



Chapter 7: Transport Layer



Introduction to Networking

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Chapter 7: Objectives

- Describe the purpose of the transport layer in managing the transportation of data in end-to-end communication.
- Describe characteristics of the TCP and UDP protocols, including port numbers and their uses.
- Explain how TCP session establishment and termination processes facilitate reliable communication.
- Explain how TCP protocol data units are transmitted and acknowledged to guarantee delivery.
- Explain the UDP client processes to establish communication with a server.
- Determine whether high-reliability TCP transmissions, or non-guaranteed UDP transmissions, are best suited for common applications.

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7.1: Transport Layer Protocols



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Role of the Transport Layer

The transport layer is responsible for establishing a temporary communication session between two applications and delivering data between them.

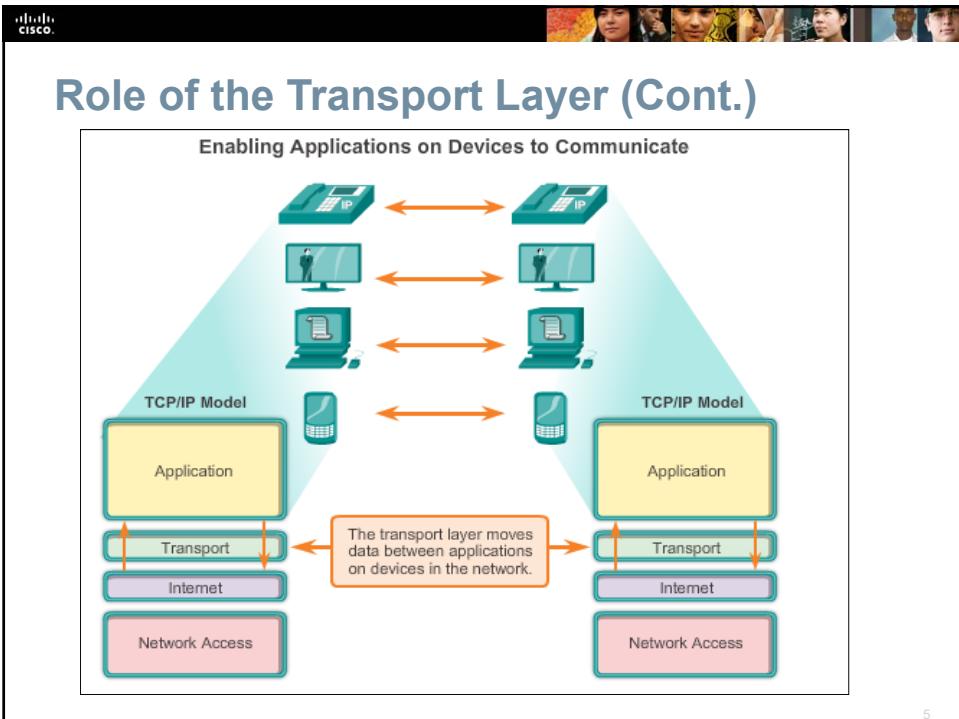
TCP/IP uses two protocols to achieve this:

- Transmission Control Protocol (TCP)
- User Datagram Protocol (UDP)

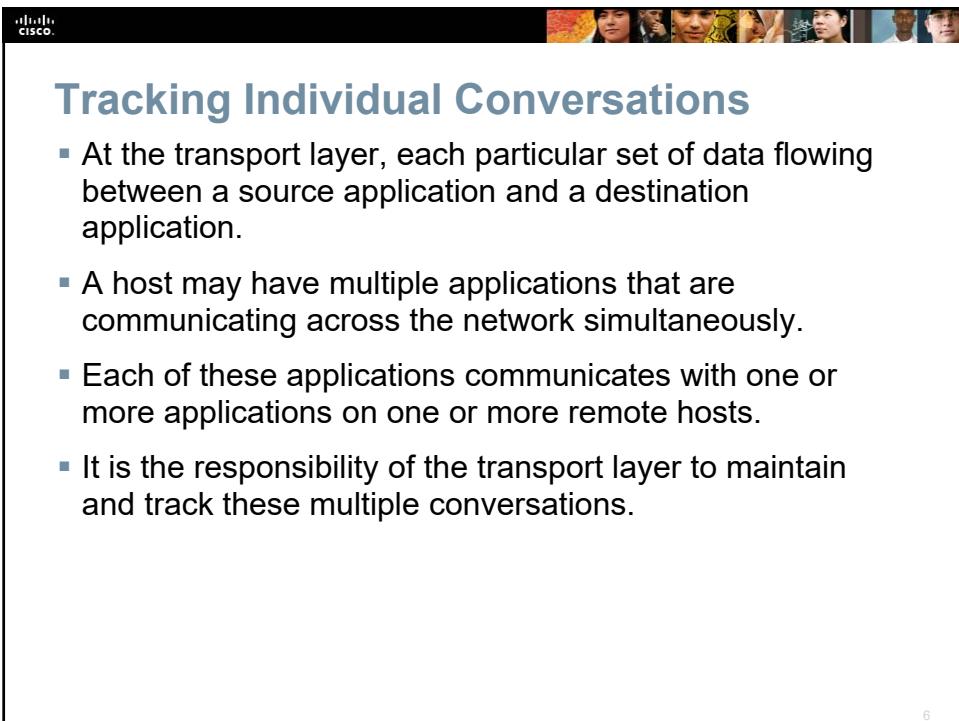
Primary Responsibilities of Transport Layer Protocols

- Tracking the individual communication between applications on the source and destination hosts
- Segmenting data for manageability and reassembling segmented data into streams of application data at the destination
- Identifying the proper application for each communication stream

