Basic terms:

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- ▶ **DeMorgan's Law**: $\neg(p \land q) = \neg p \lor \neg q$, $\neg(p \lor q) = \neg p \land \neg q$

For which values of p, q, and r is the following expression true? Give a succinct description (not the full truth table):

$$(\neg q \vee r) \wedge (p \to q) \wedge (\neg r \vee \neg p)$$

Show that the following two expressions are not logically equivalent:

$$(q \to r) \lor p$$
$$q \to (r \lor p)$$

What is the contrapositive of the statement: "If it rains, then the ground gets wet"?

Negate the statement: "If I study and sleep well, then I will pass the exam."

Use DeMorgan's Law to simplify: $\neg(p \lor \neg q)$

Which of the following are logically equivalent?

A.
$$\neg(p \land q)$$

B.
$$\neg p \lor \neg q$$

C.
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- ▶ Special notation: $\emptyset = \{\}$, the empty set.

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- ▶ Powerset (set of all subsets): $\mathbb{P}(A) = \{S \mid S \subseteq A\}$

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- ▶ Sets need not have the same "type" of element: $\{1, (3, \text{``no"}), f\}$ is a perfectly valid set.

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How would you prove that $A \nsubseteq B$? How about $A \subsetneq B$?

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- 3. Prove that $\{x^2 \mid x \in \mathbb{R}\} = \{y \ge 0 \mid y \in \mathbb{R}\}$

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- Stars and Bars: How many ways to partition n identical elements into k bins: $\binom{k+n-1}{k-1}$

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- 5. Same question but for *positive* numbers.
- 6. (Summer 2024 Review Question) Let $A = \{(a,b) \in \mathbb{R}^2 \mid a = 3 b^2\}$, $B = \{(x,y) \in \mathbb{R}^2 \mid |x| \ge 1 \text{ or } |y| \ge 1\}$. Prove that $A \subseteq B$.

Questions/Examples

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 - $f: \{2,3,4\} \to \mathbb{N}$ f(x) = 2x 1



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- ▶ So f being bijective (onto and one-to-one) means every $b \in B$ has exactly one preimage.

Which of these functions are onto, one-to-one, or both?

Signature	f(x)	Onto?	One-to-one?
$\mathbb{R} o \mathbb{R}$	x^2		
$[0,\infty) \to [0,\infty)$	x^2		
$\mathbb{N} o \mathbb{Z}$	$(-1)^x x$		
$\{2,5,6\} \rightarrow \{3,6,7,8\}$	x+1		
$\mathbb{N} \cup \{-1\} \to \mathbb{N}$	2x+1		

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- 3. Suppose we have finite sets A and B with $|A| = |B| = n \in \mathbb{N}$. Show that if $f: A \to B$ is one-to-one, then f is onto. How many such functions are there, for fixed A and B?

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- 5. Suppose that $f \circ g$ is one-to-one. Does f have to be one-to-one? Does g?

Recurrences: Review

▶ We can specify functions with a recursive formula:

$$T(n) = \begin{cases} c & n \leq B \\ \langle \text{Formula with smaller arguments to } T \rangle & n > B \end{cases}$$

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► Some examples:

	Recurrence		Closed Formula	
$T(n) = \begin{cases} $	1	$n \leq 1$	$T(n) = 2^n$	
	2T(n-1)	$n \ge 1$		
$T(n) = \begin{cases} 1 \\ 2 \end{cases}$			T(n) = 2n - 1	
	2T(n/2) + 1	n > 1	I(n) = 2n - 1	
$T(n) = \begin{cases} 0 \\ 2 \end{cases}$	/		$T(n) = n\log(n)$	
	2T(n/2) + n	n > 1	$I(n) = n \log(n)$	

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Example:
$$T(n) = \begin{cases} 0 & n \leq 1 \\ 2T(n/2) + n & n > 1 \end{cases}$$

$$T(n) & n$$

$$T(n/2) & T(n/2) & 2(n/2) = n$$

$$T(n/4) & T(n/4) & T(n/4) & T(n/4) & 4(n/4) = n$$

$$T(1) & \dots & \dots & T(1) & n(1) = n$$

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- 4. The total work is $\sum_{k=0}^{n-1} W(k) + \underbrace{L}_{\text{leaf work}}$

Recurrence Practice I

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What is T(n) for $n \geq 2$?

