Hi everyone !!

CS 374A Final Review

big bad darth vader

ACM @ UIUC

May 4, 2025

Disclaimers and Logistics

- **Disclaimer:** Some of us are CAs, but we have not seen the exam. We have no idea what the questions are. However, we've taken the course and reviewed Chandra's previous exams, so we have **suspicions** as to what the questions will be like.
- This review session is being recorded. Recordings and slides will be distributed on EdStem after the end.
- **Agenda:** We'll quickly review all topics likely to be covered, then go through a practice exam, then review individual topics by request.
    - Questions are designed to be written in the same style as Chandra ~~Kani's~~ previous exams but to be *slightly* harder, so don't worry if you don't get everything right away!
- Please let us know if we're going too fast/slow, not speaking loud enough/speaking too loud, etc.
- If you have a question anytime during the review session, please ask! Someone else almost surely has a similar question.
- We'll provide a feedback form at the end of the session.

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Induction

Template

Let x be an *arbitrary* <OBJECT>. Assume for all k s.t. k is smaller than x (by <ORDERING PROPERTY>), that $P(k)$ holds.

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If $x \neq$ <MINIMAL OBJECT>, then ..., so by IH, ..., so $P(x)$ holds.

Thus, in all cases, $P(x)$ holds.

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Thus, in all cases, $P(x)$ holds.

Some tips:

- **Always use strong induction**. All weak inductive proofs can be re-written to use strong induction with minimal changes, and the extra assumption can make your life significantly easier.
- **Write out your IH, base case, and inductive step out explicitly.** Doing so will help you avoid getting confused, and will help you avoid losing points.
- If you're performing induction on a recursive definition (strings, CFLs, etc.), generally, your inductive step will consist of one step of the recursion, and then will use IH.

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- If trying to guess whether or not a language is regular, think about memory. When processing a string through a DFA, you only need to know which state you're currently in, and do not need to look forwards/backwards in the string.
  - Implementing a DFA/NFA in code only requires $O(1)$ memory
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- If trying to guess whether or not a language is regular, think about memory. When processing a string through a DFA, you only need to know which state you're currently in, and do not need to look forwards/backwards in the string.
  - Implementing a DFA/NFA in code only requires $O(1)$ memory
  - If your checker program needs to count something without bound, the language you're checking isn't regular.
- **Regex Design Tips:** If you don't know where to start, try giving examples for strings that are in the language and strings that aren't. Look for patterns and try to build components around those patterns, then combine into something that represents the full language. Make sure to test and modify for edge cases. Explain, in English, each part of your regular expression with a short sentence. Does the explanation match the language?

DFAs/NFAs

- DFAs defined by *state set* Q, *accepting set* $A \subseteq Q$, *input alphabet* Σ, *start state* $s \in Q$, and *transition function* $\delta : Q \times \Sigma \to Q$
- NFAs allow for "trying" multiple transitions at the same time or transitioning without reading in (ϵ-transitions), accepts if there is a path to an accepting state. Transition function thereby changes to $\delta : Q \times (\Sigma \cup \{\epsilon\}) \to 2^Q$
  - Power-set construction to convert from NFA to DFA- in theory exponential-time but used in practice.
- **Tips for creating DFA/NFAs**: Break down your language into smaller patterns, and figure out what you need to store as state for each part. Make sure you clearly define all components. A drawing or transition table is just as valid as a $(Q, A, \Sigma, s, \delta)$ definition.

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Product Constructions

Given some languages $L_1, \ldots, L_n$ we want a DFA that accepts strings w satisfying $f(w \in L_1, \ldots, w \in L_n)$ where f is some logical function. Create a DFA/NFA for L using the following *rough* format:

- $Q = Q_1 \times \cdots \times Q_n$
- $\delta'(q_1, \ldots, q_n) = (\delta_1(q_1), \ldots, \delta_2(q_2))$
- $s = (s_1, \ldots, s_n)$
- $A' = \{\text{convert } f \text{ into a set expression}\}$

Practice Questions I

1. Let L be a regular language. Then, $L \cap \{0^n 1^n : n > 0\}$ is...
   (a) always regular
   (b) always irregular, but always context-free
   (c) sometimes irregular, but always context-free
   (d) sometimes non-context-free

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2. Define $\textsc{coverEvens}(w_1 w_2 \cdots w_n) = k_1 k_2 \cdots k_n$ and $\textsc{coverExponential}(w_1 w_2 \cdots w_n) = c_1 c_2 \cdots c_n$, where $k_i = \begin{cases} 1 & \text{if } i \equiv 0 \pmod{2} \\ w_i & \text{otherwise} \end{cases}$ and $c_i = \begin{cases} 1 & \text{if } \exists n \in \mathbf{Z} \text{ s.t. } i = 2^n \\ w_i & \text{otherwise} \end{cases}$.
   Which of the following is true?
   (a) If L is a regular language, then $\textsc{coverEvens}(L)$ is regular.
   (b) If $\textsc{coverEvens}(L)$ is a regular language, then L is regular.
   (c) If L is a regular language, then $\textsc{coverExponential}(L)$ is regular.
   (d) If $\textsc{coverExponential}(L)$ is a regular language, then L is regular.
   (e) Exactly two of the above are true
   (f) Exactly four of the above are true

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   (e) Exactly two of the above are true
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3. If we instead define $\textsc{UnCoverEvens}(L)$ to be $\{w : \textsc{CoverEvens}(w) \in L\}$, then would $\textsc{UnCoverEvens}(L)$ be regular for all regular L?
   (a) Yes
   (b) No

Practice Questions II

4. If L is decided by a DFA with n states, then consider the language L' consisting of all strings in L with at most 374 characters removed. Which of the following is true?
   (a) L' can be decided by a DFA with $O(n)$ states
   (b) L' can be decided by a DFA whose number of states is polynomial in n
   (c) L' can be decided by a DFA whose number of states is exponential in n
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5. Consider the language L consisting of the binary representation of all numbers congruent to 173 mod 374.
   (a) L does not have a fooling set.
   (b) L has a fooling set of size 173.
   (c) L has a fooling set of size 374.
   (d) L has a fooling set of size 375.
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   (e) L has an infinite fooling set.
6. Given a DFA M with n states, the minimum length of a string that M must accept (if $L(M)$ is non-empty) is at most:
   (a) n
   (b) $n - 1$
   (c) 2^n
   (d) $2^n - 1$
   (e) n^2

Practice Questions III

7. If L_1 is regular and L_2 is undecidable, then $L_1 \cap L_2$ is:
   (a) Always regular
   (b) Always context free
   (c) Always undecidable
   (d) Always decidable
   (e) Could be decidable or undecidable depending on L_1

