## Design

## Part A: User Environments and Exception Handling

- Exercise 1: /kern/pmap.c: mem\_init()
   Just allocate space for envs and map it to [UENVS, UENVS + PTSIZE], set permission
   PTE\_U | PTE\_P
- Exercise 2: implement functions in /kern/env.c
  Implements functions by following the hints. It take a bit more time to make load\_icode correct, following the stop in /boot/main.c

## Part B: Page Faults, Breakpoints Exceptions, and System Calls

3. Exercise 4, 5: implement trap entry in /kern/trap.c and /kern/trapentry.S \_alltraps: push required information onto stack, by the reverse order defined in Trapframe struct then call trap

Some work for the challenge:

Inspired by the implements in xv6

In trapentry.S firstly, map the interrupt number defined in inc/trap.h with and handler(named "handler"+interrupt\_number) using the marco TRAPHANDLER and TRAPHANDLER\_NOEC. Then, make a data section in trapentry.S name vectors, each element in vectors refer to a handler above.

Trap.c: trap\_init(): using the vectors defined above. Map idt elements to vectors elements using SETGATE. Except T\_BRKPT, set the dpl to 0x3(user level), others all set to 0x0(super level)

Trap.c: trap\_dispatch: Just two kind of trap number need special dispatch. One for T\_BRKPT => monitor(), one for T\_PGFLT => page\_fault\_handler()

Trap.c: page\_fault\_handler(): should panic when kernel mode page fault happens, which is conditioned by  $(tf->tf_cs \& 0x3) == 0$ 

4. Exercise 6:

Sysenter\_handler: push %edi, %ebx, %ecx, %edx, %eax in order, call syscall, then move %ebp -> %ecx, %esi -> %edx, sysexit

MSRs: add a function wrmsr in inc/x86.h then enable sep cpu by adding lines in trap\_init\_percpu. Reference: <a href="https://lwn.net/Articles/18414/">https://lwn.net/Articles/18414/</a>

lib/syscall.c: syscall():

1). first push %esp, pop %ebp to modify %ebp indirectly

2). Then make a label just after sysenter, so before sysenter get the address after Sysenter by leal. Because gcc compile will fail when try to leal a label defined After, add %= after a label to let gcc generate unique label for you.

## 5. Exercise 7:

Thisenv is not inited. Just add one line `thisenv = envs + ENVX(sys\_getenvid()); ` in lib/libmain.c

6. Exercise 8: sys\_sbrk()

First add a field called <a href="mailto:env\_break">env\_break</a> in struct Env

In kern/syscall.c: sys\_sbrk(): just increase curenv->env\_break ROUNDUP(inc, PGSIZE) And map the new allocated space with region\_alloc

7. Exercise 9: breakpoint

Trap\_dispatch is described in exercise 4, 5

Here need implement three command: si, x and c

Register the three new command by adding them into commands in kern/monitor.c Then define three functions: mon\_si, mon\_x and mon\_c

- 1). Mon\_x: print the value of address va(in hex). So things to do is just parse the string Of the address to int and get the value by dereference(\*)
- 2), mon\_c: continue executing. Set tf->tf\_eflags to tf\_efalgs & (~FL\_TF) (remove the FL\_TF bits) and return -1
  - 3). Mon\_si: execute one instruction. Add the FL\_TF bit to tf->tf\_efalgs, then print the Info required just like mon\_backtrace
  - 8. Exercise 10

User\_mem\_check:

Create a function user\_mem\_check\_page to check permission of the page where va resides. It will fail (return -1) when va >= ULIM or permission not match

Then in user\_mem\_check() just check permission of pages from va to va+len

9. Exercise 12: evilhello2

Global variables:

Struct Gatedesc \*gd

Char gdt\_pgs[2\*PGSIZE]: mapped GDT in user space

Void (\*evil func)(): the evil function

Implements the ring0\_call:

- 1). Store GDT into dp using sgdt
- 2). Map dp.pd\_base to ROUNDUP(gdt\_pgs, PGSIZE)
- 3). Setup call gate with SETCALLGATE: gdt entry gd => call\_evil\_func
- 4). Lcall GD\_TSS0, \$0