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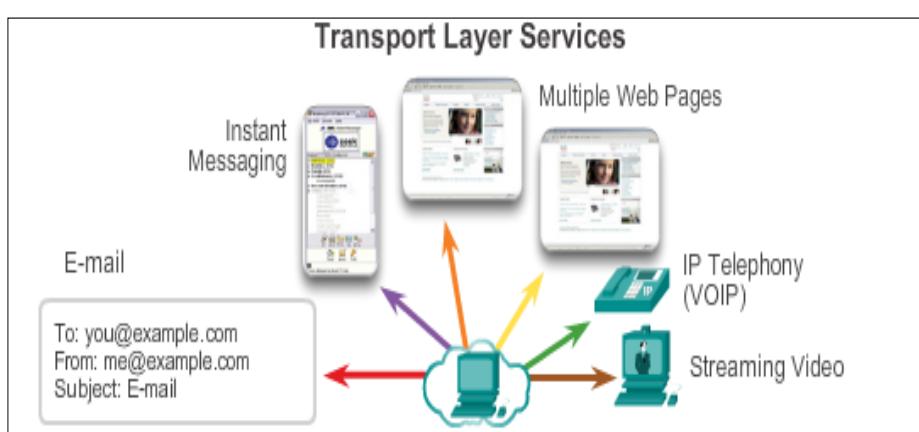
Segmenting the Data

- Most networks have a limitation on the amount of data that can be included in a single packet.
- Transport layer protocols have services that segment the application data into blocks of data that are an appropriate size.
- Header added to each segment to identify it. This header is used to track the data stream.
- At the destination, the transport layer must be able to reconstruct the pieces of data into a complete data stream.

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Conversation Multiplexing (Cont.)



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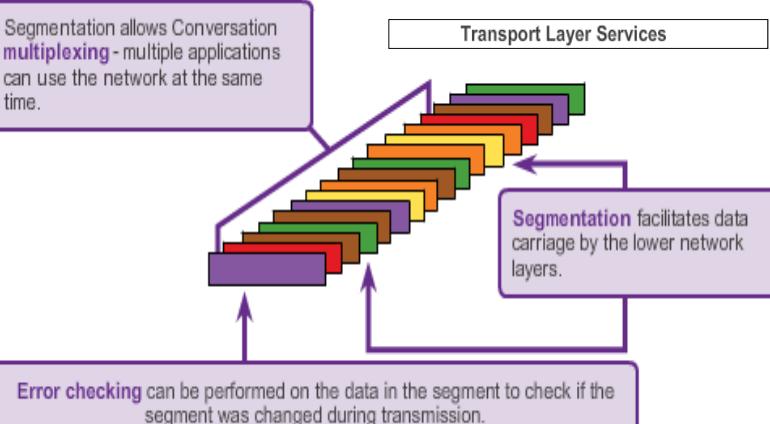
Conversation Multiplexing

- There may be many applications or services running on each host in the network.
- Enables many different communications, from many different users, to be interleaved (multiplexed) on the same network, at the same time.
- Provides the means to both send and receive data when running multiple applications.
- To pass data streams to the proper applications, the transport layer must identify the target application.
- To accomplish this, the transport layer assigns each application an identifier. This identifier is called a **port number**.

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Conversation Multiplexing (Cont.)



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Transport Layer Reliability

TCP/IP provides two transport layer protocols, **TCP and UDP**.

TCP

- Provides reliable delivery ensuring that all of the data arrives at the destination.
- Uses acknowledged delivery and other processes to ensure delivery
- Makes larger demands on the network – more overhead.

UDP

- Provides just the basic functions for delivery – no reliability.
- Less overhead.

TCP or UDP

- There is a trade-off between the value of reliability and the burden it places on the network.
- Application developers choose the transport protocol based on the requirements of their applications.

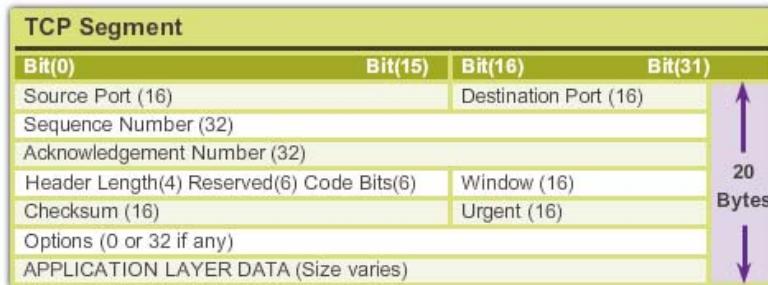
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Introducing TCP



- Connection-oriented – Creates a session between the source and destination
- Reliable delivery – Retransmits lost or corrupt data
- Ordered data reconstruction – Reconstructs numbering and sequencing of segments
- Flow control – Regulates the amount of data transmitted
- Stateful protocol – Tracks the session



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Introducing UDP

The diagram illustrates UDP characteristics and applications. It features a central computer monitor icon with four arrows pointing to different applications:

- No Ordered Delivery:** Data is reconstructed in original order. Examples: IP Telephony (VoIP) and Streaming Video.
- Unreliable Delivery:** Any segments lost are not resent. Example: Streaming Video.
- Connectionless:** No session establishment. Example: IP Telephony (VoIP).
- No Flow Control:** No congestion management. Example: Streaming Video.

Applications that use UDP:

- Domain Name System (DNS)
- Video Streaming
- VoIP

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Separating Multiple Communications

TCP and UDP use port numbers to differentiate between applications.

The diagram shows a computer icon at the top connected by a purple bracket to three application windows below: Electronic Mail, HTML Page, and Internet Chat. Arrows point from these applications to their respective protocols: POP3, HTTP, and IM. These protocols are shown above a transport layer icon. Below the transport layer are application ports and data fields. The port numbers 110, 80, and 531 are assigned to the respective application ports. A note at the bottom states: "Data for different applications is directed to the correct application because each application has a unique port number."

