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Project 1

1. Modular exponentiation:

The exponentiation performed over modulus is modular exponentiation. It is necessary to compute $x^y \bmod N$ for values of x , y , and N that are several hundred bits long. So, modular exponentiation technique helps us find the values of x , y , and N . The pseudocode for modular exponentiation is:

function modexp(x , y , N):

Input: Two n -bit integers x and N , an integer exponent y

Output: $x^y \bmod N$

$y \bmod N$

if $y = 0$: return 1

$z = \text{modexp}(x, y/2, N)$

if y is even:

return $z^2 \bmod N$

else:

return $x \cdot z^2 \bmod N$

(Reference: Book)

Now, my implementation of modular exponentiation in code(in python) is given by:

```
def mod_exp(x, y, N):
```

```
    if y == 0: return 1
```

```
    z = mod_exp(x, y // 2, N)           >>Recursively makes O(N) calls
```

```
    if y % 2 == 0:                       >>Each recursive calls make O(N^2)
```

```
        return (z * z) % N
```

```
    else:
```

```
        return x % N * (z * z) % N
```

Looking at the time complexity, this algorithm makes $O(n)$ recursive calls, and each of them takes $O(n^2)$ time, so the complexity is $O(n^3)$. So, the overall complexity is $O(N^3)$.

2. Fermat's Test algorithm for Primality test:

Fermat developed a way to test for primality using a probabilistic algorithm. This test is a primality test, giving a way to test if a number is a prime number, using Fermat's little theorem and modular exponentiation.

Fermat's Little Theorem states that if a is relatively prime to a prime number p , then $a^{p-1} \equiv 1 \bmod p$

The pseudocode for Fermat's algorithm is:

function primality(N):

Input: Positive integer N

Output: yes/no

Pick a positive integer $a < N$ at random

```

if  $a^{N-1} \equiv 1 \pmod{N}$ :
    return yes
else:
    return no

```

In a loop this algorithm is given as:

```

function primality2(N):
Input: Positive integer N
Output: yes/no
Pick positive integers  $a_1, a_2, \dots, a_k < N$  at random
if  $a_i^{N-1} \equiv 1 \pmod{N}$  for all  $i = 1, 2, \dots, k$ :
    return yes
else:
    return no

```

(Reference: Book)

My implementation of code in Python is as follows:

```

def fermat(N, k):
    for i in range(1, k):
        a = random.randint(2, N - 1)
        if mod_exp(a, N - 1, N) != 1:
            return 'composite'
    return 'prime'

```

Note that $a^{N-1} \equiv 1 \pmod{N}$

>> Making k recursive calls, requires $\log(N-1)$ complexity at each level

As seen here, the probability of error drops exponentially fast, and can be driven arbitrarily low by choosing k large enough. So, if it is prime it returns prime with a probability of $1 - \frac{1}{2^k}$. The complexity behind this algorithm: I think that since it loops for k times, the complexity must depend on k parameter and finding a^{N-1} requires $\log(N-1)$. Thus, This means that For the overall algorithm the time complexity seems to be $O(k \log^2 N)$.

3. Miller-Rabin primality test:

Miller-Rabin is also a probabilistic primality test that helps to determine if a number is composite or probably prime with some probability. In this algorithm, the process gradually processes by repeatedly taking the square root of a number that is $\equiv 1 \pmod{N}$.

My implementation of this algorithm in python is :

```

def miller_rabin(N, k):
    c=0
    d = N - 1
    while d % 2 == 0:
        c += 1
        d //= 2
    for i in range(k):

```

>>k times the first loop means it will go k loops for the worst case

```

a = random.randint(2, N-2)
b = mod_exp(a, d, N)    >>>Log (N-1) to calculate
if b!=1 and b+1 !=N:
    for j in range(1,c):
        if mod_exp(b,2, N)!= 1:  >>>repeating squaring thus the complexity is O(k(logN)^3)
            return 'composite'
        elif b!= N-1:
            break
    else:
        return 'composite'
return 'prime'

```

This algorithm uses most of the code of the Fermat's test.

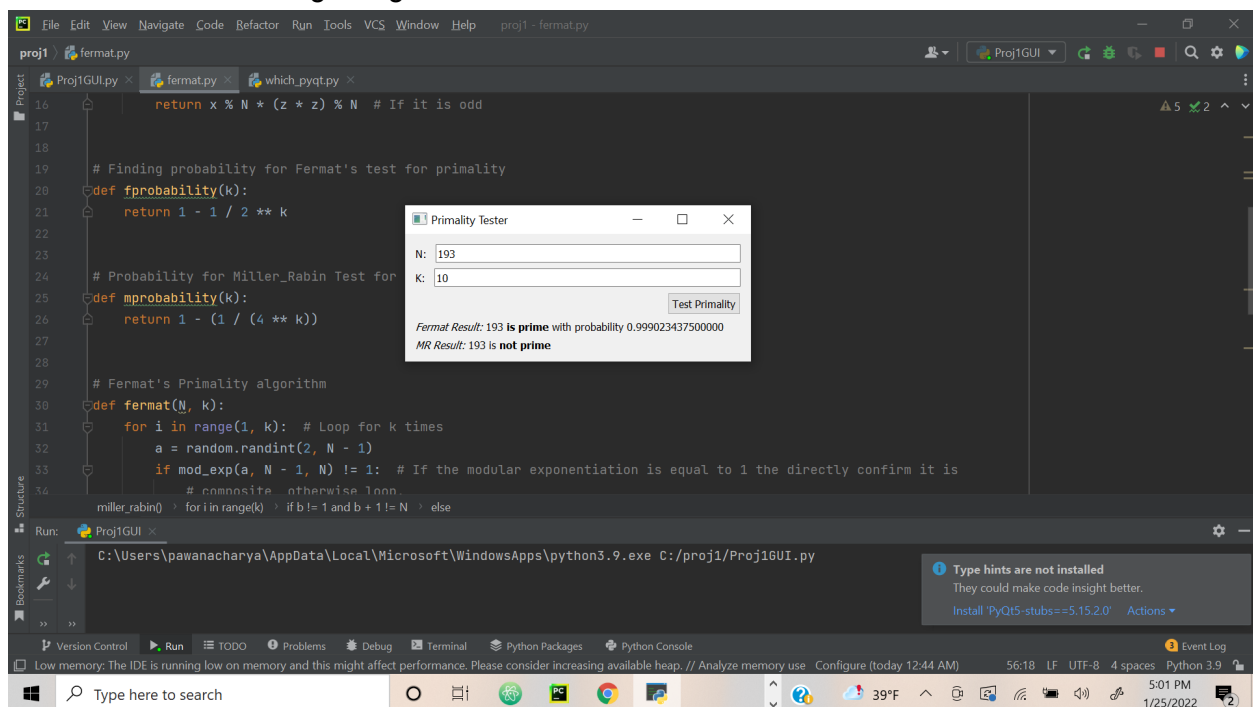
I think that the Complexity of this algorithm is : $O(k(\log N)^3)$

Agreement and Disagreement of the results on both algorithms:

For some inputs, the Rabin Miller test doesn't agree with Fermat's test. For eg: 193

But for most of the other inputs I entered, both agree with variety in probability.

One screenshot showing disagreement is :



Discussion of probability in both cases:

For composites, for at least $3/4$ of the possible choices for a , this will not be the case---either the initial test will not equal $1 \pmod{N}$.

The probability for being prime in Fermat's test is : $1 - 1/2^k$

Meaning that we can be sure that it can be prime with given probability.

Similarly, For Miller_Rabin's test: $1 - 1/4^k$.

Some screenshot of my tests are :