Answer:
$$T(n) = (c+d)2^{n-2} - d$$

Recurrence Practice III

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Answer:
$$T(n) = \frac{4}{3}(n^2 - 1) + c$$

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- ▶ In the previous slide, $T(n) = \frac{4}{3}(n^2 1) + c$ is always $\Theta(n^2)$, no matter what c is.
- Similar for $T(n) = (c+d)2^{n-2} d$ is $\Theta(2^n)$ and $T(n) = (c+2)n^{\log_2 3} 2n$ is $\Theta(n^{\log_2 3})$.

Recurrence Practice V

What is the Big-Theta running time of Shake?

```
Shake(A[1..n])
   if n < 1:
         return A
   m \leftarrow \lfloor n/2 \rfloor
   A_1 \leftarrow \text{Shake}(A[1..m])
   A_2 \leftarrow \text{Shake}(A[m+1..n])
   B \leftarrow []
   for i in 1..m:
         B.add(A_1[i])
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   return B
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Answer: $\Theta(n \log n)$

Recurrence Practice VI: HARD

$$T(n) = \begin{cases} c & n \le 1\\ T(n/2) + T(n/3) + T(n/6) + n & n > 1 \end{cases}$$

What is T(n) Big-Theta of?

Recurrence Practice VI: HARD

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Answer: $\Theta(n \log n)$

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- ▶ Goal: Prove a mathematical statement is true for all natural numbers.
- ▶ Outline: base case, inductive hypothesis, inductive step

Induction Proofs

The general strategy for proving a claim by induction is to

- (a) define the **base case(s)** and show the claim holds for them
- (b) state the **inductive hypothesis** assuming that the claim holds true for all $n < k \ (n \in \mathbb{N})$
- (c) prove that the claim holds for n = k in the rest of the **inductive step**

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- (c) prove that the claim holds for n = k in the rest of the **inductive step**
 - ▶ You can also do $n \le k$ and then n = k + 1.
 - ▶ Always use strong induction! Your inductive hypothesis must hold for all values up to k.

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Make sure to justify all the base cases that are necessary to establish the claim.

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▶ What base cases do you need to prove this claim:

$$\forall n, f_n < 2^n$$
, where $f_{n+1} = f_n + f_{n-1}$?

Induction Example I

Prove that the following holds for all natural numbers n.

$$\sum_{i=0}^{n} x^{i} = \frac{x^{n+1} - 1}{x - 1}$$

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- ▶ Base case(s): For _____, we have
- ► Inductive hypothesis: Suppose
- Inductive step: Consider $n = \underline{\hspace{1cm}}$. We want to show that

Therefore $\begin{bmatrix} & & \\ & & \end{bmatrix}$, which is what we needed to show.

Induction Example I

► Inductive step

$$\left(\sum_{i=0}^{n} x^i = \frac{x^{n+1}-1}{x-1}\right)$$

Induction Example I (Solution)

Prove that the following holds for all natural numbers n.

$$\sum_{i=0}^{n} x^{i} = \frac{x^{n+1} - 1}{x - 1}$$

- ▶ Base case(s): For n = 0, we have $\sum_{i=0}^{n} x^i = x^0 = 1$ and $\frac{x^{n+1}-1}{x-1} = \frac{x^{0+1}-1}{x-1} = \frac{x-1}{x-1} = 1$. So $\sum_{i=0}^{n} x^i = \frac{x^{n+1}-1}{x-1}$.
- ▶ Inductive hypothesis: Suppose $\sum_{i=0}^{n} x^i = \frac{x^{n+1}-1}{x-1}$ for $n = 0, 1, \dots, k$.
- ► Inductive step (next slide)

Induction Example I (Solution)

Inductive step: Consider n = k + 1. We want to show that $\sum_{i=0}^{n} x^{i} = \frac{x^{n+1}-1}{x-1}.$

