

# Homework 1

(75 points)

Name: Aima Salman

1. Order the following functions by growth rate (12 points) Indicate which of the functions grow at the same rate.

LOWEST GROWTH

$2/N$

$37 = 2^{50}$

$\text{Sqrt}(N)$

$N$

$N \log \log N$

$N \log N = N \log(N^2)$

$N \log^2 N$

$N^{1.5}$

$N^2$

$N^2 \log N$

$N^3$

$2^N$

HIGHEST GROWTH

2. Give the Big-O notation for the following expressions: (10pts, 2pt each)

a.  $2n^4 + 3n^3 - 5 = O(n^4)$

b.  $4^n - n^2 + 19 = O(4^n)$

c.  $\frac{5}{3}n = O(n)$

d.  $5n * \log(n) + 8 = O(n \log(n))$

e.  $[n(n+1)/2 + 2n] / 3 = O(n^2)$

3. For each of the following code fragments give running time analysis (Big Oh). Explain your answer (25pt, 5pts each)

```
a. sum = 0;
   for ( i=0; i < n ; i++)
       sum++;
```

The runtime of this code fragment would be  $O(n)$  because, in the worst-case scenario, the loop can run  $n$  times. This is a linearly increasing growth rate. Because there is only one loop, we use that runtime to represent the entire code segment.

```
b. sum = 0;  
    for( i = 0; i < n; i++)  
        for(j = 0; j < i ; j++)  
            sum++;
```

The runtime of this code fragment would be  $O(n^2)$  because, in the worst-case scenario, the loop can run  $n^2$  times. This is a quadratically increasing growth rate. Both the outer and inner loops can run  $n$  times each, therefore, each having run times of  $O(n)$ . Because the loops are within each other, we combine the run times by multiplying them to represent the entire code segment.

```
c. sum = 0;  
    for( i = 0; i < n; i++)  
        for( j = 0; j < i*i ; j++)  
            for( k = 0; k<j; k++)  
                sum++;
```

The runtime of this code fragment would be  $O(n^5)$  because, in the worst-case scenario, the loop can run  $n^5$  times. This is a polynomial-ly increasing growth rate. The outermost loop runs  $n$  times, the middle one runs  $n^2$  times because it squares the number from the previous loop, and the innermost loop runs  $n^2$  times as well because it runs the same amount as the previous loop ( $j$  times). From the outermost to the innermost loop, the runtimes are  $O(n)$ ,  $O(n^2)$ , and  $O(n^2)$ . Because the loops are within each other, we combine the run times by multiplying them to represent the entire code segment.

```
d. if(value < n)  
    for( i = 0; i < n; i++)  
    {  
        System.out.println(i);  
    }  
else  
    System.out.println(value);
```

The runtime of this code fragment would be  $O(n)$  because, in the worst-case scenario, the loop can run  $n$  times. This is a linearly increasing growth rate. Because there is only one loop, we use that runtime to represent the entire code segment. The if-else statements do not affect the runtime as they are  $O(1)$ .

```
e. sum2 = 0;  
    sum5 = 0;  
  
    for(i=1; i<=n/2; i++)
```

```

{
    sum2 = sum + 2;
}
for(j=1; j<=n*n; j++)
{
    sum5 = sum + 5;
}

```

The runtime of this code fragment would be  $O(n^2)$  because, in the worst-case scenario, the loop can run  $n^2$  times. This is a quadratically increasing growth rate. The first loop can run  $n/2$  times and the second one  $n^2$  times, therefore making the runtimes  $O(n)$  and  $O(n^2)$  respectively. Because the loops are not within each other, we pick the loop with the worst run time to represent the entire code segment.

4. *What is the time complexity of the below function? (10 points)*

```

void fun(int n, int arr[])
{
    int i = 0, j = 0;
    for(; i < n; ++i)
        while(j < n && arr[i] < arr[j])
            j++;
}

```

The runtime of this code would be  $O(n)$  because, in the worst-case scenario, the loop can run  $n$  times. This is a linearly increasing growth rate. Although the loops are within each other,  $j$  does not reset with each iteration of  $i$ . Therefore, it just continues to increase. So, the most that the loops can run is  $2n$ .

5. *What is the Big-O running time for this code? Explain your answer. (10 points)*

```

int i = numItems;
while (i > 0)
{
    i = i / 2; // integer division will eventually reach
zero
}

```

The runtime of this code would be  $O(\log(n))$  because, in the worst-case scenario, the loop can run  $\log(n)$  times. This is a logarithmically increasing growth rate. Each iteration of this loop looks like  $n/2^k$ , so it will be a logarithmic function when solved.

6. *An algorithm takes 0.5 ms for input size 100. How long will it take for input size 500 if the running time is the following (assume lower order terms are negligible) (8 Points)*
- Linear*

- $\frac{T(500)}{T(100)} = \frac{500}{100}$
- $\frac{T(500)}{0.5} = \frac{500}{100}$
- $T(500) = 2.50ms$

*b.  $O(N \log N)$*

- $\frac{T(500)}{T(100)} = \frac{500 \log(500)}{100 \log(100)}$
- $\frac{T(500)}{0.5} = \frac{5 \log(500)}{\log(100)}$
- $T(500) = 3.37ms$

*c. Quadratic*

- $\frac{T(500)}{T(100)} = \frac{500^2}{100^2}$
- $\frac{T(500)}{0.5} = \frac{250000}{10000}$
- $T(500) = 12.5ms$

*d. Cubic*

- $\frac{T(500)}{T(100)} = \frac{500^3}{100^3}$
- $\frac{T(500)}{0.5} = \frac{125000000}{1000000}$
- $T(500) = 62.5ms$