

Introduction to Parallel Programing with MPI Part 1

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Outline

Part 1

- Basics of parallel programming
- Basics of MPI
- Point to point communication
- Blocking vs. non-blocking calls
- Collective communication

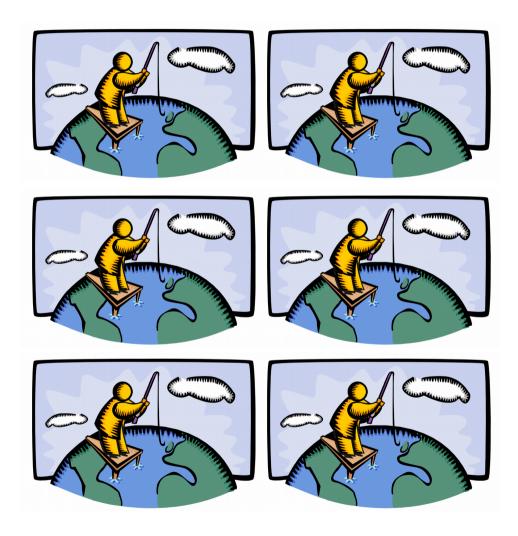
Part 2

- MPI file I/O
- Custom data types
- Communicators
- One sided communication



Parallelism

Parallelism means doing multiple things at the same time: you can get more work done at the same time





Less Fish More Fish



Terminology

- **Threads**: Execution sequences that share a single memory area (address space)
- Process: An execution sequence with its own address space
- Shared memory parallelism/ multithreading: parallelism via multiple threads
- **Distributed memory parallelism/ multiprocessing**: parallelism via multiple processes
- High Performance Computing (HPC): Combining the computing power of multiple machines to solve larger problems



Distributed Parallelism



- Suppose you want to do a jigsaw puzzle of 1000 pieces
- Lets assume it take you one hour to finish



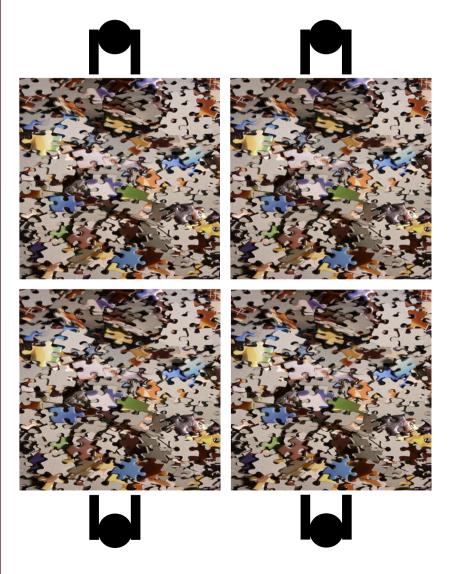
Distributed Parallelism



- Now, we divide the 1000 pieces into two tables
- Let Casey work in the other table on his half of the puzzle
- Both of you work independently: up to 2X speedup
- Communication at the end is expensive: you need to move two tables



Distributed Parallelism

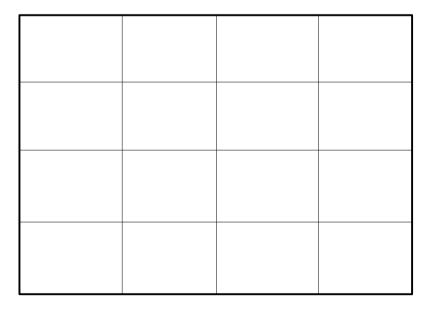


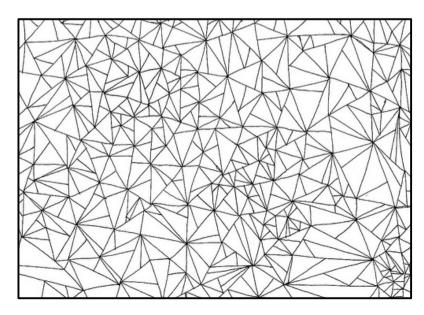
- More people joins!
- You can solve the puzzle faster
- Initial splitting of puzzle pieces may be harder (load balancing)
- Need more communication to finish the puzzle (communication overhead)