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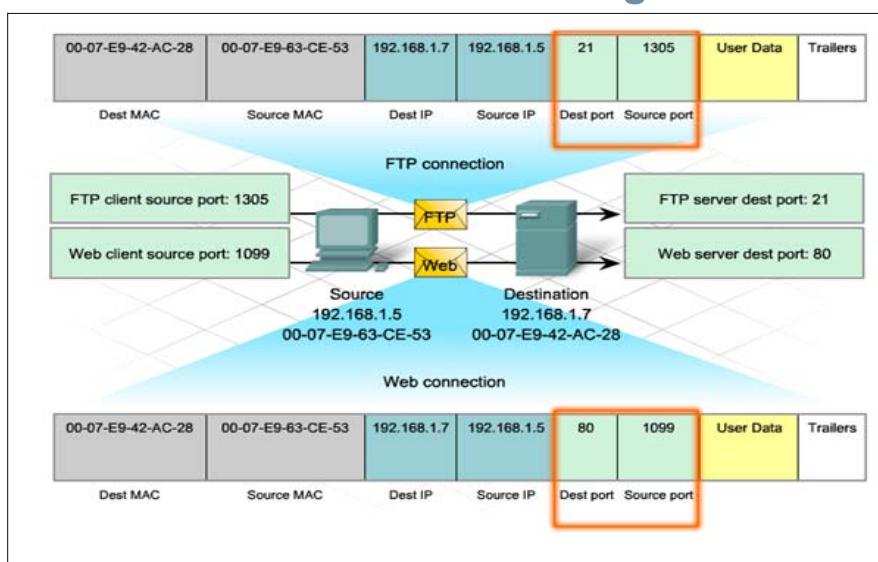
TCP and UDP Port Addressing

- When a message is delivered using either TCP or UDP, the protocols and services requested are identified by a port number.
- A port is a numeric identifier within each segment that is used to keep track of specific conversations and destination services requested.
- Every message that a host sends contains both a **source** and **destination** port. The source and destination ports are placed within the segment.
- The segments are then encapsulated within an IP packet. The IP packet contains the IP address of the source and destination.
- The combination of the source and destination IP addresses and the source and destination port numbers is known as a **socket**.
- The socket is used to identify the server and service being requested by the client.

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TCP and UDP Port Addressing



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Types of Port Numbers

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There are different types of port numbers, as shown in Figure:

Well-known Ports (Numbers 0 to 1023)

- These numbers are reserved for services and applications.
- They are commonly used for applications such as HTTP (web server), Internet Message Access Protocol (IMAP)/Simple Mail Transfer Protocol (SMTP) (email server) and Telnet.

Registered Ports (Numbers 1024 to 49151)

- These port numbers are assigned to user processes or applications.

Dynamic or Private Ports (Numbers 49152 to 65535)

- Also known as ephemeral ports, these are usually assigned dynamically to client applications when the client initiates a connection to a service.

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TCP and UDP Port Addressing (Cont.)

Port Numbers

Port Number Range	Port Group
0 to 1023	Well Known (Contact) Ports
1024 to 49151	Registered Ports
49152 to 65535	Private and/or Dynamic Ports

Registered TCP Ports:

- 1863 MSN Messenger
- 2000 Cisco SCCP (VoIP)
- 8008 Alternate HTTP
- 8080 Alternate HTTP

Well Known TCP Ports:

- 21 FTP
- 23 Telnet
- 25 SMTP
- 80 HTTP
- 110 POP3
- 194 Internet Relay Chat (IRC)
- 443 Secure HTTP (HTTPS)

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TCP and UDP Port Addressing (Cont.)

Registered UDP Ports: 1812 RADIUS Authentication Protocol 5004 RTP (Voice and Video Transport Protocol) 5040 SIP (VoIP)	Well Known UDP Ports: 69 TFTP 520 RIP
Registered TCP/UDP Common Ports: 1433 MS SQL 2948 WAP (MMS)	Well Known TCP/UDP Common Ports: 53 DNS 161 SNMP 531 AOL Instant Messenger, IRC

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TCP and UDP Port Addressing (Cont.)

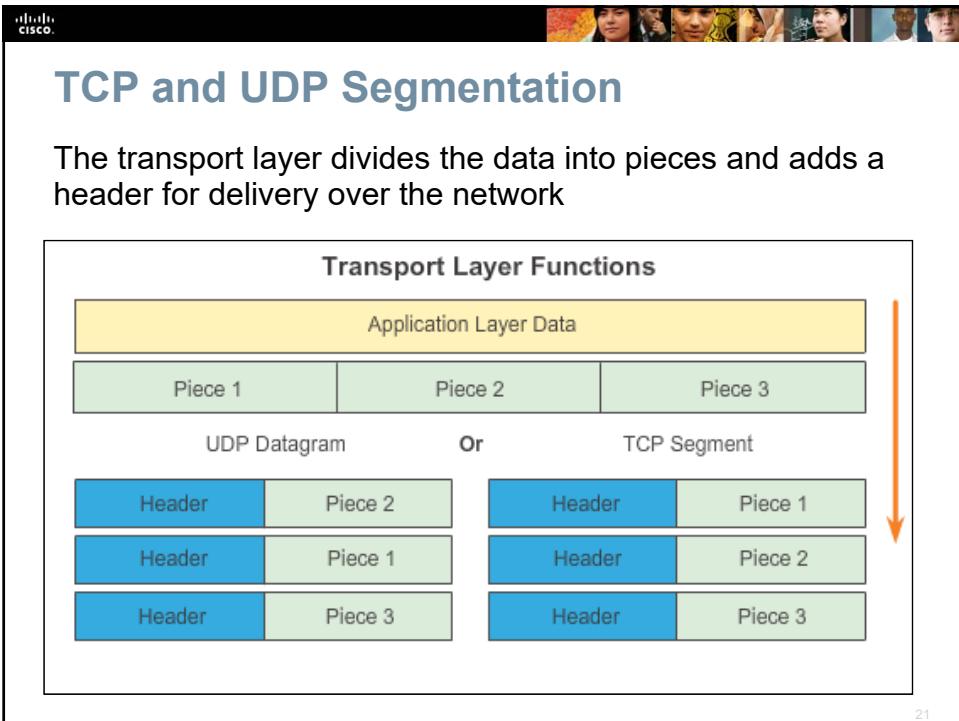
Netstat is used to examine TCP connections that are open and running on a networked host.

```
C:\>netstat
Active Connections

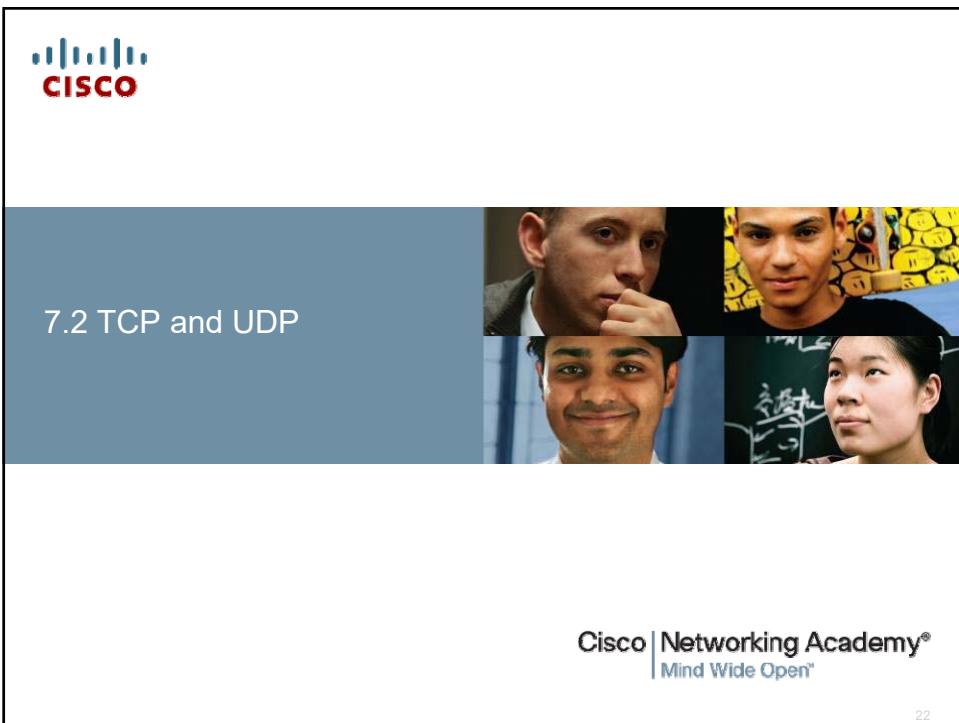
Proto  Local Address      Foreign Address          State
TCP    kenpc:3126        192.168.0.2:netbios-ssn  ESTABLISHED
TCP    kenpc:3158        207.138.126.152:http   ESTABLISHED
TCP    kenpc:3159        207.138.126.169:http   ESTABLISHED
TCP    kenpc:3160        207.138.126.169:http   ESTABLISHED
TCP    kenpc:3161        sc.msn.com:http       ESTABLISHED
TCP    kenpc:3166        www.cisco.com:http    ESTABLISHED

C:\>
```

