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The Art of Compiler Construction using C#

Ву

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CHAPTER 1

Version 0.1

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The Art of Compiler Construction using C#

Thanks to the availability of information and better tools writing a compiler has become just an excersise in software engineering. The Compilers are not difficult programs to write. The various phases of compilers are easy to understand in an independent manner. The relationship is not purely sequential. It takes some time to put phases in perspective in the job of compilation of programs.

The task of writing a compiler can be viewed in a top down fashion as follows

Parsing => Creation of Abstract Syntax Tree => Tree Traversal to generate the Object code or Recursive interpretation.

Abstract Syntax Tree

In computer science, an abstract syntax tree (AST), or just syntax tree, is a tree representation of the abstract (simplified) syntactic structure of source code written in a certain programming language. Each node of the tree denotes a construct occurring in the source code. The syntax is abstract in the sense that it does not represent every detail that appears in the real syntax. For instance, grouping parentheses is implicit in the tree structure, and a syntactic construct such as if cond then expr may be denoted by a single node with two branches. Most of you might not be aware of the fact that , programming languages are hierarchical in nature. We can model programming language constructs as classes. Trees are a natural data structure to represent most things hierarchical.

As a case in the point, let us look a simple expression evaluator. The expression evaluator will support double precision floating point value as the operands. The Operators supported are addition (+), subtraction (-), multiplication (*) and division. The Object model support Unary operators (+, -) as well. We are planning to use a composition model for modeling an expression.

In most imperative programming languages, an expression is something which u evaluate for it's value. Where as statements are something which you executes for it's effect.

let us define an abstract class for Exp

```
/// <summary>
/// Abstract for Expression evaluation
/// </summary>
abstract class Exp
{
    public abstract double Evaluate(RUNTIME_CONTEXT cont);
}
```

For the time being RUNTIME_CONTEXT is an empty class

```
/// <summary>
/// One can store the stack frame inside this class
/// </summary>
public class RUNTIME_CONTEXT
{
    public RUNTIME_CONTEXT()
    {
        }
    }
}
```

Modeling Expression

Once u have declared the interface and it's parameters , we can create a hierarchy of classes to model an expression.

```
class Exp // Base class for Expression class NumericConstant // Numeric Value class BinaryExp // Binary Expression class UnaryExp // Unary Expression
```

Take a look at the listing of NumericConstant class

```
public override double Evaluate(RUNTIME_CONTEXT cont)
{
    return _value;
}
```

Since the class is derived from Exp , it ought to implement the Evaluate method. In the Numeric Constant node , we will store a IEEE 754 double precision value. While evaluating the tree , the node will return the value stored inside the object.

Binary Expression

In a Binary Expression , one will have two Operands (Which are themselves expressions of arbitary complexity) and an Operator.

```
/// This class supports Binary Operators like + , - , / , *
public class BinaryExp : Exp
  private Exp _ex1, _ex2;
  private OPERATOR _op;
  /// <param name="a"></param>
  /// <param name="b"></param>
  public BinaryExp(Exp a, Exp b, OPERATOR op)
     _{\text{ex}1} = a;
     ex2 = b;
     _{op} = op;
  /// While evaluating apply the operator after evaluating the left and right operands
  /// <param name="cont"></param>
  public override double Evaluate(RUNTIME CONTEXT cont)
    switch (_op)
       case OPERATOR.PLUS:
         return ex1.Evaluate(cont) + ex2.Evaluate(cont);
       case OPERATOR.MINUS:
         return ex1.Evaluate(cont) - ex2.Evaluate(cont);
```

```
case OPERATOR.DIV:
    return _ex1.Evaluate(cont) / _ex2.Evaluate(cont);
    case OPERATOR.MUL:
    return _ex1.Evaluate(cont) * _ex2.Evaluate(cont);
}

return Double.NaN;
}
```

Unary Expression

In an unary expression, one will have an Operand (which can be an expression of arbitary complexity) and an Operator which can be applied on the Operand.

```
/// <summary>
/// This class supports Unary Operators like + ,- ,/ ,*
/// </summary>
public class UnaryExp: Exp
{
    private Exp _ex1;
    private OPERATOR _op;
/// <summary>
/// <param name="a"></param>
/// <param name="b"></param>
/// <param name="b"></param>
/// <param name="op"></param>
/// <param name="cont"></param>
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
///
```

```
case OPERATOR.PLUS:
    return _ex1.Evaluate(cont);
    case OPERATOR.MINUS:
    return -_ex1.Evaluate(cont);
}

return Double.NaN;
}
```

In the CallSLANG project, we will include the SLANG DOT NET assembly before composing the expression.

```
using System;
using System.Collections.Generic;
using System.Linq;
using System.Text;
using SLANG_DOT_NET; // include SLANG_DOT_NET assembly
namespace CallSLANG
  class Program
    static void Main(string[] args)
      // Abstract Syntax Tree (AST) for 5*10
      Exp e = new BinaryExp(new NumericConstant(5),
                   new NumericConstant(10),
                   OPERATOR.MUL);
      // Evaluate the Expression
      Console.WriteLine(e.Evaluate(null));
      // AST for (10 + (30 + 50))
      e = new UnaryExp(
             new BinaryExp(new NumericConstant(10),
               new BinaryExp(new NumericConstant(30),
                       new NumericConstant(50),
                  OPERATOR.PLUS),
             OPERATOR.PLUS),
```

```
OPERATOR.MINUS);

//
// Evaluate the Expression
//
Console.WriteLine(e.Evaluate(null));

//
// Pause for a key stroke
//
Console.Read();

}
}
```

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CHAPTER 2

Version 0.1

INPUT Analysis

Compilers are programs which translate source language to a target language. The Source language can be a language like C,C++ or Lisp. The potential target languages are assembly languages, object code for the microprocessors like intel x86, itanium or power pc. There are programs which translate java to C++ and Lisp to C. In such case, target language is another programming language.

Any compiler has to understand the input. Once it has analyzed the input characters, it should convert the input into a form which is suitable for further processing.

Any input has to be parsed before the object code translation. To Parse means to understand. The Parsing process works as follows

The characters are grouped together to find a token (or a word). Some examples of the tokens are '+','*',while , for , if etc. The module which reads character at a time and looks for legal token is called a lexical analyzer or Lexer. The input from the Lexer is passed into a module which identifies whether a group of tokens form a valid expression or a statement in the program. The module which determines the validity of expressions is called a parser. Rather than doing a lexical scan for the entire input , the parser requests the next token from the lexical analyzer. They act as if they are co-routines.

To put everything together let us write a small program which acts a four function calculator. The calculator is capable of evaluating mathematical expressions which contains four basic arithmetical operators, paranthesis to group the expression and unary operators. Given below is the Lexical Specifications of the calculator.

```
TOK PLUS - '+'
TOK MUL - '*'
TOK SUB - '-'
TOK DIV - '/'
TOK OPAREN - '('
TOK CPAREN - ')'
TOK DOUBLE - [0-9]+
The stuff can be converted into C# as follows
/// Enumeration for Tokens
public enum TOKEN {
   ILLEGAL TOKEN = -1, // Not a Token
   TOK PLUS = 1, // '+'
   TOK_MUL, // '*
   TOK DIV, // '/'
   TOK SUB, // '-'
   TOK_OPAREN, // '('
   TOK CPAREN, // ')'
   TOK DOUBLE, // '('
   TOK_NULL // End of string
```

The Lexical Analysis Algorithm scans through the input and returns the token associated with the operator. If it has found out a number, returns the token associated with the number. There should be another mechanism to retrieve the actual number identified.

Following pseudo code shows the schema of the lexical analyzer

```
while ( there is input ) {
    switch(currentchar) {
    case Operands:
        advance input pointer
        return TOK_XXXX;
    case Number:
        Extract the number( Advance the input )
        return TOK_DOUBLE;
        default:
        error
    }
}
```

The following C# code is a literal translation of the above algorithm.

```
using System;
using System.Collections.Generic;
using System.Ling;
using System.Text;
namespace SLANG DOT NET
  /// Enumeration for Tokens
  public enum TOKEN
    ILLEGAL TOKEN = -1, // Not a Token
    TOK PLUS = 1, // '+'
    TOK MUL, // '*'
    TOK DIV, // '/'
    TOK_SUB, // '-'
    TOK OPAREN, // '('
    TOK CPAREN, // ')'
    TOK DOUBLE, // '('
    TOK NULL // End of string
  // A naive Lexical analyzer which looks for operators , Parenthesis
  // and number. All numbers are treated as IEEE doubles. Only numbers
  // without decimals can be entered. Feel free to modify the code
  // to accomodate LONG and Double values
```

```
public class Lexer
  String IExpr; // Expression string
  int index; // index into a character
  int length; // Length of the string
  double number; // Last grabbed number from the stream
  // Ctor
  public Lexer(String Expr)
    IExpr = Expr;
    length = IExpr.Length;
    index = 0;
  // Grab the next token from the stream
  public TOKEN GetToken()
    TOKEN tok = TOKEN.ILLEGAL TOKEN;
    // Skip the white space
    while (index < length &&
    (IExpr[index] == '' || IExpr[index] == '\t'))
      index++;
    // End of string ? return NULL;
    if (index == length)
      return TOKEN.TOK_NULL;
    switch (IExpr[index])
      case '+':
         tok = TOKEN.TOK_PLUS;
         index++;
         break;
      case '-':
         tok = TOKEN.TOK_SUB;
         index++;
         break;
       case '/':
         tok = TOKEN.TOK_DIV;
         index++;
         break;
      case '*':
```

```
tok = TOKEN.TOK_MUL;
       index++;
       break;
     case '(':
       tok = TOKEN.TOK_OPAREN;
       index++;
       break;
     case ')':
       tok = TOKEN.TOK_CPAREN;
       index++;
       break;
     case '0':
    case '1':
    case '2':
     case '3':
     case '4':
     case '5':
    case '6':
    case '7':
    case '8':
    case '9':
          String str = "";
          while (index < length &&
          (IExpr[index] == '0' \parallel
          IExpr[index] == '1' ||
          IExpr[index] == '2' ||
          IExpr[index] == '3' \parallel
          IExpr[index] == '4' \parallel
          IExpr[index] == '5' ||
          IExpr[index] == '6' ||
          IExpr[index] == '7' ||
          IExpr[index] == '8' ||
          IExpr[index] == '9')
            str += Convert.ToString(IExpr[index]);
            index++;
          number = Convert.ToDouble(str);
          tok = TOKEN.TOK_DOUBLE;
       break;
    default:
       Console. WriteLine("Error While Analyzing Tokens");
       throw new Exception();
  return tok;
public double GetNumber() { return number; }
```

The Grammar

In computer science, a formal grammar (or grammar) is a set of formation rules (grammar) that describe which strings formed from the alphabet of a formal language are syntactically valid within the language. A grammar only addresses the location and manipulation of the strings of the language. It does not describe anything else about a language, such as its semantics (i.e. what the strings mean).

A context-free grammar is a grammar in which the left-hand side of each production rule consists of only a single nonterminal symbol. This restriction is non-trivial; not all languages can be generated by context-free grammars. Those that can are called context-free languages.

The Backus Naur Form (BNF) notation is used to specify grammars for programming languages, commnd line tools, file formats to name a few. The semantics of BNF is beyond the scope of this book.

Grammar of the expression evaluator

```
<Expr> ::= <Term> | Term { + | - } <Expr> <Term> ::= <Factor> | <Factor> {*|/} <Term> <Factor> ::= <number> | ( <expr> ) | {+|-} <factor>
```

There are two types of tokens in any grammar specifications. They are terminal tokens (terminals) or non terminals. In the above grammar, operators and <number> are the terminals. <Expr>,<Term>,<Factor> are non terminals. Non terminals will have at least one entry on the left side.

Conversion of Expression to the psuedo code

```
// <Expr> ::= <Term> { + | - } <Expr>
Void Expr() {
    Term();

if ( Token == TOK_PLUS || Token == TOK_SUB )
    {
        // Emit instructions
        // and perform semantic operations

        Expr(); // recurse
    }
}
```

Converstion of term to the psuedo code

```
// <Term> ::= <Factor> { * | / } <Term>
Void Term() {
    Factor();
```

```
if ( Token == TOK_MUL || Token == TOK_DIV )
{
    // Emit instructions
    // and perform semantic operations

Term(); // recurse
}
```

The following psuedo code demonstrates how to map <Factor> into code

```
// <Factor> ::= <TOK_DOUBLE>|(<expr>)| {+|-} <Factor>
//
Void Factor() {
    switch(Token)
    case TOK_DOUBLE:
    // push token to IL operand stack return
    case TOK_OPAREN:
        Expr(); //recurse
        // check for closing parenthesis and return
    case UNARYOP:
        Factor(); //recurse
        default:
        //Error
}
```

The class RDParser is derived from the Lexer class. By using an algorithm by the name Recursive descent parsing , we will evaluate the expression. A recursive descent parser is a top-down parser built from a set of mutually-recursive procedures where each such procedure usually implements one of the production rules of the grammar. Thus the structure of the resulting program closely mirrors that of the grammar it recognizes.

```
using System;
using System.Collections.Generic;
using System.Linq;
using System.Text;

namespace SLANG_DOT_NET
{
/// <summary>
///
/// </summary>
public class RDParser: Lexer
```

```
TOKEN Current_Token;
public RDParser(String str)
 : base(str)
public Exp CallExpr()
  Current_Token = GetToken();
 return Expr();
public Exp Expr()
  TOKEN 1 token;
  Exp RetValue = Term();
  while (Current_Token == TOKEN.TOK_PLUS || Current_Token == TOKEN.TOK_SUB)
    1_token = Current_Token;
    Current Token = GetToken();
    Exp e1 = Expr();
    RetValue = new BinaryExp(RetValue, e1,
      1_token == TOKEN.TOK_PLUS ? OPERATOR.PLUS : OPERATOR.MINUS);
 return RetValue;
public Exp Term()
  TOKEN l_token;
  Exp RetValue = Factor();
  while (Current_Token == TOKEN.TOK_MUL || Current_Token == TOKEN.TOK_DIV)
    1_token = Current_Token;
    Current Token = GetToken();
    Exp e1 = Term();
    RetValue = new BinaryExp(RetValue, e1,
```

```
1_token == TOKEN.TOK_MUL ? OPERATOR.MUL : OPERATOR.DIV);
 return RetValue;
public Exp Factor()
  TOKEN 1 token;
  Exp RetValue = null;
  if (Current_Token == TOKEN.TOK_DOUBLE)
    RetValue = new NumericConstant(GetNumber());
    Current Token = GetToken();
  else if (Current_Token == TOKEN.TOK_OPAREN)
    Current_Token = GetToken();
    RetValue = Expr(); // Recurse
    if (Current_Token != TOKEN.TOK_CPAREN)
      Console.WriteLine("Missing Closing Parenthesis\n");
      throw new Exception();
    Current_Token = GetToken();
  else if (Current_Token == TOKEN.TOK_PLUS || Current_Token == TOKEN.TOK_SUB)
    1_token = Current_Token;
    Current_Token = GetToken();
    RetValue = Factor();
    RetValue = new UnaryExp(RetValue,
       1_token == TOKEN.TOK_PLUS ? OPERATOR.PLUS : OPERATOR.MINUS);
  else
    Console.WriteLine("Illegal Token");
    throw new Exception();
 return RetValue;
```

Using the Builder Pattern, we will encapsulate the Parser, Lexer class activities

```
using System;
using System.Collections.Generic;
using System.Linq;
using System.Text;
namespace SLANG_DOT_NET
  /// Base class for all the Builders
  public class AbstractBuilder
  public class ExpressionBuilder: AbstractBuilder
     public string _expr_string;
     public ExpressionBuilder(string expr)
       _expr_string = expr;
     public Exp GetExpression()
       try
         RDParser p = new RDParser(_expr_string);
         return p.CallExpr();
       catch (Exception)
         return null;
```

In the CallSLang Project, the expression compiler is invoked as follows

```
using System;
using System.Collections.Generic;
using System.Linq;
using System.Text;
using SLANG_DOT_NET;

namespace CallSLANG
{
    class Program
    {
        static void Main(string[] args)
        {
            ExpressionBuilder b = new ExpressionBuilder("-2*(3+3)");
            Exp e = b.GetExpression();
            Console.WriteLine(e.Evaluate(null));

            Console.Read();
        }
    }
}
```

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CHAPTER 3

Version 0.1

STATEMENTS

The crux of the SLANG4.net can be summed up in two sentences

Expression is what you evaluate for it's value

Statement is what you execute for it's effect (on variables)

The above two maxims can be converted into a computational structure as follows

A) Expression is what you evaluate for it's value

```
/// <summary>
/// Abstract for Expression evaluation
/// Expression is what you evaluate for it's value
/// </summary>
public abstract class Exp
{
   public abstract double Evaluate(RUNTIME_CONTEXT cont);
}
```

B) Statement is what you execute for it's effect (on variables or lack of it)

```
/// <summary>
/// Statement is what you Execute for it's Effect
/// </summary>
public abstract class Stmt
{
    public abstract bool Execute(RUNTIME_CONTEXT con);
}
```

Let us implement a Print statement for the SLANG4.net compiler. The basic idea is as follows you add a class to model a statement and since the class has to inherit from the Stmt (abstract class), it ought to implement Execute Method.

```
/// <summary>
/// Implementation of Print Statement
/// </summary>

public class PrintStatement : Stmt
{
/// <summary>
/// At this point of time , Print will
/// spit the value of an Expression on the screen.
/// </summary>
private Exp_ex;
/// <summary>
```

```
/// Ctor just stores the expression passed as parameter
/// 
/// <param name="ex">
/// <param name="ex">
/// <param>
public PrintStatement(Exp ex)
{
    _ex = ex;
}

/// <summary>
/// Execute method Evaluates the expression and
/// spits the value to the console using
/// Console.Write statement.
/// 
/// <param name="con">
/// <param name="con"</p>
// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// <p
```

Let us add a PrintLine statement as well. PrintLine implementation is not different from Print statement. The only difference is it emits a new line after the expression value.

```
/// <summary>
/// Implementation of PrintLine Statement
/// </summary>
public class PrintLineStatement : Stmt
{
    private Exp _ex;

    public PrintLineStatement(Exp ex)
    {
        _ex = ex;
    }

    /// <summary>

/// Here we are calling Console.WriteLine to emit
/// an additional new line .

/// </summary>
/// <param name="con"></param>
/// /// returns> //returns>
public override bool Execute(RUNTIME_CONTEXT con)
{
        double a = _ex.Evaluate(con);
        Console.WriteLine(a.ToString());
        return true;
    }
}
```

Once we have created to classes to implement Print and PrintLine statement , we need to modify our parser (frontend) to support the statement in the language.

