011001110101010
010000010011101111100100010010100000010011110110000111
000100101010111101010110010101101010000110011110011110
1011000100001110001011010110001000110000101110111011011
01111101001100011110101011100101000001010101010011011
11010101101010001001000110100110100101000011100010100011
0100010000001001000011010100101000110001101001001100010
001010000110000000110010010110001111111111000111001001
00001111111100101110110000110001000110101111000101101010
1100100100010111101100110001111010110011000101101110001
1001001101101100010001000110011101001001000111001011001
1110010000001101101110101100101011101110111101001110
00101001101110010100010011111011110001111001011011101
11101011110100010000010000000010011100100111110100011101
```

n=57

```
001111110110010110001111101110101110101111001000011111110
110001011110010010011001101100101001011111000111100100100
000000001000001001111001111001011011000110010011110100010
0101100010110100100000001001111111001001011011110010010
1010001001000011111100001110011010001011011110001010010
10000111000100111010111101101010010111101011000010101110
011110110001011100000010100101000000110000000010101001100
00000100111110001011111100101100110110001101110000001110
0111010011010100011011111101000010100110000000000000000
11101001000110111000001100000101100110111111011110001110
01000111000001001010101101001110110001000000001101111000
1011111011001000000100001010111011001111101011001100100
111100110111010101001000011000001000111110000010001000100
00101001100111001110110000101010010101011111111110100
010011111100100010110110101111101110001001010000011110
01110110111111000000010000000000000100101110000001000110
00001010110101010011100100111111000101110100010011100
1001000001101100001110101110011111101101111100011010100
111010100000110001110011101000110101100100100011110001100
0100100001110101100000100001001111000110000011001101110
```

## Solução 103668.h

n=10

```
0100101001
0000000110
0010101010
1110010001
1110011010
0010100000
0000011101
1001111110
0000010100
0100110010
0101110000
0000001011
0101011011
0000000001
1101100000
0000001011
0000011001
0000010011
1100100001
1010110101
```

n=11

```
01101011011
00011110100
10001010010
00110011001
00000110011
01110001011
00000000000
10110110101
00100011101
00000011100
01101001001
11010101100
00100001000
00001010010
01101101011
10110011011
01101010010
00011110000
00110110110
00011011011
```

n=12

```
111011000100
010111100110
001001011001
000010110011
011001101001
001011001010
001000100111
010001100010
100001000110
011110100100
101110001001
101001010101
001011010011
110000110101
111100110010
010101100101
111000011011
110000000100
110101000010
111000100111
```

n=13

```
1111101101110
1000010100010
1110111111001
1100011000011
1011000011011
0101100010011
1011101110101
0101001011101
1101010000000
1001100010000
1100010001110
0010000111110
1001000100000
1110101011110
0011011111001
0000001100111
```

## The Merkle-Hellman Cryptosystem

```
1111010110110
0100001101010
1011010110100
0011111001011
```

n=14

```
01110000010100
11101010001110
00001110001001
01010110110111
10100011001101
10001101110010
01001111011110
10001010110111
10111011000100
11001001111100
01110101000011
00010010001110
01010100011001
01110110110011
10111000100111
10111011001010
10110001100011
00101110111001
00111000111100
11011000000000
```

n=15

```
100010010101010
011001010100001
111000010000100
001000001110001
111101110010001
010100010101011
001011101001010
000110001100111
110110110110011
110010011100001
100101100010010
110011011111001
000000011001010
110000011010101
001001001000000
000000001000100
001001101101101
011110011111101
001101001101110
111011100110001
```

n=16

```
1100111100110110
1000100001000100
0111100010000010
1001000001001000
0110111010100000
1011010000010110
1001110101110100
0100000011110101
0110110000001010
0001110010000110
1110000011000100
1111001100001010
1001111100100101
1001100101010010
0100010001000010
1100001101110110
0011110010110110
0101100001001011
1011110101110000
0001010010000010
```

n=17

```
00011000111111000
10111110000000000
00010100010010001
01010001011110001
11011000101101000
11100100111110111
00001111100100100
10011101011110011
00011001000000000
01011011000100111
10100100111001111
00001011000000000
11010110011101010
10001101011000111
10111110101101100
00000010111011000
11000101000010111
10110110101110100
10010000111101011
11101001011100111
```

## The Merkle-Hellman Cryptosystem

n=18

```
011110111110000111
000000000101001000
010100010011011010
111110100100101000
00000111011110010
000011101000111010
111101010010010101
101111000001011010
101000110100101011
010100011011100100
000100111011111000
110010111100000110
100100111110110001
101001101100110101
001010010001011000
10100111111000110
101010001011011101
011111110011001110
101110110011111001
001001001011110000
```

n=19

```
0000011001000111101
0010110001100001111
0100011010001100010
0000011001010010111
0000000001001001101
0000000101010001000
111110000001100110
0001011110101001101
1111001010001000100
0000101001010110001
1001011101011011000
0001001000110000100
1100001111110111101
0001011011101100000
1000111011110000011
0100011000101101011
0001011110010101101
1001010110010001101
0001010001011011101
1001101010101000111
```

n=20

```
01011011010010100001
00110110101110001111
11111111000011111111
10001100101001010100
11110010001000110100
10101100101010101000
10010100111011010000
00100101000000100110
00100001011101011001
10110111101111010101
11011110111010111111
01100001101001110111
01011001110010010001
10110011100110001011
01011100010000101011
11000101001001100111
00000011001001001101
01110110000101000111
00011000111001110010
1101000111010111011
```

n=21

```
110101000101001011110
000001000011101000000
000101000110101001110
001011000100101001111
011010001001011100010
011001011010010111010
000000110001001000000
110001011101001010011
101000101001011110101
111000011011110010100
100000011101101111001
111001110101010010101
100110111100101010011
010011000001101110010
111011100110001001101
011101110000110110010
111000100010011101010
000110001001010010111
010110111001100101001
011111011010000110111
```

n=22

```
1000101111110010011001
0100110100110110001010
0000110001001001100010
0001100100001100111100
```

## The Merkle-Hellman Cryptosystem

```
1100011101000011000101
111110101010101111100
1010011110010100011001
1011111010001010010001
101111101001100100101
0111010111101110001010
0101110000110111001011
0000010101000100100001
1001001110000011110101
1101111001111000110111
0001011010000101110101
1000111010100110010110
1010010010101001001001
1111000110000111110110
1000100101101000001110
1100000101100011000001
```

n=23

```
00111110011010011010010
10011101001110111101000
00010001011011001100111
10001111001001001000110
10001101110111110101000
00110011001111010011010
00000000000111010111001
10100011111000101011111
00000000000101011000001
10100100010001100111010
10000000100010100011011
11110110101110100101101
1100011111010001111110
01011011000101101110000
10001011101001100101100
10000001100011011110000
10101001100101110110111
11000110100011001010000
00000000000110010110100
00001011111100011110110
```

n=24

```
111110001100010101001001
010001001101010110101110
001010011101001100010100
011110011010110110110111
100001001000011110001100
101111011010111101010110
111100000000110010010110
110111101000110000011100
101011000011000111011011
111101000011000100110010
011000010101001000110000
00101001111101000111111
011001000101011010100101
111000100101110011000001
010111000001010110100001
111000000100000011110011
001000111011100000100110
101111101110100010111100
100111011010010011001100
010101101111001010100100
```

n=25

```
1011111000010111111000110
0100111100110101001101111
1000000101100100000100111
1011001101000001100011110
1101101000000011000010100
1010011010001000010010110
1000110011101100001010111
1111101010000100110001111
0110000110001010101100110
1110111110010110001001010
0001110110001001110111111
0111101100010100011100010
1011010101101111011110011
0010011000101110101101001
0101011011011010000000111
0000100011011100100101101
0000101100101100101010001
0111001010000000000100111
0111000010000100100011010
0000010011100100101000111
```

n=26

```
10111000110100110100100001
00110100100110011101111101
00010001101011110101110101
10100101111001010100110101
10001011101000010010001001
00101101100110000111101100
00101111010000100001011011
10000101010011000101100111
00111101011011001000000001
```

## The Merkle-Hellman Cryptosystem

```
00110100111100011000101110
01101001000001000111001000
11110011101101101011110001
0010111110100101001110010
0001001100001011111110101
0000110001010011110100101
11110100110111001100011001
0100011000001011111100101
00110001000111010010000000
10010111100110111011001011
10100010010111011010000110
```

n=27