 $\sum_{i=0}^{n} x^{i} = \frac{x^{i-1}}{x-1}.$ $\sum_{i=0}^{k+1} x^{i} = \sum_{i=0}^{k+1} x^{i} = x^{k+1} + \sum_{i=0}^{k} x^{i}.$ By the inductive hypothesis, $\sum_{i=0}^{k} x^{i} = \frac{x^{k+1}-1}{x-1}.$ So:

$$\sum_{i=0}^{k+1} x^i = x^{k+1} + \sum_{i=0}^k x^i = x^{k+1} + \frac{x^{k+1} - 1}{x - 1}$$

$$= \frac{(x - 1)x^{k+1}}{x - 1} + \frac{x^{k+1} - 1}{x - 1}$$

$$= \frac{x^{k+2} - x^{k+1} + x^{k+1} - 1}{x - 1}$$

$$= \frac{x^{k+2} - 1}{x - 1}.$$

Therefore $\sum_{i=0}^{k+1} x^i = \frac{x^{k+2}-1}{x-1}$, which is what we needed to show.



Inequality Induction

▶ Very similar to 'equality' induction! Don't overthink it! Just manipulating inequalities instead of equations, still primarily algebra

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- Very similar to 'equality' induction! Don't overthink it!
 Just manipulating inequalities instead of equations, still primarily algebra
- ➤ The only 'trick' is that you may have to simplify or change to a term to something even smaller/bigger (resp. the inequality) to make your algebra match your 'goal'. (Confused? The next example should illuminate this more clearly.)

Inequality Induction Example I

Let f_n be the *n*th Fibonacci number (i.e. $f_{n+1} = f_n + f_{n-1}$). Prove that $f_n \geq (\frac{3}{2})^{n-2}$.

Inequality Induction Example I

Let f_n be the *n*th Fibonacci number (i.e. $f_{n+1} = f_n + f_{n-1}$ where $f_1 = f_2 = 1$). Prove that $f_n \ge (\frac{3}{2})^{n-2}$.

- ▶ Base cases: For n = 1, we have $f_1 = 1$, $\left(\frac{3}{2}\right)^{1-2} = \frac{2}{3}$. $1 \ge \frac{2}{3}$. For n = 2, we have $f_2 = 1$, $\left(\frac{3}{2}\right)^{2-2} = 1$. $1 \ge 1$.
- ▶ Inductive hypothesis: Suppose that $f_n \ge \left(\frac{3}{2}\right)^{n-2}$ for n = 1, 2, ..., k-1.
- ▶ Inductive step: Consider n = k. We want to show that $f_n \ge (\frac{3}{2})^{n-2}$. $f_k = \underline{\hspace{1cm}}$

By the inductive hypothesis, we have $f_{k-1} \ge \left(\frac{3}{2}\right)^{k-3}$ and $f_{k-2} \ge \left(\frac{3}{2}\right)^{k-4}$.

So ______ Therefore $f_n \ge \left(\frac{3}{2}\right)^{n-2}$, which is what we needed to show.

Inequality Induction Example I

Inductive step: Consider n = k. We want to show that $f_n \ge (\frac{3}{2})^{n-2}$.

$$f_k = f_{k-1}^2 + f_{k-2}$$

By the inductive hypothesis, we have $f_{k-1} \ge \left(\frac{3}{2}\right)^{k-3}$ and $f_{k-2} \ge \left(\frac{3}{2}\right)^{k-4}$. So

$$f_n \ge \left(\frac{3}{2}\right)^{n-3} + \left(\frac{3}{2}\right)^{n-4}$$

$$\ge \left(\frac{3}{2}\right)^{n-4} \left(\frac{3}{2} + 1\right) = \left(\frac{3}{2}\right)^{n-4} \left(\frac{5}{2}\right)$$

$$\ge \left(\frac{3}{2}\right)^{n-4} \left(\frac{9}{4}\right) = \left(\frac{3}{2}\right)^{n-4} \left(\frac{3}{2}\right)^2$$

$$\ge \left(\frac{3}{2}\right)^{n-2}$$

Therefore $f_n \geq \left(\frac{3}{2}\right)^{n-2}$, which is what we needed to show.



Tree Induction

▶ Most important detail to remember: induct on **height of the tree**, *never* the number of nodes/leaves/etc.

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Tree Induction

- ▶ Most important detail to remember: induct on **height of the tree**, *never* the number of nodes/leaves/etc.
- ▶ Also critical: build your tree from the 'bottom-up' (adding root nodes at the top, never by adding leaves)
- ▶ May have to do casework, e.g. if node has different values depending on number/properties of children