We are going to add few more tokens to support the Statements in the SLANG.

```
public enum TOKEN

{
    ILLEGAL_TOKEN = -1, // Not a Token
    TOK_PLUS = 1, // '+'
    TOK_MUL, // '*'
    TOK_DIV, // '/'
    TOK_SUB, // '-'
    TOK_OPAREN, // '('
    TOK_CPAREN, // ')'
    TOK_DOUBLE, // 'number'
    TOK_NULL, // End of string
    TOK_NULL, // End of string
    TOK_PRINT, // Print Statement
    TOK_PRINTLN, // PrintLine
    TOK_UNQUOTED_STRING,
    TOK_SEMI // ;
}
```

In the Lexer.cs module, we add a new data structure to be used for Keyword lookup.

```
/// <summary>
/// Keyword Table Entry
/// </summary>
///
public struct ValueTable
{
   public TOKEN tok;  // Token id
   public String Value;  // Token string
   public ValueTable(TOKEN tok, String Value)
   {
      this.tok = tok;
      this.Value = Value;
   }
}
```

In the ctor of Lexer.cs , we will populate an array of ValueTables with Token and it's textual representation as given below.

```
_val = new ValueTable[2];

_val[0] = new ValueTable(TOKEN.TOK_PRINT, "PRINT");

_val[1] = new ValueTable(TOKEN.TOK_PRINTLN, "PRINTLINE");
```

We need to add a new entrypoint into the RDParser.cs class to support statements. The grammar for the SLANG at this point of time (to support statement) is as given below.

```
<stmtlist> := { statement }+

{statement} := <printstmt> | <printlinestmt>

<printstmt> := print <expr>;

<printlinestmt>:= printline <expr>;

<Expr> ::= <Term> | Term { + | - } <Expr>
<Term> ::= <Factor> | <Factor> {*|/} <Term>
<Factor>::= <number> | ( <expr> ) | {+|-} <factor>
```

The new entry point to the parser module is as follows...

```
/// <summary>
/// The new Parser entry point
/// </summary>
/// <returns></returns>
public ArrayList Parse()
{
    GetNext(); // Get the Next Token
//
// Parse all the statements
//
return StatementList();
}
```

The StatementList method implements the grammar given above. The BNF to source code translation is very easy and without much explanation it is given below

```
/// <stmmary>
/// The Grammar is
///
/// <stmtlist> := { <statement> }+
///
/// {<statement> := <printstmt> | <printlinestmt>
/// <printstmt> := print <expr >;
///
/// <printlinestmt>:= printline <expr>;
///
/// <Expr> ::= <Term> | <Term> { + | - } <Expr>
/// <Term> ::= <Factor> | <Factor> {*//} <Term>
/// <Factor> ::= <number> | ( <expr> ) | {+|-} <factor>
///
/// </summary>
```

```
/// <returns></returns>
private ArrayList StatementList()
{
    ArrayList arr = new ArrayList();
    while (Current_Token != TOKEN.TOK_NULL)
    {
        Stmt temp = Statement();
        if (temp != null)
        {
            arr.Add(temp);
        }
     }
     return arr;
}
```

The method Statement queries the statement type and parses the rest of the statement.

```
/// This Routine Queries Statement Type
/// to take the appropriate Branch...
/// Currently, only Print and PrintLine statement
/// are supported..
/// if a line does not start with Print or PrintLine ..
/// an exception is thrown
private Stmt Statement()
  Stmt retval = null;
  switch (Current Token)
     case TOKEN.TOK PRINT:
       retval = ParsePrintStatement();
       GetNext();
       break;
     case TOKEN.TOK PRINTLN:
       retval = ParsePrintLNStatement();
       GetNext();
       break;
     default:
       throw new Exception("Invalid statement");
       break;
  return retval;
/// Parse the Print Staement .. The grammar is
/// PRINT <expr>;
    Once you are in this subroutine, we are expecting
    a valid expression (which will be compiled) and a
    semi collon to terminate the line..
    Once Parse Process is successful, we create a PrintStatement
/// Object..
private Stmt ParsePrintStatement()
```

```
GetNext();
  Exp a = Expr();
  if (Current_Token != TOKEN.TOK_SEMI)
     throw new Exception("; is expected");
  return new PrintStatement(a);
/// Parse the PrintLine Staement .. The grammar is
/// PRINTLINE <expr>;
/// Once you are in this subroutine, we are expecting
/// a valid expression ( which will be compiled ) and a
/// semi collon to terminate the line..
/// Once Parse Process is successful, we create a PrintLineStatement
/// Object..
private Stmt ParsePrintLNStatement()
  GetNext();
  Exp a = Expr();
  if (Current Token != TOKEN.TOK SEMI)
    throw new Exception("; is expected");
  return new PrintLineStatement(a);
```

Finally in the callSlang Project, I have invoked these routines to demonstrate how everything is put together.

```
Stmt s = obj as Stmt;
s.Execute(null);
}

/// summary>
///
// summary>
static void TestSecondScript()
{
string a = "PRINTLINE -2*10;" + "\r\n" + "PRINTLINE -10*-1;\r\n PRINT 2*10;\r\n";
RDParser p = new RDParser(a);
ArrayList arr = p.Parse();
foreach (object obj in arr)
{
Stmt s = obj as Stmt;
s.Execute(null);
}
/// summary>
/// </ri>
/// summary>
/// param name="args"></param>
static void Main(string[] args)
{
// TestFirstScript();
TestSecondScript();
Console.Read();
}
}
```

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CHAPTER 4

Version 0.1

Types, Variables and Assignment Statement

In this step, we will support data types and variables to SLANG. Assignment statement will also be implemented in this step.

The language supports only three data types viz

NUMERIC STRING BOOLEAN

Let us add an Enum for data types

```
/// <summary>
/// Type info enumerations
/// </summary>
public enum TYPE_INFO
{
    TYPE_ILLEGAL = -1, // NOT A TYPE
    TYPE_NUMERIC, // IEEE Double precision floating point
    TYPE_BOOL, // Boolean Data type
    TYPE_STRING, // String data type
}
```

Every variable will have a name , type and a slot for storing it's value in the symbol table. Moreover , Functions return SYMBOL_INFO.

```
/// <summary>
/// Symbol Table entry for variable
/// using Attributes, one can optimize the
/// storage by simulating C/C++ union.
/// </summary>
public class SYMBOL_INFO
{

public String SymbolName; // Symbol Name
public TYPE_INFO Type; // Data type
public String str_val; // memory to hold string
public double dbl_val; // memory to hold double
public bool bol_val; // memory to hold boolean
}
```

The next step is to modify the Exp and Stmt interface to reflect the variable support.

```
/// <summary>
/// In this Step , we add two more methods to the Exp class
/// TypeCheck => To do Type analysis
/// get_type => Type of this node
/// </summary>
abstract class Exp
{
   public abstract SYMBOL_INFO Evaluate(RUNTIME_CONTEXT cont);
   public abstract TYPE_INFO TypeCheck(COMPILATION_CONTEXT cont);
   public abstract TYPE_INFO get_type();
}
```

The class RunTime context contains the Symbol Table during interpretation.

```
/// A Context is necessary for Variable scope...will be used later
public class RUNTIME_CONTEXT
  /// Symbol Table for this context
  private SymbolTable m_dt;
  /// Create an instance of Symbol Table
  public RUNTIME CONTEXT()
     m_dt = new SymbolTable();
  /// Property to retrieve Table
  public SymbolTable TABLE
     get
       return m_dt;
     set
       m_dt = value;
```

```
/// A Context is necessary for Variable scope...will be used later
public class COMPILATION_CONTEXT
  /// Symbol Table for this context
  private SymbolTable m_dt;
  /// Create an instance of Symbol Table
  public COMPILATION_CONTEXT()
    m_dt = new SymbolTable();
  /// Property to retrieve Table
  public SymbolTable TABLE
    get
      return m dt;
    set
      m dt = value;
```

Let us write the code for BooleanConstant node. This will store TRUE or FALSE value.

```
/// <summary>
/// <summary>
/// Node for Boolean Constant ( TRUE, FALSE }
/// Value
/// </summary>
public class BooleanConstant : Exp
{
/// <summary>
/// Info Field
```

```
private SYMBOL INFO info;
/// Ctor
/// <param name="pvalue"></param>
public BooleanConstant(bool pvalue)
  info = new SYMBOL_INFO();
  info.SymbolName = null;
  info.bol val = pvalue;
  info.Type = TYPE INFO.TYPE BOOL;
/// Evaluation of boolean will given the value
/// <param name="local"></param>
/// <param name="global"></param>
public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
  return info;
/// <param name="global"></param>
public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
  return info. Type;
/// <summary>
public override TYPE_INFO get_type()
  return info. Type;
```

The next thing which we should support is NumericConstant. This will store a IEEE 754 double precision floating point value.

```
/// <summary>
/// </summary>
public class NumericConstant : Exp
{
/// <summary>
/// Info field
/// </summary>
private SYMBOL_INFO info;
```

```
/// <summary>
public NumericConstant(double pvalue)
  info = new SYMBOL_INFO();
  info.SymbolName = null;
  info.dbl val = pvalue;
  info.Type = TYPE_INFO.TYPE_NUMERIC;
}
/// Evaluation of boolean will given the value
/// <param name="global"></param>
public override SYMBOL_INFO Evaluate(RUNTIME_CONTEXT cont)
  return info;
/// <summary>
/// <param name="local"></param>
/// <param name="global"></param>
public override TYPE_INFO TypeCheck(COMPILATION CONTEXT cont)
  return info. Type;
/// <summary>
public override TYPE_INFO get_type()
  return info. Type;
```

The AST node for storing String Literal is as given below

```
/// <summary>
/// To Store Literal string enclosed
/// in Quotes
/// </summary>
public class StringLiteral : Exp
```

```
/// info field
private SYMBOL_INFO info;
/// Ctor
/// <param name="pvalue"></param>
public StringLiteral(string pvalue)
  info = new SYMBOL INFO();
  info.SymbolName = null;
  info.str val = pvalue;
  info.Type = TYPE_INFO.TYPE_STRING;
/// Evaluation of boolean will given the value
/// <param name="local"></param>
/// <param name="global"></param>
public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
  return info;
/// <param name="global"></param>
public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
  return info. Type;
public override TYPE INFO get type()
  return info. Type;
```

Addding support to variable is an involved activity. The code has been extensively commented to explain the rationale.

```
/// <summary>
/// Node to store Variables
/// The data types supported are
/// NUMERIC
```

```
STRING
/// BOOLEAN
/// The node store only the variable name, the
/// associated data will be found in the
/// Symbol Table attached to the
    COMPILATION CONTEXT
public class Variable: Exp
  private string m name; // Var name
  TYPE_INFO _type;
                       // Type
  /// this Ctor just stores the variable name
  public Variable(SYMBOL_INFO inf)
    m_name = inf.SymbolName;
  /// Creates a new symbol and puts into the symbol table
  /// and stores the key (variable name)
  /// <param name=" value"></param>
  public Variable(COMPILATION_CONTEXT st, String name, double _value)
    SYMBOL INFO s = new SYMBOL INFO();
    s.SymbolName = name;
    s.Type = TYPE INFO.TYPE NUMERIC;
    s.dbl val = value;
    st.TABLE.Add(s);
    m_name = name;
  /// Creates a new symbol and puts into the symbol table
  /// and stores the key (variable name)
  /// <param name=" value"></param>
  public Variable(COMPILATION CONTEXT st, String name, bool value)
    SYMBOL INFO s = new SYMBOL INFO();
    s.SymbolName = name;
    s.Type = TYPE_INFO.TYPE_BOOL;
    s.bol val = value;
    st.TABLE.Add(s);
    m name = name;
  /// Creates a new symbol and puts into the symbol table
     and stores the key (variable name)
```

```
/// <param name="st"></param>
/// <param name="name"></param>
/// <param name=" value"></param>
public Variable(COMPILATION CONTEXT st, String name, string value)
  SYMBOL INFO s = new SYMBOL INFO();
  s.SymbolName = name;
  s.Type = TYPE_INFO.TYPE_STRING;
  s.str_val = _value;
  st.TABLE.Add(s);
  m name = name;
/// Retrieves the name of the Variable (method version)
public string GetName()
  return m_name;
/// Retrieves the name of the Variable (property version)
public string Name
  get
    return m_name;
  set
    m_name = value;
/// To Evaluate a variable , we just need to do a lookup
/// in the Symbol table ( of RUNTIME CONTEXT )
/// <param name="st"></param>
/// <param name="glb"></param>
public override SYMBOL_INFO Evaluate(RUNTIME_CONTEXT cont)
  if (cont.TABLE == null)
    return null;
  else
    SYMBOL_INFO a = cont.TABLE.Get(m_name);
    return a;
```

```
/// Look it up in the Symbol Table and
/// return the type
/// <param name="local"></param>
/// <param name="global"></param>
public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
  if (cont.TABLE == null)
    return TYPE_INFO.TYPE_ILLEGAL;
  else
    SYMBOL_INFO a = cont.TABLE.Get(m_name);
    if (a != null)
      _type = a.Type;
      return _type;
    return TYPE_INFO.TYPE_ILLEGAL;
    this should only be called after the TypeCheck method
    has been invoked on AST
public override TYPE_INFO get_type()
  return _type;
```

At this point of time Expression hierarchy (without operators) looks like as follows

```
class Exp
class BooleanConstant
class NumericConstant
class StringLiteral
class Variable
```

Once we have created nodes to represents constants of the type which we are planning to support , we created a variable node. The next challenge is to add support for the operators. Till now , I had UnaryExp and BinaryExp. For clarity , I will have classes like Plus (+), Minus (-), Div (/) and Mul (*) for BinaryExp and will have classes UnaryPlus (+), UnaryMinus (-) for Unary operators

The first operator to be supported is Binary +

```
/// the node to represent Binary +
public class BinaryPlus: Exp
  /// Plus has got a left expression (expl )
  /// and a right expression...
  /// and a Associated type information
  private Exp exp1, exp2;
  TYPE INFO type;
  /// <param name="e1"></param>
  public BinaryPlus(Exp e1, Exp e2)
    \exp 1 = e1; \exp 2 = e2;
  public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
    SYMBOL INFO eval left = exp1.Evaluate(cont);
    SYMBOL INFO eval right = exp2.Evaluate(cont);
    if (eval_left.Type == TYPE_INFO.TYPE_STRING &&
       eval right. Type == TYPE INFO. TYPE STRING)
```

```
SYMBOL_INFO ret_val = new SYMBOL_INFO();
    ret val.str val = eval left.str val + eval right.str val;
    ret_val.Type = TYPE INFO.TYPE STRING;
    ret_val.SymbolName = "";
    return ret val;
  else if (eval_left.Type == TYPE_INFO.TYPE_NUMERIC &&
    eval right. Type == TYPE INFO. TYPE NUMERIC)
    SYMBOL INFO ret val = new SYMBOL INFO();
    ret val.dbl val = eval left.dbl val + eval right.dbl val;
    ret val.Type = TYPE INFO.TYPE NUMERIC;
    ret_val.SymbolName = "";
    return ret_val;
  else
    throw new Exception("Type mismatch");
/// <summary>
/// <param name="local"></param>
/// <param name="global"></param>
public override TYPE_INFO TypeCheck(COMPILATION CONTEXT cont)
  TYPE INFO eval left = exp1.TypeCheck(cont);
  TYPE INFO eval right = exp2.TypeCheck(cont);
  if (eval_left == eval_right && eval_left != TYPE_INFO.TYPE_BOOL)
    _type = eval_left;
    return _type;
  else
    throw new Exception("Type mismatch failure");
public override TYPE_INFO get_type()
  return _type;
```

Where as Evaluate routine for StringLiteral , NumericConstant , BooleanConstant and Variable just involves returning the SYMBOL_INO , in the case of Operators things are bit evolved...

```
public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
  SYMBOL INFO eval left = exp1.Evaluate(cont);
  SYMBOL INFO eval right = exp2.Evaluate(cont);
  if (eval left.Type == TYPE INFO.TYPE STRING &&
    eval right.Type == TYPE INFO.TYPE STRING)
    SYMBOL INFO ret val = new SYMBOL INFO();
    ret val.str val = eval left.str val + eval right.str val;
    ret val.Type = TYPE INFO.TYPE STRING;
    ret_val.SymbolName = "";
    return ret val;
  else if (eval left.Type == TYPE INFO.TYPE NUMERIC &&
    eval right.Type == TYPE INFO.TYPE NUMERIC)
    SYMBOL INFO ret val = new SYMBOL INFO();
    ret_val.dbl_val = eval_left.dbl_val + eval_right.dbl_val;
    ret val.Type = TYPE INFO.TYPE NUMERIC;
    ret_val.SymbolName = "";
    return ret val;
  else
    throw new Exception("Type mismatch");
```

In the above code snippet, Left and Right expressions are evaluated and the types are queried. In our compiler, operations involving numerics and strings are successful only if all the operands are of the same type.