```
000011100111000110110001011
011001011010111110100011001
010001100111111100000101001
100100010011000110000011111
100111001010010101111100000
110000110010011101001110111
000011011100001110110011001
010111101101100100101101101
000101001110010111011111111
010100011000111011100000110
001011101101010100111001011
101011100011111000001010110
111101011110101100011001010
110011100011100001010100100
101111101010101010000100011
110100100110011000000000101
011100101011100110010110101
0001010110011110101010110
100000000100111101101101100
101100001011001110000110001
```

n=28

```
1101100101000110000010010110
0011000111100111111110011001
0111110111101001010111101010
0101001111001101000110101011
0000101000110111000010010000
0111001001011101100111101101
0001101001010010111110000111
0000011101000100000001100001
000000000101111111001010010
0111110110001010000100010010
0101010011010101000111111010
001000001011000111111110001
1001010101011010000111100001
1000010001011101110011001100
1010101100011101101100010000
0100011101110010111011100100
0110100110110110100110010000
0110011010110001111010011011
1011000001001100101010001001
0110011100000101010110010001
```

n=29

```
11110011000011101111100101001
01011001101101011100010101011
11100100010111111011000101100
10101101000011010110110011001
10100100011111110001101100110
00001110011011111010111001011
11001110001110001011000101010
10000011110001101110110001001
00111001001011111110110011110
01100100000011001010000110110
11011100101111101110111011001
11100100101111111001001100001
11101110110110010010110100010
01101101111000000001100101001
10100010111001100011011010011
11101110011111001100011100100
10100000100100011101111100011
00110011000010100110000110000
10011110101101100010110000101
01001110000110001110110110100
```

n=30

```
111001111001011011010001110101
0000111101000111100000110100100
001000011101001001010110111011
010000011001001010111011101011
10100111101101001110111101110
01111100110100101111001000110
10010010111010011011110000011
111001000000110111100111101000
01011110000100111011111010011
11001010011011111010011101000
001001011100110011000110111101
001000111100011101001001000011
01011001101001011111110010000
10010111011100100010011110010
```

## The Merkle-Hellman Cryptosystem

```
001110010101100111000011001010
1100010000110101010111001100
01010110110000101110011100000
000000010100110000100010111010
00100101100111100101101101000
0110101000000001110010110100
```

n=31

```
000101001100100010110000011001
100100100101011001110100101111
010011011010110011100111010111
0000110011110010001001000010001
011000101011111010000001010101
1111110110000010010010101100010
100101111010001101111001010000
100001001010110010101010111101
0011100000100011101101011000101
100011010101101110010100110001
0000110001101101100011100101100
1111101000110110100101010110100
101001001001111011111010100000
100110010101100001001011101011
101010100011001001111100101100
0100000100101000000101111011100
101000001111111000001000110001
111110100001010111010010110111
11011010101100111010110110110
1010000010001100110010101001000
```

n=32

```
111111111111100111001110100111
0101011110010111011110001010100
1011011000011101001011010011110
01000001011000000010110110010110
0110110000111001111010000010000
10100011111001011010101110000011
0000000111111100000011010100111
00000100000101000101000001111100
1100111010101111010010100110001
10000101110111000111001101000100
00100011101101011011110100011000
0000100101100000100111000000101
10011100110111010000001011010100
010011101000001101100011101110
11101000100001100100101110110001
1001111001001011101110101110100
11000110110010001000001011000110
11101111001101010100011101110110
1001111001011101010110001111100
1011010110011000110011001111101
```

n=33

```
100010111011101010100001110101001
10001001101111010101011111101001
100110110000011000010010110110110
0011010111010111101010100011111
110011010001111010111000000111101
101001101011001110011100010111001
11111111101111000101010000000111
000100001100110101111000111001011
111010000011001010100100010100111
00011110101100111101010111000110
110111011110001000010011111110010
001101111101101001010001110110101
00011100010001100000000010100111
001001010101001000011000010000110
010011010010011110011111010110111
110011100011000010101101110000100
00001101110101000011111011110110
1001000111011111010100101101111
011010111000010101100011000000001
1111011000010010011110001001101
```

n=34

```
011100010100100100111111011011110
110110110111010100101111110001110
100100010101010110001001101011101
111011110001001010000111111010100
000111010111011000001010010111011
0010101110110100010011101000101100
0100011001011001110111100011011010
11011101111000000111010111111110
110000110110001111011000100110111
0101100110011110011100110100111101
0111110101000011100010000001011100
0101111000101000001001010010000101
010110100110011101001001101011100
011101010100011111100011011101111
1111010011000101001111100110100000
1111011010010010001101000111001001
1111010001000010010101100000101100
101001100101111100101111110100100
101001101001100000111100111101111
```

## The Merkle-Hellman Cryptosystem

0101010000111010011001001000011001

n=35

0010101010001101011000011011111101  
000110001011111001001010110001101  
11000010010110011011110000111100101  
00011011010001000110101010001101100  
0010110110110011010101000010011000010  
11101100111000011001010101110011011  
0100111101010000011110101000111000  
010000010101011101110100110100100  
1011011101110000110111111001111100  
10111010001010101001000001001100111  
1110101000110101101111111111100000  
11001101110111001110110001010000010  
00011011010101101110001110110000100  
01011000101100011010100110010000110  
1010010011000010100000000101110000  
0110011001000000101111010111001001  
01111100011000011010000101011000101  
00001000111001011010100111000110111  
11001011110011000101100111001000010  
11100100100101010010100101011011010

n=36

010000000010010000111001110101001010  
110111000101111011001010001011110110  
110000010100000100001011100111011101  
111010101011111101000111001110011101  
111000000100101100101100110001011110  
001001001011101011101010101110010000  
000001000001011001110010111001010000  
000010110011100110000010000000010010  
101110111101111110001111011000000010  
1100100001001101010110010100111101  
000101011000011101000010011100001111  
101100010111011100001101110101100000  
111100011000110001000101001101001001  
1011010101101000001010110100110011  
001101101100000111000110110100010101  
001101100010001000110010110001111100  
111111000100101111001011010000011000  
001000101101000000101001011110110101  
001110111001011000000001110101100010  
000100011011101110010101011010011000

n=37

1101000000110011001010100011000100100  
1100110111001111000110001001110000101  
1011010000111100001100000001110000001  
101101000100000101110100111001000000  
0111010011110111010111100001000010100  
0100111011111100010111000110100110100  
001100010011011101011000001010110100  
101000100100001100000111100111111110  
1010011110100110110010110111011011011  
0111111110010011011101111000100101001  
1111001100110010001001010010111011010  
001100111010011110000111011101100100  
0100011010100000101101110010101100010  
110100111101101011111011000000100010  
1100001110000001011001101100011011001  
0011100110011000101000110010100101011  
1010110010111001001010100011101000010  
1110100000111010100000011011110110101  
1011000010010111001101000000010100100  
1010000000110111100110100001010100001

n=38

11100000011000100001011000010100001011  
11010010010011010101000100010101111011  
111100100000100110010100101011001110  
110101111100001111110001101010101111  
11000101001001110101111100011011100000  
0010010000010000001010000101111001001  
0110010010111101001011110110011010000  
0100010010111111101100101001010001110  
01111010101000110010010110000111100001  
00001100110100101101111111000010000101  
010001100001110111110011111111000110  
10111011111001100101011110000110101101  
10010010001010101001100001010011101001  
0010110111010111101110001011001011000  
11011101111010011101111011010010100100  
11111111011111000100000001010100011110  
11000100110110011100100010111010001010  
10110001001100111110111111000000110010  
11111000001011100010100111011110110010  
10001011010101110011100011100100101010

n=39

110101100000110001010110100101011010010  
011100101001111000000100100011010010101



## The Merkle-Hellman Cryptosystem

```
1011010110111100100000110111110001001
101001001001110110100101110101111101100
1101001111100100111000110001001100111
10110101110001010000001101110010000000
1011011100100011000110001111001111110
01101100110111000110010011010111000011
10010010100100000000110110010000110001
000000100101110111000101001101101110101
111101101101001000111110100001001110011
001101001001111100111010100111110000010
11101011010100001010000000100110101110
000000011100100000011101100011010010100
011010101000010000110100110110100010110
010001111111010000101101101110010001100
1100100111010101100011110001111101001
110111010011010001001100101101010110111
100010100001000011000101001100110111010
001011110001111010011100111000101001010
```

n=40