Practice Questions III

7. If L_1 is regular and L_2 is undecidable, then $L_1 \cap L_2$ is:
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   (b) $L \cup \{\epsilon\}$ is context-free
   (c) $L \cap R$ is context-free for any regular language R
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   (e) $\{w\#w \mid w \in L\}$ is context-free

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9. Given a regular expression of length n, the equivalent minimum DFA might have
   9.1 $O(n)$ states
   9.2 $O(n^2)$ states
   9.3 $O(2^n)$ states
   9.4 $O(n!)$ states
   9.5 Always exactly n states

Practice Questions IV

1. Find a regexes for the following languages:
   (a) $\{\theta^a b^b \theta^c \mid a \geq 0 \text{ and } b \geq 0 \text{ and } c \geq 0 \text{ and } a \equiv b + c \pmod 2\}$

Practice Questions IV

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   (a) $\{\theta^a b^b \theta^c \mid a \geq 0 \text{ and } b \geq 0 \text{ and } c \geq 0 \text{ and } a \equiv b + c \pmod 2\}$
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2. Formally define DFAs/NFAs that accepts the following languages:
   (a) All strings whose ninth-to-last symbol is `0`, or equivalently, the set

$$\{x \circ z \mid x \in \Sigma^* \text{ and } z \in \Sigma^8\}.$$

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   (b) All strings w such that

   $$(\#(\texttt{0}, w) \bmod 3) + (\#(\texttt{1}, w) \bmod 7) = (|w| \bmod 4).$$

Recursion

- **Definition:** Reducing the problem to a smaller instance of itself, where eventually we can terminate in a base case.
  - Think: If we have a problem of size n, we want to continuously reduce to a problem smaller than n.
  - Example: Tower of Hanoi

Template

procedure AMAZINGRECURSIVEALGO(n)
    if $n ==$ [some base case] **then**
        return [value]
    else
        return `AmazingRecursiveAlgo`$(n-1)$

- Similar to **induction**!

Recursion: Runtime Analysis

- **General Form:**

$$T(n) = \underbrace{r}_{\text{\# of subproblems}} \cdot \overbrace{T\left(\frac{n}{c}\right)}^{\text{work at each subproblem}} + \underbrace{f(n)}_{\text{work at current level}}$$

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$$\begin{aligned}
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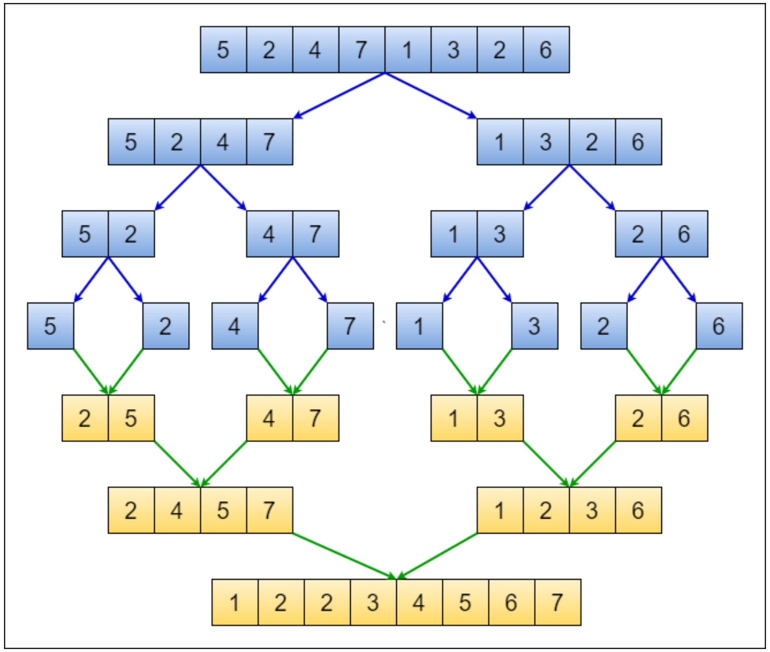
  - **Intuition:** If each level contains more work than the level below it, then the root level will dominate. If each level contains the same amount of work, then we have $log_c n$ levels with $f(n)$ work. If each level contains less work than the work below it, then the leaf nodes will dominate.

Divide and Conquer Algos: Merge Sort

- **Purpose:** Sort an arbitrary array.
- **Time Complexity**: $O(n \log n)$
- **Intuition:** Three phases: (a) split the array in half, (b) sort each side, (c) merge the sorted halves by repeatedly comparing smallest elements on each side not yet inserted.

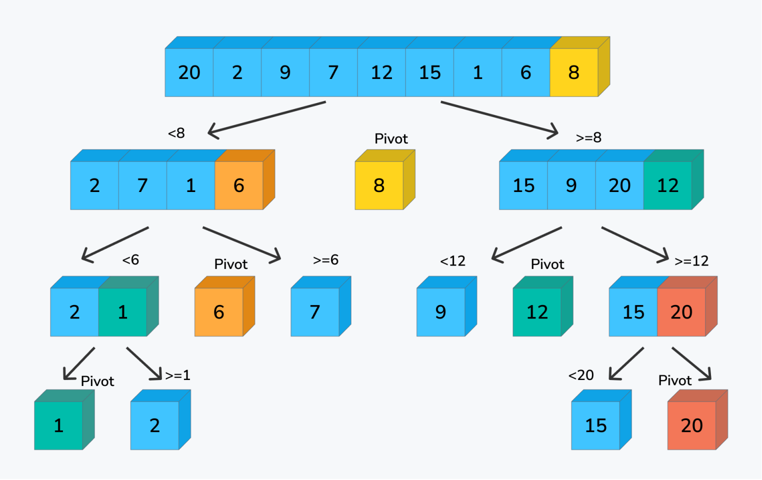


Divide and Conquer Algos: Quickselect

- **Purpose:** Get the n^{th} smallest element in an arbitrary array.
- **Time Complexity**: Avg: $O(n)$ | Worst; $O(n^2)$, ($O(n)$ with MoM)
- **Intuition**: Pick a pivot P with a value P_V and rearrange the array such that all the elements that are less than P_V are to the left of P and all the elements that are greater than P_V are to the right of P, just like quick select. If the length of the array of elements that are less than P_V is greater than n, then we know that the n^{th} smallest element is to the left of P and we recurse on the left subarray. Otherwise, we know that the n^{th} smallest element is to the right of P and we recurse on the right subarray.
  - **Why the poor worst case performance?**
  - Again, because we can get unlucky and pick the worst possible pivot at every step.
  - We can guarantee linear performance with a better pivot-picking algorithm such as MedianOfMedians
    - Finds element that larger than $\frac{3}{10}$ and smaller than $\frac{7}{10}$ of the array's elements.
    - Runs in $O(n)$ time