The routine TypeCheck is similar to Evaluate. Only difference is that TypeCheck updates the type information of the nodes in a Recursive manner. The routine get_type is only valid once you have called TypeCheck routine.

BinaryMinus is similar to BinaryPlus. The only difference is only Numerics can be subtracted.

```
class BinaryMinus: Exp
  private Exp exp1, exp2;
  TYPE_INFO _type;
 public BinaryMinus(Exp e1, Exp e2)
    \exp 1 = e1; \exp 2 = e2;
  public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
    SYMBOL INFO eval left = exp1.Evaluate(cont);
    SYMBOL INFO eval right = exp2.Evaluate(cont);
    if (eval left.Type == TYPE INFO.TYPE NUMERIC &&
      eval right. Type == TYPE INFO. TYPE NUMERIC)
      SYMBOL INFO ret val = new SYMBOL INFO();
      ret_val.dbl_val = eval_left.dbl_val - eval_right.dbl_val;
      ret_val.Type = TYPE_INFO.TYPE_NUMERIC;
      ret val.SymbolName = "";
      return ret_val;
    else
      throw new Exception("Type mismatch");
 /// <summary>
 /// <param name="local"></param>
  public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
    TYPE INFO eval left = exp1.TypeCheck(cont);
    TYPE INFO eval right = exp2. TypeCheck(cont);
    if (eval left == eval right && eval left == TYPE INFO.TYPE NUMERIC)
      _type = eval left;
      return _type;
    else
```

```
{
    throw new Exception("Type mismatch failure");
}

public override TYPE_INFO get_type()
{
    return _type;
}
```

Multiplication and Division operators are only valid for Numeric Types. If you have understood the implementation of BinaryPlus , the Mul and Div operators are easy to follow

```
class Mul: Exp
 private Exp exp1, exp2;
 TYPE_INFO_type;
 public Mul(Exp e1, Exp e2)
    \exp 1 = e1; \exp 2 = e2;
  public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
    SYMBOL INFO eval left = exp1.Evaluate(cont);
    SYMBOL_INFO eval_right = exp2.Evaluate(cont);
    if (eval_left.Type == TYPE_INFO.TYPE_NUMERIC &&
      eval right.Type == TYPE INFO.TYPE NUMERIC)
      SYMBOL_INFO ret_val = new SYMBOL_INFO();
      ret val.dbl val = eval left.dbl val * eval right.dbl val;
      ret_val.Type = TYPE_INFO.TYPE_NUMERIC;
      ret val.SymbolName = "";
      return ret val;
    else
      throw new Exception("Type mismatch");
```

```
/// <summary>
 /// </summary>
 /// <param name="local"></param>
 /// <param name="global"></param>
 public override TYPE_INFO TypeCheck(COMPILATION_CONTEXT cont)
    TYPE INFO eval_left = exp1.TypeCheck(cont);
    TYPE_INFO eval_right = exp2.TypeCheck(cont);
    if (eval_left == eval_right && eval_left == TYPE_INFO.TYPE_NUMERIC)
      _type = eval_left;
      return _type;
    else
      throw new Exception("Type mismatch failure");
 public override TYPE_INFO get_type()
    return _type;
/// <summary>
class Div: Exp
 private Exp exp1, exp2;
  TYPE_INFO _type;
 public Div(Exp e1, Exp e2)
    \exp 1 = e1; \exp 2 = e2;
 public override SYMBOL_INFO Evaluate(RUNTIME_CONTEXT cont)
    SYMBOL_INFO eval_left = exp1.Evaluate(cont);
    SYMBOL_INFO eval_right = exp2.Evaluate(cont);
    if (eval_left.Type == TYPE_INFO.TYPE_NUMERIC &&
```

```
eval right.Type = TYPE INFO.TYPE NUMERIC)
    SYMBOL INFO ret val = new SYMBOL INFO();
    ret val.dbl val = eval left.dbl val / eval right.dbl val;
    ret_val.Type = TYPE_INFO.TYPE_NUMERIC;
    ret_val.SymbolName = "";
    return ret_val;
  else
    throw new Exception("Type mismatch");
/// <param name="local"></param>
/// <param name="global"></param>
public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
  TYPE INFO eval left = exp1.TypeCheck(cont);
  TYPE_INFO eval_right = exp2.TypeCheck(cont);
  if (eval left == eval right && eval left == TYPE INFO.TYPE NUMERIC)
    _type = eval_left;
    return _type;
  else
    throw new Exception("Type mismatch failure");
public override TYPE_INFO get_type()
  return _type;
```

UnaryPlus and UnaryMinus is also similar to the implementation of other operators. Both these operators are only applicable for Numeric data type.

```
/// the node to represent Unary +
class UnaryPlus : Exp
  /// <summary>
  /// Plus has got a right expression (exp1)
  /// and a Associated type information
  private Exp exp1;
  TYPE INFO type;
  /// <param name="e2"></param>
  public UnaryPlus(Exp e1)
    exp1 = e1;
  /// <summary>
  public override SYMBOL_INFO Evaluate(RUNTIME_CONTEXT cont)
    SYMBOL_INFO eval_left = exp1.Evaluate(cont);
    if (eval_left.Type == TYPE_INFO.TYPE_NUMERIC)
      SYMBOL INFO ret val = new SYMBOL INFO();
      ret val.dbl val = eval left.dbl val;
      ret_val.Type = TYPE_INFO.TYPE_NUMERIC;
      ret_val.SymbolName = "";
      return ret val;
    else
      throw new Exception("Type mismatch");
  /// <param name="local"></param>
  /// <param name="global"></param>
  public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
```

```
TYPE_INFO eval_left = exp1.TypeCheck(cont);
    if (eval left == TYPE INFO.TYPE NUMERIC)
       _type = eval_left;
      return _type;
    else
      throw new Exception("Type mismatch failure");
  public override TYPE_INFO get_type()
    return _type;
/// the node to represent Unary -
class UnaryMinus : Exp
 /// <summary>
  /// Plus has got a right expression (exp1)
  /// and a Associated type information
  private Exp exp1;
  TYPE_INFO _type;
  /// <param name="e1"></param>
  /// <param name="e2"></param>
  public UnaryMinus(Exp e1)
    exp1 = e1;
  /// <param name="cont"></param>
  public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
    SYMBOL_INFO eval_left = exp1.Evaluate(cont);
    if (eval left.Type == TYPE INFO.TYPE NUMERIC)
      SYMBOL_INFO ret_val = new SYMBOL_INFO();
```

```
ret val.dbl val = -eval left.dbl val;
    ret_val.Type = TYPE_INFO.TYPE_NUMERIC;
    ret_val.SymbolName = "";
    return ret val;
  else
    throw new Exception("Type mismatch");
/// <param name="local"></param>
/// <param name="global"></param>
public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
  TYPE INFO eval left = exp1.TypeCheck(cont);
  if (eval left == TYPE INFO.TYPE NUMERIC)
     _type = eval_left;
    return _type;
  else
    throw new Exception("Type mismatch failure");
public override TYPE_INFO get_type()
  return _type;
```

The statement related nodes are moved to a seperate module by the name AstForStatements.

In this step, we have added support for Variable Declaration and Assignment statement.

The AST for Variable declaration is given below

```
/// Compile the Variable Declaration statements
public class VariableDeclStatement: Stmt
  SYMBOL INFO m inf = null;
  Variable var = null;
  public VariableDeclStatement(SYMBOL INFO inf)
    m inf = inf;
  /// </summary>
  public override SYMBOL INFO Execute(RUNTIME CONTEXT cont)
    cont.TABLE.Add(m inf);
    var = new Variable(m inf);
    return null;
```

In the parser, before we create VariableDeclStatement node, we need to insert the variable's SYMBOL_INFO into the SymbolTable of the COMPILATION_CONTEXT. The VariableDeclStatement node just stores the variable name in the Variable AST.

While Executing the VariableDeclStatement , a Variable is created in the Symbol table of RUNTIME CONTEXT.

Both Compilation Context (COMPILATION_CONTEXT) and Run time Context (RUNTIME_CONTEXT) just contains references to respective symbol tables.

The AST for Assignment statement is given below

```
/// <summary>
/// Assignment Statement
/// <summary>
public class AssignmentStatement : Stmt
{
    private Variable variable;
    private Exp expl;

    public AssignmentStatement(Variable var, Exp e)
    {
        variable = var;
        exp1 = e;

    }

    public AssignmentStatement(SYMBOL_INFO var, Exp e)
    {
        variable = new Variable(var);
        exp1 = e;

    }

    public override SYMBOL_INFO Execute(RUNTIME_CONTEXT cont)
    {
        SYMBOL_INFO val = exp1.Evaluate(cont);
        cont.TABLE.Assign(variable, val);
        return null;
    }
}
```

At this point of time, AST for Statements is as shown below

```
class Stmt
class VariableDeclStatement
class AssignmentStatement
class PrintStatement
class PrintLineStatement
```

The class SymbolTable is just a vector of name/value pair. The source code of the SymbolTable is given below.

```
/// <summary>
/// Symbol Table for Parsing and Type Analysis
/// </summary>
public class SymbolTable
{
/// <summary>
```

```
/// private data structure
private System.Collections.Hashtable dt = new Hashtable();
/// Add a symbol to Symbol Table
/// <param name="s"></param>
public bool Add(SYMBOL INFO s)
  dt[s.SymbolName] = s;
  return true;
/// Retrieve the Symbol
public SYMBOL_INFO Get(string name)
  return dt[name] as SYMBOL INFO;
/// Assign to the Symbol Table
/// <param name="value"></param>
public void Assign(Variable var, SYMBOL_INFO value)
  value.SymbolName = var.GetName();
  dt[var.GetName()] = value;
/// Assign to a variable
/// <param name="var"></param>
public void Assign(string var, SYMBOL INFO value)
  dt[var] = value;
```

The class CsyntaxErrorLog and CsemanticErrorLog (in SupportClasses.cs) is meant for error logging while the compilation process is going on.....

Let us go back to Lexical Analysis stage once again. This time we have added lot of new keywords to the language and Token set has become bit larger than the previous step.

```
public enum TOKEN
     ILLEGAL TOKEN = -1, // Not a Token
     TOK PLUS = 1, // '+'
     TOK_MUL, // '*'
     TOK DIV, // '/'
     TOK SUB, // '-'
     TOK OPAREN, // '('
     TOK CPAREN, // ')'
     TOK NULL, // End of string
     TOK PRINT, // Print Statement
     TOK PRINTLN, // PrintLine
     TOK UNQUOTED STRING, // Variable name, Function name etc
     TOK SEMI, //;
     //----- Addition in Step 4
     TOK_VAR_NUMBER, // NUMBER data type
TOK_VAR_STRING, // STRING data type
TOK_VAR_BOOL, // Bool data type
TOK_NUMERIC, // [0-9]+
TOK_COMMENT, // Comment Token ( presently not used )
     TOK_BOOL_TRUE, // Boolean TRUE
TOK_BOOL_FALSE , // Boolean FALSE
     TOK_STRING, // String Literal
     TOK ASSIGN
                            // Assignment Symbol =
```

We have also moved couple of routines and state variables to Lexer class.

The two notable addition are

```
/// <summary>
/// Current Token and Last Grabbed Token
/// </summary>
protected TOKEN Current_Token; // Current Token
protected TOKEN Last_Token; // Penultimate token
```

Since we have added support for string type..., we need to support string literals (or the last grabbed string) in the lexical analyzer...

```
/// <summary>
/// last_str := Token assoicated with
/// </summary>

public String last_str; // Last grabbed String
```

We need to update the keyword table with additional key words supported by the compiler

```
///
// Fill the Keywords
//
//
keyword[0] = new ValueTable(TOKEN.TOK_BOOL_FALSE, "FALSE");
keyword[1] = new ValueTable(TOKEN.TOK_BOOL_TRUE, "TRUE");
keyword[2] = new ValueTable(TOKEN.TOK_VAR_STRING, "STRING");
keyword[3] = new ValueTable(TOKEN.TOK_VAR_BOOL, "BOOLEAN");
keyword[4] = new ValueTable(TOKEN.TOK_VAR_NUMBER, "NUMERIC");
keyword[5] = new ValueTable(TOKEN.TOK_PRINT, "PRINT");
keyword[6] = new ValueTable(TOKEN.TOK_PRINTLINE");
```

The Parsing of the statements starts from Parse Routine of RDParser.cs

```
/// <summary>
/// The new Parser entry point
/// </summary>
/// <returns></returns>

public ArrayList Parse(COMPILATION_CONTEXT ctx)
{

GetNext(); // Get the Next Token
//
// Parse all the statements
//
return StatementList(ctx);
}
```

Any variable encountered during the Parse Process will be put into the symbol table associated with Compilation Context.