```
1010011100011111111101011010001001000111
1100100000111011000110000010111011101110
0111011100101001110011100110110101011011
0010110111000100110010000010101100000100
1110101110101011101010000000101110101001
1110110010000010001011100011011001001101
1110100100100101101100010101010111100100
0001001100011011001011010001010011001110
0010101110011010111010110011010001000010
0001111000100101110101110000100010000101
0101110001011010011010111110001101110111
1101111000010111100100010000111110001100
1111000010100000100100001011001110000110
0111011110011110001001111001101110000010
11010011111000110101111101001111110101
1100010000000011110001110000101101010010
0000100000000011101100101101011111010100
100011000001011001110000000011111111000
0001101110101000110000010000000110001000
1011101111000011011011101100001011010011
```

n=41

```
01000010100000110010110101111010110111001
01100110100001110111110100001001011000000
00101011000110011000011110010011100011110
1100101101001000001111111111000000010100
1010011111000001100100110111110101000001
1111001111001001010111110000110111101111
1011100100011111110110001110101110110011
0100010111010011010110011101010101111101
11011101111000101110001100000000110111011
01011111010000101001110110100001001101111
00000001000111100111111111101011000110000
00110110001110001100111111110010010111001
10001110011010110101000110100010000100000
01110110011111110011110101011100111100100
10010101011001111000110000111110101011111
10100010011110110100011100100000101101101
0101111101100111000011110000010011001000
01101000010100100000010101101100010101111
01110111011100101001101100001010001100010
01100011010111101110111110010000011110000
```

n=42

```
001101000000000101000010101111101010110001
011000111000001011010111110010101100000111
000011101110101011100100111110001110011100
00110100110000111001000010010011111110011
000001000111111110100110100101101010111111
1011100101010011011101110110110011001101
101111000111001101001010101111110011000110
10001001010101011001110001101011010111001
000100010010101000101100010001101011011100
111110101101011100011111001101011110110100
11011011001111011001111111010001010001110
100110000111010101100101001100110011010110
100011110011110000001011101110110011010001
1000111001111001001010110110101101001010
000011010010011010110111110011111001001111
001100010100110111000000001001111100110111
110110001001010001001001011101101101100110
110011110000001101010100001001101100101110
101000000101110100110011111110010110110100
0100000110111000100110101101001111000001
```

n=43

```
100010111001100110111111111111011010000110
1000011100101101101000100010101001001110000
1010010001111011100100011101100000010100101
0101101100010100101110000110001111000111001
0001110010101101100101100010001111100110100
1000011111000101110000000101001011100011000
0000111110011111101111101000000111010110111
```

## The Merkle-Hellman Cryptosystem

```
0000001000011110011101011001000100101000001
1010010011010000010011100010111000011100000
1010111000110001110010101000100001011101111
111100110110110111100000101001010011100000
010010101000101000100110000011111010100111
1100101101100000111000110001011011110000110
0110101000111011101011010000011011000010110
1011010101001100110110100110001000001011101
0000000110001101100001011000110010101111110
0010011101010011001111000010001000110010111
011100000100001110110111111000011110100111
001001111011001101000111111110111000100010
1000011010100111110011110110010010010010111
```

n=44

```
0010000011111000110001010011010010111011110
11101110011101000101010111100010000111010
010101000011111010011110100110001011101010
0011101100010110000011100110101011010011100
010010001001110010111011110011000010111110
100101111101101111100010010101110110101111
0110010100010111011110110110100110011000100
00000111001001111111001111110101011101110
101000010111001010111001111000101111011011
01111001001101011101100011001010110010000110
0000111010010110011001000010100100101111101
00000100011100010111000110110110111101110001
011010111011000110101101001111111111000110
10001111111001100000100110110000100011001111
01011001100001101110110000101101000011000110
0001000110110111101100011001110100110100110
11101100000111010001000000101000001011011011
11100011111011000011001011010000111010111111
10000100100000010001011111011100100110110110
10011000101011011001100001100001100100001001
```

n=45

```
111101110011010011001110011010000100011001101
011110011111001110001111100110101011001100101
110100101111101010101001110000001111011001010
001111010011100111011000000000011100000100101
100011111011100010110000100001100101111110110
100001000000010010011000110101100100101100110
111100100011001111000111001100110010110011001
000010101110110001011000001001010000111000000
00011000111100000110001011000111111110001110
100110001010111110000011100100110100001010010
00101110000010100010100001001000011111110111
10000010111110111110100110010111001110111100
000000001000000110010000110010111011110110110
001011011001000000010001010001101101110011100
11001011111110111001010011011110110101110001
011101111100101101011110100010100110001010010
110000100110001100111110011111001010100111101
11001100101110011000000010000110111100001001
111111111100001011110010111011011000111000100
011101011001010011000011110111101111100011001
```

n=46

```
1010111000001010100111110100110101110101110111
110000110010111010101110100000000000101011100
0101000000111110101011101100001001100101011011
1001011111010111000011110001100001101110000111
01101011101100101010101110100010101100110100
1011011110111011000010001001101110011111101010
101001101011001001100101011001110000100010010
0100101110100000001101101110000101000010110100
1110111100001101111100110100001101000001000100
110000100111010111111011111000101001001100111
010111101111100001101010000110101010001110001
1010010010010100101110011100010001000101110101
1111100000011010000001100111111011010110110110
1110001011100001000101100001111100100011010011
0110100101100011100001000100110000110101000110
01011110101110011101010100010100111000111000
1000001001000101110110000100101100111100111101
0011110010110011111010010001001101100010111111
11100010101000001011001000000000111011100011
0111011101011001011101001010111001010001100111
```

n=47

```
00110100100010010111010010001011111111101000110
11111110100000110101001010010011100100011101111
111101000100110101001010100011011010001100010
10100010110011110100011011100110011100110110
00101101101010111101101100010011110101110110111
0010011101100110000110010101001110010111101010
001101000110000110110010110110110010010111000
10010000010110110011111010000110001100100011001
00011011101100111001111010100101001001101000100
1010110011111011110101010011011111010100010010
0001011100110001011110100100111100101101001000
1100000000100100010100011101101000101001001
```

## The Merkle-Hellman Cryptosystem

```
11010101010111100101101001011101001001010011000
101001001110000101111110100001000010011000000
10101111011000110001110100101010011011110010
1100000010101010010100010101001011101010000100
1011101011101111011011110011000101110000010
111100010111001010010000100000101010010111000
11000000010011001100010010100101110100101110
1100100101110010001011100111011110011100100011
```

n=48

```
00100010010001101001101110000001101110001111110
11001000000110000000110101100111011000110110110
001010010000000100011110010001010100100011001010
000011111001111011001111010100010000010001011110
1110001110110101000010010011010101000111010100
010100010010011011101000110110111101101101110001
111110000101111000111001010111011110100111010110
0110101110011111011000011101111101000011100011
100110100100110110111001111010001111001011100100
001101010111100000110100101011000110010010010010
000100000001111001011000010110010111101110000111
111000110110000101100110000110111101100111011100
001111011100010001000001110001010000111000111100
101001010110010000011100110011110001110001010111
10111111100111011011001100111110011001000101010
11001001011010110111001110100110001010111001010
010010110001001001010100100110100100101010011000
11011111011111011100000101001010100001110010100
0111000110101110101011111110000101111000110010
100100001011100110101101001001010100000011100001
```

n=49

```
011011001101001001000100111011000100001111110100
1100000010100001100001101101000010000110001100110
0001001110010111011010111101010011100011001110000
010010000001010001101100011101101110011111110001
1011011010110011101110001011001110100110011101001
0000110000110111000000101100000101011101011011001
010010100001010101111011110001100001000111101
101110100100100011111100100110001010110110011101
1100010001100111010110011100110100000001000110011
010111011111010110011000110011101111010101100111
001001011000101001011110110010100010111001111101
0100110100001111010001010111100000010110010110110
0110111111000101111011000010010101011110010110111
010101011011110001110111101010000111101101000010
0100001111000001110011110010010100000001110001110
0110010001101000110111010101010011110001011101100
1101000001101100110001100111011101011001001110101
1001101111110011001111111000100011011100010000001
0111111000100111010000110011101110100110010101010
0101010001100000111110111011001101101001100011
```

n=50

