

2016

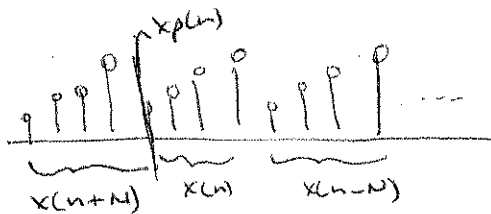
(5)

Relationship between $x(n)$ and $x_p(n)$

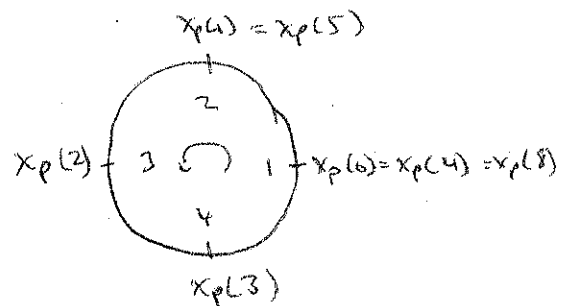
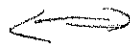
$$x(n) = \{1, 2, 3, 4\} \Rightarrow \begin{array}{c} \text{1} \quad \text{2} \quad \text{3} \quad \text{4} \\ \text{1} \quad \text{2} \quad \text{3} \end{array} \quad L=4$$

now, say $x(n)$ is one period of a signal $x_p(n)$ that has period $N \geq L$. How do we picture this?

Case 1: $N=L$



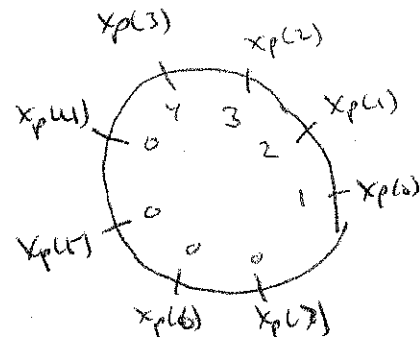
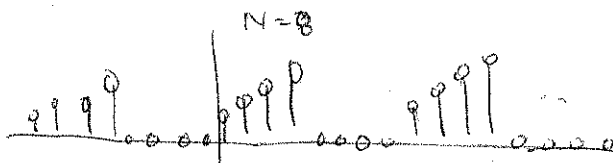
equivalent picture



↑
This is a handy way to think of the periodic signal. Time goes counter-clockwise.

Case 2: $N > L$

here, we "pad" $x(n)$ with zeros



why not always $N=L$?

→ sometime have 2 signals of different lengths

→ see example 7.1.2 in book / class slides

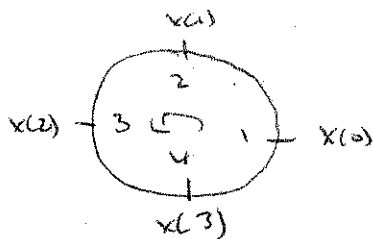
(6)

Shifting periodic sequence

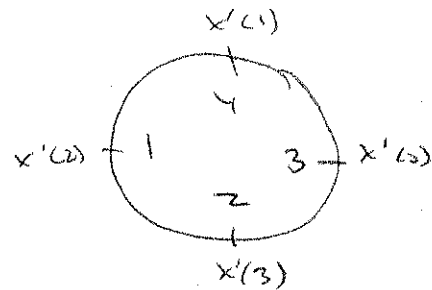
→ For linear time shift (used before) samples just slide over.

→ for a periodic time shift, new signals rotate in:

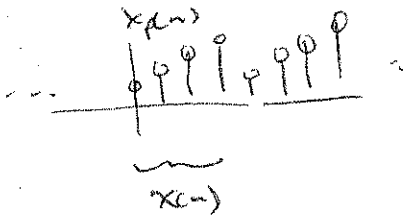
start with:



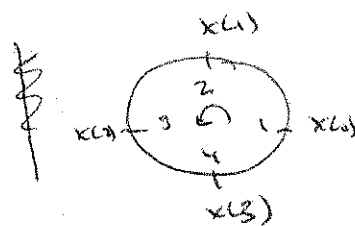
periodic shift
→ 2



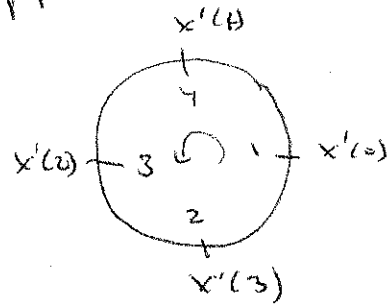
$$x'(n) = \{3, 4, 1, 2\}$$



Time reversal - see class ppt



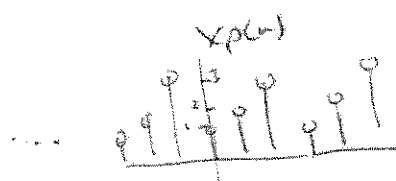
reversal



① Circular convolution lecture - CCL

Review from last lecture:

