# **Assessment 3 - Computer Science**

Make sure that you always use notations consistent with lecture notes.

Marks: 15 marks, weighted 15% of your final mark.

Submission format: Adobe PDF (preferred) or MS Word Doc.

Submit to: Moodle.

Late submission penalty: No late submission allowed.

The deadline for assignment 3 is: Mon 5th, August 10:00 pm/

# Q1. Binary Search Trees (BST) (3 marks total)

- (1) (1 mark) Starting from an empty binary search tree what is the final binary search tree constructed by inserting the following ements in the specified order: {20, 40, 13, 4, 89, 21, 65, 15, 28, 46, 26, 36, 8,14,18}.
- (2) (1 mark) Draw a BST with the same elements from (q1.1) but has a maximum height of 3 instead.
- (3) (1 mark) Using the tree you built from (q.1.1), perform removing node {21} and draw the resulting BSV.

# Q2. Decision Trees (4 marks total)

Outlook	Temperature	Wind	Play Tennis		
Sunny	Hot	Weak	No		
Sunny	Hot	Strong	No		
Overcast	Hot	Weak	Ses		
Rain	Mild	Weak	Yes		
Rain	Cool	Weak	Yes		
Rain	Cool	Strong	No		
Overcast	Cool	Weak	Yes		
Sunny	Mild	Strong	Yes		
Sunny	O WIGHT O	Strong	Yes		
Overcast	Mild	Strong	No		

1) (3 marks) Computing the conditional entropy and the information gain for each predictor using the ID3 algorithm, give the steps.

(2) (1 mark) Build a decision tree using the computed information gain for the predictors.

## Q3. Bloom Filter (4 marks total)

UNSW Estate Management needs to check if a license plate might have Parking Permits quickly. It needs to check a large number of license plates every day. To boost efficiency, they employ a Bloom filter. They have a Bloom filter of size m = 10 and k = 3. The hash functions that they employ are publicly listed below:  $h_1$ ,  $h_2$ , and  $h_3$ .

$$h_1(plate\_num) = \left(\sum_{n=1}^{l} (c_n - 1)\right) mod m$$

$$h_2(plate\_num) = \left(\sum_{n=1}^{l} (c_n + h_1(plate\_num))\right) mod m$$

 $h_3(plate\_num) = (h_1(plate\_num) + h_2(plate\_num))$  m

### Notes:

- i. *plate\_num* denotes a single plate number entry.
- ii. *l* denotes the length of the plate number.
- iii.  $c_n$  denotes the n<sup>th</sup> character of a **plate\_num** e.g., for "VWX901",  $c_1 = 'V'$ ,  $c_2 = 'W'$ ,  $c_3 = 'X'$ ,  $c_4 = 9'$ ,  $c_5 = '0'$ , and  $c_6 = '1$ .
- iv. If  $c_n$  is not a number, consult the following table below.

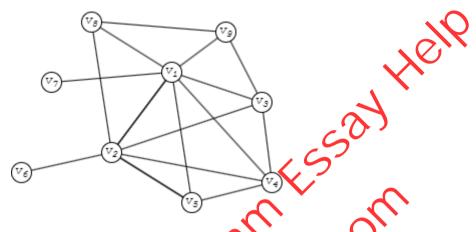
Α	В	С	D	Ε	F	G	Н	I	7	K	L	М
$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$			$\uparrow$	10	ノ	$\uparrow$	$\uparrow$	1
0	1	2	3	4	5	6	7	8	9	10	11	12
N	0	Р	Q	R	S	7	C)	V	W	Х	Y	Z
$\uparrow$	$\uparrow$	$\uparrow$	,1C	<b>&gt;</b>	0	ĴC		$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$
13	14	15	-16	17	18	19	20	21	22	23	24	25
	A	A B	N O P  ↑ ↑ ↑ ↑  N ↓ ↑ ↑	N O P Q  ↑ ↑ ↑ ↑ ↑	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$							

#### Tasks:

- (1) (2 marks) Given an empty bloom filter, initialize the bloom filter by inserting the following stream of 4 plate numbers with Parking Permit from  $S_1$  into the Bloom filter.  $S_1 = \{NSW 17 U CMD256", "XUV971", "KBM899"\}$ . Using the bloom filter to check plate numbers in  $S_2 = \{MNO121", "QRS247"\}$ , give the steps.
- (2) (1 mark) If the bloom filter size m=100 and there are 15 plates to initialize the bloom filter, what is the optimal number of hash functions? Compute the new probability of false positive.
- (3) mark) What can we do to decrease the false positive rate?

## Q4. K-cores (4 marks total)

The k-core of a graph is the *maximal* subgraph such that every vertex in the k-core has a degree of at least k. We provide the following graph G with 9 vertices.



- (1) (1 mark) Draw the 3-core of G.
- (2) (1 mark) Add one edge to the original graph 6 to make a 4-core. Justify your answer.
- (3) (2 marks) Another subgraph netric k-truss, is similar to k-core, the definition of k-truss is as follows.

**Definition**: A k-truss is a subgraph where every edge is part of at least k-2 triangles within the subgraph.

Note: In a graph G=(V,E), let V be the set of vertices and E be the set of edges. A triangle is a set of three vertices  $\{v1,v2,v3\}\subseteq V$  such that  $(v1,v2)\in E$ ,  $(v2,v3)\in E$ , and  $(v1,v3)\in E$ .

Draw the 4 truss of G, justify your answer.