```
00011100011000000101100100111010011011010001010111
01111100110010000111100011100000111000101100011000
01010101100100011000100000010111110001000011110110
00011010111101001111100000010110101011000111110011
0110100000110000111000011101101011010010101000000
101010110111010100100011110010011000011101101100001
10100100001010011000001011110101100011011110000010
01111100011010111001101000001001110110010001101100
11000110101111111001100111101010110100000001000
01000101111011100110010010111100001101101010000111
0100010000011100011010111110010100101000001111110
1110110010011011011101000101110110100000001110100
11101110100100110110000111011011101000101100111
10010111110101000100011010110100100000100111100111
00100000110000001000011001010110001110101101001101
10101000001100010010101010011011110100000001111
011101111110000011011110100101100101001000001011000
0000001000110011101010000000100001001000000010010
0011111001011101010010010100101111100110011110100
0111000101101011000010100011001011011101000111111
```

n=51

```
001000001001100001000110000100110001100011000011011
00100010010111100101101111001110110100111011100110
10100000101001101100110010001111011001001010101011
11111110011101110101000111011010101100111111001
00110100010100101110100000111100001000000111101111
000000101100000100110000010011101001010111101111001
100111000011101001101010010000111000101110111101000
0000011000101101001010000101110110111111101010001
100001100100001000010111011001000001010001010011000
010010100010001001110000011010100011011110101000011
111001110000011010110001001100001010101010000101100
0100111011000111000011001011001101101111111010000
111101110010110000111010101010111011101001110010101
110000110100100011001100011110010100100011110100100
000101001000001110111100011100001010111001111010101
000001011101010101110011100101110000000011101100111
0110110111110000001001010011100111111001011010110001
```

## The Merkle-Hellman Cryptosystem

111010001110000110000110010001111001110011001111000  
101011010101100011101010000101001111100001000100010  
010011000010110100011101010010010100011111010010111

n=52

1101110001101100110001100100100100110101101001000111  
100001011101111110101100010011011000110101101011000  
0111101101011101000011001110010101100101001001111000  
1011000110110001010100000111001100111010000001101110  
01011010000011011011010111000001111100000100000000  
1100001010111011111010001000110010001100010111101  
1101111110101011001100111011100101110001110110000100  
1110100001111101110000100100100101101100010011101010  
111001000011100001011001110101101000100101101101000  
0101001000011100101001010011001011111011101101000  
11010000010010000101110111101011000011000000010110  
1100101110011000001100110000100100100101110101100110  
111101011110010101110011100001100100100100110001110  
1101011001001110101010011101101000100001011111100  
1101100000111010110111001111111110011011110010011  
010101101111000101110101010100011000100010110100110  
111101111111011100100011001110011001010001000110101  
001101100110100011101101001111001000110101010101101  
100001100001011101000000010110000000110001010110110  
0110110110001111100000111100000011011100111011100011

n=53

10110110100011011001011000111101011110010100110000011  
0101001111011100001100111101111101100100111110010000  
0001110011101110010100011101001111100011001111101001  
0011011010010111101001111100011110100011011110010111  
011010101010111111110110101010011110011001000011111  
00110010011000111001011010000010001101111101110101  
000101011101011001110011110010100101100010110110010  
100101110111001010110011110001111100101011100001110  
11011100011010111100011101000011001101110001101101111  
000001011101010000100000000001011111010011100101110  
0000110100001110101101011111010111110010111000100100  
110101001101110000111110111110011011100100100000010  
010001011100111100110111110110000001000011010100011  
00001111001101000010010001100000000010101110100011111  
00001110001001110011000000010010100011001101100000110  
0000011100111101011101011000101011001001011010010100  
11001011010010101001001100111111001000010110111101  
0111001011110001110001000111011100001001111000101010  
00100111011100011110000001100100000110010111001100011  
11011101101001100000110111000000110010100101010001110

n=54

011101100010010010000110110110001100000001000101011001  
010001111011111111111100110011001000110101011000000  
11010010010101101010011110110100101011101100110100110  
1001011111001111010111000000000101111001011100110010  
1101101000001010110101000100111101011011110110111110  
01011011010001111000000101011100110100011111101111  
0011110100001111101001101110011010011111010001100011  
10011010010000011100110000101010000010101100100011100  
00001011001011001010000001101101111000101011000001010  
1111100100101100000010101111010111110111000001110110  
0000000101110111110110100011100000111000110111101111  
11000111110000110100111000111010010010101111011101  
01010000011010011001111101101000001000000000100101100  
100111001101011110001000000110110010001100010101100110  
00110101011011111101110010101011100010101110110101000  
111110110011000000101100100111000111011001011101100111  
11011001001110100010010001010100001010001100000111111  
00100100111001111100100100100101010001010111000000000  
1100111000011000010101111000010100111101110111101101  
010010111101111011111100100010100111011010001100101100

n=55

0111010100111110001110110100110101001110101001110100010  
0101100101001010110101011100111100000010100100011010000  
1101011010010111011100001110000010001000110011001010010  
0101111000011100111011000011010101100100110010110110110  
1011000101010101101110100010001110110000101010111011  
0110110101100010110111001011010101101011101110101011011  
1011000100101111100100010110101000011101001110100010011  
011000101111110111110010001000011101011100011110111101  
110100111101011110001111000110101111100010110010010010  
0101011111001000000101110001101110010100000101101011111  
10100100100000001000010001011000100100100111000111001  
011011110001100110100111001111000010000101001110110000  
0001110011010001101000001010111001110001110101110111000  
01111100011011010011101100011001110000111001111001110  
0110100011110101111100100101011101111100101110000010  
111000110110100111011010101111000000010001101000111101  
0100101010000101000110111001101100110111001001100101  
11010011011100111100010100011100011011100010001000111  
0110000011011110011100101100100110101010111100110101111  
01001011111100101011101001111010110101001001010100000

n=56

## The Merkle-Hellman Cryptosystem

```
101001111100010000100111101010100001101011111110110000
11001001010001111001101000010111010010000110011101100111
00111001100111101100111001000101001000101000001010001001
00100100010100101010000101010100010100011101111100011101
0001001110001100001010101000101001111010000001100110011
01111100011000010111011100110011100111100101010100000001
0011100101101010111011110110110001001100111001011110111
00010010011101011010001001111101101010110000010111010011
010011100001100110111110010010000100111111111110001100
01110001110101110101001001000111100101011001000011111010
111000011000100100111000111100100000101111110010111011
000110001110001100111111100000000010100110110101101101
10001010010111001010010000101001000000010011001100110011
00000001110110010100010101110100010000100010110110001011
11000111110100111010110011001110000010110101101011010000
01001001111001100111100101110010111100011010100010010101
000110110100111011110101000001000000011100011010100100
11000100000111011110000100001000101011111000100101010001
00101011110111101000010101111000101000100000011111001
10001111111110010001011000011100101101111000110010000001
```

n=57

```
110011110101101110110000000100000111011100101100111100100
101010011101001010011001111111100001001100011101000101101
0111010111111000000000001100110000111001101011111101010
100101011011101101110011110100100101000001100101000111110
000111100110101001100100100010000010100000110001011001010
100110111010111011000111001110111010111010000000010100110
0010111001000111000101010000001000011000110011111110111
10000110000011010110100001011011110011010111111011111001
110011111110010100000011100101001110010001000111101000110
001011111010101010110011001011110100101001100101111100011
111101010101111000011000100101101101000100001010011001011
1110100000000001000010001100111011100011010100101110011000
10000101110111011100000101011010110100111110111100101100
111000110101011011100011010110010101101100010100100101010
101010101110010000111011100111110010010000100100011011011
101100110110000101110101110011000100000001010100001100010
001000001110011100100001111001011000111010110010101111000
001011000010011111111101010100101110010100001110111101011
1010000101011010100011011000100100101101011100111100101
000010010010010000010010110011011111100000100100000111
```

# Código Desenvolvido

## Código Desenvolvido em C

