

CS 161 HW 5

1) $P \Rightarrow \neg Q : \neg P \vee \neg Q$

$Q \Rightarrow \neg P : \neg Q \vee \neg P$

	P	Q	$P \Rightarrow \neg Q$	$Q \Rightarrow \neg P$
0	T	T	X	X
1	T	F	✓	✓
2	F	F	✓	✓
3	F	T	✓	✓

Let d be $P \Rightarrow \neg Q$ and B be $Q \Rightarrow \neg P$. Because $M(d) = M(B) = \{1, 2, 3\}$, the two sentences are equivalent.

$$P \Leftrightarrow \neg Q : \begin{matrix} P \Rightarrow \neg Q : \neg P \vee \neg Q \\ \neg Q \Rightarrow P : Q \vee P \end{matrix} : (\neg P \vee \neg Q) \wedge (Q \vee P)$$

	P	Q	$P \Leftrightarrow \neg Q$	$((P \wedge \neg Q) \vee (\neg P \wedge Q))$
0	T	T	X	X
1	T	F	✓	✓
2	F	F	X	X
3	F	T	✓	✓

Let d be $P \Leftrightarrow \neg Q$ and B be $((P \wedge \neg Q) \vee (\neg P \wedge Q))$. Because $M(d) = M(B) = \{1, 3\}$, the two sentences are equivalent.

2)

- a. $\neg(\text{Smoke} \Rightarrow \text{fire}) \vee (\neg \text{Smoke} \Rightarrow \neg \text{fire})$
 $\neg(\neg \text{Smoke} \vee \text{fire}) \vee (\text{Smoke} \vee \neg \text{fire})$
 $(\text{Smoke} \wedge \neg \text{fire}) \vee (\text{Smoke} \vee \neg \text{fire})$

	Smoke	fire	$(\text{Smoke} \Rightarrow \text{fire}) \Rightarrow (\neg \text{Smoke} \Rightarrow \neg \text{fire})$
0	T	T	✓
1	T	F	✓
2	F	T	✗
3	F	F	✓

Because the set of worlds that the sentence is valid in is not all possible worlds, but not no worlds either, the sentence is neither valid nor unsat.

- b. $\neg(\text{Smoke} \Rightarrow \text{fire}) \vee ((\text{Smoke} \vee \text{heat}) \Rightarrow \text{fire})$
 $\neg(\neg \text{Smoke} \vee \text{fire}) \vee (\neg(\text{Smoke} \vee \text{heat}) \vee \text{fire})$
 $(\text{Smoke} \wedge \neg \text{fire}) \vee ((\neg \text{Smoke} \wedge \neg \text{heat}) \vee \text{fire})$

	Smoke	fire	heat	$(\text{Smoke} \Rightarrow \text{fire}) \Rightarrow ((\text{Smoke} \vee \text{heat}) \Rightarrow \text{fire})$
0	T	T	T	✓
1	T	T	F	✓
2	T	F	T	✓
3	T	F	F	✓
4	F	T	T	✓
5	F	T	F	✓
6	F	F	T	✗
7	F	F	F	✓

Because the set of worlds that the sentence is valid in is not all possible worlds, but not empty, the sentence is neither valid nor unsat.

c) LHS: $(\text{Smoke} \wedge \text{heat}) \Rightarrow \text{fire}$

$$\neg (\text{Smoke} \wedge \text{heat}) \vee \text{fire}$$

RHS: $((\text{Smoke} \Rightarrow \text{fire}) \vee (\text{heat} \Rightarrow \text{fire}))$

$$((\neg \text{Smoke} \vee \text{fire}) \vee (\neg \text{heat} \vee \text{fire}))$$

① $\neg (\text{Smoke} \wedge \text{heat}) \vee \text{fire} \Rightarrow ((\neg \text{Smoke} \vee \text{fire}) \vee (\neg \text{heat} \vee \text{fire}))$

① $\neg (\neg (\text{Smoke} \wedge \text{heat}) \vee \text{fire}) \vee ((\neg \text{Smoke} \vee \text{fire}) \vee (\neg \text{heat} \vee \text{fire}))$

① $(\text{Smoke} \wedge \neg \text{fire}) \vee (\neg \text{Smoke} \vee \text{fire} \vee \neg \text{heat})$