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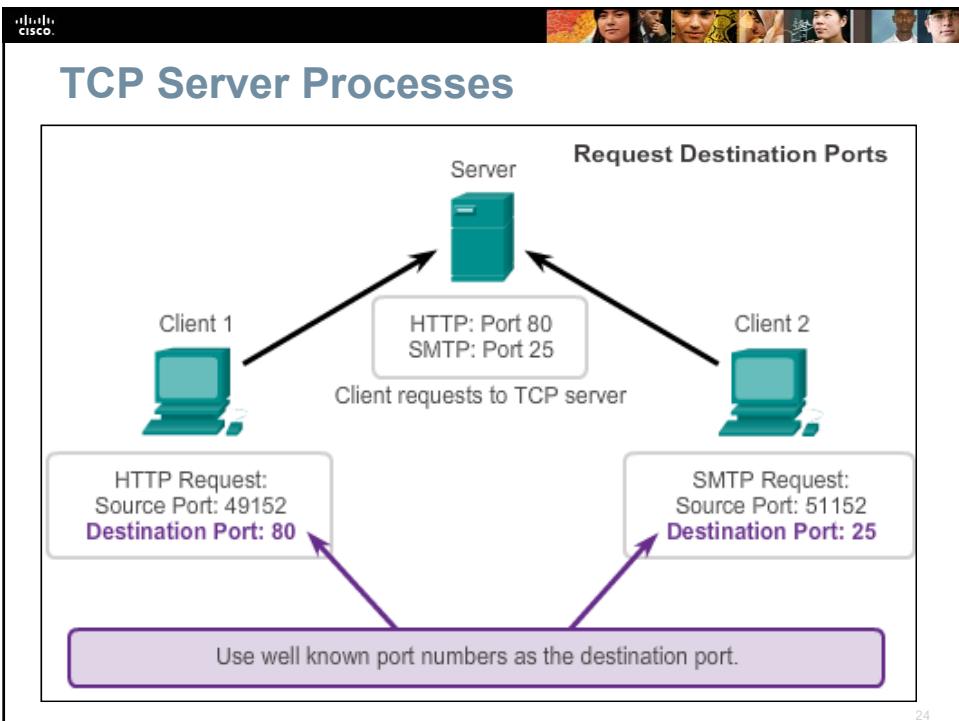
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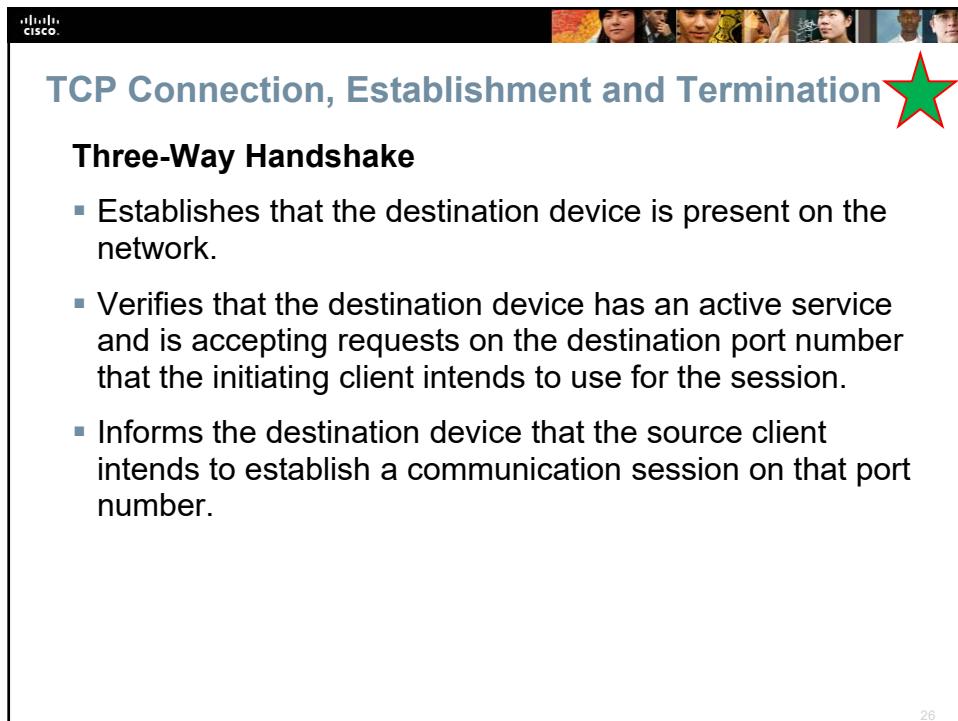
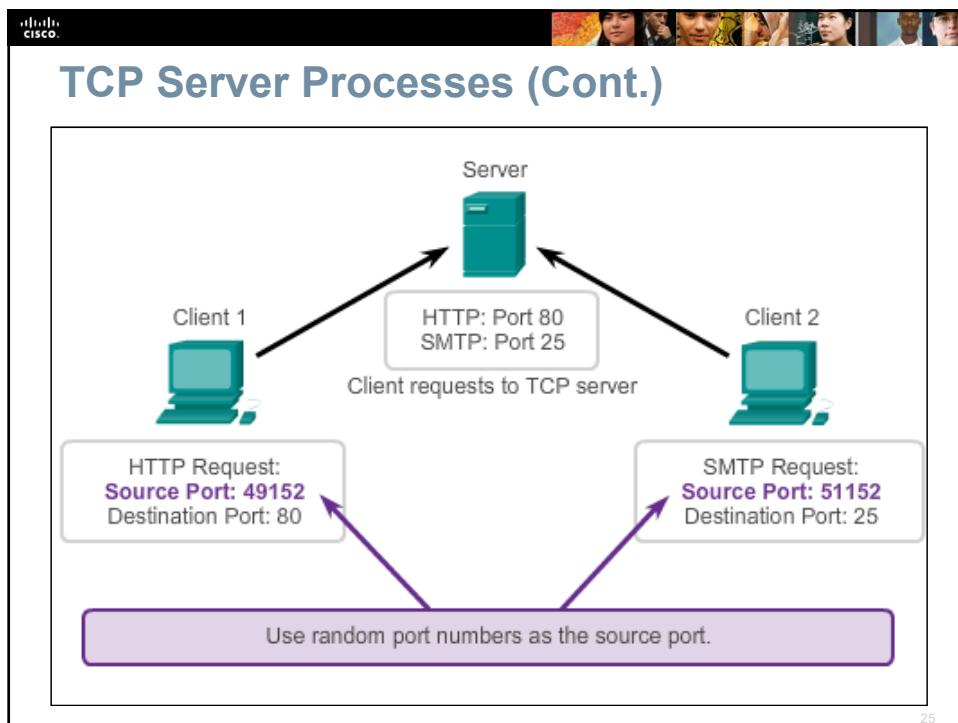
TCP Server Processes