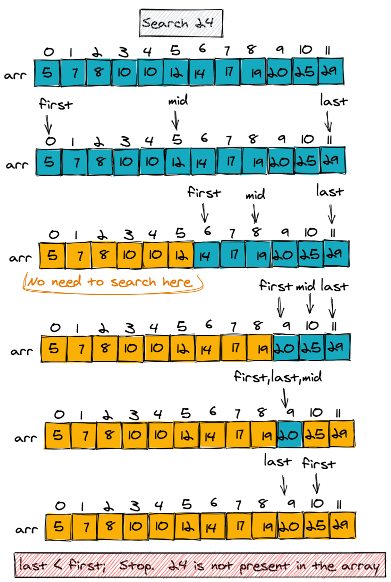
Divide and Conquer Algos: Quicksort

- **Purpose:** Sort an arbitrary array.
- **Time Complexity**: Avg: $O(n \log n)$ | Worst: $O(n^2)$ ($O(n \log n)$ deterministic with quickselect partitioning)
- **Intuition**: Pick a pivot and rearrange the array such that all the elements that are less than the pivot value are to the left of the pivot value and all the elements that are greater than the pivot value are to the right of the pivot value. Then sort each side.
  - **Why the poor worst case performance?**
  - Because we can get unlucky and pick the worst possible pivot at every step.



Divide and Conquer Algos: Binary Search

- **Purpose:** Find the existence of an element in a sorted array
- **Time Complexity**: $O(\log n)$
- **Intuition**: Say we are trying to find the value n. Pick the middle element M in the array. If $n > M$, the element must be to the right of n and we recurse on the right. Otherwise, we recurse on the left.



Backtracking

- Technique to methodically explore the solutions to a problem via the reduction to said problem to a smaller variant of itself, a.k.a **recursion**.
- Intuitively, think of the problem space as a maze that we are trying to find the exit of. For each path, you would traverse until you reach a dead end, at which point you **back track** to try a different path.
- To find recurrence, think ”What information about a subset of my current problem space would be really nice to know?”

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$$\text{LIS}(i, j) = \begin{cases} 0 & \text{if } i = 0 \\ \text{LIS}(i-1, j) & \text{if } A[i] \geq A[j] \\ \max \begin{cases} \text{LIS}(i-1, j) \\ 1 + \text{LIS}(i-1, i) \end{cases} & \text{else} \end{cases}$$

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- "What is the length of a longest increasing subsequence in an arbitrary array?"

$$\mathrm{LIS}(i,j) = \begin{cases} 0 & \text{if } i = 0 \\ \mathrm{LIS}(i-1,j) & \text{if } A[i] \geq A[j] \\ \max \begin{cases} \mathrm{LIS}(i-1,j) \\ 1 + \mathrm{LIS}(i-1,i) \end{cases} & \text{else} \end{cases}$$

This kind of sucks; we're redoing computation that we've already done! What if instead, we computed all the subproblems beforehand, wrote down the solutions, then did the recursion?

Dynamic Programming

- It’s backtracking, but we compute all of the subproblems iteratively.
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- **How to solve a DP:**
  - Identify how we can take advantage of a recursive call on a smaller subset of the input space.
  - Identity base cases
  - Identity recurrences (they should cover all possible cases at each step)

Dynamic Programming

Let's look at the LIS example from before: "What is the length of a longest increasing subsequence in an arbitrary array?"