```

1 //
2 // AED, November 2021
3 //
4 // Solution of the first practical assignment (subset soma problem)
5 //
6 // Place your student numbers and names here
7 // Student 1: André Butuc (nMec: 183538)
8 // Student 2: Gonçalo Silva (nMec: 183668)
9 //
10 #if __STDC_VERSION__ < 199901L
11 # error "This code must be compiled in c99 mode or later (-std=c99)" // to handle the unsigned long long data type
12 #endif
13 #ifndef STUDENT_H_FILE
14 # define STUDENT_H_FILE "103538.h"
15 #endif
16
17 // Para manipular os algoritmos que são corridos.
18 #define IterativeBruteForce // -> Ativa iterative_brute_force se descomentado
19 #define RecursiveBruteForce // -> Ativa recursive_brute_force se descomentado
20 #define CleverRecursiveBruteForce // -> Ativa clever_recursive_brute_force se descomentado
21 #define HorowitzSahni // -> Ativa horowitz_sahni se descomentado
22 #define SchroeppeShamir // -> Ativa schroeppe_shamir se descomentado
23
24 // Para manipular as versões HorowitzSahni e SchroeppeShamir
25 #define QuickSort // -> Ativa versão QuickSort em Horowitz e Schroeppe se descomentado
26 #define MergeSort // -> Ativa versão MergeSort em Horowitz e Schroeppe se descomentado
27
28 // Para manipular se há ou não escrita em ficheiro.
29 #define WriteFile // -> Ativa escrita em ficheiro dos tempos de execução, se descomentado.
30 #define PrintTime // -> Ativa a impressão em terminal dos tempos de execução, se descomentado.
31 #define PrintOnlySolution // -> Imprime só as combinações no Iterativo,
32 // não ativar quando não estiver a correr o IterativeBruteForce ou não ativar
33 // quando não está ativada o OutputSolutionsRC ou OutputSolutionHS
34
35 // Para manipular se há impressão ou não das soluções no terminal.
36 #define OutputSolutionsRC // -> Ativar para ver as soluções dos algoritmos recursive_brute_force ou clever_recursive_brute_force, nunca ativar com outros algoritmos.
37 #define OutputSolutionsHS // -> Ativar para ver as soluções dos algoritmos horowitz_sahni ou schroeppe_shamir, nunca ativar com outros algoritmos.
38
39
40 // include files
41 #include <stdio.h>
42 #include <stdlib.h>
43 #include "../P02/elapsed_time.h"
44 #include STUDENT_H_FILE
45
46 typedef struct {
47     integer_t mascara;
48     integer_t soma;
49 } mascara_struct_t;
50
51 typedef struct {
52     integer_t mascara;
53     integer_t soma;
54     int indice1;
55     int indice2;
56 } heapStruct_t;
57
58 /*
59 Seria uma boa ideia colocar este algoritmo no relatório de forma a dar entender ao professor
60 o nosso algoritmo piloto que nos permitiu estudar o problema e começar a produzir soluções
61 É possível calcular as somas até n=27.
62 É também interessante mostrar o gráfico desta solução porque o best-case e o worst-case irão ser bastantes diferentes.
63 Por exemplo: Para n=28; Best-Case: 1.4 segundos; Worst-Case: > 15 segundos;
64
65 */
66 // -----iterative_brute_force----- //
67
68 int iterative_brute_force(int n, integer_t p[n], integer_t desired_soma)
69 {
70     integer_t test_soma;
71     for (int comb=0; comb<(1<<n); comb++)
72     {
73         test_soma=0;
74         for (int bit=0; bit<n; bit++)
75         {
76             if ((comb & (1 << bit)) != 0)
77             {
78                 test_soma += p[bit];
79             }
80         }
81         if (test_soma==desired_soma)
82         {
83             for (int bit=0; bit<n; bit++)
84             {
85                 if ((comb & (1 << bit)) != 0)
86                 {
87                     printf("%d ", bit);
88                 }
89             }
90             printf("\n");
91             return 1;
92         }
93     }
94     return 0;
95 }
96 // -----

```

# The Merkle-Hellman Cryptosystem

O algoritmo `recursive_brute_force`, como o nome indica encontra a soma recursivamente. Em cada valor de `p` há um "branching": o bit do valor de `p` fica a 1 ou fica a 0. Em cada ramo, a soma parcial (`partial_soma`) é incrementada (no ramo do bit igual a 1) ou mantém-se igual (no ramo do bit igual a 0). De seguida é incrementado o índice do array `p` e é chamada de novo a função `recursive_brute_force`.

As condições de paragem da função recursiva são:

- `if(partial_soma == desired_soma)`, que significa que foi encontrada a solução. Retornando 1 irá desencadear o desenvolvimento das sucessivas chamadas prévias da função. Sendo que o bit da folha que desencadeou a terminação da função com sucesso irá ser armazenado no array `b` e assim sucessivamente até à raiz.
- `if(current_index == n)`, que significa, caso verdadeiro, que a recursividade chegou ao fim do array `p` e como já não há mais elementos ou combinações a somar e como ainda não foi encontrada solução, retorna o valor 0 que por sua vez indica que não foi encontrada a solução=soma.

\*/

```
// -----recursive_brute_force----- //
```

```
int recursive_brute_force(int n, integer_t p[n], int b[n], integer_t desired_soma, int current_index, integer_t partial_soma)
```

```
{  
    if(partial_soma == desired_soma)  
        return 1;  
    if(current_index == n)  
        return 0;  
    if (recursive_brute_force(n, p, b, desired_soma, current_index + 1, partial_soma + 0 * p[current_index]) != 0){  
        b[current_index]=0;  
        return 1;  
    }  
    if (recursive_brute_force(n, p, b, desired_soma, current_index + 1, partial_soma + 1 * p[current_index]) != 0){  
        b[current_index]=1;  
        return 1;  
    }  
    return 0;  
}
```

```
// ----- //
```

/\*

O algoritmo `clever_recursive_brute_force`, é semelhante ao algoritmo `recursive_brute_force`, no entanto, apresenta mais dois argumentos de entrada que servem calcular uma soma a contar dos últimos índices para os primeiros.

É feita esta soma simultânea de baixo para cima e de cima para baixo do array `p` para que quando os `current_index` e `current_last_index` se "encontrarem", ou seja, quando forem iguais, podemos avaliar se é possível realmente chegar à soma-solução ou se o nosso array de valores tem valores demasiado pequenos para alguma vez chegar à soma solução.

As condições de paragem da função recursiva são:

- `if(partial_soma == desired_soma)`, que significa que foi encontrada a solução. Retornando 1 irá desencadear o desenvolvimento das sucessivas chamadas prévias da função. Sendo que o bit da folha que desencadeou a terminação da função com sucesso irá ser armazenado no array `b` e assim sucessivamente até à raiz.
- `if(current_index == n)`, que significa, caso verdadeiro, que a recursividade chegou ao fim do array `p` e como já não há mais elementos ou combinações a somar e como ainda não foi encontrada solução, retorna o valor 0 que por sua vez indica que não foi encontrada a solução=soma.

(Até agora iguais ao algoritmo de `recursive_brute_force`.)

- `if(partial_soma > desired_soma)`, a `partial_soma` já é superior à soma-solução logo nunca irá ser alcançada porque os valores de `p`

```
203 int partition(mascara_struct_t arr[], int low, int high) {  
204     mascara_struct_t pivot = arr[high];  
205     int i = (low - 1);  
206     for (int j = low; j <= high - 1; j++) {  
207         if (arr[j].soma < pivot.soma) {  
208             i++;  
209             swap(&arr[i], &arr[j]);  
210         }  
211     }  
212     swap(&arr[i + 1], &arr[high]);  
213     return (i + 1);  
214 }  
215  
216  
217 void quicksort(mascara_struct_t arr[], int low, int high) {  
218     if (low < high) {  
219         int pi = partition(arr, low, high);  
220         quicksort(arr, low, pi - 1);  
221         quicksort(arr, pi + 1, high);  
222     }  
223 }  
224  
225 void insertion_sort(mascara_struct_t *data, int first, int one_after_last)  
226 {  
227     int i, j;  
228     for(i = first + 1; i < one_after_last; i++)  
229     {  
230         mascara_struct_t tmp = data[i];  
231         for(j = i; j > first && tmp.soma < data[j - 1].soma; j--) {  
232             data[j] = data[j - 1];  
233         }  
234         data[j] = tmp;  
235     }  
236 }  
237  
238 void merge_sort(mascara_struct_t *data, int first, int one_after_last)  
239 {  
240     int i, j, k, middle;  
241     mascara_struct_t *buffer;  
242     if(one_after_last - first < 40)  
243         insertion_sort(data, first, one_after_last);  
244     else  
245     {  
246         middle = (first + one_after_last) / 2;  
247         merge_sort(data, first, middle);  
248         merge_sort(data, middle, one_after_last);  
249         buffer = (mascara_struct_t *)malloc((size_t)(one_after_last - first) * sizeof(mascara_struct_t)) - first; // no error check!  
250         i = first;  
251         j = middle;  
252         k = first;
```

# The Merkle-Hellman Cryptosystem