The Logic of the StatementList is as follows,

while there is more statements
Parse and Add Statements to the ArrayList

```
/// The Grammar is
/// <stmtlist> := { <statement> }+
/// {<statement> := <printstmt> | <printlinestmt>
/// <printstmt> := print <expr >;
/// <vardeclstmt> := STRING <varname>; |
     NUMERIC <varname>; |
            BOOLEAN <varname>;
/// <printlinestmt>:= printline <expr>;
/// <Expr> ::= <Term> | <Term> { + | - } <Expr>
/// <Term> ::= <Factor> | <Factor> {*|/} <Term>
/// <Factor>::= <number> | ( <expr> ) | {+|-} <factor>
     <variable> | TRUE | FALSE
private ArrayList StatementList(COMPILATION CONTEXT ctx)
  ArrayList arr = new ArrayList();
  while (Current Token != TOKEN.TOK NULL)
    Stmt temp = Statement(ctx);
    if (temp != null)
       arr.Add(temp);
  return arr;
```

The Statement method just queries the statement type and calls the appropriate Parse Routines

```
/// <summary>
/// This Routine Queries Statement Type
/// to take the appropriate Branch...
/// Currently, only Print and PrintLine statement
/// are supported..
/// if a line does not start with Print or PrintLine ..
/// an exception is thrown
/// </summary>
/// <returns></returns>
private Stmt Statement(COMPILATION_CONTEXT ctx)
```

```
Stmt retval = null;
switch (Current Token)
  case TOKEN.TOK_VAR_STRING:
  case TOKEN.TOK VAR NUMBER:
  case TOKEN.TOK_VAR_BOOL:
    retval = ParseVariableStatement(ctx);
    GetNext();
    return retval;
  case TOKEN.TOK PRINT:
    retval = ParsePrintStatement(ctx);
    GetNext();
    break;
  case TOKEN.TOK PRINTLN:
    retval = ParsePrintLNStatement(ctx);
    GetNext();
    break;
  case TOKEN.TOK UNQUOTED STRING:
    retval = ParseAssignmentStatement(ctx);
    GetNext();
    return retval;
  default:
    throw new Exception("Invalid statement");
return retval;
```

The Source code of the ParseVariableDeclStatement is as given below

```
/// <summary>
/// Parse Variable declaration statement
/// </summary>
/// <param name="type"></param>

public Stmt ParseVariableDeclStatement(COMPILATION_CONTEXT ctx)
{

//--- Save the Data type
    TOKEN tok = Current_Token;

// --- Skip to the next token , the token ought
// to be a Variable name ( UnQouted String )
    GetNext();

if (Current_Token == TOKEN.TOK_UNQUOTED_STRING)
{
    SYMBOL_INFO symb = new SYMBOL_INFO();
    symb.SymbolName = base.last_str;
    symb.Type = (tok == TOKEN.TOK_VAR_BOOL)?
    TYPE_INFO.TYPE_BOOL: (tok == TOKEN.TOK_VAR_NUMBER)?
    TYPE_INFO.TYPE_NUMERIC: TYPE_INFO.TYPE_STRING;

//-------- Skip to Expect the SemiColon
```

```
GetNext();
  if (Current Token == TOKEN.TOK SEMI)
    // ----- Add to the Symbol Table
    // for type analysis
    ctx.TABLE.Add(symb);
    // ----- return the Object of type
    // ----- VariableDeclStatement
    // This will just store the Variable name
    // to be looked up in the above table
    return new VariableDeclStatement(symb);
  else
    CSyntaxErrorLog.AddLine("; expected");
    CSyntaxErrorLog.AddLine(GetCurrentLine(SaveIndex()));
    throw new CParserException(-100, ", or; expected", SaveIndex());
else
  CSyntaxErrorLog.AddLine("invalid variable declaration");
  CSyntaxErrorLog.AddLine(GetCurrentLine(SaveIndex()));
  throw new CParserException(-100, ", or; expected", SaveIndex());
```

Assignment statement is easy to parse as the required ground work has already been done...!

```
/// <summary>
/// Parse the Assignment Statement
/// <variable> = <expr>
/// </summary>
/// <parentment="pb">
/// <parentment="pb"
// <parentment="pb">
/// <parentment="pb"
// <
```

```
//---- The next token ought to be an assignment
// expression....
GetNext();
if (Current Token != TOKEN.TOK ASSIGN)
  CSyntaxErrorLog.AddLine("= expected");
  CSyntaxErrorLog.AddLine(GetCurrentLine(SaveIndex()));
  throw new CParserException(-100, "= expected", SaveIndex());
//----- Skip the token to start the expression
// parsing on the RHS
GetNext();
Exp exp = Expr(ctx);
//---- Do the type analysis ...
if (exp.TypeCheck(ctx) != s.Type)
  throw new Exception("Type mismatch in assignment");
// ----- End of statement (;) is expected
if (Current_Token != TOKEN.TOK_SEMI)
  CSyntaxErrorLog.AddLine("; expected");
  CSyntaxErrorLog.AddLine(GetCurrentLine(SaveIndex()));
  throw new CParserException(-100, "; expected", -1);
// return an instance of AssignmentStatement node..
// s => Symbol info associated with variable
// exp => to evaluated and assigned to symbol info
return new AssignmentStatement(s, exp);
```

Parsing Expressions

The grammar for expression is given below

```
<Expr> ::= <Term> | <Term> { + | - } <Expr> <Term> ::= <Factor> | <Factor> {*|/} <Term> <Factor>::= <number> | ( <expr> ) | {+|-} <factor> <variable> | TRUE | FALSE
```

Let us take a look at the Expression routine, ie the top most expression parsing routine at this point of time... (In future, when logical expressions and relational expressions are added, we modify the

grammar)

The Term routine handles the mul and the div operators

```
/// <summary>
/// <ferm> ::= <Factor> | <Factor> | *|/ <ferm>
/// </summary>
public Exp Term(COMPILATION_CONTEXT ctx)
{
    TOKEN 1_token;
    Exp RetValue = Factor(ctx);

    while (Current_Token == TOKEN.TOK_MUL || Current_Token == TOKEN.TOK_DIV)
{
        | 1_token = Current_Token;
        | Current_Token == GetToken();

        | Exp el = Term(ctx);
        | if (l_token == TOKEN.TOK_MUL)
        | RetValue = new Mul(RetValue, el);
        | else
        | RetValue = new Div(RetValue, el);
        | return RetValue;
        | }

        | return RetValue;
        | }
```

The factor routine is where we handle Variables, unary Operations, Constants etc....

```
/// <Factor>::= <number> | ( <expr> ) | {+|-} <factor>
/// <variable> | TRUE | FALSE
public Exp Factor(COMPILATION CONTEXT ctx)
  TOKEN 1 token;
  Exp RetValue = null;
  if (Current_Token == TOKEN.TOK_NUMERIC)
    RetValue = new NumericConstant(GetNumber());
    Current Token = GetToken();
  else if (Current Token == TOKEN.TOK STRING)
    RetValue = new StringLiteral(last str);
    Current Token = GetToken();
  else if (Current Token = TOKEN.TOK BOOL FALSE ||
       Current Token == TOKEN.TOK BOOL TRUE)
    RetValue = new BooleanConstant(
      Current_Token == TOKEN.TOK_BOOL_TRUE ? true : false);
    Current Token = GetToken();
  else if (Current_Token == TOKEN.TOK_OPAREN)
    Current Token = GetToken();
    RetValue = Expr(ctx); // Recurse
    if (Current Token != TOKEN.TOK CPAREN)
      Console.WriteLine("Missing Closing Parenthesis\n");
      throw new Exception();
    Current Token = GetToken();
  else if (Current Token == TOKEN.TOK PLUS || Current Token == TOKEN.TOK SUB)
    1 token = Current Token;
    Current Token = GetToken();
    RetValue = Factor(ctx);
    if (1 token == TOKEN.TOK PLUS)
      RetValue = new UnaryPlus(RetValue);
    else
      RetValue = new UnaryMinus(RetValue);
```

In the CallSlang Project, this is how we need to invoke the Script....

```
StreamReader sr = new StreamReader(filename);
  string programs2 = sr.ReadToEnd();
  //----- Creates the Parser Object
  // With Program text as argument
  RDParser pars = null;
  pars = new RDParser(programs2);
  // Create a Compilation Context
  COMPILATION CONTEXT ctx = new COMPILATION CONTEXT();
  // Call the top level Parsing Routine with
  // Compilation Context as the Argument
  ArrayList stmts = pars.Parse(ctx);
  // if we have reached here , the parse process
  // is successful... Create a Run time context and
  // Call Execute statements of each statement...
  RUNTIME CONTEXT f = new RUNTIME CONTEXT();
  foreach(Object obj in stmts )
    Stmt s = obj as Stmt;
    s.Execute(f);
/// <param name="args"></param>
static void Main(string[] args)
  if (args == null ||
     args.Length != 1)
    Console.WriteLine("CallSlang <scriptname>\n");
    return;
  TestFileScript(args[0]);
  //----- Wait for the Key Press
  Console.Read();
```

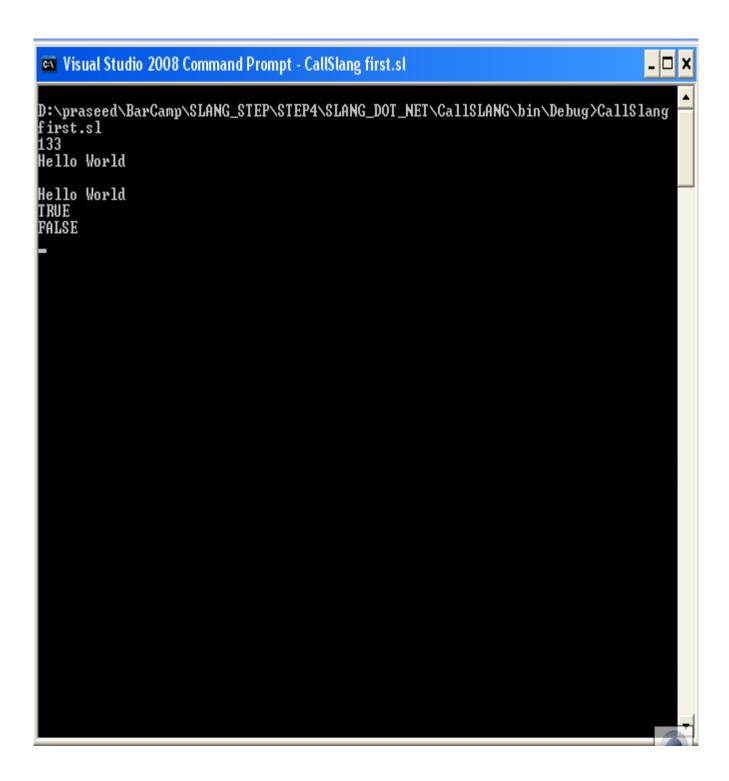
```
First.sl (Slang script)
// A simple SLANG script
// Sample #1
NUMERIC a; // Declare a Numeric variable
a = 2*3+5*30 + -(4*5+3); // Assign
PRINTLINE a; // Dump a
//---- String concatenation
PRINT "Hello " + "World";
//----- Write a new line
PRINTLINE "";
//----string data type
STRING c;
c = "Hello"; // assignment to string
//---- assignment and concatenation
C = C + "World";
PRINTLINE c;
//----- boolean variable
BOOLEAN d;
d= TRUE;
PRINTLINE d;
d= FALSE;
PRINTLINE d;
```

Second.sl

Third.sl

You can go to the command line and invoke the CallSlang executable as follows....a screen snapshot is given in the next page

CallSlang first.sl CallSlang second.sl CallSlang third.sl



SLANG4.NET

htttp://slangfordotnet.codeplex.com

The Art of Compiler Construction using C#

Ву

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CHAPTER 5

Version 0.1

Compilation To a .NET executable!

Till now, the Slang Compiler was interpreting the statements and expressions. In this step we will try to compile the source into .net IL (.net intermediate Language)

In this step, we are not going to make any change into the SLANG grammar. With Future (support for Logical Expression, Relational Expression, Control structures and Procedures) in mind, we will re factor the code.

.NET IL uses a Stack based machine model. Your Intel x86 machines are register based. Stack based machine was touted as the future of machines in the early 1980s. What this basically means , the .NET CLR has got a evaluation stack for processing expressions and functions....

Read more about Reflection. Emit, CLR and IL to understand this.

To give a quick feel of what is involed, let us think about the expression 2*(3+4)

This expression can be decomposed into

```
Acc = ADD 3 and 4 (Acc stands for accumulator)
Acc = Acc MUL 2;
```

if we are using a stack as an accumulator (which .net does) this can be written as

```
PUSH 3 // now stack contains [3]
PUSH 4 // now stack constains [34]
ADD // pop 4, pop 3, add, push result
// Now stack contains [7]
PUSH 2 // Now stack conains [72]
MUL // pop 2, pop 7, multiply, push the result
// now stack contains the result [14]
```

when the same stuff was compiled into .net ...this is what we got

This out put was generated using ILDASM.exe utility. This will help to disassemble .net executables.

To generate .NET executable, in the System.Reflection namespace there are some classes which we need to be familiar.

```
AssemblyBuilder
ModuleBuilder
TypeBuilder
MethodBuilder
```

From the name we can infer that these classes are based on GOF Builder pattern. The .NET StringBuilder class is another notable class which is based on Builder pattern. Consult Wikipedia or a decent book on design patterns to learn more about Builder Pattern. Later, we will be using Builder Pattern to Build (!) Procedure and Module from Program text.

An Assembly (.NET exe or .net DLL) is composed of Collection of Modules. In most cases , we will have one module per assembly. Module will have collection of Types (classes and structs). Each Type will have collection of method. The hierarchical relationonship between these entities are given below.

```
Hierarchy is as follows..

Assembly

Module

Type

Method
```

For each entity there is a corresponding builder available. Please include System.Reflection and System.Reflection.Emit to get access to these classes. How this has been used is given below

```
//
// Get The App Domain
//
//
//
AppDomain _app_domain = Thread.GetDomain();
AssemblyName _asm_name = new AssemblyName();
//
// One can give a strong name , if we want
//
_asm_name.Name = "MyAssembly";
//
// Save the Exe Name
//
_name = name;
//
// Create an instance of Assembly Builder
//
_asm_builder = AppDomain.CurrentDomain.DefineDynamicAssembly(
_asm_name,
_AssemblyBuilderAccess.RunAndSave);
//
// Create a module builder , from AssemblyBuilder
//
_module_builder = _asm_builder.DefineDynamicModule("DynamicModule1", _name, false);
```

```
//
// Create a class by the name MainClass..
// We compile the statements into a static method
// of the type MainClass .. the entry point will
// be called Main
//
_type_builder = _module_builder.DefineType("MainClass");
```

Now one can use Typebuilder instance to add methods to the type.

All the Executable creation logic has been included in a class by the name ExeGenerator.....

```
/// ExeGenerator - Takes care of the creation of
/// .NET executable...
public class ExeGenerator
  /// Hierarchy is as follows..
  /// Assembly
        Module
          Type
            Method
  /// Refer to Reflection.Emit documentation for
  /// more details on creation of .NET executable
  AssemblyBuilder _asm_builder = null;
  ModuleBuilder _ module_builder = null;
  TypeBuilder type builder = null;
  /// Name of the Executable
  string name = "";
  /// Program to be Compiled...
  TModule p = null;
  /// Ctor which takes Executable name
      as parameter
  public ExeGenerator(TModule p, string name)
    // The Program to be compiled...
```

```
_{p} = p;
  // Get The App Domain
  AppDomain app domain = Thread.GetDomain();
  AssemblyName _asm_name = new AssemblyName();
  // One can give a strong name, if we want
  _asm_name.Name = "MyAssembly";
  // Save the Exe Name
  _name = name;
  // Create an instance of Assembly Builder
  _asm_builder = AppDomain.CurrentDomain.DefineDynamicAssembly(
     _asm_name,
    AssemblyBuilderAccess.RunAndSave);
  // Create a module builder, from AssemblyBuilder
  module builder = asm builder.DefineDynamicModule("DynamicModule1", name, false);
  // Create a class by the name MainClass..
  // We compile the statements into a static method
  // of the type MainClass .. the entry point will
  // be called Main
  // ExeGenerator will be called from TModule.Compile method
  // We will add methods to the type MainClass as static method
  _type_builder = _module_builder.DefineType("MainClass");
/// return the type builder....
public TypeBuilder type_bulder
  get
    return type builder;
public void Save()
```

```
//
// Note :- Call this (Save ) method only after
// Compilation of All statements....
_type_builder.CreateType();

//
// Retrieve the Entry Point from TModule....
//
MethodBuilder mb = _p._get_entry_point("MAIN");

if (mb != null)
{
//
// Here we will set the Assembly as a Console Application...
// We will also set the Entry Point....
//
_ asm_builder.SetEntryPoint(mb, PEFileKinds.ConsoleApplication);
}
//
// Write the Resulting Executable...
//
_ asm_builder.Save(_name);
}
```

The Above class will be called from Tmodule.CreateExecutable (more about Tmodule later)

```
public bool CreateExecutable(string name)
{
    //
    // Create an instance of Exe Generator
    // ExeGenerator takes a TModule and
    // exe name as the Parameter...
    _exe = new ExeGenerator(this,name);
    // Compile The module...
    Compile(null);
    // Save the Executable...
    _exe.Save();
    return true;
}
```

To encapsulate the activities to be undertaken to Compile a Procedure to a static method, we have created a new class by the name, DNET EXECUTABLE GENERATION CONTEXT

For Compiling each Subroutine, an instance of this class will be created. This class contains

a) ArrayList to store local variables...

```
/// <summary>
/// Auto (Local) Variable support
/// Stores the index return by DefineLocal
/// method of MethodBuilder
/// </summary>
```

```
private ArrayList variables = new ArrayList();
```

b) Reference to an ILGenerator Object.

```
// <summary>
/// ILGenerator Object
/// </summary>
private ILGenerator ILout;
```

c) SymbolTable for doing type analysis

```
/// <summary>
/// Symbol Table for storing Types and
/// doing the type analysis
/// </summary>
SymbolTable m_tab = new SymbolTable();
```

At this point of time, we are not supplying Procedure name to this class. We do not support user defined subroutines in this step. This class only generates a method by the name MAIN.