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```
procedure LIS-ITERATIVE(A[1..n]):
    A ← [1 ... n][1 ... n]
    for all i ← 1 ... n do
        for all j ← i ... n do
            if A[i] ≤ A[j] then
                LIS[i][j] = 1
            else
                LIS[i][j] = 0
    for all i ← 1 ... n do
        for all j ← 2 ... n do
            if A[i] ≥ A[j] then
                LIS[i][j] = LIS[i − 1, j]
            else
                LIS[i][j] = max { LIS[i − 1, j]
                                { LIS[i − 1, i] + 1
    return LIS[n, n]
```

Graphs

- **Definition:** A set of vertices V connected by a set of edges E. Individual edges are notated as (u, v), where $u, v \in V$.
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Graph Algorithms: Traversal

- **BFS:**
  - **Purpose:** Reachability, Shortest Path (unweighted graph)
  - **Implementation details:** Add your neighbours to a **queue**, pop from the queue to get next node
  - **Runtime:** $O(V + E)$

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Graph Algorithms: Shortest Path

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    - $d(v, k)$ is the shortest-walk distance from s to v using at most k edges
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  - **Runtime:** $O(mn)$
- **Floyd-Warshall**
  - **Purpose:** APSP, yes negative edge weights
  - **Implementation:** Dynamic Programming recurrence:
    - $d(u, v, i)$ is the shortest-path distance from u to v only going through vertices $1 \ldots i$.
    - $d(u, v, i) = \min(d(u, v, i), d(u, i, i-1) + d(i, v, i-1))$
  - **Runtime:** $O(n^3)$

Graph Algorithms: MSTs

3 main algorithms:

- **Prim-Jarnik**: Keep a priority queue for edges to be added to the tree. Start the tree at some arbitrarily selected root vertex. When adding a vertex, add all of its neighbors to the queue. Runtime: $O(|E| \log |V|)$, $O(|V| \log |V| + |E|)$ using Quake heaps.

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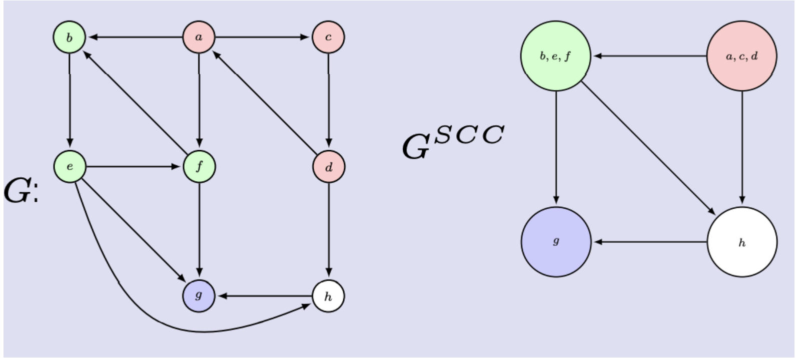
- **Borůvka**: No fancy data structures! Just find smallest edge going out of each vertex, then contract all edges that you selected! Runtime: $O(|E| \log |V|)$
- Faster (but way more complicated algorithms) exist. **Yao** (1975): $O(|E| \log \log |V|)$ with a modification of Borůvka's (using linear-time median selection). **Karger-Klein-Tarjan** (1995): $O(|E|)$ in expectation, **Chazelle** (2000): $O(|E|\alpha(|V|, |E|))$ deterministic

Graph Algorithms: SCC

SCC-Finding Algorithms (Tarjan's, Kosaraju's)

- **Purpose:** To identify (and collapse) SCCs in a (directed) graph
- **Runtime:** $O(V + E)$
- **Returns:** A metagraph that has one node for each SCC.



Graph Algorithms: Longest Path

Longest path in a Directed Acyclic Graph (DAG)

- **Purpose:** To find the longest simple path (no repeating vertices) by weight in a graph which is guaranteed to be a DAG[1].
- **Runtime[2]:** $O(V + E)$
- **Returns:** The sum of the weights of the longest path in the DAG.

---

[1] Finding the longest path in other types of graphs is at least NP-hard.
[2] This is a relatively straight-forward DP on a DAG problem if you wish to derive it.

Graph Problems: General Stuff

How to solve graph problems:

1. Identify type of problem (Reachability, Shortest Path, SCC)
2. Construct new graph
   - Add sources/sinks
   - Add vertices via $V' = V \times \{\text{some set}\}$ (Useful for tracking states)
   - Add vertices via $E' = E \times \{\text{some set}\}$ (Useful for allowing/prohibit certain behaviour)
3. Apply some stock algorithm **(DO NOT MODIFY THE ALGORITHMS - MODIFY THE INPUTS!)**
4. Draw connection between how to result of the algorithm upon the new graph relates to the solution of the original question.

Practice Problems I

1. Given a graph, consider the problem of finding the vertex that is reachable by the most other vertices. This problem is:
   (a) Solvable in $O(m+n)$ time
   (b) Solvable in $O(m \log n)$ time
   (c) Solvable in $O(mn)$ time
   (d) Solvable in polynomial time, but not any of the runtimes above
   (e) Solvable in exponential time, but not polynomial time

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2. Given an unsorted list, we want to print out the $\sqrt{n}$ smallest elements in sorted order.
   (a) We can do this in $O(\sqrt{n})$ operations
   (b) We can do this in $O(\sqrt{n} \log n)$ operations
   (c) We can do this in $O(n)$ operations
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   (e) We can do this in $O(n^{1.5})$ operations

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   (e) We can do this in $O(n^{1.5})$ operations
3. Define the binary operator @ s.t. $a \,@\, b = \frac{a+b}{2}$. Given an expression $a \,@\, b \,@\, c \,@\, \cdots$, finding the evaluation order that maximizes the total value
   (a) Can be done in $O(n^2)$ time via DP
   (b) Can be done in $O(n^3)$ time via DP
   (c) Can be done in $O(n^4)$ time via DP
   (d) Cannot be done in polynomial time, but can be done in exponential time
   (e) Can be done in $O(n!)$ time, and no faster algorithm is possible

Practice Problems II

4. Given a directed graph G with edges and some vertices s and t, some of which are negative, finding the shortest simple path from s to t using at most 374 negative edges (faster is better)...
   (a) Can be done in $O(m \log n)$ time using Dijkstra's, but only if G has no negative cycles
   (b) Can be done in $O(m \log n)$ time using Dijkstra's, even if G has a negative cycle
   (c) Can be done in $O(mn)$ time using Bellman-Ford, but only if G has no negative cycles
   (d) Can be done in $O(mn)$ time using Bellman-Ford, even if G has a negative cycle

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   (d) Can be done in $O(mn)$ time using Bellman-Ford, even if G has a negative cycle
5. Given a connected DAG G with a single sink t and weighted edges, for a given vertex s, if for a path P, $s(P)$ is the sum of all edge values for that path, computing the sum of the $s(P)$ over *all* $s \to t$ paths...
   (a) Can be calculated for all $v \in G$ in $O(n)$ time
   (b) Can be calculated for a *single* $s \in G$ in $O(n)$ time, but requires $O(n^2)$ time for all possible $s \in G$
   (c) Requires $O(n^2)$ time for a single $s \in G$, but can be calculated in $O(n^2)$ time for all possible $s \in G$
   (d) Requires $O(n^2)$ time for a single $s \in G$, and requires $O(n^3)$ time for all $s \in G$
   (e) None of the above are true

Practice Problems III

6. Solve the recurrence $T(n) = T\left(\frac{3n}{4}\right) + T\left(\frac{n}{4}\right) + O(n^2)$, where $T(n) = O(1)$ when $n \leq 374$.
   (a) $O(n)$