```
213 while(k < one_after_last)
214     if(i == one_after_last || (i < middle && data[i].soma <= data[j].soma))
215         buffer[k++] = data[i++];
216     else
217         buffer[k++] = data[j++];
218 for(i = first; i < one_after_last; i++)
219     data[i] = buffer[i];
220 free(buffer + first);
221 }
222 }
223 // ----- //
224
225 /*
226 void soma_all_subsets(int n, integer_t p[n], mascara_struct_t result[1 << n], int level, integer_t mascara, integer_t subsoma, int idx)
227
228 O algoritmo void soma_all_subsets é um algoritmo recursivo que procura encontrar todas as somas possíveis e as correspondentes máscaras
229 das somas. Parte do mesmo princípio dos algoritmos recursive_brute_force e clever_recursive_brute_force, no entanto como sabemos que a função terá que percorrer todo
230 o vetor p, a única condição de paragem é if(current_index == n) que significa que o current_index foi incrementado até ter chegado ao último índice do array p.
231 No entanto, como sabemos que teremos que guardar todas as "máscaras" das somas, quando a função retorna todas devem ser guardadas.
232 Devido a este fator desenvolvemos uma estrutura que guarda tanto a soma como a máscara:
233 typedef struct {
234     integer_t mascara;
235     integer_t soma;
236 } mascara_struct_t;
237
238 E agora nas sucessivas chamadas recursivas já não é avaliado o valor de retorno, sendo este sempre void, portanto o armazenamento não é feito por indexação
239 num array que guarda os bits como nos algoritmos anteriores, mas é feito com o auxílio da variável mascara que é representada por bits, que é inicialmente 0.
240 Em cada chamada recursiva é feita uma operação bit-wise OR do valor da mascara com o valor de 1 bit deslocado para a esquerda de current_index bits, caso
241 o bit do elemento de p seja "um".
242 Vamos supor que temos o valor inicial da mascara 0 e que o elemento que estamos a avaliar tem índice de 5 no vetor p e vamos colocar este elemento com bit 1,
243 ou seja, participa na soma parcial. Então mascara | (1 << 5) => mascara | 100000 => 000000 | 100000 = 100000 (o valor da mascara é estendido com 0 mais significativos
244 até se igual à dimensão com a qual estamos a fazer a operação bit-wise OR, o mesmo acontece inversamente).
245
246 Para não haver conflito quando quisermos guardar elementos no vetor somas (que é o vetor de resultados), quando há o branching do ramo que iguala o bit do elemento a ser
247 avaliado a 0 (não participa na soma parcial) é incrementado o índice dos índices ímpares e o ramo que iguala o bit do elemento a ser avaliado a 1 (participa na soma parcial)
248 é incrementado o índice dos índices pares.
249
250 Quando o índice chega ao final do array p a recursão pára e todos os dados são guardados nos campos da sua estrutura na posição index no vetor soma.
251
252 */
253
254 void soma_all_subsets(int n, integer_t p[n], mascara_struct_t somas[1 << n], int current_index, integer_t mascara, integer_t partial_soma, int index){
255     if(current_index == n)
256     {
257         somas[index].soma = partial_soma;
258         somas[index].mascara = mascara;
259         return;
260     }
261     soma_all_subsets(n, p, somas, current_index + 1, mascara, partial_soma, 2*index);
262     soma_all_subsets(n, p, somas, current_index + 1, mascara | (1 << current_index), partial_soma + p[current_index], 2*index+1);
263 }
```

```
264
265 // type == 1 -> min_heap    type == 2 -> max_heap
266
267 void push(heapStruct_t h[], heapStruct_t elem, int* h_size, int type)
268 {
269     int i = 0;
270     for (i = *h_size; i > 0 && ((h[(i - 1) / 2].soma > elem.soma && type == 1) || (h[(i - 1) / 2].soma < elem.soma && type == 2)); i = (i - 1) / 2)
271     {
272         h[i] = h[(i - 1) / 2];
273     }
274     h[i] = elem;
275     (*h_size)++;
276 }
277
278 // type == 1 -> min_heap    type == 2 -> max_heap
279
280 heapStruct_t pop(heapStruct_t h[], int* h_size, int type)
281 {
282     int i, c;
283     heapStruct_t elem = h[0];
284
285     (*h_size)--;
286     for (i = 0; i + 2 + 1 <= *h_size; i = c) {
287         c = 2 * i + 1;
288         if (c < *h_size && ((h[c].soma > h[c + 1].soma && type == 1) || (h[c].soma < h[c + 1].soma && type == 2)))
289             c++;
290         if ((h[c].soma < h[*h_size].soma && type == 1) || (h[c].soma > h[*h_size].soma && type == 2))
291             h[i] = h[c];
292         else
293             break;
294     }
295     h[i] = h[*h_size];
296     return elem;
297 }
298 // ----- //
299
300 // ----- horowitz_sahni ----- //
301
302 int horowitz_sahni(int n, integer_t p[n], integer_t *solucao_mascara, integer_t desired_sum, int current_index, integer_t partial_soma){
303     int n1 = n/2;
304     int n2 = n - n1;
305     integer_t p1[n1];
306     integer_t p2[n2];
307     int length_a = 1 << n1;
308     int length_c = 1 << n2;
309     mascara_struct_t *a = malloc(length_a * sizeof(mascara_struct_t));
310     mascara_struct_t *c = malloc(length_c * sizeof(mascara_struct_t));
311
312 }
```



# The Merkle-Hellman Cryptosystem

```
353 for (int i=0; i<n; i++){
354     if (i < n/2){
355         p1[i] = p[i];
356     }
357     else {
358         p2[i - (n/2)] = p[i];
359     }
360 }
361
362 soma_all_subsets(n1, p1, a, 0, 0, 0, 0);
363 soma_all_subsets(n2, p2, c, 0, 0, 0, 0);
364
365 #ifdef MergeSort
366 merge_sort(a, 0, length_a);
367 merge_sort(c, 0, length_c);
368 #endif
369
370 #ifdef QuickSort
371 quicksort(a, 0, length_a-1);
372 quicksort(c, 0, length_c-1);
373 #endif
374
375 int indice_a = 0;
376 int indice_c = length_c - 1;
377 integer_t soma = 0;
378 integer_t solucao = 0;
379 int ret = 0;
380
381 while (indice_a < length_a && indice_c >= 0) {
382     soma = a[indice_a].soma + c[indice_c].soma;
383     if (soma == desired_sum) {
384         solucao = a[indice_a].mascara | (c[indice_c].mascara << n1);
385         ret = 1;
386         break;
387     }
388     if (soma < desired_sum) {
389         indice_a++;
390     }
391     if (soma > desired_sum){
392         indice_c--;
393     }
394 }
395
396 *solucao_mascara = solucao;
397 free(a);
398 free(c);
399 return ret;
400 }
401
```