By inspecting the Ctor of DNET_EXECTUABLE_GENERATION_CONTEXT, we can get a fair idea of the usage of this class...

```
public DNET EXECUTABLE GENERATION CONTEXT(TModule program,
                           Procedure proc,
                           TypeBuilder bld)
    // All the code in the Source module is compiled
    // into this procedure
    _proc = proc;
    // TModule Object
    _program = program;
    // Handle to the type (MainClass )
     bld = bld;
    // The method does not take any procedure
    System. Type [] s = null;
    // Return type is void
    System.Type ret_type = null;
    // public static void Main()
    _methinfo = _bld.DefineMethod("Main",
       MethodAttributes.Public | MethodAttributes.Static,
       ret type, s);
    // We have created the Method Prologue
    // Get the handle to the code generator
    ILout = methinfo.GetILGenerator();
```

```
}
```

Creataion of Local Variables is not a difficult task. The ILGenerator class has got a method called DeclareLocal for that. The following code snippet will explain the logic involved..

```
public int DeclareLocal(System.Type type)
{
    // It is possible to create Local ( auto )
    // Variables by Calling DeclareLocl method
    // of ILGenerator... this returns an integer
    // We store this integer value in the variables
    // collection...
    LocalBuilder lb = ILout.DeclareLocal(type);
    // Now add the integer value associated with the
    // variable to variables collection...
    return variables.Add(lb);
}
```

The entire source code of DNET EXECUTABLE GENERATION CONTEXT class is given below

```
DNET EXECUTABLE GENERATION CONTEXT is for generating
    CLR executable
public class DNET EXECUTABLE GENERATION CONTEXT
  /// ILGenerator Object
  private ILGenerator ILout;
  /// Auto (Local) Variable support
  /// Stores the index return by DefineLocal
  /// method of MethodBuilder
  private ArrayList variables = new ArrayList();
  /// Symbol Table for storing Types and
     doing the type analysis
  SymbolTable m tab = new SymbolTable();
  /// CLR Reflection.Emit.MethodBuilder
  MethodBuilder methinfo = null;
  /// CLR Type Builder ( useful for creating
  /// classes in the run time
  TypeBuilder bld=null;
  /// Procedure to compiled
```

```
Procedure _proc = null;
/// <summary>
/// Program to be compiled...
TModule program;
/// <param name="prog"></param>
/// <param name="p"></param>
public DNET EXECUTABLE GENERATION CONTEXT(TModule program,
                         Procedure proc,
                         TypeBuilder bld)
  // All the code in the Source module is compiled
 _proc = proc;
  // into this procedure
  // TModule Object
  _program = program;
  // Handle to the type (MainClass )
  _{\rm bld} = {\rm bld};
  // The method does not take any procedure
  System.Type[] s = null;
  // Return type is void
  System.Type ret_type = null;
  // public static void Main()
  _methinfo = _bld.DefineMethod("Main",
    MethodAttributes.Public | MethodAttributes.Static,
    ret_type, s);
  // We have created the Method Prologue
  // Get the handle to the code generator
  ILout = _methinfo.GetILGenerator();
/// <summary>
/// </summary>
public string MethodName
  get
    return proc.Name;
```

```
/// <summary>
public MethodBuilder MethodHandle
  get
    return _methinfo;
/// <summary>
public TypeBuilder TYPEBUILDER
  get
    return _bld;
/// <summary>
/// </summary>
public SymbolTable TABLE
  get
    return m_tab;
/// <summary>
public ILGenerator CodeOutput
  get
    return ILout;
/// </summary>
/// <param name="type"></param>
public int DeclareLocal(System.Type type)
```

```
{
    // It is possible to create Local ( auto )
    // Variables by Calling DeclareLocl method
    // of ILGenerator... this returns an integer
    // We store this integer value in the variables
    // collection...
    LocalBuilder Ib = ILout.DeclareLocal(type);
    // Now add the integer value associated with the
    // variable to variables collection...
    return variables.Add(lb);
}

/// <summary>
/// <param name="s">
/// <param name="s">
/// <returns> 
/// returns> 
/// returns 
/// returns 
// return variables[s] as LocalBuilder;
}
```

COMPILING EXPRESSIONS

To support compilation into .NET IL, we have an addional method appropriately named Compile to the Exp base class.

```
/// <summary>
/// In this Step , we add two more methods to the Exp class
/// TypeCheck => To do Type analysis
/// get_type => Type of this node
/// </summary>
public abstract class Exp
{

public abstract SYMBOL_INFO Evaluate(RUNTIME_CONTEXT cont);
public abstract TYPE_INFO TypeCheck(COMPILATION_CONTEXT cont);
public abstract TYPE_INFO get_type();
/// <summary>
/// Added in the STEP 5 for .NET IL code generation
/// </summary>
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// 
/// Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont);
}
```

This additional method takes a DNET_EXECUTABLE_GENERATION_CONTEXT as the parameter.

To generate Code for Loading a numericthis what we one has to do

The code emitted for a BooleanConstant

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    //
    // Retrieve the IL Code generator and Emit
    // LDC_I4 => Load Constant Integer 4
    // We are planning to use a 32 bit long for Boolean
    // True or False
    cont.CodeOutput.Emit(OpCodes.Ldc_I4, (info.bol_val) ? 1 : 0);
    return true;
}
```

The code emitted for a StringLiteral

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    //
    //    //
    // For string emit
    // LDSTR => Load String
    //
    cont.CodeOutput.Emit(OpCodes.Ldstr , info.str_val);
    return true;
}
```

The Code emitted for loading a variable reference on to the stack

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    //
    // Retrieve the Symbol information from the
    // Symbol Table. Symbol name is the key here..
    //
```

```
SYMBOL_INFO info = cont.TABLE.Get(m_name);

//

// Give the Position to retrieve the Local Variable

// Builder.

//

LocalBuilder lb = cont.GetLocal(info.loc_position);

//

// LDLOC => Load Local... we need to give

// a Local Builder as parameter

//

cont.CodeOutput.Emit(OpCodes.Ldloc, lb);

return true;

}
```

The Code emitted for BinaryPlus (+) needs to query the operand to generate the appropriate code. If the Left and the right expression is numeric... it needs to Generate Add opcode. There is no appropriate operator for String Concatenation. One needs to use String.Concat method to do that....

```
public override bool Compile(DNET EXECUTABLE GENERATION CONTEXT cont)
      // Compile the Left Expression
      exp1.Compile(cont);
      // Compile the Right Expression
      exp2.Compile(cont);
      // Emit Add instruction
      if (_type == TYPE_INFO.TYPE_NUMERIC)
        cont.CodeOutput.Emit(OpCodes.Add);
      else
        // This is a string type..we need to call
        // Concat method..
        Type [] str2 = {
                  typeof(string),
                  typeof(string)
                };
        cont.CodeOutput.Emit(OpCodes.Call,
          typeof(String).GetMethod("Concat", str2));
      return true;
```

For this routine to work correctly, one needs to call TypeCheck every time to update the type information on all nodes.. This polymorphic behavior forced us to call TypeCheck every time we finished parsing an expression.

```
private Stmt ParsePrintStatement(ProcedureBuilder ctx)
{
    GetNext();
    Exp a = Expr(ctx);
    //
    // Do the type analysis ...
    //
    a.TypeCheck(ctx.Context);
    if (Current_Token != TOKEN.TOK_SEMI)
    {
        throw new Exception("; is expected");
    }
    return new PrintStatement(a);
}
```

The code generation of Subtract Operation is as follows

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    exp1.Compile(cont);
    exp2.Compile(cont);
    cont.CodeOutput.Emit(OpCodes.Sub);
    return true;
}
```

The code emitted for UnaryMinus is as given below

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    // Compile the expression
    exp1.Compile(cont);
    //
    // Negate the value on the top of the
    // stack
    //
    cont.CodeOutput.Emit(OpCodes.Neg);
    return true;
}
```

COMPILING STATEMENTS

To support the Compilation of statement into IL an additional method has been added to the Stmt base class

```
public abstract class Stmt
{
    public abstract SYMBOL_INFO Execute(RUNTIME_CONTEXT cont);
    //
    // Added in the Step 5 for .net IL compilation
    public abstract bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont);
}
```

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)

{

// Compile the Expression

// The Output will be on the top of stack
exp1.Compile(cont);

//

// Generate Code to Call Console.Write

//

System.Type typ = Type.GetType("System.Console");
Type[] Parameters = new Type[1];

TYPE_INFO tdata = exp1.get_type();

if (tdata == TYPE_INFO.TYPE_STRING)
Parameters[0] = typeof(string);
else if (tdata == TYPE_INFO.TYPE_NUMERIC)
Parameters[0] = typeof(double);
else
Parameters[0] = typeof(bool);
cont.CodeOutput.Emit(OpCodes.Call, typ.GetMethod("Write", Parameters));
return true;
}
```

The Code emitted for VariableDeclStatement is as follows

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    //
    // Retrieve the type from the SYMBOL_INFO
    //
    System.Type type = (m_inf.Type == TYPE_INFO.TYPE_BOOL)?
        typeof(bool): (m_inf.Type == TYPE_INFO.TYPE_NUMERIC)?
        typeof(double): typeof(string);
    //
    // Get the offset of the variable
    //
    int s = cont.DeclareLocal(type);
    // Store the offset in the SYMBOL_INFO
    //
    m_inf.loc_position = s;
    //
    // Add the variable into Symbol Table..
    //
    cont.TABLE.Add(m_inf);
    return true;
}
```

The code for the Assignment statement is as given below

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    if (!exp1.Compile(cont))
    {
        throw new Exception("Compilation in error string");
    }
    SYMBOL_INFO info = cont.TABLE.Get(variable.Name);
    LocalBuilder lb = cont.GetLocal(info.loc_position);
    cont.CodeOutput.Emit(OpCodes.Stloc, lb);
    return true;
}
```

To support additional language constructs, we had two additional base classes in this project.

The class Procedure is to model a FUNCTION. At this point we compile all the statement into a Function. In future, we will support User defined Function.

```
/// <summary>
/// Abstract base class for Procedure
/// All the statements in a Program ( Compilation unit )
/// will be compiled into a PROC
/// </summary>
public abstract class PROC
{
   public abstract SYMBOL_INFO Execute(RUNTIME_CONTEXT cont);
   public abstract bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont);
}
```

The Concrete class Procedure implements in this interface.

Let us look at the Compile method of the Procedure class. Procedure is nothing but a collection of staements.

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    foreach (Stmt e1 in m_statements)
    {
        e1.Compile(cont);
    }
    cont.CodeOutput.Emit(OpCodes.Ret);
    return true;
}
```

The Execute method of the Procedure is very similar to Compile method. Instead of Calling Compile method , Execute will delegate the call to Execute mathod of statement.

```
public override SYMBOL_INFO Execute(RUNTIME_CONTEXT cont)
{
    foreach (Stmt stmt in m_statements)
        stmt.Execute(cont);
    return null;
}
```

Another interface defined is to support the concept of a module. Module (at this point of time) is just collection of statements. The collection of statements will be compiled into a Procedure. And the Procedure will embedded in a module.

When we support user defined Functions, the meaning of module will be changed a bit. Then the Module will become collection of Functions. All the Functions will be compiled as static methods of a single type. One can easily extend the compiler to support classes.

```
/// <summary>
/// A bunch of statement is called a Compilation
/// unit at this point of time... STEP 5
/// In future, a Collection of Procedures will be
/// called a Compilation unit
///
/// Added in the STEP 5
/// </summary>
public abstract class CompilationUnit
{
public abstract SYMBOL_INFO Execute(RUNTIME_CONTEXT cont);
public abstract bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont);
}
```

Since there is a preexisting type called Module available in the .net class hierarchy, we will use the name Tmodule instead.

```
/// A CodeModule is a Compilation Unit ..
    At this point of time ..it is just a bunch
    of statements...
public class TModule: CompilationUnit
  /// A Program is a collection of Procedures...
  /// Now , we support only global function...
  private ArrayList m_procs=null;
  /// List of Compiled Procedures....
  /// At this point of time..only one procedure
  /// will be there....
  private ArrayList compiled_procs = null;
  /// class to generate IL executable...
  private ExeGenerator _exe = null;
  /// Ctor for the Program ...
  /// <param name="procedures"></param>
  public TModule(ArrayList procs)
     m_procs = procs;
```

```
public bool CreateExecutable(string name)
      // Create an instance of Exe Generator
      // ExeGenerator takes a TModule and
      // exe name as the Parameter...
      exe = new ExeGenerator(this,name);
      // Compile The module...
      Compile(null);
      // Save the Executable...
      exe.Save();
      return true;
    public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
      compiled procs = new ArrayList();
      foreach (Procedure p in m procs)
        DNET EXECUTABLE GENERATION CONTEXT con = new
DNET EXECUTABLE GENERATION CONTEXT(this,p, exe.type bulder);
        compiled procs.Add(con);
        p.Compile(con);
      return true;
    public override SYMBOL INFO Execute(RUNTIME CONTEXT cont)
      Procedure p = Find("Main");
      if(p!=null)
        return p.Execute(cont);
      return null;
    public MethodBuilder get entry point(string funcname)
      foreach (DNET EXECUTABLE GENERATION CONTEXT u in compiled procs)
        if (u.MethodName.Equals( funcname))
          return u.MethodHandle;
```

```
return null;

}

public Procedure Find(string str)
{
  foreach (Procedure p in m_procs)
  {
    string pname = p.Name;
    if (pname.ToUpper().CompareTo(str.ToUpper()) == 0)
        return p;
  }
  return null;
}
```

RETURN OF THE BUILDERS

To create Tmodule and Procedure Object , we need to Parse the statemeents. We can create the Objects only when we have finished the Parsing. We need to accumulate the requisite Objects before we create Procedure. Builder Pattern is a nice way to organize the code to do these kind of stuff.

```
/// <summary>
/// Base class for all the Builder
///
///
///
/// AbstractBuilder
/// TModuleBuilder
/// ProcedureBuilder
/// /summary>
public class AbstractBuilder
{
}
```

The Source code of the Module Builder is given below

```
/// <summary>
/// A Builder for Creating a Module
/// </summary>
class TModuleBuilder : AbstractBuilder
{
/// <summary>
/// Array of Procs
/// </summary>
private ArrayList procs;
```

```
/// Array of Function Prototypes
  /// not much use as of now...
  private ArrayList protos=null;
  /// Ctor does not do much
  public TModuleBuilder()
    procs = new ArrayList();
    protos = null;
 /// Add Procedure
  /// <param name="p"></param>
  public bool Add(Procedure p)
    procs.Add(p);
    return true;
  /// Create Program
  /// </summary>
  public TModule GetProgram()
    return new TModule(procs);
  public Procedure GetProc(string name)
    foreach (Procedure p in procs)
      if (p.Name.Equals(name))
         return p;
    return null;
}
```

The source code of the ProcedureBuilder is given below...

```
public class ProcedureBuilder : AbstractBuilder
  /// Procedure name ..now it is hard coded
  /// to MAIN
  /// </summary>
  private string proc_name = "";
  /// Compilation context for type analysis
  COMPILATION_CONTEXT ctx = null;
  /// Procedure does not take any argument..
  ArrayList m_formals = null;
  /// Array of Statements
  ArrayList m stmts = new ArrayList();
  /// Return Type of the procedure
  TYPE INFO inf = TYPE INFO.TYPE ILLEGAL;
  /// <param name="name"></param>
  /// <param name=" ctx"></param>
  public ProcedureBuilder(string name, COMPILATION CONTEXT ctx)
    ctx = ctx;
    proc_name = name;
  /// <param name="info"></param>
  public bool AddLocal(SYMBOL_INFO info)
    ctx.TABLE.Add(info);
    return true;
  /// <param name="e"></param>
  public TYPE_INFO TypeCheck(Exp e)
```

```
return e.TypeCheck(ctx);
/// </summary>
/// <param name="st"></param>
public void AddStatement(Stmt st)
  m_stmts.Add(st);
/// <param name="strname"></param>
public SYMBOL_INFO GetSymbol(string strname)
  return ctx.TABLE.Get(strname);
/// Check the function Prototype
/// <param name="name"></param>
public bool CheckProto(string name)
  return true;
public TYPE_INFO TYPE
  get
    return inf;
  set
    inf = value;
public SymbolTable TABLE
  get
```

```
return ctx.TABLE;
public COMPILATION_CONTEXT Context
  get
    return ctx;
/// <summary>
public string Name
  get
    return proc_name;
  set
    proc name = value;
public Procedure GetProcedure()
  Procedure ret = new Procedure(proc_name,
      m_stmts, ctx.TABLE, inf);
  return ret;
```

The Parse Process in the RDParser modules starts from the DoParse Method

```
/// <summary>
/// </summary>
/// <returns></returns>

public TModule DoParse()
{

ProcedureBuilder p = new ProcedureBuilder("MAIN", new COMPILATION_CONTEXT());

ArrayList stmts = Parse(p);
```

```
foreach (Stmt s in stmts)
{
    p.AddStatement(s);
}

Procedure pc = p.GetProcedure();

prog.Add(pc);
    return prog.GetProgram();
}
```

In this version all the methods in the Parser takes a ProcedureBuilder object as parameter...