Important Tree Terminology

- **binary**: each node has 0, 1, or 2 children
- \triangleright *n*-ary: each node has between 0 and *n* children
- \triangleright full: each node has strictly either 0 or n children
- ► **complete**: every level, except possibly the last, is completely filled
- ▶ perfect: full and complete; all levels filled, all internal nodes have 2 children, all leaves at same depth

Tree Induction Example

Define a Filbert tree to be a binary tree containing 2D points such that:

- \triangleright Each leaf node contains (3, 1), (-2, -5), or (2,2).
- An internal node with one child labeled (a, b) has label (a + 1, b 1).
- An internal node with two children labeled (x, y) and (a, b) has label (x + a, y + b).

Prove that the point in the root node of any Filbert tree is on or below the line x = y.

Tree Induction Example

Prove that the point in the root node of any Filbert tree is on or below the line x = y.

Tree Induction Example (Solution)

Prove that the point in the root node of any Filbert tree is on or below the line x = y.

Proof by induction on h, where h is the height of the tree.

- ▶ Base case(s): For Filbert tree where h = 0, the root node is a leaf and so contains (3, 1), (-2, -5), or (2,2), all of which are on or below the line x = y.
- ▶ Inductive hypothesis: Suppose that the point in the root node of any Filbert tree is on or below the line x = y for trees of height h = 0, 1, ..., k 1 $(k \ge 1)$.

Tree Induction Example (Solution)

▶ Inductive step: Let T be a Filbert tree of height k. There are 2 cases.

Case 1: The root of T has one child subtree, whose root contains (a, b). The root of T contains (a + 1, b - 1). By the inductive hypothesis, (a, b) is on or below x = y, i.e. $b \le a$. Since $b \le a$, $b - 1 \le a + 1$, so this point is on or below x = y.

Case 2: The root of T has two child subtrees, whose roots contain (x, y) and (a, b). Then the root of T contains (x+a, y+b). By the inductive hypothesis, $y \le x$ and $b \le a$. So $y+b \le x+a \implies (x+a, y+b)$ is on or below x=y. In all cases the root node contains a point on or below x=y, which is what we needed to show.

(Another) Inequality Induction Example

Prove by induction that for any two lists of nonnegative numbers (x_1, \ldots, x_n) and (y_1, \ldots, y_n) ,

$$\left(\sum_{i=1}^{n} x_i^2\right) \left(\sum_{i=1}^{n} y_i^2\right) \ge \left(\sum_{i=1}^{n} x_i y_i\right)^2$$

You may use the AM-GM inequality: For any real numbers $a,b \geq 0, \frac{a+b}{2} \geq \sqrt{ab}$.

(Another) Inequality Induction Example

$$\left(\sum_{i=1}^n x_i^2\right) \left(\sum_{i=1}^n y_i^2\right) \ge \left(\sum_{i=1}^n x_i y_i\right)^2, \frac{\frac{a+b}{2}}{2} \ge \sqrt{ab}$$

Graphs: Review

▶ (Abstract Def) An (undirected) graph is a tuple G = (V, E) where V is any set and $E \subseteq \{\{u, v\} : u, v \in V\}$.

Graphs: Review

- ▶ (Abstract Def) An (undirected) graph is a tuple G = (V, E) where V is any set and $E \subseteq \{\{u, v\} : u, v \in V\}$.
- ▶ (Usable Def) A graph is a set of vertices and edges between them.

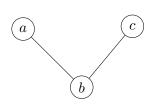


Figure: A graph with vertices $V = \{a, b, c\}$ and edges $E = \{\{a, b\}, \{b, c\}\}$.

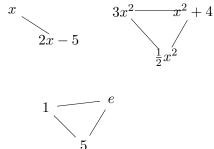
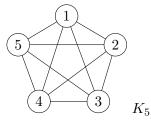


Figure: A graph with edges if nodes are Big-Theta of each other