- Each application process running on the server is configured to use a port number, either by default or manually by a system administrator.
- An individual server cannot have two services assigned to the same port number within the same transport layer services.
- There can be many simultaneous ports open on a server, one for each active server application.

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TCP Three-Way Handshake – Step 1

Step 1: The initiating client requests a client-to-server communication session with the server

TCP 3-Way Handshake (SYN)

A protocol analyzer shows initial client request in frame 10

TCP segment in this frame shows:

- SYN flag set to validate an Initial Sequence Number
- Randomized sequence number valid (relative value is 0)
- Random source port 1061
- Well-known destination port is 80 (HTTP port) indicates web server (httpd)

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TCP Three-Way Handshake – Step 2

Step 2: The server acknowledges the client-to-server communication session and requests a server-to-client communication session.

TCP 3-Way Handshake (SYN, ACK)

A protocol analyzer shows server response in frame 11

Frame 11: 52 bytes on wire (496 bits), 62 bytes captured
 Ethernet II, Src: cisco_63:74:a0 (00:0f:24:63:74:a0), Dst: 10.1.1.1 (0.0.0.0)
 Internet Protocol Version 4, Src: 192.168.254.254 (192.168.254.254), Dst: 10.1.1.1 (0.0.0.0)
 Transmission Control Protocol, Src Port: http (80), Dst Port: http (80)
 Source port: http (80)

• ACK flag set to indicate a valid Acknowledgement number
 • Acknowledgement number response to initial sequence number as relative value of 1
 • SYN flag set to indicate the Initial Sequence Number for the server to client session
 • Destination port number of 1061 corresponding to the clients source port
 • Source port number of 80 (HTTP) indicating the web server service (httpd)

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TCP Three-Way Handshake – Step 3

Step 3: The initiating client acknowledges the server-to-client communication session.

TCP 3-Way Handshake (ACK)

No.	Time	Source	Destination
13	16.303490	10.1.1.1	192.168.254.254
11	16.304896	192.168.254.254	10.1.1.1
12	16.304925	10.1.1.1	192.168.254.254
13	16.305153	10.1.1.1	192.168.254.254
14	16.307875	192.168.254.254	10.1.1.1