We saw that a periodic signal $x_p(n)$ could be thought of as a circular sequence:



$$x_p(n) = x_p(n+4) = \dots$$



$$x_p(n) = x_p(n+1) = x_p(n+2) = \dots$$

$$x_p(n) = x_p(n+3) = \dots$$

math to go with this:

$$x_p(n) = x(n \bmod N) \\ \equiv x(n)_N$$

Examples:

$$0 \bmod 3 = 0$$

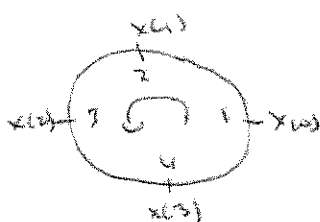
$$2 \bmod 3 = 2$$

$$3 \bmod 3 = 0$$

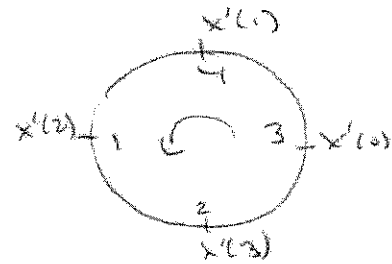
$$4 \bmod 3 = 1$$

etc.

Review: periodic shift



periodic shift
by
2



(1- ppt)

+ time reversal (PPT)

check time-reversal $x' = x(N-n)$, $0 \leq n \leq N-1$

$$\text{if } N=4, \quad x'(0) = x(4-0) = x(0)$$

$$x'(1) = x(4-1) = x(3)$$

$$x'(2) = x(4-2) = x(2)$$

etc.

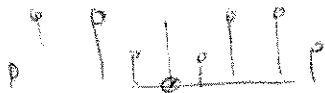
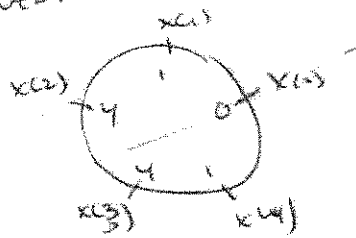
② CCL

Properties: even and odd

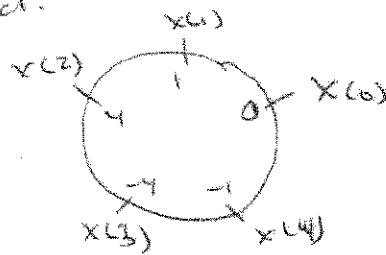
even signal: $x(N-n) = x(n)$ $0 \leq n \leq N-1$

odd signal: $x(N-n) = -x(n)$

even:



odd:



DFT properties - Part 7.2, table 7.2

Just as w/ CTFT & DTFT, we can define properties of DFT. Most are very familiar, but have some adjustment for circular periodic nature of DFT

Symmetry properties : For real $x(n)$

$$X(N-k) = X^*(k) \quad \text{"circular Hermitian symmetry"}$$

$$|X(N-k)| = |X(k)| \quad \text{circular even}$$

$$\angle X(N-k) = -\angle X(k) \quad \text{circular odd.}$$

where do these come from:

- real & even $x(n) \rightarrow$ Real, even $X(k)$

- real & odd $x(n) \rightarrow$ Imag, odd $X(k)$

③ DFT

Properties - ① time reversal

If $x(n) \leftrightarrow X(k)$, then

$$x((-n))_N = x(N-n) \leftrightarrow X((-k))_N = X(N-k)$$

(like DTFT: $x(-n) \leftrightarrow X(-\omega)$)

Proof:

$$\text{DFT}\{x(N-n)\} = \sum_{n=0}^{N-1} x(N-n) e^{-j2\pi kn/N} \quad \text{put in PPT}$$

② Circular time shift

if $x(n) \leftrightarrow X(k)$, then

$$x((n-l))_N \leftrightarrow X(k) e^{-j2\pi kl/N}$$

proof: see PPT

③ similarly, circular freq shift

$$x(n) e^{j2\pi ln/N} \leftrightarrow X((k-l))_N$$

④ Parseval's: $\sum_{k=0}^{N-1} |x(n)|^2 = \frac{1}{N} \sum_{k=0}^{N-1} |X(k)|^2$

⑤ the big one: convolution/multiplication duality
(= PPT)

Final result $x_1(n) \otimes x_2(n) \leftrightarrow X_1(k) X_2(k)$

look at on-line applets for intuition - also

PoM 7.2.2

Doing circular convolution

(N)

(4) eu

time domain

- locate $x_1(m)$ ~~on~~ on N -point circle
- " ~~$x_2(m)$~~ $x_2((-m))_N$ " " " "
- for $n=0, 1, N-1$
 - rotate x_2 by one point CCW
 - multiply
 - add

freq domain (DFT)