```
402 // ----- //
403
404 // -----schroepel_shamir----- //
405
406 int schroepel_shamir (int n, integer_t p[], integer_t desired_soma, integer_t *solucao_mascara) {
407
408     int n1 = n/4 + (n%4)/2;
409     int n2 = n/4;
410     int n3 = n/4 + (n%4)/2;
411     int n4 = n-n1-n2-n3;
412
413     integer_t p1[n];
414     integer_t p2[n2];
415     integer_t p3[n3];
416     integer_t p4[n4];
417
418     for (int i=0; i<n; i++){
419         if (i < n1){
420             p1[i] = p[i];
421         }
422         else if (i < n1+n2){
423             p2[i - n1] = p[i];
424         }
425         else if (i < n1+n2+n3){
426             p3[i - n1 - n2] = p[i];
427         }
428         else {
429             p4[i - n1 - n2 - n3] = p[i];
430         }
431     }
432
433     int length_a1 = 1 << n1;
434     int length_a2 = 1 << n2;
435     int length_b1 = 1 << n3;
436     int length_b2 = 1 << n4;
437
438     mascara_struct_t * a1 = malloc((length_a1) * sizeof(mascara_struct_t));
439     mascara_struct_t * a2 = malloc((length_a2) * sizeof(mascara_struct_t));
440     mascara_struct_t * b1 = malloc((length_b1) * sizeof(mascara_struct_t));
441     mascara_struct_t * b2 = malloc((length_b2) * sizeof(mascara_struct_t));
442
443     soma_all_subsets(n1, p1, a1, 0, 0, 0, 0);
444     soma_all_subsets(n2, p2, a2, 0, 0, 0, 0);
445     soma_all_subsets(n3, p3, b1, 0, 0, 0, 0);
446     soma_all_subsets(n4, p4, b2, 0, 0, 0, 0);
447
448     #ifdef MergeSort
449     merge_sort(a1, 0, length_a1);
450
```

# The Merkle-Hellman Cryptosystem

```
451 merge_sort(a2, 0, length_a2);
452 merge_sort(b1, 0, length_b1);
453 merge_sort(b2, 0, length_b2);
454 #endif
455
456 #ifdef QuickSort
457 quicksort(a1, 0, length_a1 - 1);
458 quicksort(a2, 0, length_a2 - 1);
459 quicksort(b1, 0, length_b1 - 1);
460 quicksort(b2, 0, length_b2 - 1);
461 #endif
462
463 heapStruct_t min_heap[length_a1];
464 int n_min_heap = 0;
465
466 heapStruct_t max_heap[length_b1];
467 int n_max_heap = 0;
468
469 for (int i = 0; i < length_a1; i++) {
470     heapStruct_t soma = {
471         .mascara = a1[i].mascara | (a2[0].mascara << n1),
472         .soma = a1[i].soma + a2[0].soma,
473         .indice1 = i,
474         .indice2 = 0
475     };
476     push(min_heap, soma, &n_min_heap, 1);
477 }
478
479 for (int i = 0; i < length_b1; i++) {
480     heapStruct_t soma = {
481         .mascara = b1[i].mascara | (b2[length_b2 - 1].mascara << n3),
482         .soma = b1[i].soma + b2[length_b2 - 1].soma,
483         .indice1 = i,
484         .indice2 = length_b2 - 1
485     };
486     push(max_heap, soma, &n_max_heap, 2);
487 }
488
489 int ret = 0;
490 integer_t solucao = 0;
491
492 while (n_max_heap > 0 && n_min_heap > 0) {
493     integer_t partial_soma = min_heap[0].soma;
494     partial_soma += max_heap[0].soma;
495
496     if (partial_soma == desired_soma) {
497         solucao = min_heap[0].mascara | (max_heap[0].mascara << (n1+n2));
498         ret = 1;
499         break;
500     }
```

```
501 }
502
503 if (partial_soma < desired_soma) {
504     heapStruct_t vmin = pop(min_heap, &n_min_heap, 1);
505     vmin.indice2++;
506     if (vmin.indice2 < length_a2) {
507         heapStruct_t w = {
508             .mascara = a1[vmin.indice1].mascara | (a2[vmin.indice2].mascara << n1),
509             .soma = a1[vmin.indice1].soma + a2[vmin.indice2].soma,
510             .indice1 = vmin.indice1,
511             .indice2 = vmin.indice2
512         };
513         push(min_heap, w, &n_min_heap, 1);
514     }
515 }
516
517 if (partial_soma > desired_soma) {
518     heapStruct_t vmax = pop(max_heap, &n_max_heap, 2);
519     vmax.indice2--;
520
521     if (vmax.indice2 >= 0) {
522         heapStruct_t w = {
523             .mascara = b1[vmax.indice1].mascara | (b2[vmax.indice2].mascara << n3),
524             .soma = b1[vmax.indice1].soma + b2[vmax.indice2].soma,
525             .indice1 = vmax.indice1,
526             .indice2 = vmax.indice2
527         };
528         push(max_heap, w, &n_max_heap, 2);
529     }
530 }
531 }
532
533 *solucao_mascara = solucao;
534 free(a1);
535 free(a2);
536 free(b1);
537 free(b2);
538 return ret;
539 }
540 // ----- //
541
542 // -----main program----- //
543 int main(void)
544 {
545     fprintf(stderr, "Program configuration:\n");
546     fprintf(stderr, "  min_n ..... %d\n", min_n);
547     fprintf(stderr, "  max_n ..... %d\n", max_n);
548     fprintf(stderr, "  n_somas ..... %d\n", n_sums);
549     fprintf(stderr, "  n_problems .. %d\n", n_problems);
550 }
```

# The Merkle-Hellman Cryptosystem

```
551 fprintf(stderr, " integer_t ... %d bits\n", 8 * (int)sizeof(integer_t));
552
553 #ifdef WriteFile
554 FILE *fpb = NULL;
555 fpb = fopen("test.txt", "w");
556 #endif
557 for(int i = 0; i < n_problems; i++)
558 {
559     int n = all_subset_sum_problems[i].n;
560     if(n > max_n)
561         continue;
562     #if defined(OutputSolutions) || defined(PrintOnlySolution)
563     printf("n=%d\n", n);
564     #endif
565     integer_t *p = all_subset_sum_problems[i].p;
566     for(int j = 0; j < n_sums; j++)
567     {
568         integer_t desired_soma = all_subset_sum_problems[i].sums[j];
569         #if defined(HorowitzSahni) || defined(SchroepelShamir)
570         integer_t b = 0;
571         #endif
572
573         #if defined(RecursiveBruteForce) || defined(CleverRecursiveBruteForce)
574         int b1[];
575         for (int l=0; l < n; l++){
576             b[l]=0;
577         }
578         #endif
579         int success;
580         if (n==min_n){
581             #ifdef WriteFile
582             char stringb[20];
583             sprintf(stringb, "%d", n);
584             fprintf(fpb, "%s ", stringb);
585             #endif
586
587             double time = cpu_time();
588             #ifdef IterativeBruteForce
589             success=iterative_brute_force(n, p, desired_soma);
590             #endif
591
592             #ifdef RecursiveBruteForce
593             success=recursive_brute_force(n, p, b, desired_soma, 0, 0);
594             #endif
595
596             #ifdef CleverRecursiveBruteForce
597             success=clever_recursive_brute_force(n, p, b, desired_soma, 0, 0, n-1, 0);
598             #endif
599 }
```

```
600 #ifdef HorowitzSahni
601     success=horowitz_sahni(n, p, &b, desired_soma, 0, 0);
602 #endif
603
604 #ifdef SchroepelShamir
605     success=schroepel_shamir(n, p, desired_soma, &b);
606 #endif
607
608 if (success==1){
609     ;
610     #if defined(PrintOnlySolution)
611     printf("Solução encontrada, n=%d.\n", n);
612     #endif
613
614     #ifdef OutputSolutionsHS
615     for (int i = 0; i < n; i++) {
616         printf("%s", b & i ? "1" : "0");
617         b = b >> 1;
618     }
619     printf("\n");
620     #endif
621
622     #ifdef OutputSolutionsRC
623     for (int i=0; i<n; i++){
624         printf("%d", b[i]);
625     }
626     printf("\n");
627     #endif
628 }
629 else {
630     printf("Solution not found.\n");
631     return 0;
632 }
633
634 time = cpu_time()-time;
635 #if defined(PrintTime) && !defined(PrintOnlySolution)
636 printf("time needed= %f\n", time);
637 #endif
638 #ifdef WriteFile
639 char tempob[50];
640 sprintf(tempob, "%f\n", time);
641 fprintf(fpb, "%s", tempob);
642 #endif
643 }
644 }
645 printf("\n");
646
647 }
```

```

648 }
649 #ifdef WriteFile
650 fclose(fpb);
651 #endif
652 return 1;
653 }
654
655 // ----- //
```

## Código Desenvolvido em MatLab