```
private ArrayList StatementList(ProcedureBuilder ctx)
{
    ArrayList arr = new ArrayList();
    while (Current_Token != TOKEN.TOK_NULL)
    {
        Stmt temp = Statement(ctx);
        if (temp != null)
        {
            arr.Add(temp);
        }
    }
    return arr;
}
```

In the previous step , these routines were taking a COMPILATION_CONTEXT as paramter. In this step , this object is encapsulated in ProcedureBuilder object.

Now onwards, we will have three projects in the solution

```
SLANG_DOT_NET => Compilation Engine
SLANGCOMPILE => Calls the Compilation Engine to generate .NET executable
SLANGINTERPRET => Calls the Compilation Engine to interpret the statements
```

Let us take a look at the source code of SLANGCOMPILE

```
using System;
using System.Collections;
using System.Linq;
using System.Text;
using System.IO;
using SLANG_DOT_NET;
namespace SLANGCOMPILE
{
    class Caller
    {
        /// <summary>
        /// Driver routine to call the program script
```

```
static void TestFileScript(string filename)
  if (filename == null)
    return;
  // ----- Read the contents from the file
  StreamReader sr = new StreamReader(filename);
  string programs2 = sr.ReadToEnd();
  sr.Close();
  sr.Dispose();
  //----- Creates the Parser Object
  // With Program text as argument
  RDParser pars = null;
  pars = new RDParser(programs2);
  TModule p = null;
  p = pars.DoParse();
  // Now that Parse is Successul...
  // Create an Executable...!
  if (p.CreateExecutable("First.exe"))
    Console.WriteLine("Creation of Executable is successul");
    return;
/// <param name="args"></param>
static void Main(string[] args)
  if (args == null ||
     args.Length != 1)
    Console.WriteLine("SLANGCOMPILE <scriptname>\n");
  TestFileScript(args[0]);
  //----- Wait for the Key Press
  Console.Read();
```

```
using System;
using System.Collections.Generic;
using System.Ling;
using System. Text;
using System.Collections;
using System.IO;
using SLANG_DOT_NET;
namespace CallSLANG
  class Caller
    /// Driver routine to call the program script
    static void TestFileScript(string filename)
      if (filename == null)
           return;
      // ----- Read the contents from the file
      StreamReader sr = new StreamReader(filename);
      string programs2 = sr.ReadToEnd();
      //----- Creates the Parser Object
      // With Program text as argument
      RDParser pars = null;
       pars = new RDParser(programs2);
       TModule p = null;
      p = pars.DoParse();
      // Now that Parse is Successul...
      // Do a recursive interpretation...!
      RUNTIME CONTEXT f = new RUNTIME CONTEXT(p);
      SYMBOL_INFO fp = p.Execute(f);
    /// </summary>
    /// <param name="args"></param>
    static void Main(string[] args)
      if (args == null ||
         args.Length != 1)
```

Let us how our compiler has generated code for some script....

Using the ILDASM utility, here is what we get from our compiler...

```
.method public static void Main() cil managed
  .entrypoint
 // Code size
                    35 (0x23)
  .maxstack 3
 11_0000: 1dc.r8
                     2.
 1_0009: 1dc.r8
                     3.
 11_0012: 1dc.r8
 11_001b: add
 11_001c: mul
 11_001d: call
                    void [mscorlib]System.Console::writeLine(float64)
 11_0022: ret
} // end of method MainClass::Main
```

The next script will check string Concatentaion

```
PRINT "Hello World" + " SL ";
```

The output generated by our compiler was as follows

```
.method public static void Main() cil managed
```

The next script tests the working of unary expressions.

The code generated by the compiler is as follows

```
.method public static void Main() cil managed
 .entrypoint
 // Code size
                   40 (0x28)
 .maxstack 2
 .locals init (float64 v_0,
         string V_1)
 11_0000: 1dc.r8
 11_0009: neg
 11_000a: neg
 11_000b: neg
 1_000c: stloc.0
 11_000d: 1dloc.0
 11_000e: 1dc.r8
 11_0017: mul
 11_0018: 1dc.r8
                    10.
 11_0021: add
 11_0022: call
                    void [mscorlib]System.Console::WriteLine(float64)
```

```
IL_0027: ret
} // end of method MainClass::Main
```

This script test allmost all the features of the compiler (implemented at this point of time)

```
// A simple SLANG script
// Sample #1
NUMERIC a; // Declare a Numeric variable
a = 2*3+5*30 + -(4*5+3); // Assign
PRINTLINE a; // Dump a
//---- String concatenation
PRINT "Hello " + "World";
//----- Write a new line
PRINTLINE "";
//----string data type
STRING c;
c = "Hello"; // assignment to string
//---- assignment and concatenation
C = C + "World";
PRINTLINE c;
//-----boolean variable
BOOLEAN d;
d= TRUE;
PRINTLINE d;
d= FALSE;
PRINTLINE d;
```

The output is

```
.method public static void Main() cil managed
 entrypoint
 // Code size
                    156 (0x9c)
 .maxstack 4
 .locals init (float64 v_0,
         string V_1,
         bool V_2)
 1_0000: 1dc.r8
                    2.
 1_0009: 1dc.r8
                     3.
 11_0012: mul
 1_0013: 1dc.r8
                     5.
 11_001c: 1dc.r8
                    30.
 1L_0025: mul
 11_0026: 1dc.r8
                    4.
 1_002f: 1dc.r8
                    5.
 1_0038: mul
 11_0039: 1dc.r8
                     3.
 1_0042: add
 11_0043: neg
 11_0044: add
 11_0045: add
 1_0046: stloc.0
 11_0047: 1dloc.0
 11_0048: call
                    void [mscorlib]System.Console::WriteLine(float64)
 1_004d: 1dstr
                     "неПо"
 11_0052: 1dstr
                     "world"
 11_0057: call
                    string [mscorlib]System.String::Concat(string,
                                                       string)
                    void [mscorlib]System.Console::Write(string)
 1L 005c: call
 11_0061: 1dstr
 1_0066: call
                    void [mscorlib]System.Console::WriteLine(string)
 11_006b: 1dstr
                     "неПо"
 1_0070: stloc.1
 1_0071: ldloc.1
 11_0072: 1dstr
                     "world"
 11_0077: call
                    string [mscorlib]System.String::Concat(string,
                                                       string)
 1_007c: stloc.1
 1_007d: 1dloc.1
 11_007e: call
                    void [mscorlib]System.Console::WriteLine(string)
 11_0083: 1dc.i4
                    0x1
 1,0088: stloc.2
```

```
IL_0089: Idloc.2
IL_008a: call void [mscorlib]System.Console::WriteLine(bool)
IL_008f: Idc.i4 0x0
IL_0094: stloc.2
IL_0095: Idloc.2
IL_0096: call void [mscorlib]System.Console::WriteLine(bool)
IL_009b: ret
} // end of method MainClass::Main
```

SLANG4.NET

htttp://slangfordotnet.codeplex.com

The Art of Compiler Construction using C#

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CHAPTER 6

Version 0.1

Supporting Control Structures

In this step , we will support Control Structures like IF statement and While Statement. Before we begin , we need to extend the Exp hierarchy with Relational Expression and Logical Expressions.

Let us extend the list of tokens with Additonal Tokens for Relational expressions , Logical Expressions and Control structures

```
•
```

```
// ----- Added in Step 6
// for relational operator support
                 // '=='
TOK EQ,
                 // '<>'
TOK NEQ,
TOK GT,
                 // '>'
                 // '>='
TOK GTE,
TOK_LT,
                 // '<'
                 // '<='
TOK_LTE,
TOK AND,
                  // '&&'
TOK OR,
                 // '||'
TOK NOT,
                  // '!'
//----- Added in Step 6 for
// Control structures support
TOK IF,
                   // IF
TOK_THEN,
                   // Then
TOK_ELSE,
                  // Else Statement
TOK ENDIF,
                  // Endif Statement
TOK WHILE,
                   // WHILE
TOK WEND
                   // Wend Statement
```

Additonal Keywords supported include

```
//------ Added in the step 6
//------ To support control structures

keyword[7] = new ValueTable(TOKEN.TOK_IF, "IF");
keyword[8] = new ValueTable(TOKEN.TOK_WHILE, "WHILE");
keyword[9] = new ValueTable(TOKEN.TOK_WEND, "WEND");
keyword[10] = new ValueTable(TOKEN.TOK_ELSE, "ELSE");
keyword[11] = new ValueTable(TOKEN.TOK_ENDIF, "ENDIF");
keyword[12] = new ValueTable(TOKEN.TOK_THEN, "THEN");
```

We need to create a new enum for Relational Operators.

Supporting RELATIONAL EXPRESSIONS

Now we have got everything ready to write a class for Relational Expression.

```
public class RelationExp: Exp
{

/// <summary>
/// Which Operator
/// <summary>
RELATION_OPERATOR m_op;
/// <summary>
/// Left and Right Expression
/// <summary>
private Exp ex1, ex2;
/// <summary>
/// Type of this node
/// </summary>
TYPE_INFO_type;
///
/// Operand Types .. if operands are string
/// we need to generate call to String.Compare
/// method...
///
TYPE_INFO_optype;
}
```

The Evaluate routine will give good insight into the implementation of Relational Expression. The basic idea is to Evaluate the Left Expression (Exp1) and Right Expression (Exp2). Once the evaluation is over, one can query the Operand Type (The Operands need to be of same type) and perform the appropriae action.

For Boolean and String type only '==' and '<>' are supported. The TypeChecking routine makes sure that only these operands are permited for Boolean and String Type.

```
/// <summary>
/// The logic of this method is obvious...
/// Evaluate the Left and Right Expression...
/// Query the Type of the expressions and perform
/// appropriate action
/// </summary>
/// <param name="cont"></param>
/// <returns></returns>
```

```
public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
  SYMBOL INFO eval left = ex1.Evaluate(cont);
  SYMBOL INFO eval right = ex2.Evaluate(cont);
  SYMBOL INFO ret val = new SYMBOL INFO();
  if (eval left.Type == TYPE INFO.TYPE NUMERIC &&
    eval_right.Type == TYPE_INFO.TYPE_NUMERIC)
    ret val.Type = TYPE INFO.TYPE BOOL;
    ret val.SymbolName = "";
    if (m op == RELATION OPERATOR.TOK EQ)
      ret_val.bol_val = eval_left.dbl_val == eval_right.dbl_val;
    else if (m op == RELATION OPERATOR.TOK NEQ)
      ret val.bol val = eval left.dbl val != eval right.dbl val;
    else if (m_op == RELATION_OPERATOR.TOK_GT)
      ret val.bol val = eval left.dbl val > eval right.dbl val;
    else if (m op == RELATION OPERATOR.TOK GTE)
      ret val.bol val = eval left.dbl val >= eval right.dbl val;
    else if (m op == RELATION OPERATOR.TOK LTE)
      ret val.bol val = eval left.dbl val <= eval right.dbl val;
    else if (m op == RELATION OPERATOR.TOK LT)
      ret val.bol val = eval left.dbl val < eval right.dbl val;
    return ret val;
  else if (eval left.Type == TYPE INFO.TYPE STRING &&
    eval_right.Type == TYPE_INFO.TYPE_STRING)
    ret val.Type = TYPE INFO.TYPE BOOL;
    ret val.SymbolName = "";
    if(m_op == RELATION_OPERATOR.TOK_EQ)
      ret_val.bol_val = ( String.Compare(
          eval left.str val,
          eval right.str val) == 0)? true:false;
    else if (m op == RELATION OPERATOR.TOK NEQ)
      ret val.bol val = String.Compare(
          eval left.str val,
          eval right.str val) != 0;
    else
      ret val.bol val = false;
```

```
return ret_val;

}
if (eval_left.Type == TYPE_INFO.TYPE_BOOL &&
    eval_right.Type == TYPE_INFO.TYPE_BOOL)
{

ret_val.Type = TYPE_INFO.TYPE_BOOL;
    ret_val.SymbolName = "";

if (m_ op == RELATION_OPERATOR.TOK_EQ)
    ret_val.bol_val = eval_left.bol_val == eval_right.bol_val;
    else if (m_ op == RELATION_OPERATOR.TOK_NEQ)
    ret_val.bol_val = eval_left.bol_val != eval_right.bol_val;
    else
    {
        ret_val.bol_val = false;
    }
    return ret_val;
}

return null;
}
```

Let us take a look at TypeCheck routine of RelationalExp class.

```
/// <summary>
       Recursively check the type and bubble up the type
   /// information to the top...
   public override TYPE INFO TypeCheck(COMPILATION CONTEXT cont)
     TYPE_INFO eval_left = ex1.TypeCheck(cont);
     TYPE INFO eval right = ex2.TypeCheck(cont);
     if (eval_left != eval_right)
       throw new Exception("Wrong Type in expression");
     if (eval left == TYPE INFO.TYPE STRING &&
        (!(m op == RELATION OPERATOR.TOK EQ ||
         m_op == RELATION_OPERATOR.TOK_NEQ)))
       throw new Exception("Only == amd != supported for string type ");
     if (eval_left == TYPE_INFO.TYPE_BOOL &&
        (!(m_op == RELATION_OPERATOR.TOK_EQ ||
         m op == RELATION OPERATOR.TOK NEQ)))
       throw new Exception("Only == amd != supported for boolean type ");
```

```
// store the operand type as well
_optype = eval_left;
_type = TYPE_INFO.TYPE_BOOL;
return _type;
}
```

To Compile a RelationalExpression involving String Operands (== and <>), we need to generate calls to String.Compare method.

```
private bool CompileStringRelOp(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
    // Compile the Left Expression
    ex1.Compile(cont);
    // Compile the Right Expression
    ex2.Compile(cont);
    // This is a string type..we need to call
    // Compare method..
     Type [] str2 = {
                 typeof(string),
                 typeof(string)
           };
       cont.CodeOutput.Emit(OpCodes.Call,
       typeof(String).GetMethod("Compare", str2));
       if (m op == RELATION OPERATOR.TOK EQ)
         cont.CodeOutput.Emit(OpCodes.Ldc I4, 0);
         cont.CodeOutput.Emit(OpCodes.Ceq);
       else
         // This logic is bit convoluted...
         // String.Compare will give 0, 1 or -1
         // First we will check whether the stack value
         // is zero..
         // This will put 1 on stack ..if value was zero
         // after string.Compare
         // Once again check against zero ...it is equivalent
         // to negation
         cont.CodeOutput.Emit(OpCodes.Ldc I4, 0);
         cont.CodeOutput.Emit(OpCodes.Ceq);
         cont.CodeOutput.Emit(OpCodes.Ldc_I4, 0);
         cont.CodeOutput.Emit(OpCodes.Ceq);
```

```
return true;}
```

The code generation for Boolean and Numeric values are similar. in the parsing phase itself, we flag error if some operands are applied on Booleans.

```
public override bool Compile(DNET EXECUTABLE GENERATION CONTEXT cont)
    if ( optype == TYPE INFO.TYPE STRING)
       return CompileStringRelOp(cont);
    // Compile the Left Expression
    ex1.Compile(cont);
    // Compile the Right Expression
    ex2.Compile(cont);
    if (m op == RELATION OPERATOR.TOK EQ)
       cont.CodeOutput.Emit(OpCodes.Ceq);
    else if (m op == RELATION OPERATOR.TOK GT)
       cont.CodeOutput.Emit(OpCodes.Cgt);
    else if (m op == RELATION OPERATOR.TOK LT)
       cont.CodeOutput.Emit(OpCodes.Clt);
    else if (m op == RELATION OPERATOR.TOK NEQ)
       // There is no IL instruction for !=
       // We check for the equivality of the
       // top two values on the stack ...
       // This will put 0 ( FALSE ) or 1 (TRUE)
       // on the top of stack...
       // Load zero and check once again
       // Check == once again...
       cont.CodeOutput.Emit(OpCodes.Ceq);
       cont.CodeOutput.Emit(OpCodes.Ldc I4, 0);
       cont.CodeOutput.Emit(OpCodes.Ceq);
    else if (m_op == RELATION_OPERATOR.TOK_GTE)
       // There is no IL instruction for >=
       // We check for the < of the
       // top two values on the stack ...
       // This will put 0 ( FALSE ) or 1 (TRUE)
       // on the top of stack...
       // Load Zero and
       // Check == once again...
       cont.CodeOutput.Emit(OpCodes.Clt);
       cont.CodeOutput.Emit(OpCodes.Ldc I4, 0);
       cont.CodeOutput.Emit(OpCodes.Ceq);
```

```
else if (m_op == RELATION_OPERATOR.TOK_LTE)
{
    // There is no IL instruction for <=
    // We check for the > of the
    // top two values on the stack ...
    // This will put 0 ( FALSE ) or 1 (TRUE)
    // on the top of stack...
    // Load Zero and
    // Check == once again...

cont.CodeOutput.Emit(OpCodes.Cgt);
    cont.CodeOutput.Emit(OpCodes.Ldc_I4, 0);
    cont.CodeOutput.Emit(OpCodes.Ceq);
}

return true;
}
```

Supporting LOGICAL EXPRESSIONS

SLANG is planning to support AND (&&) , OR (\parallel) and NOT (!) Operators. AND AND OR are binary operators and NOT is a unary operator. We will be having a seperate class for NOT.