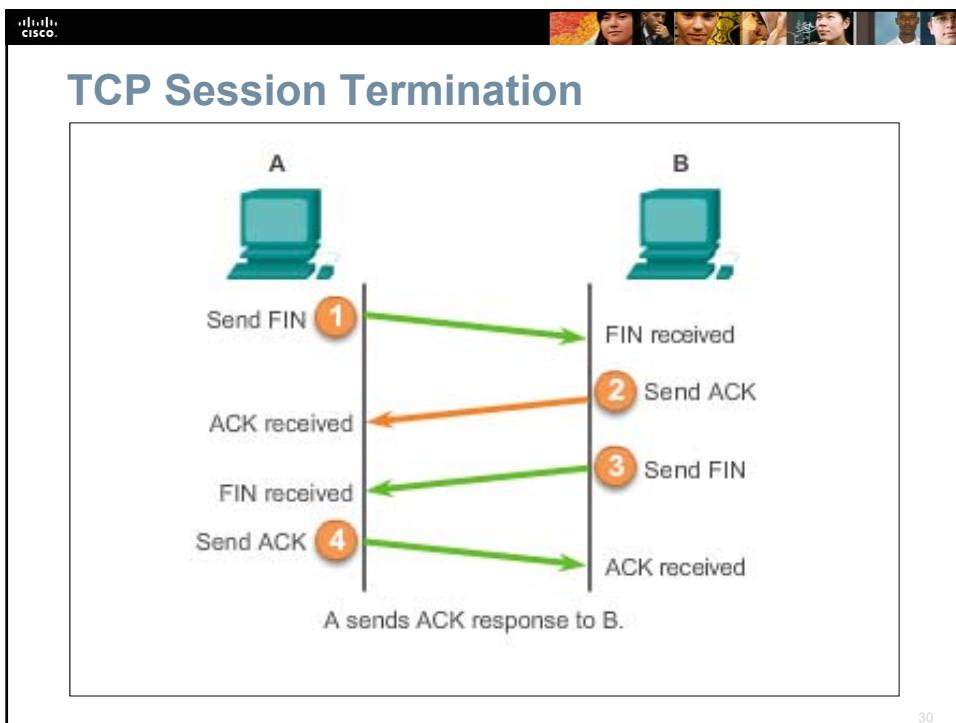
Frame 12: 54 bytes on wire (432 bits), 54 bytes captured
 Ethernet II, Src: vmware_kbe:62:88 (00:50:56:be:62:88)
 Internet Protocol Version 4, Src: 10.1.1.1 (10.1.1.1)
 Transmission Control Protocol, Src Port: k'osk (1061)

A protocol analyzer shows client response to session in frame 12

The TCP segment in this frame shows:

- ACK flag set to indicate a valid Acknowledgement number
- Acknowledgement number response to initial sequence number as relative value of 1
- Source port number of 1061 to corresponding
- Destination port number of 80 (HTTP) indicating the web server service (httpd)

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TCP Reliability – Ordered Delivery

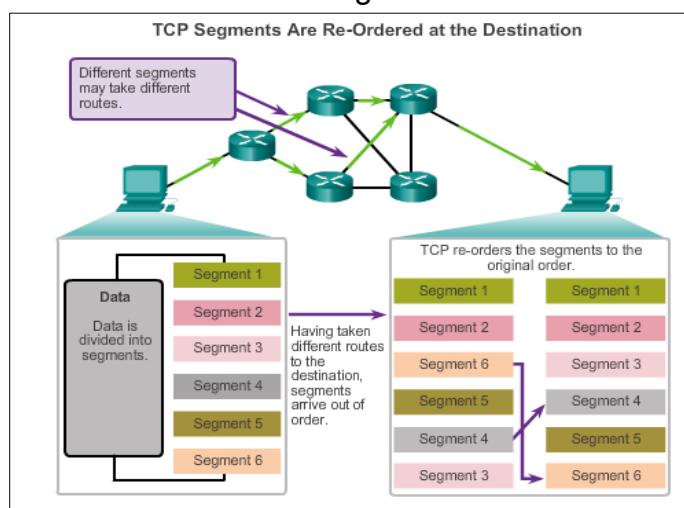
- When services send data using TCP, segments may arrive at their destination out of order.
- During session setup, an initial sequence number (ISN) is set.
- As data is transmitted during the session, the sequence number is incremented by the number of bytes that have been transmitted.
- This data byte tracking enables each segment to be uniquely identified and acknowledged. Missing segments can be identified.
- Segment sequence numbers enable the reliability by indicating how to reassemble and reorder received segments, as shown in the figure.

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TCP Reliability – Ordered Delivery

Sequence numbers are used to reassemble segments into their original order.



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Acknowledgement and Window size

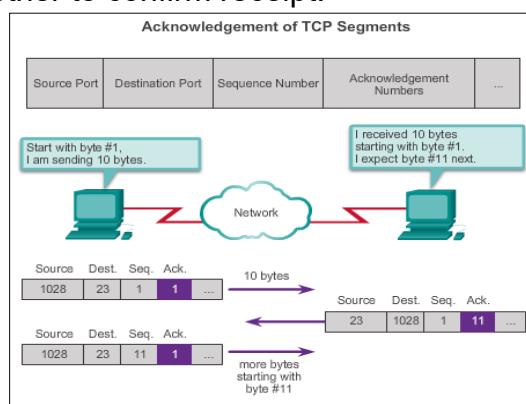
- TCP ensures that each segment reaches its destination.
The sequence (SEQ) number and acknowledgement (ACK) number are used together to confirm receipt of the bytes of data.
- The SEQ number indicates the relative number of bytes that have been transmitted in this session.
- TCP uses the ACK number sent back to the source to indicate the next byte that the receiver expects to receive.
- The source is informed that the destination has received all bytes indicated by the ACK number.
- The amount of data that a source can transmit before an acknowledgement must be received is called the **window size**, which enables the management of lost data and flow control.

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Acknowledgement and Window Size

The sequence number and acknowledgement number are used together to confirm receipt.



The window size is the amount of data that a source can transmit before an acknowledgement must be received.

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TCP Flow Control- Window Size and Acknowledgements

- Flow control helps maintain the reliability of TCP transmission by adjusting the rate of data flow between source and destination for a given session.
- Flow control is accomplished by limiting the amount of data segments forwarded by requiring acknowledgments of receipt prior to sending more.
- TCP uses window sizes to attempt to manage the rate of transmission to the maximum flow that the network and destination device can support, while minimizing loss and retransmissions.

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TCP Flow Control- Window Size and Acknowledgements

TCP Segment Acknowledgement and Window Size			
Sender		Receiver	
Window size = 3000			
Sequence number 1	1500 bytes	Receive 1 - 1500	
Sequence number 1501	1500 bytes	Receive 1501 - 3000	
Receive Acknowledge		Acknowledgement number 3001	
Sequence number 3001	1500 bytes	Receive 3001 - 4500	
Sequence number 4501	1500 bytes	Receive 4501 - 6000	
Receive Acknowledge		Acknowledgement number 6001	

The **window size** determines the number of bytes sent before an acknowledgement is expected.

The **acknowledgement** number is the number of the next expected byte.