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   (e) $O(n^3)$
7. The problem of determining if a graph G contains a triangle (cycle of length 3) that includes a specific vertex v can be solved in
   (a) $O(m)$
   (b) $O(n^2)$
   (c) $O(nm)$
   (d) $O(n^3)$
   (e) Is NP-Complete

Practice Problems IV

8. Which of the following is true? (If more than one statement is true, pick the strongest).
   (a) A MST will *never* contain the maximum-weight edge in a cycle, and will *always* contain the minimum-weight edge
   (b) A MST *may* contain the maximum-weight edge in a cycle, but will *always* contain the minimum-weight edge
   (c) A MST will *never* contain the maximum-weight edge in a cycle, but *may* contain the minimum-weight edge
   (d) A MST *may* contain the maximum-weight edge in a cycle, and *may* contain the minimum-weight edge
9. Given a graph G whose edges are colored either red or blue,
   (a) We can find a cycle where at least 1/3 of the edges are blue in $O(n+m)$ time
   (b) We can find a cycle where at least 1/3 of the edges are blue in $O(m \log n)$ time
   (c) We can find a cycle where at least 1/3 of the edges are blue in $O(nm)$ time
   (d) We can find a cycle where at least 1/3 of the edges are blue in $O(m^k)$ time for some $k > 2$
   (e) Finding a cycle with 1/3 of the edges being blue is NP-hard

how ba-a-a-ad can i be

We are attempting to write a text formatter. Ideally, we want our formatter to format our text reasonably; we wouldn't want our formatter to just put each word on a new line, for example. Fortunately, we have a function badness(n), where given n characters of whitespace, it will return to us a integer of "badness". Given that each line can hold at most M character s and an array $L[1 \cdots n]$, where $L[i]$ is the length of the ith word, describe an algorithm that finds the smallest possible total badness of a given sequence of words.

the skyway sucks

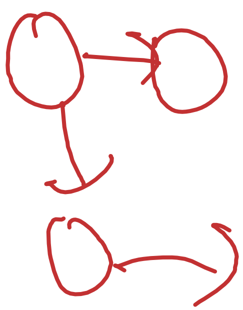
$G = V, E$

$G' = V \times \{0 \ldots k/20\}$

edge → $(u, i) \to (v, i)$

$(u, i) \to (v, i-1) \to wt = wt + 1$

You're at your office in downtown Chicago, and you're trying to get to your friend's apartment elsewhere in the city. You're given the map of Chicago as a directed graph *G*, with streets as edges and intersections as vertices. Each edge is weighted with the travel time (in minutes) that it takes to travel the length of the road. Some edges are annotated as moving bridges, which are only usable for certain periods (notated as intervals in minutes after you start). You can also bribe the bridge operator into moving the bridge and letting you cross by slipping him a crisp \$20, in which case you can cross after 1 minute, but you only have \$*k*. Describe and analyze an algorithm to calculate the quickest route that you can go.



Dij-upd:

if not bridge

dec-key(t + e.wt)

if bridge:

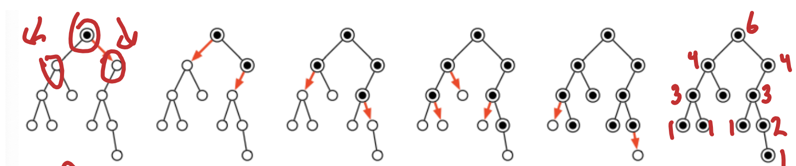
dec-key(min_start + e.wt)

max(t, minimal entry > t)

tell me more, tell me more, did you get very far?

Suppose we need to distribute a message to all the nodes in a given binary tree. Initially, only the root node knows the message. In a single round, each node that knows the message is allowed (but not required) to forward it to at most one of its children. Describe and analyze an algorithm to compute the minimum number of rounds required for the message to be delivered to all nodes in the tree. For example, given the following tree as input, your algorithm should return the integer 5.

topo sort



A message being distributed through a binary tree in five rounds.

$$\text{Mess}(R) = \min\begin{cases} \max(1+M(R.L),\ 2+M(R.R)) \\ \max(1+M(R.R),\ 2+M(R.L)) \end{cases} \text{ if 2 children}$$

$1 + R.\text{child}$ if 1 child

0 if leaf

$U(v)$

P and NP

- A **decision problem** is a problem with a true/false answer. (yes/no, etc.)
- **P** is the set of decision problems with a polynomial-time solver.
- **NP** is the set of decision problems with a polynomial-time *nondeterministic* solver.
- Alternatively, NP is the set of decision problems with a polynomial-time *certifier* for "true" answers, given a polynomial-size *certificate*.
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For example, consider the yes/no problem of deciding whether a graph $G = (V, E)$ has a path containing all its vertices. (Hamiltonian Path)

- If you were given the path already ($O(V)$ length) as a certificate, you could certify that the answer is "yes" in polynomial time.
- Therefore, this problem is in NP.

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  - Intuitively, with an NP problem, we can verify a "yes" answer quickly if we have the solution in front of us.

For example, consider the yes/no problem of deciding whether a graph $G = (V, E)$ has a path containing all its vertices. (Hamiltonian Path)

- If you were given the path already ($O(V)$ length) as a certificate, you could certify that the answer is "yes" in polynomial time.
- Therefore, this problem is in NP.

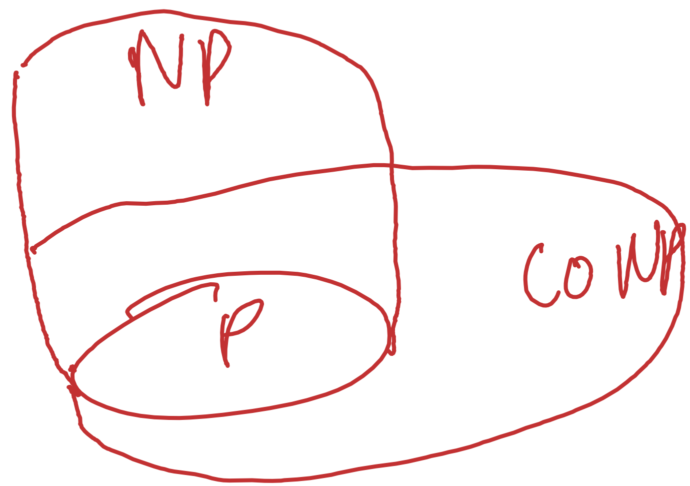
Formally, an algorithm C is a certifier for problem X when $s \in X$ if and only if there exists string t such that $C(s, t) =$ true.

- t here is a "certificate."
- We can show X is NP by providing this information, and showing C is polynomial-time and t is polynomial-size (with respect to the size of the input s).

co-NP

- **co-NP** is the set of decision problems X whose complements $\overline{X}$ are in NP.
- Alternatively, NP is the set of decision problems with a polynomial-time certifier for **"false"** answers, given a polynomial-size certificate.
- For example, the problem of deciding whether a graph *doesn't* have a Hamiltonian path is in co-NP.

co-NP isn't on your skillset, but be aware that this is *not* the same thing as NP.



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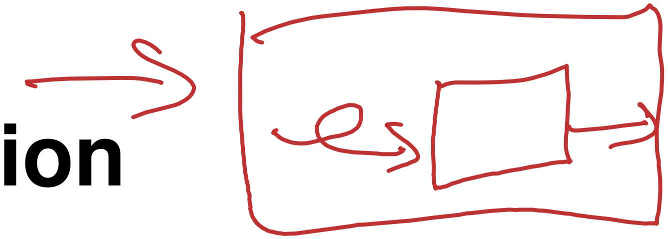

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Make sure you're going in the right direction!

If you're trying to prove that a problem is NP-hard or undecidable, you need to reduce **from** an NP-hard/undecidable problem **to** the problem you want to prove is hard (in other words, show that an oracle for your problem can be used to solve an NP-hard/undecidable problem). The most common mistake on exams is reducing in the wrong direction.

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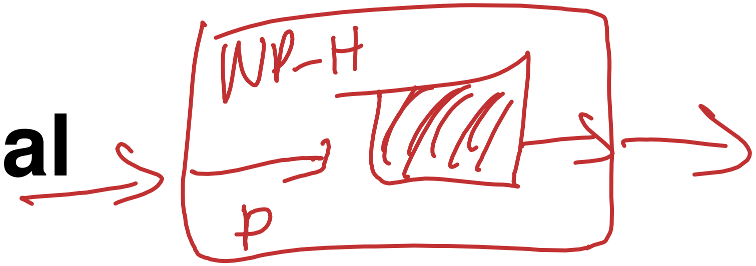

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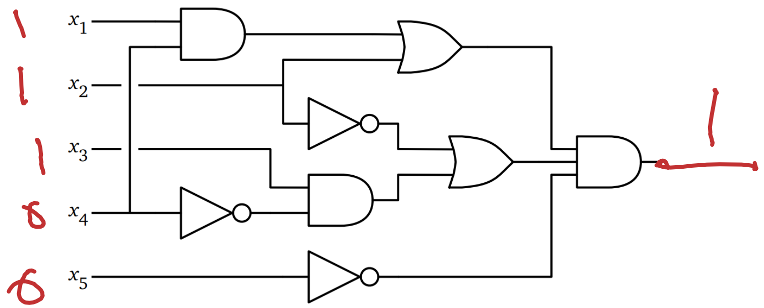
Template- Reduction

Assume that there exists an oracle function *B* which runs in [TIME CONSTRAINT]. Thus, we can solve *A* as follows:

```
procedure A(input):
    Do some preprocessing to create instances of problem B
    outputs ← B(generated inputs)
    Do some postprocessing on outputs to get the correct answer for A
```

A Tour of NP-Hard Problems: CircuitSAT and 3SAT

- **CircuitSAT**: The "original" NP-complete problem. Given a boolean circuit, is there a set of inputs that makes it return `true`?



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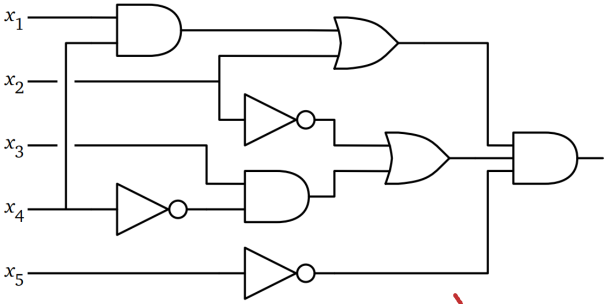

- **3SAT**: Given a boolean formula of the form $(a \vee b \vee c) \wedge (\overline{a} \vee d \vee e) \wedge \cdots$, is there an assignment to the input variables that makes it return `true`?

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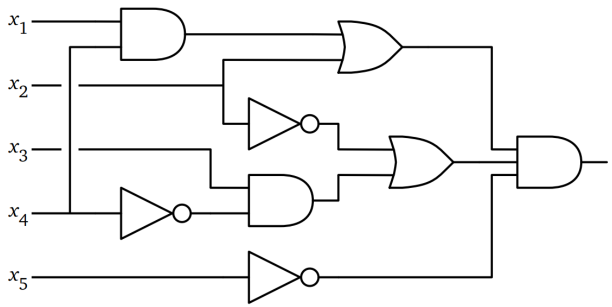

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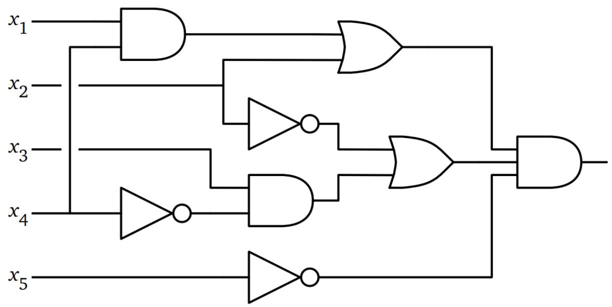

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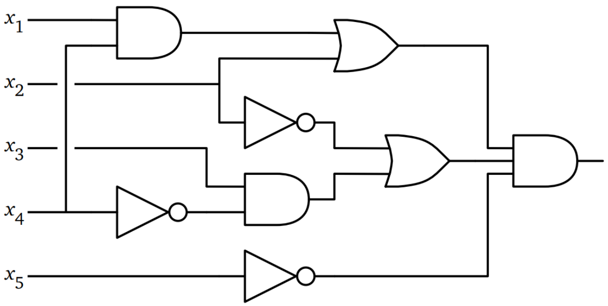

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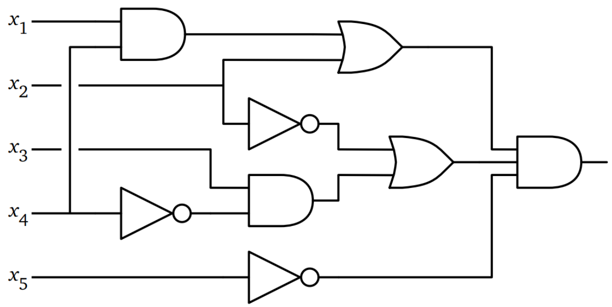

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Be careful with *k*-SAT variants!

While k-SAT for $k \geq 3$ is NP-complete, there is a polynomial-time algorithm for 2SAT. (Using strongly connected components!)

A Tour of NP-Hard Problems: CircuitSAT and 3SAT

$a, \bar{a}$ $(a \vee b \vee c) \wedge$

Consider the problem **MajSAT**: Clauses now consist of 5 literals, and you must satisfy at least 3 literals in each clause. Is **MajSAT** in NP, NP-hard, both, or neither? Prove why by either stating an algorithm or providing a reduction.

Maj(a, b, c, d, e)

(a, b, c, t, t)

Maj(t, t, t, t, t)

proc 3S:

for clause $(a \vee b \vee c)$

c' ← maj(a, b, c, t, t)

add Maj(t, t, t, t, t)

Ora(p)

return

A Tour of NP-Hard Problems: Max{Clique, IndSet}, MinVertexCover

- **MaxClique**: Given a graph G and positive integer h, can we find a K_h subgraph in G (i.e. a set of h nodes where each one has an edge to every other)?

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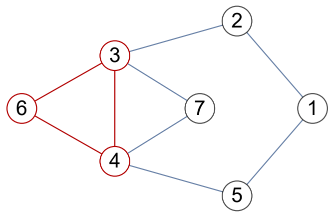

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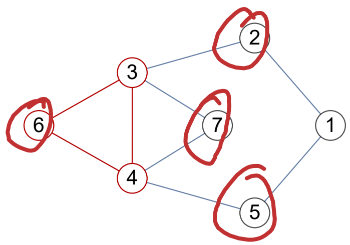

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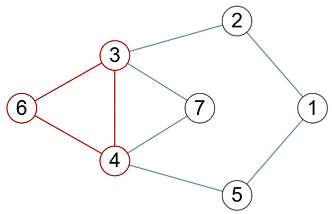

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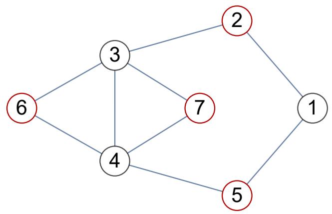

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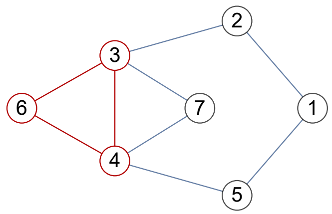

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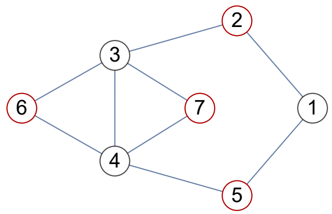

- **MinVertexCover**: Given a graph G and positive integer h, can we find a set of h nodes so that all edges have at least one endpoint chosen?

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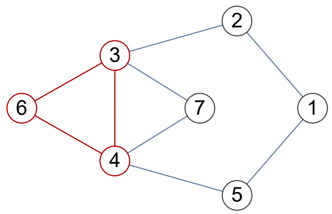

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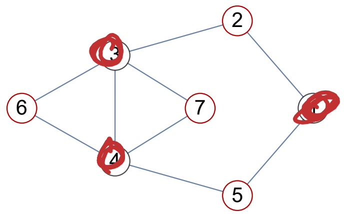

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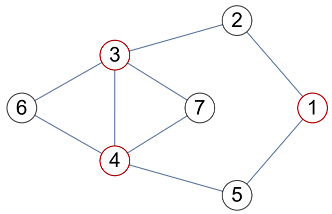


A Tour of NP-Hard Problems: Max{Clique, IndSet}, MinVertexCover

ACM is writing their review session for CS/ECE 374B MT3. While making slides, each CA writes 2 problems, either alone or in collaboration with other CAs. Since all of the CAs all have inflated egos, they won't show up to the review session unless one of the problems that they worked on is in the review session. Show that determining whether we can run a review session with at most *k* problems is NP-complete.