```
class LogicalExp : Exp
{
    /// <summary>
    ///    //summary>
    TOKEN m_op;
    ///    //summary>
    /// Operands
    ///    //summary>
    private Exp ex1, ex2;
    ///    //    //summary>
    ///    //    //summary>
    Type of the node...
    ///    //summary>
    TYPE_INFO_type;
}
```

The Type checking logic makes sure that only Boolean types are given to Logical Operators as operands.

```
public override TYPE_INFO TypeCheck(COMPILATION_CONTEXT cont)
{
    TYPE_INFO eval_left = ex1.TypeCheck(cont);
    TYPE_INFO eval_right = ex2.TypeCheck(cont);

// The Types should be Boolean...
// Logical Operators only make sense
```

```
// with Boolean Types

if (eval_left == eval_right &&
      eval_left == TYPE_INFO.TYPE_BOOL )
{
      _type = TYPE_INFO.TYPE_BOOL;
      return_type;
    }
    else
{
      throw new Exception("Wrong Type in expression");
    }
}
```

The Evaluate Routine just apply the requisite operator after the evaluation of Left and Right Expression. The Type checking routine sees to it that both the operands are Boolean.

```
public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
    SYMBOL INFO eval left = ex1.Evaluate(cont);
    SYMBOL_INFO eval_right = ex2.Evaluate(cont);
    if (eval left.Type == TYPE INFO.TYPE BOOL &&
      eval right.Type == TYPE INFO.TYPE BOOL)
      SYMBOL INFO ret val = new SYMBOL INFO();
      ret val.Type = TYPE INFO.TYPE BOOL;
      ret_val.SymbolName = "";
      if (m op == TOKEN.TOK AND)
        ret_val.bol_val = ( eval_left.bol_val && eval_right.bol_val);
      else if (m op == TOKEN.TOK OR)
        ret val.bol val = (eval left.bol val || eval right.bol val);
      else
        return null;
      return ret val;
    return null;
```

The Compilation of Logical And and Logical Or to IL is simple. It is very similar to code generation for any other operator.

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    ex1.Compile(cont);
    ex2.Compile(cont);
    if (m_op == TOKEN.TOK_AND)
        cont.CodeOutput.Emit(OpCodes.And);
```

```
else if (m_op == TOKEN.TOK_OR)
     cont.CodeOutput.Emit(OpCodes.Or);

return true;
}
```

Compiling LOGICAL NOT

Logical Not (!) is a unary operator. We have created a special class for this operator.

```
/// <summary>
 /// Logical!
 class LogicalNot : Exp
   private Exp ex1;
    TYPE_INFO _type;
   /// <summary>
   /// <param name="op"></param>
   public LogicalNot(Exp e1)
      ex1 = e1;
   /// </summary>
   public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
      SYMBOL INFO eval left = ex1.Evaluate(cont);
      if (eval_left.Type == TYPE_INFO.TYPE_BOOL)
        SYMBOL INFO ret val = new SYMBOL INFO();
        ret_val.Type = TYPE_INFO.TYPE_BOOL;
        ret_val.SymbolName = "";
        ret_val.bol_val = !eval_left.bol_val;
        return ret_val;
     else
          return null;
```

```
/// <summary>
public override TYPE_INFO TypeCheck(COMPILATION_CONTEXT cont)
  TYPE INFO eval left = ex1.TypeCheck(cont);
  if (
    eval_left == TYPE_INFO.TYPE_BOOL)
     _type = TYPE_INFO.TYPE_BOOL;
    return _type;
  else
    throw new Exception("Wrong Type in expression");
/// <summary>
/// </summary>
public override bool Compile(DNET EXECUTABLE GENERATION CONTEXT cont)
  ex1.Compile(cont);
  // Check whether top of the stack is 1 (TRUE)
  // Check Whether the previous operation was successful
  // Functionally equivalent to Logical Not
  // Case Top of Stack is 1 (TRUE)
  // Top of Stack \Rightarrow [1]
  // LDC_I4 => [11]
  // CEQ \implies [1]
  // LDC_I4 \Rightarrow [10]
  // CEQ \Rightarrow [0]
  // Case Top of Stack is 0 (FALSE)
  // Top of Stack \Rightarrow [0]
  // LDC_I4 \Rightarrow [01]
  // CEQ \Rightarrow [0]
  // LDC_I4 => [00]
  // CEQ => [1]
  cont.CodeOutput.Emit(OpCodes.Ldc_I4, 1);
  cont.CodeOutput.Emit(OpCodes.Ceq);
  cont.CodeOutput.Emit(OpCodes.Ldc I4, 0);
  cont.CodeOutput.Emit(OpCodes.Ceq);
  return true;
```

```
/// <summary>
/// </summary>
/// <returns></returns>
public override TYPE_INFO get_type()
{
    return _type;
}
```

In the Lexical Analyzer , the code is written for supporing the new operators

```
case "!":
  tok = TOKEN.TOK_NOT;
  index++;
  break;
case '>':
  if (IExpr[index + 1] == '=')
    tok = TOKEN.TOK GTE;
    index += 2;
  else
    tok = TOKEN.TOK_GT;
    index++;
  break;
case '<':
  if(IExpr[index + 1] == '=')
    tok = TOKEN.TOK LTE;
    index += 2;
  else if (IExpr[index + 1] == '>')
    tok = TOKEN.TOK_NEQ;
    index += 2;
  else
    tok = TOKEN.TOK LT;
    index++;
  break;
case '=':
  if(IExpr[index + 1] == '=')
    tok = TOKEN.TOK_EQ;
    index += 2;
  else
```

```
tok = TOKEN.TOK_ASSIGN;
    index++;
  break;
case '&':
  if(IExpr[index + 1] == '\&')
    tok = TOKEN.TOK_AND;
    index += 2;
  else
    tok = TOKEN.ILLEGAL_TOKEN;
    index++;
  break;
case ":
  if(IExpr[index + 1] == '|')
    tok = TOKEN.TOK_OR;
    index += 2;
  else
    tok = TOKEN.ILLEGAL_TOKEN;
    index++;
  break;
case '/':
  if(IExpr[index + 1] == '/')
    SkipToEoln();
    goto re_start;
  else
    tok = TOKEN.TOK_DIV;
    index++;
  break;
```

Extending the Object Model to support IF STATEMENT

The grammar for IF statement is

This can be easily converted to a class which will look as follows....

```
class IfStatement : Stmt
{
    /// <summary>
    /// cond expression
    /// the type ought to be boolean
    /// </summary>
    private Exp cond;
    /// <summary>
    /// List of statements to be
    /// <summary>
    private ArrayList stmnts;
    /// <summary>
    /// List of statements to be
    /// <summary>
    private ArrayList ond is false
    /// </summary>
    private ArrayList else_part;
}
```

The Evaluate Routine of if is as given below

```
/// <summary>
/// Interpret the if statement
/// <param name="cont"></param>
public override SYMBOL_INFO Execute(RUNTIME_CONTEXT cont)
  // Evaluate the Condition
  SYMBOL_INFO m_cond = cond.Evaluate(cont);
  // if cond is not boolean..or evaluation failed
  if (m_cond == null ||
    m_cond.Type != TYPE_INFO.TYPE_BOOL)
    return null;
  if (m cond.bol val == true)
    // if cond is true
    foreach (Stmt rst in stmnts)
       rst.Execute(cont);
  else if (else_part != null)
    // if cond is false and there is
```

```
// else statement ..!
foreach (Stmt rst in else_part)
rst.Execute(cont);

}
return null;
}
```

The if statement generates code as given below

```
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
       Label true label, false label;
      // Generate Label for True
       true_label = cont.CodeOutput.DefineLabel();
       // Generate Label for False
       false_label = cont.CodeOutput.DefineLabel();
      // Compile the expression
       cond.Compile(cont);
      // Check whether the top of the stack contain
      // 1 ( TRUE)
      cont.CodeOutput.Emit(OpCodes.Ldc I4, 1);
       cont.CodeOutput.Emit(OpCodes.Ceq);
      // if False, jump to false label ...
      // ie to else part
       cont.CodeOutput.Emit(OpCodes.Brfalse, false label);
       foreach (Stmt rst in stmnts)
         rst.Compile(cont);
       // Once we have reached here...go
       // to True label...
       cont.CodeOutput.Emit(OpCodes.Br, true label);
      // Place a Label here...if the condition evaluates
      // to false , jump to this place..
       cont.CodeOutput.MarkLabel(false label);
       if (else_part != null)
         foreach (Stmt rst in else part)
           rst.Compile(cont);
```

```
}
}
//
// Place a label here...to mark the end of the
// IF statement
cont.CodeOutput.MarkLabel(true_label);
return true;
}
```

The While statement is very similar to IF statement.

```
class WhileStatement : Stmt
   private Exp cond;
    private ArrayList stmnts;
   /// <param name="s"></param>
    public WhileStatement(Exp c, ArrayList s)
      cond = c;
      stmnts = s;
    /// <param name="cont"></param>
    public override SYMBOL INFO Execute(RUNTIME CONTEXT cont)
   Test:
      SYMBOL_INFO m_cond = cond.Evaluate(cont);
      if (m cond == null || m cond.Type != TYPE INFO.TYPE BOOL)
        return null;
      if (m cond.bol val != true)
        return null;
      SYMBOL_INFO tsp = null;
      foreach (Stmt rst in stmnts)
        tsp = rst.Execute(cont);
        if (tsp != null)
```

```
return tsp;
  goto Test;
/// </summary>
/// <param name="cont"></param>
public override bool Compile(DNET EXECUTABLE GENERATION CONTEXT cont)
  Label true label, false label;
  true label = cont.CodeOutput.DefineLabel();
  false_label = cont.CodeOutput.DefineLabel();
  cont.CodeOutput.MarkLabel(true label);
  cond.Compile(cont);
  cont.CodeOutput.Emit(OpCodes.Ldc I4, 1);
  cont.CodeOutput.Emit(OpCodes.Ceq);
  cont.CodeOutput.Emit(OpCodes.Brfalse, false label);
  foreach (Stmt rst in stmnts)
    rst.Compile(cont);
  cont.CodeOutput.Emit(OpCodes.Br, true label);
  cont.CodeOutput.MarkLabel(false label);
  return true;
```

Modifiying the Parser to support Control Structures

The advent of While and IF automatically requires statement blocks. The StatementList routine should Parse uptil a Block is finished...

At this point of time, Blocks can end when we reach

WEND ENDIF ELSE <END OF INPUT>

```
/// <summary>
/// The Grammar is
///
/// <stmts> := { stmt }+
/// {stmt} := <vardeclstmt> |
/// <printstmt>|<assignmentstmt>|
```

```
<ifstmt>| <whilestmt> |
/// <vardeclstmt> ::= <type> var name;
/// <printstmt> := PRINT <expr>;
/// <assignmentstmt>:= <variable> = value;
/// <ifstmt>::= IF <expr> THEN <stmts> [ ELSE <stmts> ] ENDIF
/// <whilestmt>::= WHILE <expr> <stmts> WEND
/// <type> := NUMERIC | STRING | BOOLEAN
/// <expr> ::= <BExpr>
/// <BExpr> ::= <LExpr> LOGIC OP <BExpr>
/// <LExpr> ::= <RExpr> REL OP <LExpr>
/// <RExpr> ::= <Term> ADD OP <RExpr>
/// <Term>::= <Factor> MUL OP <Term>
/// <Factor> ::= <Numeric> |
           <String> |
           TRUE |
           FALSE |
            "(" <expr> ")" |
            {+|-|!} <Factor>
   <LOGIC OP> := '&&' | '||'
   <REL_OP> := '>' |' < '|' >=' |' <=' |' <>' |' =='
    <MUL OP> := '*' |' /'
/// <ADD OP> := '+' |' -'
private ArrayList StatementList(ProcedureBuilder ctx)
  ArrayList arr = new ArrayList();
  while (
      (Current Token != TOKEN.TOK ELSE) &&
      (Current Token != TOKEN.TOK ENDIF) &&
      (Current Token != TOKEN.TOK WEND) &&
      (Current_Token != TOKEN.TOK_NULL)
    Stmt temp = Statement(ctx);
    if (temp != null)
      arr.Add(temp);
  return arr;
```

In the Statement Method we need to add following code as well..

```
case TOKEN.TOK_IF:
    retval = ParseIfStatement(ctx);
    GetNext();
    return retval;

case TOKEN.TOK_WHILE:
```

```
retval = ParseWhileStatement(ctx);
GetNext();
return retval;
```

The Source code of ParselFStatement is given below

```
/// <summary>
   /// </summary>
   /// <param name="pb"></param>
   public Stmt ParseIfStatement(ProcedureBuilder pb)
      GetNext();
     ArrayList true_part = null;
      ArrayList false part = null;
      Exp exp = BExpr(pb); // Evaluate Expression
      if (pb.TypeCheck(exp) != TYPE_INFO.TYPE_BOOL)
        throw new Exception("Expects a boolean expression");
      if (Current Token != TOKEN.TOK THEN)
        CSyntaxErrorLog.AddLine(" Then Expected");
        CSyntaxErrorLog.AddLine(GetCurrentLine(SaveIndex()));
        throw new CParserException(-100, "Then Expected", SaveIndex());
      GetNext();
      true part = StatementList(pb);
      if (Current Token == TOKEN.TOK ENDIF)
        return new IfStatement(exp, true_part, false_part);
      if (Current_Token != TOKEN.TOK_ELSE)
        throw new Exception("ELSE expected");
      GetNext();
      false part = StatementList(pb);
      if (Current Token != TOKEN.TOK ENDIF)
```

```
throw new Exception("END IF EXPECTED");
}
return new IfStatement(exp, true_part, false_part);
}
```

The source code of ParseWhileStatement is given below

```
public Stmt ParseWhileStatement(ProcedureBuilder pb)
{
    GetNext();
    Exp exp = BExpr(pb);
    if (pb.TypeCheck(exp) != TYPE_INFO.TYPE_BOOL)
    {
        throw new Exception("Expects a boolean expression");
    }

ArrayList body = StatementList(pb);
    if ((Current_Token != TOKEN.TOK_WEND))
    {
        CSyntaxErrorLog.AddLine("Wend Expected");
        CSyntaxErrorLog.AddLine(GetCurrentLine(SaveIndex()));
        throw new CParserException(-100, "Wend Expected", SaveIndex());
    }

    return new WhileStatement(exp, body);
}
```

Given below is some sample script and corresponding IL code generated by the compiler....