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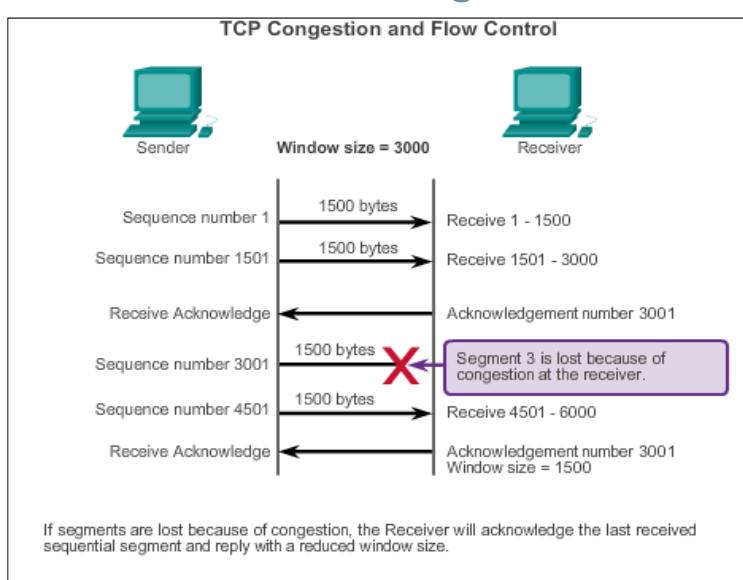
TCP Flow Control – Congestion Avoidance

- When network resources are constrained, TCP can reduce the window size to require that received segments be acknowledged more frequently.
- As shown in the figure, if a receiving host has congestion, it may respond to the sending host with a segment that specifies a reduced window size.
- After a period of transmission with no data losses or constrained resources, the receiver begins to increase the window field, which reduces the overhead on the network.

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TCP Flow Control – Congestion Avoidance



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UDP Low Overhead vs. Reliability

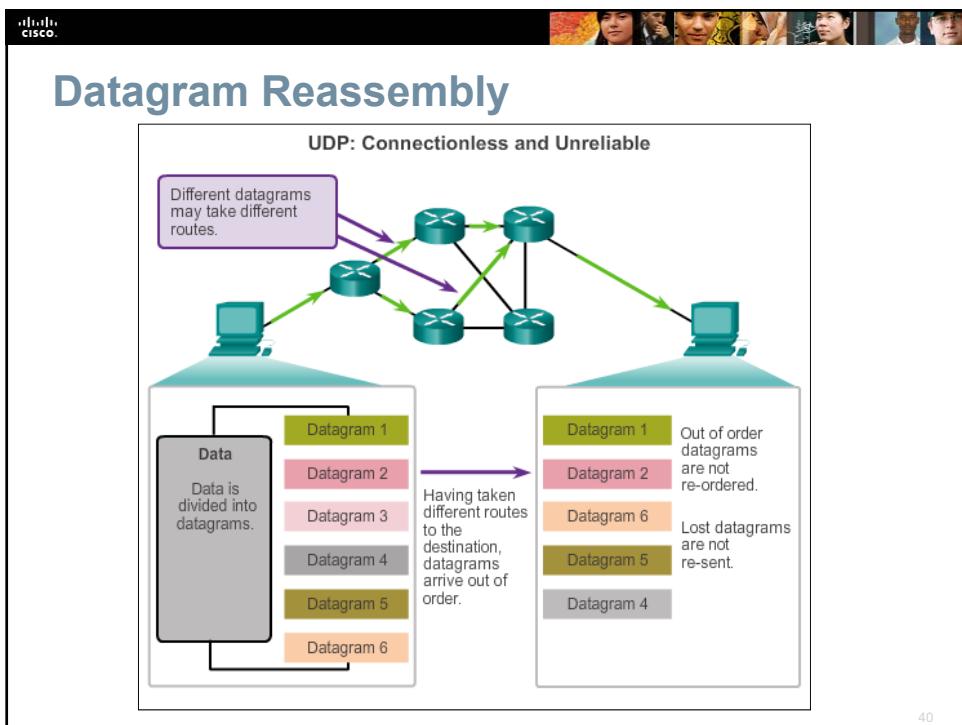
UDP

- Simple protocol that provides the basic transport layer function
- Used by applications that can tolerate small loss of data
- Used by applications that cannot tolerate delay

Used by

- DNS
- Simple Network Management Protocol (SNMP)
- Dynamic Host Configuration Protocol (DHCP)
- Trivial File Transfer Protocol (TFTP)
- IP telephony or VoIP
- Online games

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UDP Server and Client Processes

- UDP-based server applications are assigned well-known or registered port numbers.
- UDP client process randomly selects port number from range of dynamic port numbers as the source port.

Clients Sending UDP Requests

Server

Client 1

Client 2

Client 1 DNS Request:
Source Port 49152
Destination Port 53

Client 2 RADIUS User Authentication Request:
Source Port 51152
Destination Port 1812

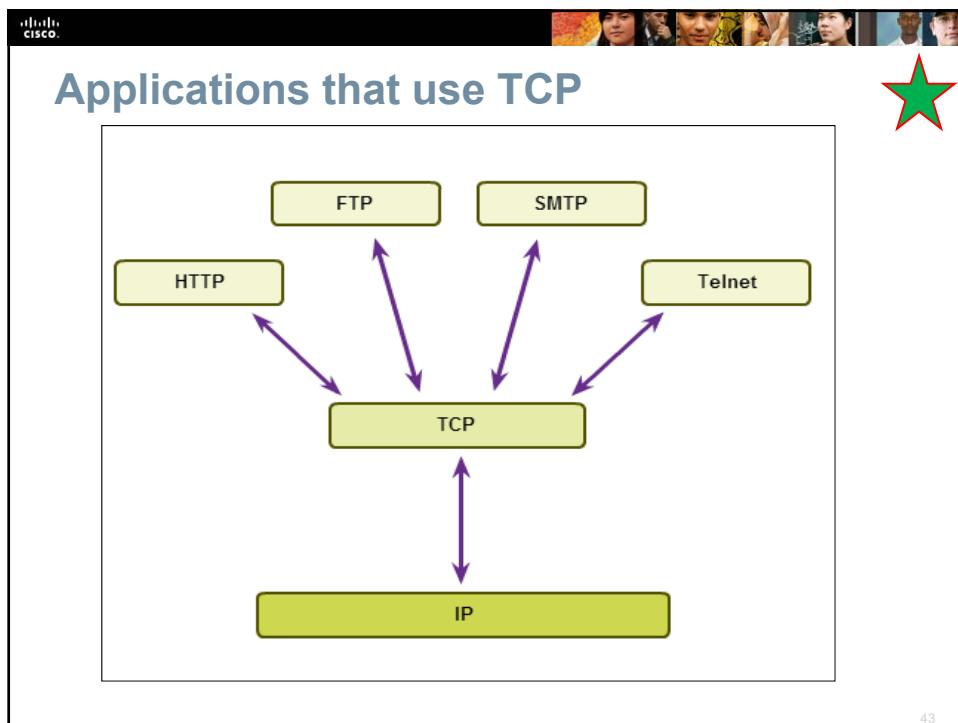
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Applications that use TCP