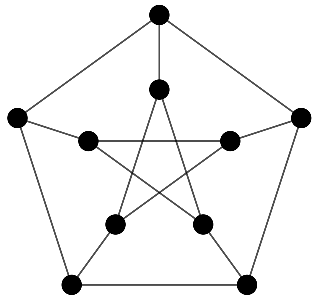
MinVtxCover:

edge -> CA

Vtx -> problem

A Tour of NP-Hard Problems: Graph Coloring

- Given an (undirected) graph, can we color the nodes with at most *k* colors so that no two vertices that share an edge are of the same color?



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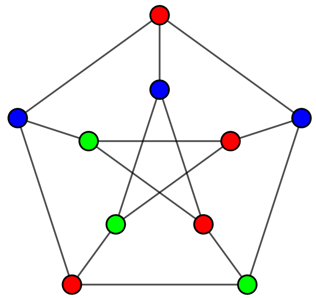

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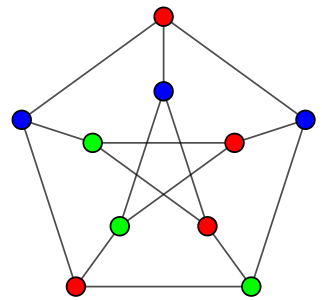

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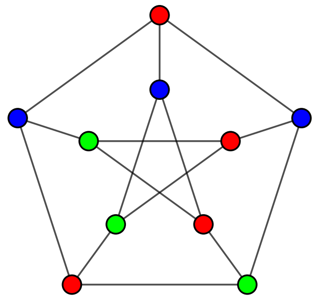

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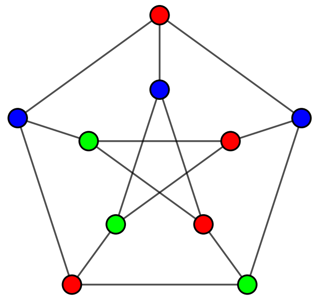

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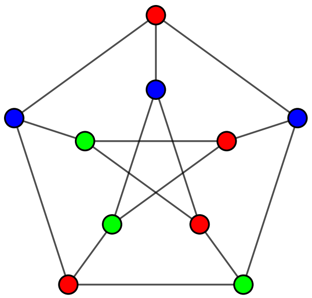

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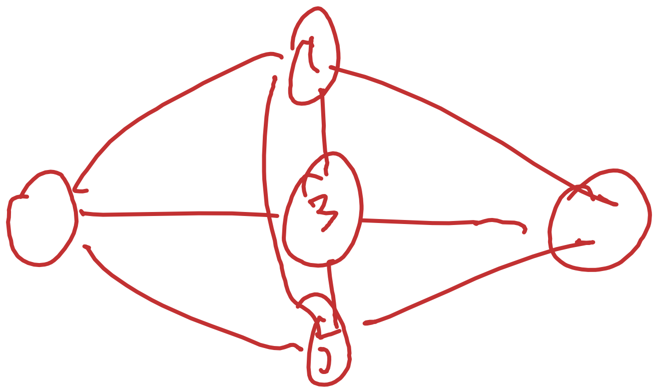
Be careful with *k*-coloring variants!

While k-coloring for $k \geq 3$ is NP-complete, you can find whether a graph is bipartite (2-colorable) using DFS.

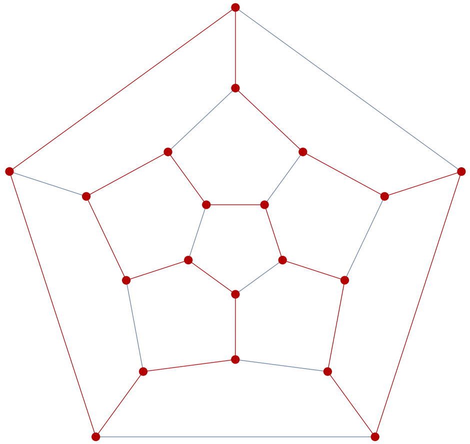
A Tour of NP-Hard Problems: Graph Coloring

Consider the problem **Safe7Color**, which asks you to color a graph with 7 colors, such that it is a violation if there is an edge $u \leftrightarrow v$ where $c(u)$ and $c(v)$ differ by 0 or 1 (mod 7). Is this problem NP-hard?



A Tour of NP-Hard Problems: Hamiltonian Paths and Cycles

- A **Hamilton Path** is a path that goes through each vertex *exactly* once. Likewise, a **Hamiltonian Cycle** is a cycle that goes through each node *exactly* once.
  - Every graph with a Hamiltonian cycle has a Hamiltonian path, but not every graph with a Hamiltonian path has a Hamiltonian cycle.



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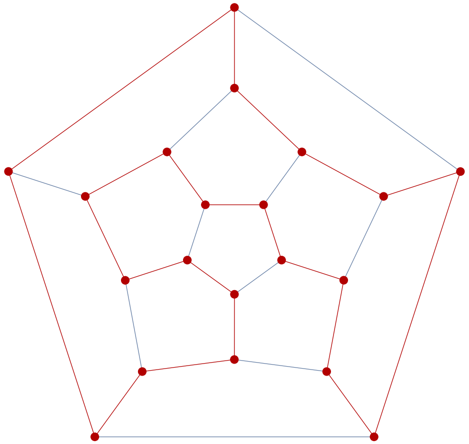

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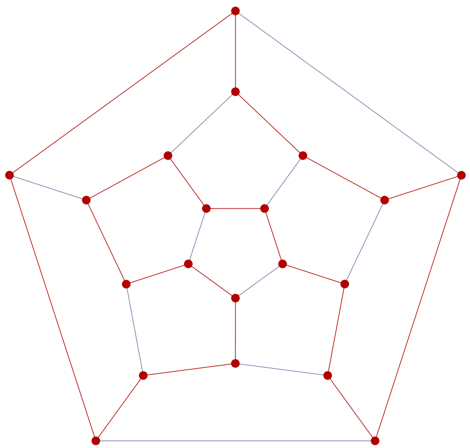

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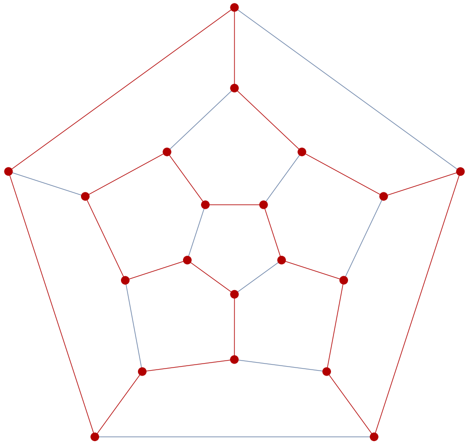
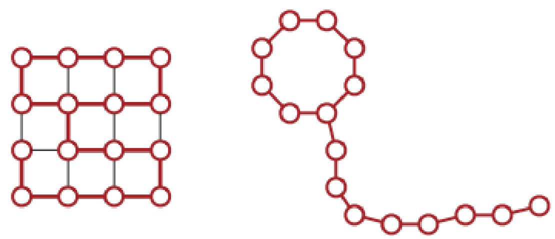
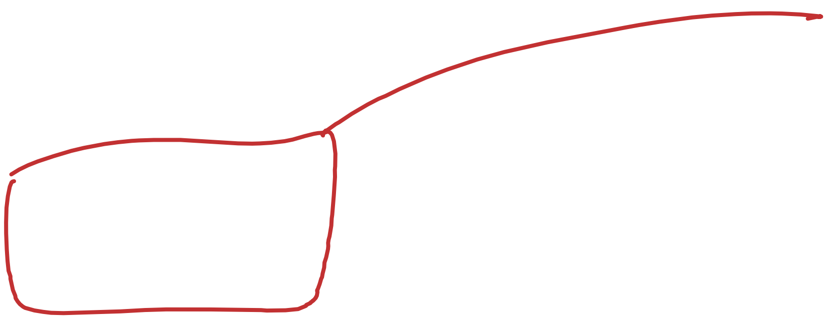


- Consider reducing from **HamPath** or **HamCycle** if. . .
  - You're given a graph, and you're asked to find a sequence of vertices
  - You have a resource pool, and you want to use up everuthing

A Tour of NP-Hard Problems: Hamiltonian Paths and Cycles

A **balloon graph** of size ℓ is a cycle of length ℓ attached to a path of length ℓ, where the cycle and the path are disjoint except for the connecting vertex. Show that it is NP-hard to determine whether a graph has a balloon subgraph of size at least k.





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- **Checkers**: given a $n \times n$ checkerboard, is there a move that captures at least k checkers?

Undecidability

- A language is **decidable** if there exists an algorithm which always returns `true` to all inputs in L and `false` to inputs not in L
  - If we can only return `true` to all inputs in L and either return `false` *or* infinite-loop for all other inputs, the language is merely **acceptable**.

Theorem (Turing, 1936)

The language `Halt`*:* $\{(f, w) : \text{ the function } f \text{ does not infinite loop on input } w\}$ *is undecidable.*