Load Balancing

- Distribute the work evenly among workers
 - Processes should not wait for others to finish: less efficient
- If a problem scales linearly, load balancing is easy
- Load balancing can be easy/hard depending on the problem





Easy

Hard



Communication Overhead

- Input data, parameter initialization, etc. need to be communicated to each node
- Outputs generated by each node should be gathered
- Communications during the calculation stage
 - Information about "boundaries" (eg. weather simulations)
 - N-body codes needs position of all other particles to calculate the total force on one particle
- There are ways to programatically minimize communication overhead in addition to use faster interconnects (hardware)



Why bother with all this?

- This seems like a lot of extra work!
- Why bother?
 - Done right, HPC can run your code much faster (few hours compared to week(s) on your desktop!) → get more done
 - You can solve bigger problems → better, exciting science
 - Today's HPC will be your desktop in 10 to 15 years! → stay ahead of your time



MPI The Message Passing Interface



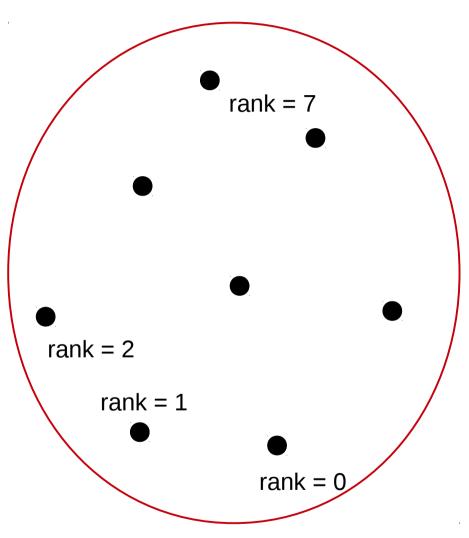
What is MPI?

- MPI is a language-independent communications protocol
 - Implements the framework (software) for communication between processes
- MPI is an API (Application Programming Interface)
 - Only the structure and behavior of each routine is defined
 - Implementation of routines are platform specific
- Different implementations: OpenMPI, MPICH, MVAPICH2, Intel MPI, ...
- MPI Consists of a header file, set of library routines, and a runtime environment
- An MPI code usually has server (master) and client (slave) sections
 - Only one node executes server sections
 - All nodes except server runs client sections



MPI basics

- Start and stop
 - MPI_Init
 - MPI_Finalize
- Environment awareness
 - MPI_Comm_size
 - MPI_Comm_rank



Communicator (MPI_COMM_WORLD)



MPI Program Structure (F90)

```
program my_mpi_test
                            header file
   IMPLICIT NONE
   include (mpif.h"
   [other includes]
   integer :: my rank, n procs, mpi err
   [other declarations]
   call MPI_Init(mpi_err)
   call MPI Comm Rank(MPI_COMM_WORLD, my_rank, mpi_err)
   call MPI Comm size (MPI COMM WORLD, n procs, mpi err)
   [do your work here]
   call MPI Finalize(mpi err)
end program my mpi test
```



MPI Program Structure (C)

```
#include <stdio.h>
#include "mpi.h"
[other includes]
                                                  Optional list of
                                                  arguments to
Int main (int argc, char* argv[])
                                                  be sent to all
                                                  processes
   int my rank, n procs, mpi err;
   mpi err = MPI Init(&argc, &argv);
   mpi err = MPI Comm Rank(MPI COMM WORLD, &my rank);
   mpi err = MPI Comm size(MPI COMM WORLD, &n procs);
   [do your work here]
   mpi_err = MPI_Finalize();
}
```