- find $X_1(k)$ for $k=0, 1, \dots, N-1$
- " $X_2(k)$ " " "
- find $X_3(k) = X_1(k) X_2(k)$, $k=0, \dots$
- $x_3(n) = F^{-1}\{X_3(k)\}$ for $n=0, 1, \dots, N-1$

$$x = [1 \ 2]$$

$$y = [\cancel{1} \ \cancel{2} \ 4 \ \cancel{5} \ \cancel{6}]$$

mod 11

$$x \circledast y = 13 \ 10 \ 13$$

~~conv~~

$$x * y = 3 \ 10 \ 13 \ 10$$

$$\text{IP } x = [1 \ 2]$$

$$y = [3 \ 4 \ 5]$$

~~draw out~~

a) draw out circles

b) do computation

c) how long is linear conv $x * y$

(5) cll

7.3 Use of DFT in linear filtering

We want to find $y(n) = h(n) * x(n)$
 \uparrow
regular convolution

we know

$$1) Y(\omega) = H(\omega) X(\omega)$$

$$\text{so } y(n) = F^{-1}(H(\omega) X(\omega))$$

we also know that if $x(n)$ is length L ,
 $h(n)$ is length M , then

$$y = h * x \text{ is of length } L + M - 1 //$$

* thus if we want to find $Y(\omega)$, we need a
number of points in the DFT $N \geq (L + M - 1)$

- using fewer points will mean we've undersampled
 $Y(\omega)$. Then, $y(n) = F^{-1}(Y(\omega))$ will have
time-domain aliasing.

- using enough samples means $h \circledast x = h * x$

~~if~~ If we handle this ok, we can use DFT to do linear
convolution as follows:

1) given $x(n)$ of length L
 $h(n)$ of length M

2) set $N = L + M - 1$

3) zero-pad $x(n)$ so it's of length N ($M-1$ zeros)
" " h " " " " ($L-1$ zeros)

4) take N -point DFT's of x, h

$$5) y(n) = \text{DFT}^{-1}\{Y(\omega)\}$$

6 CCL

- go through on-line examples
- go thru example on board

DFT wrap / FFT Lecture

Topic 1: review derivation of $\tilde{X}_1(k) \tilde{X}_2(k) \leftrightarrow X_1 \circledast X_2$
(in PPT)

Review last time

→ Algorithm for circular convolution: place X_1 and $X_2((-n))_N$ on circle, rotate X_2 counter-clockwise, multiply, add

→ avoiding time domain aliasing

- if $x(n)$ is length L , $h(n)$ is length M , then
linear convolution output is length $L+M-1 = N$

- thus, the DFT of $y(n) = x(n) * h(n)$ must have
 $N \geq L+M-1$

- thus, we must zero-pad x and h to avoid
time-domain aliasing; then,

$$X \circledast h = x * h \quad \text{for 1st } N \text{ points}$$

- ~~finally~~ $y(n) = \text{DFT}^{-1}(\tilde{X}(k) \tilde{H}(k))$ in frequency domain

DFT-based filtering of long sequences

If $x(n)$ is a long sequence, h is short, then:

- 1) break $x(n)$ into short blocks, length L
- 2) use DFT to filter each block
- 3) paste things back together

(2)

Overlap-add strategy (Figure is in PPT)algorithm: for block $\alpha = 1, 2, 3, \dots$ 1) take a length L block of data2) zero pad with $M-1$ ~~not~~ zeros ~~(N-M)~~~~so~~ so for first block,

$$x_{\alpha} = [x(1), \dots, x(L), \underset{\substack{\uparrow \\ M-1 \text{ zeros}}}{0, 0, \dots, 0}]$$

3) zero pad $h(n)$ so it's length N

$$h_{\alpha} = [h(1), h(2), \dots, h(M), \underset{\substack{\uparrow \\ L-1}}{0, 0, \dots, 0}]$$

4) find $\tilde{Y}_{\alpha}(k) = \sum_{\alpha}(k) H_{\alpha}(k)$ ← no subscript, h always the same.

$$\text{and } \tilde{y}_{\alpha} = \text{IDFT}^{-1} \left\{ \sum_{\alpha}(k) H(k) \right\}$$

5) Problem: the last $M-1$ samples of $\tilde{y}_{\alpha}(n)$ do not include the effects of the first $M-1$ samples in next block.solution: add last $M-1$ points of $\tilde{y}_{\alpha}(n)$ to first $M-1$ points of $\tilde{y}_{\alpha+1}(n)$ Overlap-~~add~~^{save} strategy (figure is in PPT)algorithm for block $\alpha = 1, 2, 3, \dots$ 1) take length L data block2) put previous $(M-1)$ points ~~at~~ at beginning x
(for $\alpha=1$, pre-pad w/ $M-1$ zeros)3) Thus input length $N = L + M - 1$; ~~for length~~ do an N -point DFT- note we'd really want a $L + (M-1) \times 2 = L + 2M - 2$ length filter; since we are short, first $M-1$ points are time aliased

(3)

this N -point DFT gives us

$$\tilde{Y}_\alpha(k) = \sum_{\alpha} X_\alpha(k) H(k), \quad N = M+L-1 \text{ points}$$

$$\tilde{y}_\alpha(n) = \text{IDFT} \left\{ \sum_{\alpha} X_\alpha(k) H(k) \right\}$$

\uparrow
 $M+L-1$

4) because we used a DFT that was too short, the first $M-1$ points of $\tilde{y}_\alpha(n)$ are aliased.

5) Thus we discard the first $M-1$ points, keep rest, and continue.

Both overlap-add and overlap-sum process L data points in each block using an $N = L+M-1$ DFT. Thus, choice is really personal preference.