```

%% iterative_brute_force
fileID=fopen('iterative_brute_force_103668.txt','r');

tline=fgetl(fileID);
n_array_ib=zeros(1,21);
time_array=zeros(1,20*21);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    disp(time);
    if ~ismember(n,n_array_ib)
        n_array_ib(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;

    tline=fgetl(fileID);
end

time_array_media_ib=zeros(1,21);
time_array_best_ib=zeros(1,21);
time_array_worst_ib=zeros(1,21);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    % if time_array(i)<best

```

```

for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_ib(count_tm)=media;
        time_array_best_ib(count_tm)=best;|
        time_array_worst_ib(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end

```

---

```

%% recursive brute_force
fileID=fopen('recursive_brute_force_103530.txt','r');
tline=fgetl(fileID);
n_array_rb=zeros(1,26);
time_array=zeros(1,20*26);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    disp(time);
    if ~ismember(n,n_array_rb)
        n_array_rb(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;

    tline=fgetl(fileID);
end

time_array_media_rb=zeros(1,26);
time_array_best_rb=zeros(1,26);
time_array_worst_rb=zeros(1,26);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);

```

---

```
for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_rb(count_tm)=media;
        time_array_best_rb(count_tm)=best;
        time_array_worst_rb(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end
```

---

```

%% clever_recursive_brute_force
fileID=fopen('clever_recursive_brute_force_103530.txt','r');
tline=fgetl(fileID);
n_array_cb=zeros(1,26);
time_array=zeros(1,20*26);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    disp(time);
    if ~ismember(n,n_array_cb)
        n_array_cb(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;
    tline=fgetl(fileID);
end
time_array_media_cb=zeros(1,25);
time_array_best_cb=zeros(1,25);
time_array_worst_cb=zeros(1,25);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    if time_array(i)>worst

```



```

for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_cb(count_tm)=media;
        time_array_best_cb(count_tm)=best;
        time_array_worst_cb(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end

```

```

%% quicksort horowitz
fileID=fopen('quicksort_horowitz_103668.txt','r');
tline=fgetl(fileID);
n_array_qh=zeros(1,48);
time_array=zeros(1,20*48);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    if ~ismember(n,n_array_qh)
        n_array_qh(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;

    tline=fgetl(fileID);
end
time_array_media_qh=zeros(1,48);
time_array_best_qh=zeros(1,48);
time_array_worst_qh=zeros(1,48);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    if time_array(i)>worst

```

```

for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_qh(count_tm)=media;
        time_array_best_qh(count_tm)=best;
        time_array_worst_qh(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end

```

```

%% merge sort horowitz
fileID=fopen('mergesort_horowitz_103668.txt','r');
tline=fgetl(fileID);
n_array_mh=zeros(1,48);
time_array=zeros(1,20*48);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    if ~ismember(n,n_array_mh)
        n_array_mh(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;

    tline=fgetl(fileID);
end
time_array_media_mh=zeros(1,48);
time_array_best_mh=zeros(1,48);
time_array_worst_mh=zeros(1,48);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    if time_array(i)>worst

```

```

for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_mh(count_tm)=media;
        time_array_best_mh(count_tm)=best;
        time_array_worst_mh(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end

```

```

%% quicksort Schroepfel
fileID=fopen('quicksort_schroepfel_103668.txt','r');
tline=fgetl(fileID);
n_array_qs=zeros(1,48);
time_array=zeros(1,20*48);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    if ~ismember(n,n_array_qs)
        n_array_qs(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;

    tline=fgetl(fileID);
end
time_array_media_qs=zeros(1,48);
time_array_best_qs=zeros(1,48);
time_array_worst_qs=zeros(1,48);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    if time_array(i)>worst

```

```

for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_qs(count_tm)=media;
        time_array_best_qs(count_tm)=best;
        time_array_worst_qs(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end

```

---

```

%% merge sort Schroepfel
fileID=fopen('mergesort_schroepfel_103668.txt','r');
tline=fgetl(fileID);
n_array_ms=zeros(1,48);
time_array=zeros(1,20*48);
count=1;
count_t=1;
while ischar(tline)
    n=str2double(tline(1:2));
    time=str2double(tline(4:end));
    if ~ismember(n,n_array_ms)
        n_array_ms(count)=n;
        count=count+1;
    end
    time_array(count_t)=time;
    count_t=count_t+1;

    tline=fgetl(fileID);
end
time_array_media_ms=zeros(1,48);
time_array_best_ms=zeros(1,48);
time_array_worst_ms=zeros(1,48);
count_20=1;
count_tm=1;
soma=0;
media=0;
worst=time_array(1);
best=time_array(1);
for i=1:length(time_array)
    if time_array(i)>worst

```

---



```

for i=1:length(time_array)
    if time_array(i)>worst
        worst=time_array(i);
    end
    if time_array(i)<best
        best=time_array(i);
    end
    if count_20==20
        count_20=1;
        soma=soma+time_array(i);
        media=soma/20;
        time_array_media_ms(count_tm)=media;
        time_array_best_ms(count_tm)=best;
        time_array_worst_ms(count_tm)=worst;
        worst=time_array(i);
        best=time_array(i);
        soma=0;
        media=0;
        count_tm=count_tm+1;
        continue
    end
    soma=soma+time_array(i);
    count_20=count_20+1;
end

```

## The Merkle-Hellman Cryptosystem

```
%% all algorithms graph
figure(1);
ib(1) = semilogy(n_array_ib,time_array_media_ib,"-","DisplayName","iterative brute force average","Color",[0.4 0.5 0.8]);
hold on;
ib(2) = semilogy(n_array_ib,time_array_best_ib,"-","DisplayName","iterative brute force best","Color",[0.4 0.5 0.3]);
ib(3) = semilogy(n_array_ib,time_array_worst_ib,"-","DisplayName","iterative brute force worst","Color",[0.5 0.1 0.7]);

rb(1) = semilogy(n_array_rb,time_array_media_rb,"-","DisplayName","brute force average","Color",[0 1 0]);
rb(2) = semilogy(n_array_rb,time_array_best_rb,"-","DisplayName","brute force best","Color",[0 1 1]);
rb(3) = semilogy(n_array_rb,time_array_worst_rb,"-","DisplayName","brute force worst","Color",[0.9 0 0.6]);

cb(1) = semilogy(n_array_cb,time_array_media_cb,"-","DisplayName","clever recursive brute force average","Color",[0 0.5 1]);
cb(2) = semilogy(n_array_cb,time_array_best_cb,"-","DisplayName","clever recursive brute force best","Color",[1 0.5 0.5]);
cb(3) = semilogy(n_array_cb,time_array_worst_cb,"-","DisplayName","clever recursive brute force worst","Color",[0 0 0]);

qh(1) = semilogy(n_array_qh,time_array_media_qh,"-","DisplayName","quick horowitz average","Color",[1 0 0]);
qh(2) = semilogy(n_array_qh,time_array_best_qh,"-","DisplayName","quick horowitz best","Color",[1 1 0]);
qh(3) = semilogy(n_array_qh,time_array_worst_qh,"-","DisplayName","quick horowitz worst","Color",[0.1 0 0.9]);

mh(1) = semilogy(n_array_mh,time_array_media_mh,"-","DisplayName","merge horowitz average","Color",[0.1 0.5 0.1]);
mh(2) = semilogy(n_array_mh,time_array_best_mh,"-","DisplayName","merge horowitz best","Color","black");
mh(3) = semilogy(n_array_mh,time_array_worst_mh,"-","DisplayName","merge horowitz worst","Color",[1 1 0.2]);

qs(1) = semilogy(n_array_qs,time_array_media_qs,"-","DisplayName","quick Schroeppele average","Color",[0.1 0.5 0.1]);
qs(2) = semilogy(n_array_qs,time_array_best_qs,"-","DisplayName","quick Schroeppele best","Color",[0.3 0.6 1]);
qs(3) = semilogy(n_array_qs,time_array_worst_qs,"-","DisplayName","quick Schroeppele worst","Color",[0.3 1 1]);

ms(1) = semilogy(n_array_ms,time_array_media_ms,"-","DisplayName","merge Schroeppele average","Color",[0.8 0.4 0]);
ms(2) = semilogy(n_array_ms,time_array_best_ms,"-","DisplayName","merge Schroeppele best","Color",[0.8 0 0]);
ms(3) = semilogy(n_array_ms,time_array_worst_ms,"-","DisplayName","merge Schroeppele worst","Color",[0.4 0 0.9]);

grid on;
title('Titulo para o gráfico');
hold off;
lgd=legend;
```