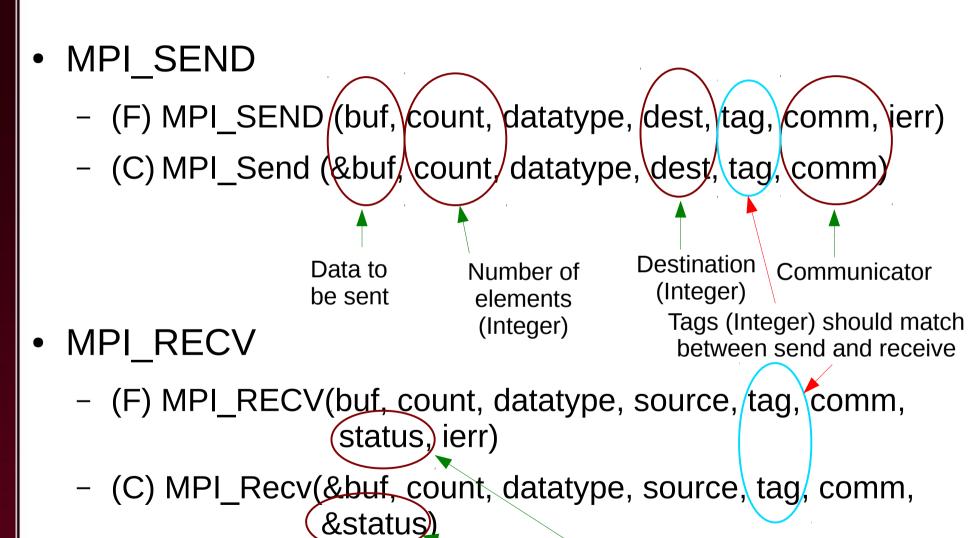
Compiling

- GNU, Intel, and PGI OpenMPI and MVAPICH2 compilers available at HPC
- Use modules to load selected compiler
 - Eg: module load gnu-openmpi
- MPI compiler wrapper scripts are used for compiling

Language	Script	Compiler	
С	mpicc	gcc/icc/pgcc	
C++	mpiCC mpic++ mpicxx	g++/icpc/pgCC	
Fortran	mpif90	gfortran/ifort/pgf90	
	mpif77	g77	



Point to point calls



Stores source & tag information



Point to point calls

MPI_SENDRECV

- (F) MPI_SENDRECV(sendbuf, sendcount, sendtype, dest, sendtag, recvbuf, recvcount, recvtype, source, recvtag, comm, status, ierr)
- (C) MPI_Sendrecv(&sendbuf, sendcount, sendtype, dest, sendtag, &recvbuf, recvcount, recvtype, source, recvtag, comm, &status)
- Can be used with individual SEND/RECV calls or another SENDRECV command



MPI Data types

Type of data sent or received should be explicitly specified

C Data Types		Fortran Data Types	
MPI_CHAR	char	MPI_CHARACTER	character
MPI_INT	int	MPI_INTEGER	integer
MPI_FLOAT	float	MPI_REAL	real
MPI_DOUBLE	double	MPI_DOUBLE_PRECISION	double precision

 Above is not a complete list. MPI standard also allows creation of user defined data types.



MPI Hello World

Codes available at

https://github.com/prasadhmd/MPIworkshop/tree/master/part1



Determinism

- MPI is nondeterministic: the arrival order of messages sent from two processes, A and B, to a third process, C, is not defined
- **Source** variable determines from which process a message is received. Using **MPI_ANY_SOURCE** as the source will let it receive from any sender
- Tag variable specifies which message from a given source to receive. Using MPI_ANY_TAG as the tag will let it receive messages with any tag



Deadlock Conditions

```
if (rank == 0) {
    MPI_Send(..., 1, tag1, MPI_COMM_WORLD);
    MPI_Recv(..., 1, tag2, MPI_COMM_WORLD, &status);
} else if (rank == 1) {
    MPI_Send(..., 0, tag2, MPI_COMM_WORLD);
    MPI_Recv(..., 0, tag1, MPI_COMM_WORLD, &status);
}
```