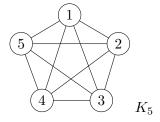
Special Graphs

▶ The graph K_n has n nodes and an edge between every pair of vertices:

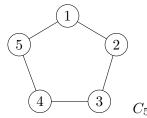


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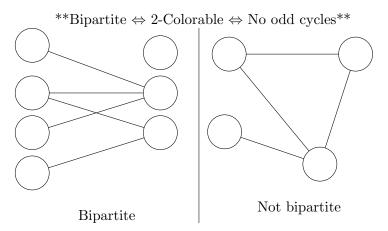


▶ The graph C_n has n nodes in a single cycle.



Bipartite

A graph is bipartite if its vertices can be divided into disjoint sets L and R such that every edge is between L and R.



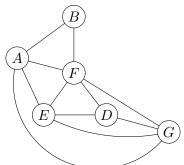
Paths, Walks, and Cycles

A walk in a graph is a sequence of vertices $(v_1, v_2, ..., v_k)$ and edges $(e_1, e_2, ..., e_{k-1})$ where each edge connects the two vertices on either end of it $(e_i = \{v_i, v_{i+1}\})$.

A path is a walk that doesn't repeat vertices.

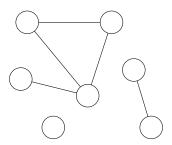
A cycle is a path where the start and end vertices are connected with an edge $(\{v_n, v_1\} \in E)$.

The distance between two vertices is the length of the shortest path between them.



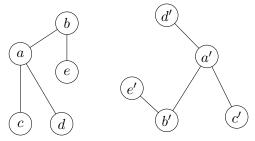
Connected Components

A connected component of a graph is a maximal set of vertices where each pair has a connecting walk. The graph below can be split into 3 connected components.



Isomorphisms

▶ Two graphs G = (V, E) and G' = (V', E') are isomorphic if there is a bijection $f: V \to V'$ such that $\{v_1, v_2\} \in E$ if and only if $\{f(v_1), f(v_2)\} \in E'$. (What type of relation is this?)

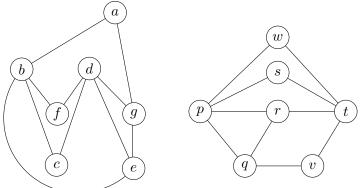


▶ For which n is K_n bipartite? C_n ?

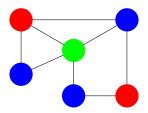
- ▶ For which n is K_n bipartite? C_n ?
- ightharpoonup For a graph with n vertices, what is an upper bound on the number of distinct paths?

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- \triangleright (SU24) Are these graphs isomorphic? What is the diameter of the left graph? What is the distance between a and c?

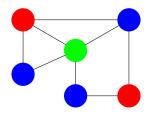


▶ A (proper) k-coloring of a graph G = (V, E) is an function $f: V \to \{1, 2, ..., k\}$ such that for every edge $\{u, v\} \in E, f(u) \neq f(v)$.



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- The chromatic number of a graph G is $\chi(G)$, the smallest k such that G is k-colorable.



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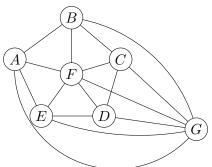
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- ▶ If G has K_n as a subgraph, $\chi(G) \ge n$.
- ▶ If G can be colored with k vertices, then $\chi(G) \leq k$.

▶ What is an upper bound for $\chi(G)$, given that G has n vertices?

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- ► (SU24) What is the chromatic number of the graph below? Prove it.



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- ightharpoonup (Past Examlet) In a tree of height h, which best describes the diameter of the tree?
 - (a) h (b) 2h (c) $\leq 2h$ (d) h+1 (e) $\leq h$

Tree Induction

Given a tree T and a function $w: T \to \mathbb{R}_+$, the weighted sum of the tree is $\sum_{v \in T} w(v) 2^{-h(v)}$, where h(v) is the depth of the node.

A weight function is fair if, for every node $v \in T$,

- 1. If the node has one child c, w(u) = w(c)
- 2. If the node has two children c_1 and c_2 , $w(u) = w(c_1) + w(c_2)$

Let T be a tree and r be its root. Prove by (strong) induction that for any fair weight function w, the weighed sum of the tree is no more than 2w(r).