② $(\neg \text{Smoke} \vee \text{fire} \vee \neg \text{heat}) \Rightarrow \neg (\text{Smoke} \wedge \text{heat}) \vee \text{fire}$

② $\neg (\neg \text{Smoke} \vee \text{fire} \vee \neg \text{heat}) \vee (\neg (\text{Smoke} \wedge \text{heat}) \vee \text{fire})$

② $(\neg (\neg \text{Smoke} \vee \text{fire}) \wedge \text{heat}) \vee (\neg \text{Smoke} \vee \neg \text{heat} \vee \text{fire})$

② $(\text{Smoke} \wedge \neg \text{fire} \wedge \text{heat}) \vee \neg \text{Smoke} \vee \neg \text{heat} \vee \text{fire}$

$$[(\text{Smoke} \wedge \neg \text{fire}) \vee (\neg \text{Smoke} \vee \text{fire} \vee \neg \text{heat})] \wedge$$

$$[(\text{Smoke} \wedge \neg \text{fire} \wedge \text{heat}) \vee \neg \text{Smoke} \vee \neg \text{heat} \vee \text{fire}]$$

	Smoke	fire	heat	sentence
0	T	T	T	✓
1	T	T	F	✓
2	T	F	T	✓
3	T	F	F	✓
4	F	T	T	✓
5	F	T	F	✓
6	F	F	T	✓
7	F	F	F	✓

Because the set of all worlds the sentence is valid in is all possible worlds, the sentence is valid.

3)

a. $\Delta = \{ \text{mythical} \Rightarrow \neg \text{mortal}, \neg \text{mythical} \Rightarrow \text{mortal} \wedge \text{mammal}, \neg \text{mortal} \vee \text{mammal} \Rightarrow \text{horned}, \text{horned} \Rightarrow \text{magical} \}$

b. mythical: a $(\neg a \vee \neg b) \wedge (a \vee (b \wedge c)) \wedge (\neg(\neg b \vee c) \vee d) \wedge$
 mortal: b $(\neg d \vee e)$
 mammal: c
 horned: d $(\neg a \vee \neg b) \wedge ((a \vee b) \wedge (a \vee c)) \wedge ((b \wedge \neg c) \vee d) \wedge$
 magical: e $(\neg d \vee e)$

$(\neg a \vee \neg b) \wedge (a \vee b) \wedge (a \vee c) \wedge (b \vee d) \wedge (\neg c \vee d) \wedge (\neg d \vee e)$

c. Δ IS CNF from part b

α is a

$\Delta \models \alpha$?

$\Delta \wedge \neg \alpha$ unsat?

0) $\neg a \vee \neg b$

1) $a \vee b$

2) $a \vee c$

3) $b \vee d$

4) $\neg c \vee d$

5) $\neg d \vee e$

6) $\neg a$

7) b

8) c

9) d

10) e

Δ

$\neg \alpha$

1,6

2,6

4,8

5,9

No more rules can be applied, and no contradictions have been found. Thus, $\Delta \wedge \neg \alpha$ is SAT, and Δ cannot be used to prove the unicorn is mythical.

α is e

0) $\neg a \vee \neg b$

1) $a \vee b$

2) $a \vee c$

3) $b \vee d$

4) $\neg c \vee d$

5) $\neg d \vee e$

6) $\neg e$

$\neg \alpha$

7) $\neg d$

8) $\neg c$

9) a

10) $\neg b$

11) d

0,6

4,7

2,8

0,9

3,10

$\neg a$ and 11 contradict.

$\neg d$ and d never true, so

$\Delta \wedge \neg \alpha$ is unsat and

$\Delta \models \alpha$. Thus, we can use Δ to prove the unicorn is magical.

3)

c. continued

homed {

- α is d
- 0) $\neg a \vee \neg b$
- 1) $a \vee b$
- 2) $a \vee c$
- 3) $b \vee d$
- 4) $\neg c \vee d$
- 5) $\neg d \vee e$
- 6) $\neg d$

} Δ

} $\neg \alpha$

7) b 3, 6
 8) $\neg a$ 0, 7
 9) c 2, 8
 10) d 4, 9

6 and 10 contradict,
 so $\Delta \wedge \neg \alpha$ is unsat
 and $\Delta \models \alpha$. Thus,
 we can use Δ to
 prove the unicorn
 is homed.