- Above code will NEVER complete! SEND from server will wait for a RECV request from rank 1 process while the it waits for a RECV request from server
- Following is one possible fix

```
if (rank == 0) {
    MPI_Send(..., 1, tag1, MPI_COMM_WORLD);
    MPI_Recv(..., 1, tag2, MPI_COMM_WORLD, &status);
} else if (rank == 1) {
    MPI_Recv(..., 0, tag1, MPI_COMM_WORLD, &status);
    MPI_Send(..., 0, tag2, MPI_COMM_WORLD);
}
```



MPI Communication types

Blocking Calls

- A call will not "return" until it is "complete"
 - A blocking call will pause the program execution on the sender process(es) until the receiving process(es) receive the complete message
- "Safe" but can slow down the program execution

Non-Blocking Calls

- Only instructs MPI to carry out the communication: No waiting for the data transfer to be complete
- "Fast" but synchronization can be tricky



Non-Blocking point to point calls

Immediate send

- (F) MPI_ISEND(buf, count, datatype, dest, tag, comm, request, ierr)
- (C) MPI_Isend(&buf, count, datatype, dest, tag, comm, &request)

request (Integer type in Fortran, MPI_Request type in C): a handle for the communication to retrieve the status later

Immediate receive

- (F) MPI_IRECV(buf, count, datatype, source, tag, comm, request, ierr)
- (C) MPI_Irecv(&buf, count, datatype, source, tag, comm, &request)



program ring implicit none include 'mpif.h'

Integer :: numtasks, rank, next, prev, buf(2), ierr Integer :: stats(MPI_STATUS_SIZE,4), reqs(4)

call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)

prev = rank - 1
next = rank + 1
if (rank == 0) prev = numtasks - 1
if (rank == numtasks - 1) next = 0

Exchange data between nearest neighbors

call MPI_IRECV(buf(1), 1, MPI_INTEGER, prev, 1, MPI_COMM_WORLD, req.(2), ierr) call MPI_IRECV(buf(2), 1, MPI_INTEGER, next, 2, MPI_COMM_WORLD, reqs(2), ierr)

call MPI_ISEND(rank, 1, MPI_INTEGER, prev, 2, MPI_COMM_WORLD, reqs(3), ierr) call MPI_ISEND(rank, 1, MPI_INTEGER, next, 1, MPI_COMM_WORLD, reqs(4), ierr)

calMPI_WAITALL(4, reqs, stats, ierr)

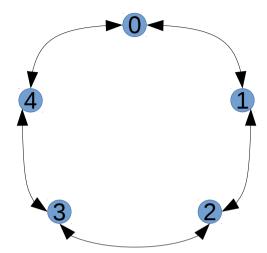
call MPI_FINALIZE(ierr)

end program ring

wait until all transfers are completed this is a **BLOCKING** command

Codes available at

https://github.com/prasadhmd/MPIworkshop/tree/master/part1





Non-Blocking tests

 It is sometimes useful to have a non-blocking way to check the status of a communication

```
(F) MPI_TEST(request, flag, status, ierr)
```

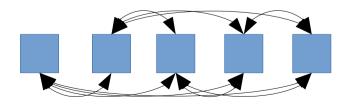
- (C) MPI_Test(&request, &flag, &status) False/True depending on in/completion

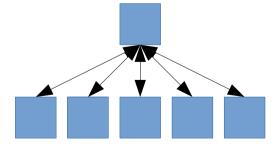
 MPI_TESTALL, MPI_TESTANY are two variants of MPI_TEST which checks an array of communications



Collective Communication

- Sometimes it is necessary to communicate between groups of processes
 - Data Movement: transfer initial values, gather individual results, etc.
 - Collective Computation: gather partial results and combine to get the final result
 - Synchronization: keep everyone in the same page
- Types of collective calls
 - One to all and all to one
 - All to all