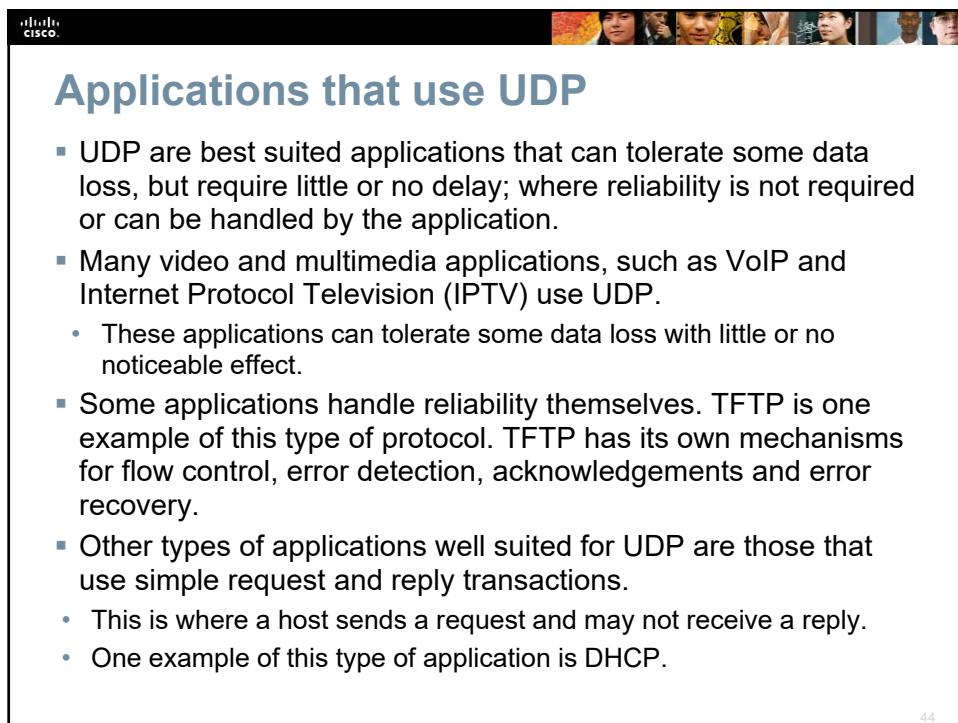
- TCP best suited for applications that need reliable transport and can tolerate some delay.
- As shown in the figure, some examples of well-known applications that use TCP include:
 - Hypertext Transfer Protocol (HTTP)
 - File Transfer Protocol (FTP)
 - Simple Mail Transfer Protocol (SMTP)
 - Telnet



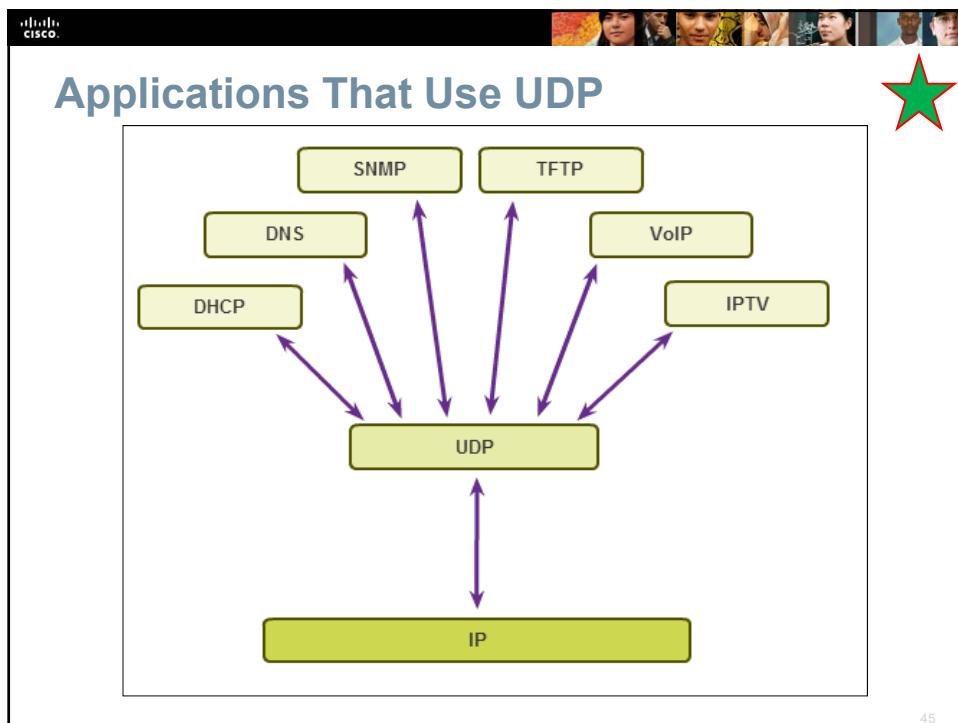
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Chapter 7: Summary

In this chapter, you learned:

- The role of the transport layer is to provide three main services: multiplexing, segmentation and reassembly, and error checking. It does this by:
 - Dividing data received from an application into segments.
 - Adding a header to identify and manage each segment.
 - Using the header information to reassemble the segments back into application data.
 - Passing the assembled data to the correct application.
- How TCP and UDP operate and which popular applications use each protocol.
- Ports provide a “tunnel” for data to get from the transport layer to the appropriate application at the destination.

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