4) Figure 1)

decomposable: yes, vars on each side of AND gates are different

determinism: no, top or gate inputs aren't mut. excl. (A=true,

$$① (\neg A \wedge B) \vee (\neg B \wedge A) : \checkmark$$

B=false,

$$② (C) \vee (\neg D \wedge \neg C) : \checkmark$$

C=true,

$$③ (\neg A \wedge \neg B) \vee (A) : \checkmark$$

D=false)

$$④ (C \wedge \neg D) \vee (D \wedge \neg C) \checkmark$$

$$⑤ ① \wedge ② \vee ③ \wedge ④$$

smooth: No, with ②, inputs on rightside aren't all present on RHS (D).

Figure 2)

decomposable: yes, vars in each side of ANDs are diff.

determinism: no, or gate ① is same on each side (A=false,

$$① (\neg A \wedge B) \vee (\neg A \wedge B) : X$$

B=true

$$② (C \wedge D) \vee (\neg D \wedge \neg C)$$

leads to

$$③ (\neg A \wedge B) \vee (\neg A \wedge B)$$

both sides

$$④ (C \wedge \neg D) \vee (D \wedge \neg C)$$

being true)

$$⑤ ① \wedge ② \vee ③ \wedge ④$$

smooth: yes, in each or gate, ①-⑤, all vars present in either side are also pres. on other side.

5)

\rightarrow XOR

a.

	A	B	$(\neg A \wedge B) \vee (\neg B \wedge A)$
0	T	T	X
1	T	F	✓
2	F	T	✓
3	F	F	X

$$WMC = w(A, \neg B) + w(\neg A, B)$$

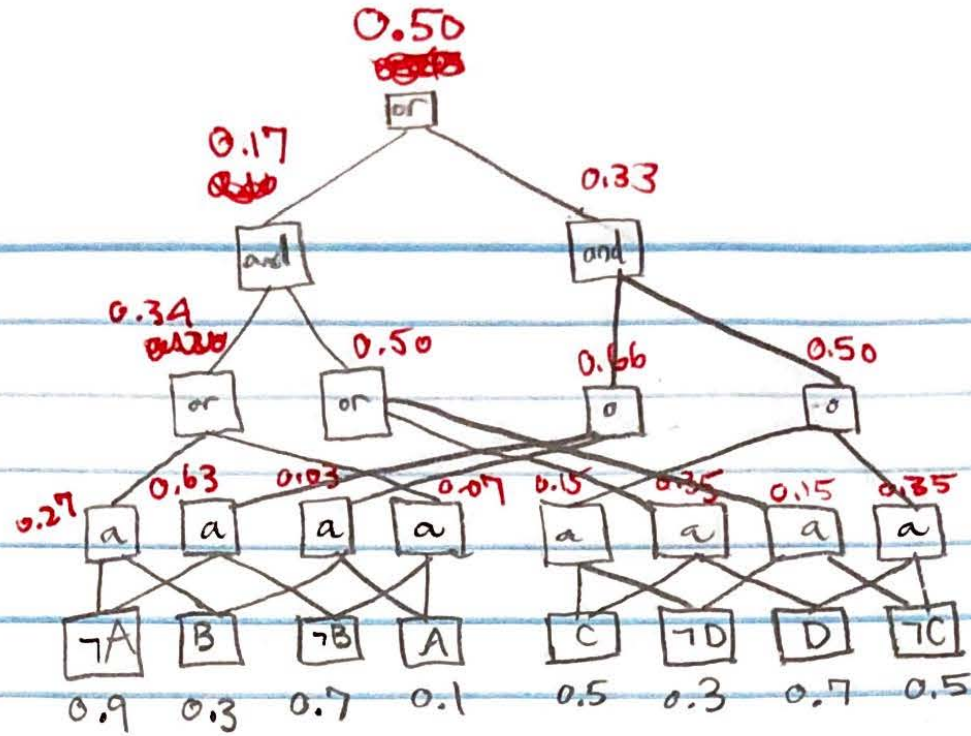
$$= w(A) \cdot w(\neg B) + w(\neg A) \cdot w(B)$$

$$= 0.1 \cdot 0.7 + 0.9 \cdot 0.3$$

$$= 0.34$$

- b. The count on the root is same as the WMC for the formula.
The same calculations are performed.

c.



$$WMC = 0.50$$