Collective Communication



sender

- (F) MPI_BCAST(buffer, count, datatype, root, comm, ierr)
- (C) MPI_Bcast(&buffer, count, datatype, root, comm)



2 2 2

- Reduction
 - (F) MPI_REDUCE(sendbut) recvbut, count, datatype, op, root, comm, ierr)
 - (C) MPI_Reduce(&sendbuf) &recvbuf, count, datatype op, root, comm)

partial result



+

3 5 1 2 4 15 1 5 1 1 2 1 4

operation: MPI_SUM, MPI_PROD, MPI_MIN, ...

final result



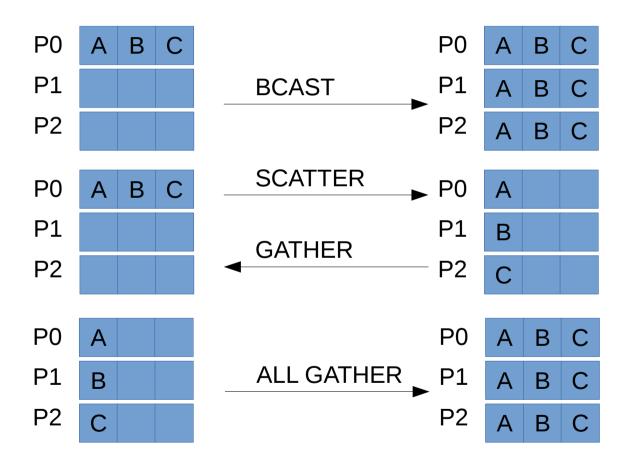
MPI Reduction Example

Codes available at

https://github.com/prasadhmd/MPIworkshop/tree/master/part1



Collective Communication





Collective Communication

- Collective communication calls can be faster than a set of send-receive calls
- Very useful for data parallelism
 - Same set of instructions operate on different sub sets of data
- There is significant communication overhead
- These are blocking calls
 - Non-blocking collective calls are also available



MPI Reduce example MPI Scatter/Gather example

Codes available at

https://github.com/prasadhmd/MPIworkshop/tree/master/part1



Parallel Programming Strategies



Parallel Programming Strategies

- Client-server
- Data Parallelism
- Task Parallelism
- Pipeline



Client-Server

- Server (master) decides what clients (slaves) do
- Embarrassingly parallel: The problem can easily be broken into roughly equal amounts of work per process and has very little communication (low communication overhead)
- Has near linear speedup and easy to program
- Eg: Monte Carlo methods widely used to simulate a physical phenomena or calculate an integral
 - Randomly generate large number of samples (realizations) of a phenomenon/equation and take the average over all samples
 - Simulation stops when the average value converges



Client-Server example Calculating Pi



Data Parallelism

- Each process does exactly same operations on a unique subset of data
- Most scientific problems involve calculus: solving differential equations etc.
- Numerically solving these equations over a large domain is very common
- Data parallelism can be applied to parallelize this type of problems
- Eg: CFD, Heat transfer, Weather prediction, etc.



Task Parallelism

- Each process does different operations on exactly same set of data
- Task parallelism is a widely used technique
- N body problem
 - N objects interacting with each other via forces: stars under gravity, molecules under electrostatic force etc.
 - Send properties of each object to all processes and let each process find the total force on a subset of particles
 - After each time step, use MPI_Allreduce
 - Applications: Cosmology, structural biology, machine learning



Pipeline Parallelism

- Each process does its work, passes its set of data to next process and receives next set of data from previous process
- All processes are connected to form a data pipeline
- Every process execute same tasks and results are passed to the "next" process and more data is received from the "previous" process
- Workers can be connected in a circular (closed) loop or linear (open)
- Matrix multiplication can be done in a pipeline