OneToHundred.sl

```
PRINTLINE I;
 I = I + 1;
WEND
//----- Print all even numbers
I=0:
WHILE ( I <= 100 )
 PRINTLINE I;
 | = | + 2|
WEND
// ----- Testing String Comparison
STRING s1;
STRING s2;
s1="Hello";
s2="Hell"+"o";
if (s1 == s2) then
 PRINTLINE "S1 and S2 are equal";
endif
IF!(s1 <> s2) then
  PRINTLINE "Tested Logical Not";
endif
```

The IL code generated by the compiler is as follows

```
.method public static void Main() cil managed
 .entrypoint
 // Code size
                   254 (Oxfe)
 .maxstack 8
 .locals init (float64 v_0,
        string V_1,
         string V_2)
 11_0000: 1dc.r8
                  0.0
 1_0009: stloc.0
 1_000a: 1dloc.0
 11_000b: 1dc.r8
                    100.
 11_0014: cqt
 11_0016: 1dc.i4
                   0x0
 11_001b: ceq
 11_001d: 1dc.i4
                   0x1
 11_0022: ceq
 1_0024: brfalse 1_0040
```

```
IL 0029: 1dloc.0
11 002a: call
                  void [mscorlib]System.Console::writeLine(float64)
11_002f: 1dloc.0
11_0030: 1dc.r8
                   1.
1L 0039: add
1 003a: stloc.0
1 1 003b: br
                   1 ■ 000a
IL 0040: 1dc.r8
                   0.0
1,0049: stloc.0
11_004a: 1dloc.0
11_004b: 1dc.r8
                   100.
11_0054: cgt
11_0056: 1dc.i4
                  0x0
11_005b: ceq
11_005d: 1dc.i4
                  0x1
11_0062: ceq
1,0064: brfalse 1,0080
11.0069: 1dloc.0
11 006a: call
                  void [mscorlib]System.Console::WriteLine(float64)
11_006f: Idloc.0
11.0070: 1dc.r8
                   2.
11_0079: add
1_007a: stloc.0
1 1 007b: br
                   1L 004a
                   "неПо"
11.0080: 1dstr
1_0085: stloc.1
                   "HeTT"
11_0086: 1dstr
11_008b: 1dstr
                   "o"
11_0090: call
                  string [mscorlib]System.String::Concat(string,
                                                     strina)
1 0095: stloc.2
11_0096: Idloc.1
11_0097: 1dloc.2
11_0098: call
                  int32 [mscorlib]System.String::Compare(string,
                                                     string)
11_009d: 1dc.i4
                  0x0
11_00a2: ceq
11_00a4: 1dc.i4
                  0x1
11_00a9: ceq
■ 00ab: brfalse
                   11 00bf
11_00b0: 1dstr
                   "S1 and S2 are equal"
                  void [mscorlib]System.Console::WriteLine(string)
11_00b5: call
IL 00ba: br
                   1 ■ 00bf
1_00bf: 1dloc.1
11_00c0: 1dloc.2
11 00c1: call
                  int32 [mscorlib]System.String::Compare(string,
                                                     string)
11_00c6: 1dc.i4
                  0x0
11_00cb: ceq
11_00cd: 1dc.i4
                  0x0
11_00d2: ceq
11. 00d4: 1dc.i4
                  0x1
```

```
IL_00d9: ceq
IL_00db: ldc.i4  0x0
IL_00e0: ceq
IL_00e2: ldc.i4  0x1
IL_00e7: ceq
IL_00e9: brfalse  IL_00fd
IL_00ee: ldstr  'Tested Logical Not''
IL_00f3: call  void [mscorlib]System.Console::WriteLine(string)
IL_00f6: ret
} // end of method MainClass::Main
```

SLANG4.NET

htttp://slangfordotnet.codeplex.com

The Art of Compiler Construction using C#

Ву

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CHAPTER 7

Version 0.1

FUNCTIONS

This step will support Sub routine call. This will make the language useful for doing some tasks which were diffcult to acheive earlier.

The Grammar for Function is

We need to add one more class to Stmt hierarchy. The purpose is to implment the Return Value.

<retstmt> := Return <expr>;

```
}
/// <summary>
/// </param name="cont"></param>
/// <returns></returns>
public override bool Compile(DNET_EXECUTABLE_GENERATION_CONTEXT cont)
{
    if (m_el != null)
    {
        m_el.Compile(cont);
    }
    cont.CodeOutput.Emit(OpCodes.Ret);
    return true;
}
```

We need to Add an additional node to Exp class heirarchy to model Function call.

```
/// <summary>
 /// The node to model Function Call
 /// in the Expression hierarchy...
 class CallExp: Exp
   /// <summary>
  /// Procedure Object
   Procedure m proc;
   /// ArrayList of Actuals
  ArrayList m actuals;
   /// <summary>
   /// procedure name ...
   string _procname;
   /// <summary>
   /// Is it a Recursive Call?
   bool isrecurse;
   /// Return type of the Function
   TYPE_INFO _type;
   /// Ctor to be called when we make a ordinary
   /// subroutine call
   /// <param name="proc"></param>
   /// <param name="actuals"></param>
   public CallExp(Procedure proc, ArrayList actuals)
     m_proc = proc;
```

```
m actuals = actuals;
/// Ctor to implement Recursive sub routine
/// <param name="actuals"></param>
public CallExp(string name, bool recurse, ArrayList actuals)
   procname = name;
  if (recurse)
     isrecurse = true;
  m_actuals = actuals;
  // For a recursive call Procedure Address will be null
  // During the interpretation time we will resolve the
  // call by look up...
  // m_proc = cont.GetProgram().Find(_procname);
  // This is a hack for implementing one pass compiler
  m proc = null;
public override SYMBOL INFO Evaluate(RUNTIME CONTEXT cont)
  if (m_proc != null)
    // This is a Ordinary Function Call
    RUNTIME_CONTEXT ctx = new RUNTIME_CONTEXT(cont.GetProgram());
    ArrayList lst = new ArrayList();
    foreach (Exp ex in m actuals)
       lst.Add(ex.Evaluate(cont));
    return m proc.Execute(ctx, lst);
  else
    // Recursive function call...by the time we
    // reach here..whole program has already been
    // parsed. Lookup the Function name table and
    // resolve the Address
    m_proc = cont.GetProgram().Find(_procname);
    RUNTIME_CONTEXT ctx = new RUNTIME_CONTEXT(cont.GetProgram());
```

```
ArrayList lst = new ArrayList();
    foreach (Exp ex in m_actuals)
       lst.Add(ex.Evaluate(cont));
    return m_proc.Execute(ctx, lst);
/// <param name="cont"></param>
public override TYPE_INFO TypeCheck(COMPILATION_CONTEXT cont)
  if (m_proc != null)
    _type = m_proc.TypeCheck(cont);
  return _type;
/// <summary>
/// <returns></returns>
public override TYPE_INFO get_type()
  return _type;
/// <summary>
/// <param name="cont"></param>
public override bool Compile(DNET EXECUTABLE GENERATION CONTEXT cont)
  if (m_proc == null)
    // if it is a recursive call..
    // resolve the address...
    m_proc = cont.GetProgram().Find(_procname);
  string name = m proc.Name;
  TModule str = cont.GetProgram();
  MethodBuilder bld = str._get_entry_point(name);
```

```
foreach (Exp ex in m_actuals)
{
    ex.Compile(cont);
}
cont.CodeOutput.Emit(OpCodes.Call, bld);
return true;
}
```

To support the function call , we need to modify the Parser a bit to support Function Invocation.

```
else if (Current Token == TOKEN.TOK UNQUOTED STRING)
        String str = base.last_str;
        if (!prog.IsFunction(str))
          // if it is not a function..it ought to
          // be a variable...
          SYMBOL INFO inf = ctx.GetSymbol(str);
          if (inf == null)
             throw new Exception("Undefined symbol");
          GetNext();
          return new Variable(inf);
        // P can be null, if we are parsing a
        // recursive function call
        Procedure p = prog.GetProc(str);
        // It is a Function Call
        // Parse the function invocation
        Exp ptr = ParseCallProc(ctx, p);
        GetNext();
        return ptr;
```

The ParseCallProc subroutine takes care of the compilation of Subroutine call.

```
/// <summary>
///
/// </summary>
/// <param name="pb"></param>
```

```
/// <param name="p"></param>
public Exp ParseCallProc(ProcedureBuilder pb, Procedure p)
  GetNext();
  if (Current_Token != TOKEN.TOK_OPAREN)
    throw new Exception("Opening Parenthesis expected");
  GetNext();
  ArrayList actualparams = new ArrayList();
  while (true)
    // Evaluate Each Expression in the
    // parameter list and populate actualparams
    // list
    Exp exp = BExpr(pb);
    // do type analysis
    exp.TypeCheck(pb.Context);
    // if , there are more parameters
    if (Current Token == TOKEN.TOK_COMMA)
       actualparams.Add(exp);
       GetNext();
       continue;
    if (Current Token != TOKEN.TOK CPAREN)
       throw new Exception("Expected paranthesis");
    else
      // Add the last parameters
       actualparams.Add(exp);
      break;
  // if p is null, that means it is a
  // recursive call. Being a one pass
  // compiler , we need to wait till
  // the parse process to be over to
  // resolve the Procedure.
  if(p!=null)
    return new CallExp(p, actualparams);
  else
    return new CallExp(pb.Name,
               true, // recurse!
```

```
actualparams);
```

Parsing Function Defenition....

The actual routine to parse a subroutine is given in the ParseFunction method of the RDParser class

```
/// <summary>
  /// Parse A Single Function.
   ProcedureBuilder ParseFunction()
     // Create a Procedure builder Object
     ProcedureBuilder p = new ProcedureBuilder("", new COMPILATION CONTEXT());
     if (Current Token != TOKEN.TOK FUNCTION)
       return null;
     GetNext();
     // return type of the Procedure ought to be
     // Boolean , Numeric or String
     if (!(Current Token == TOKEN.TOK VAR BOOL ||
       Current Token == TOKEN.TOK VAR NUMBER ||
       Current Token = TOKEN.TOK VAR STRING))
       return null;
     //----- Assign the return type
     p.TYPE = (Current Token == TOKEN.TOK VAR BOOL)?
       TYPE INFO.TYPE BOOL: (Current Token == TOKEN.TOK VAR NUMBER)?
       TYPE INFO.TYPE NUMERIC: TYPE INFO.TYPE STRING;
     // Parse the name of the Function call
     GetNext();
     if (Current Token != TOKEN.TOK UNQUOTED STRING)
       return null;
     p.Name = this.last_str; // assign the name
     // ----- Opening parenthesis for
     // the start of <paramlist>
     GetNext();
     if (Current_Token != TOKEN.TOK_OPAREN)
       return null;
     //---- Parse the Formal Parameter list
     FormalParameters(p);
```

```
if (Current_Token != TOKEN.TOK_CPAREN)
    return null;

GetNext();

// ------- Parse the Function code
ArrayList lst = StatementList(p);

if (Current_Token != TOKEN.TOK_END)
{
    throw new Exception("END expected");
}

// Accumulate all statements to
// Procedure builder
//
foreach (Stmt s in lst)
{
    p.AddStatement(s);
}

return p;
}
```

The Method Formal Paramaters parse the formal parameter list and add to the Function Prototype list in the Tmodule class.

In the function given below, a and b are called Formal Paramters...

```
FUNCTION BOOLEAN TEST( NUMERIC a , STRING b )
PRINTLINE a;
PRINTLINE b;
RETURN true;
END
```

The above function might be called with actual parameters (can be expressions) as follows

```
TEST(10,"Hello World")
```

We need to bind 10 to a and "Hello World" to b before execution of the code..

```
/// <summary>
/// </summary>
/// <param name="pb"></param>
void FormalParameters(ProcedureBuilder pb)
{

if (Current_Token != TOKEN.TOK_OPAREN)
throw new Exception("Opening Parenthesis expected");
GetNext();
```

```
ArrayList lst_types = new ArrayList();
while (Current Token = TOKEN.TOK VAR BOOL ||
  Current_Token == TOKEN.TOK_VAR_NUMBER ||
  Current Token = TOKEN.TOK VAR STRING)
  SYMBOL INFO inf = new SYMBOL INFO();
  inf.Type = (Current_Token == TOKEN.TOK_VAR_BOOL)?
    TYPE INFO.TYPE BOOL: (Current Token == TOKEN.TOK VAR NUMBER)?
    TYPE_INFO.TYPE_NUMERIC: TYPE_INFO.TYPE_STRING;
  GetNext();
  if (Current Token != TOKEN.TOK UNQUOTED STRING)
    throw new Exception("Variable Name expected");
  inf.SymbolName = this.last str;
  lst types.Add(inf.Type);
  pb.AddFormals(inf);
  pb.AddLocal(inf);
  GetNext();
  if (Current Token != TOKEN.TOK COMMA)
    break;
  GetNext();
prog.AddFunctionProtoType(pb.Name, pb.TYPE, lst_types);
return;
```

To Parse Return Statement , we need to a block in the Statement method of RDParse class as given below

```
case TOKEN.TOK_RETURN:
    retval = ParseReturnStatement(ctx);
    GetNext();
    return retval;
```

The Parser entry point is changed a bit to support Function call

```
/// <summary>
/// The new Parser entry point
/// </summary>
/// <returns ></returns >
public TModule DoParse()
{
    try
    {
        GetNext(); // Get The First Valid Token
        return ParseFunctions();
    }
    catch (Exception e)
    {
        Console.WriteLine("Parse Error -----");
        Console.WriteLine(e.ToString());
        return null;
    }
}
```

A module is nothing but a list of functions. The above statement is computationally equivalent to

```
public TModule ParseFunctions()
{
    while (Current_Token == TOKEN.TOK_FUNCTION)
    {
        ProcedureBuilder b = ParseFunction();
        Procedure s = b.GetProcedure();

        if (s == null)
        {
                  Console.WriteLine("Error While Parsing Functions");
                 return null;
        }

        prog.Add(s);
        GetNext();
    }

/// Convert the builder into a program
///
// return prog.GetProgram();
}
```

The Formal Grammar for SLANG at this point of time is as follows

```
<Module> ::= {<Procedure>}+;
<Procedure>::= FUNCTION <type> func name '(' arglist ')'
              <stmts>
         FND
<type> := NUMERIC | STRING | BOOLEAN
arglist ::= '(' {} ')' | '(' <type> arg name [, arglist ] ')'
<stmts> := { stmt }+
{stmt} := <vardeclstmt> | <printstmt>|<printlnstmt>
       <assignmentstmt>|<callstmt>|<ifstmt>|
       <whilestmt> | <returnstmt>
<vardeclstmt> ::= <type> var name;
<printstmt> := PRINT <expr>;
<assignmentstmt>:= <variable> = value;
<ifstmt>::= IF <expr> THEN <stmts> [ ELSE <stmts> ] ENDIF
<whilestmt>::= WHILE <expr> <stmts> WEND
<returnstmt>:= Return <expr>
<expr> ::= <BExpr>
<BExpr> ::= <LExpr> LOGIC OP <BExpr>
<LExpr> ::= <RExpr> REL_OP <LExpr>
<RExpr> ::= <Term> ADD OP <RExpr>
<Term>::= <Factor> MUL OP <Term>
<Factor> ::= <Numeric> | <String> | TRUE | FALSE | <variable> | '(' <expr> ')' | {+|-|!}
<Factor> | <callexpr>
<callexpr> ::= funcname '(' actuals ')'
<LOGIC OP> := '&&' | '||'
<REL OP> := '>' |' < '|' >=' |' <=' |' <>' |' =='
<MUL OP> := '*' |' /'
<ADD OP> := '+' |' -'
```

Here are some sample programs according to the above grammar

```
ENDIF
return 0;

END

FUNCTION BOOLEAN MAIN()
NUMERIC d;
d= Quad(1,0-6,9);

IF ( d == 0 ) then
PRINT "No Roots";
ELSE
IF ( d == 1 ) then
PRINT "Discriminant is zero";
ELSE
PRINT "Two roots are available";
ENDIF
ENDIF
ENDIF
```

```
//
// A simple Slang Script ... can only run with
// STEP 7 or above
//
// This code will print Fibonacci series between 1
// and 1000
FUNCTION BOOLEAN MAIN()
 NUMERIC newterm;
 NUMERIC prevterm;
 NUMERIC currterm;
 currterm = 1;
 prevterm = 0;
 newterm = currterm + prevterm;
 PrintLine newterm;
while ( newterm < 1000 )
  prevterm = currterm;
  currterm = newterm;
  newterm = currterm + prevterm;
```

```
PrintLine newterm;
wend
END
```

```
// Recursive Fibonacci series routine
//
FUNCTION NUMERIC FIB( NUMERIC n )
   IF ( n <= 1 ) then
      return 1;
   ELSE
     RETURN FIB(n-1) + FIB(n-2);
   ENDIF
END
//
// Main routine
//
FUNCTION BOOLEAN MAIN()
NUMERIC d;
d=0;
While ( d <= 10 )
 PRINTLINE FIB(d);
  d = d+1;
Wend
END
```

```
//
FUNCTION NUMERIC FACT( NUMERIC d )
  IF ( d <= 0 ) THEN
     return 1:
  ELSE
     return d*FACT(d-1);
  ENDIF
END
//
// Entry Point
//
FUNCTION BOOLEAN MAIN()
NUMERIC d;
d=0;
While ( d <= 10 )
  PRINTLINE FACT(d);
  d = d+1;
Wend
END
```

The Compiler/Interpreter can be used as follows... Copy all the scripts to a particular directory...copy the Slanginterpret.exe and SlangCompile.exe to the same folder. Go to Visual studio command prompt

To compile scripts to .net executable

SLANGCOMPILE <scriptname>

This will produce first.exe ... you can execute the file by typing first.exe at the command prompt.

Eg :-SlangCompile onetohundred.sl SlangCompile Fibrec.sl

To interpret the script

SLANGINTERPRET <scriptname>

The above scripts have run successfully. There might be some bugs in the code...that is given as a HOMEWORK for you. !!!!!!