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Let $\mathcal{L}$ *be any set of languages that satisfies the following conditions:*

- *There is a Turing machine Y such that Accept(Y)* $\in \mathcal{L}$ *.*
- *There is a Turing machine N such that Accept(N)* $\notin \mathcal{L}$ *.*

Then, the language $\mathit{AcceptIn}(\mathcal{L}) \leftarrow \{\langle M \rangle \mid \mathit{Accept}(M) \in \mathcal{L}\}$ *is undecidable.*

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  3. Abuse the fact that you can put code into a function to derive a contradiction.

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For each of the following languages, either show that they are decidable by describing an algorithm that decides them, or show that they are undecidable by reduction and by Rice's theorem when possible.

(a) AcceptsRegular $= \{\langle M \rangle : M\text{'s accept set is regular}\}$

Halt (M, w):

m'(x):

run M on w

ret x has prime len

ret ¬AccReg(M')

Rice's

Y

acc everyth.

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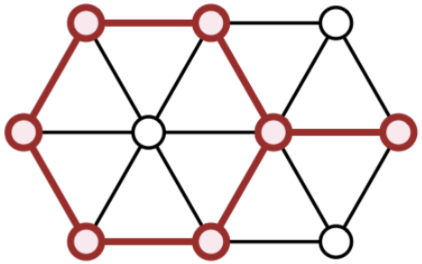
(f) NeverLeft $= \{\langle M, w \rangle : M \text{ never moves left on input } w\}$

twiangles are scawy

A subset S of vertices in an undirected graph G is a called triangle-free if, for every triple of vertices $u, v, w \in S$,at least one of the three edges uv, uw, vw is absent from G. Prove that find the size of teh largest triangle-free subset of vertices in a given undirected graph is NP-hard.



A triangle-free subset of 7 vertices and its induced edges.
This is ***not*** the largest triangle-free subset in this graph.

$G=(V,E)$ $\subset$ $\subseteq$

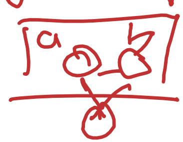
$G'=(V',E') \Rightarrow G$ copy

$+\ \forall u,v \in V,\ u \neq v$, we add a vertex w_i and add corresponding edges

Let's define W as the set of vertices we added to get from $G \rightarrow G'$

Lemma 1: Let $S \subseteq G'$, where $S = A \cup W$. S is a triangle-free subset in G' $\Leftrightarrow$ A is an independent set in G

independent $\Rightarrow$ triangle. Suppose A is independent, but S is NOT triangle-free



triangle $\Rightarrow$ independent: Suppose S is triangle-free but A is NOT independent.

Lemma 1: Let $S \subset G'$, where $S = A \cup W$.
Then A independent set in G $\iff$ S triangle-free in G'

Goal: Reduce from Largest Triangle Free to Maximal Ind. Set.

1) Let S be the maximal triangle free subset of G'. By Lemma, we have $S = A \cup W$, where A is ind. set $\in G$. Thus, the maximal independent set must at least be big as $|A|$.
Proof by contradiction: Suppose $\exists A' \in G$ s.t. $|A'| > |A|$ and A' is independent in G. Then, $\exists S' = A' \cup W$, where S' is triangle-free and $|S'| > |S|$. But, we assumed that S is maximal! Oops.

2) Let A be the maximal ind. set in G. By Lemma, $\exists S$ s.t. S is triangle free in G'. Thus, the maximal triangle-free subset in G' is at least as big as $|S|$. Proof by contradiction! Suppose $\exists S'$ s.t. S' is triangle free in G', and $|S'| > |S|$. Thus, $\exists A'$ where $S' = A' \cup W$, and $|A'| > |A|$. This is a contradiction! Oops.
Thus, S is the maximal triangle free subset in G'.

Feedback

- Please fill out the feedback form: `go.acm.illinois.edu/374a_final_feedback`