Simulation Results of Channel Capacity using Parity Bits for Error Correction

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1 Introduction

Claude Shannon was a major figure in creating the mathematical discipline of information theory. This paper analyzes how emperical results line up with the statements made by this mathematical model of information. More formally, we will be analyzing the effect of using parity bits for error correction and how well they are at approaching the formal limit on the conveyance of data.

2 Mathematical Description

The official description of converyable information is the mutual information between to system of events. Let A be the source alphabet. Let B be the reception alphabet. Let Q be the transmission matrix (IE the probability of receiving $b_j \in B$ given a transmission of $a_i \in A$. Let these entries be represented as $q_{i,j}$ respectively. For our example we will only consider the simplier channel known as the binary symettric channel. Thus

$$Q = \begin{pmatrix} q_{0,0} & q_{0,1} \\ q_{0,1} & q_{1,1} \end{pmatrix} = \begin{pmatrix} p & 1-p \\ 1-p & p \end{pmatrix}$$

Given such a channel, the probability of bit errors can be represented by a Bernoulli Trial.

The mutual information between two system of events (definition given elsewhere) is defined as following:

$$I(A,B) = \sum_{i=1}^{m} \sum_{j=1}^{n} Pr(A_i \cup B_j) \log \left(\frac{Pr(A_i \cup B_j)}{Pr(A_i)Pr(B_j)} \right)$$

... =
$$\sum_{i=1}^{m} Pr(A_i) \sum_{j=1}^{n} q_{i,j} \log \left(\frac{q_{i,j}}{\sum_{t=1}^{m} Pr(A_t) q_{t,j}} \right)$$

It has been proved elsewhere that the maximal channel capacity (maximum I(A,B) based on $Pr(A_i)$ – probability that a source letter is transmitted) for a BSC is when $Pr(A_1 = 0) = Pr(A_2 = 1) = \frac{1}{2}$. Furthermore for these values,

$$I(A,B) = p * log(2*p) + (1-p) * log(2*(1-p))$$

For example, if p = .99 then I(A, B) = .919.

3 Problem Statement

To test this model of channel capacity and mutual information, a Matlab program was written that would satisfy the important requirements listed above $(Pr(A_i) = \frac{1}{2}, Q)$ as described above). The necessary error correction was implemented using parity bits. The Matlab code is attached below.

```
PRECISION = 100;
         RETRY_COUNT = 10000;
         %first avg transmission
        %second error rate of transmission
       %third avg channel capacity
         A = zeros(3, 100);
         A = zeros(2, PRECISION);
        temp = 0;
10
         for i = 1:PRECISION
11
              for j = 1:RETRY_COUNT
12
                   temp = transmit(generate_random_msg(10,4), i/PRECISION, 20);
13
                   if temp == -1
14
                        A(2,i) = A(2,i) + 1;
15
                   else
16
                        A(1,i) = A(1,i) + temp;
17
                    end
               end
19
               %compute avg ignoring failed transmissions
              temp = RETRY_COUNT - A(2,i);
21
               if temp == 0
               % 200 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 100 = 10
23
                   A(1,i) = A(1,i)/(RETRY_COUNT-A(2,i));
25
               end
         end
27
28
         %compute channel cap graphics_toolkit
         %min bytes computed by testing for zero error
30
         min_bytes = transmit(generate_random_msg(10,4), 0, 20);
         %remember to subtract parity bytes because not in orginal msq
32
         A(3,:) = (ones(1,PRECISION)*(min_bytes-10))./A(1,:);
         figure
35
       %bits of info needed for transmission
36
         a = subplot(3,1,1);
         plot(A(1,:))
38
         ylabel(a,'Number of Bits Needed for Transmission')
         xlabel(a, 'Error Chance (1-Success Percentage)')
         %number of errors in transmission
         b = subplot(3,1,2);
42
         plot(A(2,:))
        ylabel(b,'Number of Errant Messages')
         xlabel(b, 'Error Chance (1-Success Percentage)')
       %channel capacity
       c = subplot(3,1,3);
47
        plot(A(3,:))
         ylabel(c,'Channel Capacity')
       xlabel(c,'Error Chance (1-Success Percentage)')
```

Listing 1: Main Simulator

```
function bit_count = transmit(msg, err_prob, retry_lim)
     msg = insert_parity_bit(msg);
     num_transmissions = 0;
     transmitted_msg = 0;
     for i = 1:size(msg)(1)
       transmitted_msg = transmit_msg(msg(i,:),err_prob);
       %check rough equality with parity bit
       %guarentees proper transmission for one or less errors
       temp_c = 0;
       while transmitted_msg(1) != parity(transmitted_msg(2:end)) && temp_c < retry_lim
11
         num_transmissions = num_transmissions + 1;
         temp_c = temp_c + 1;
13
         transmitted_msg = transmit_msg(msg(i,:),err_prob);
15
       %msg would fail using given parity scheme and error probality
16
       if isequal(transmitted_msg, msg(i,:)) == false
17
         bit_count = -1;
         return
19
       end
20
       %always one transmission
       num_transmissions = num_transmissions + 1;
22
23
24
     %num trans multiplied by the size of packet transmitted
25
     bit_count = num_transmissions * size(msg)(2);
26
   end
```

Listing 2: Transmission Simulator

```
function msg = transmit_msg(msg_in, err_prob)
      msg = zeros(size(msg_in));
      rnd_num = 0;
      for i = 1:size(msg_in)(2)
        rnd_num = rand();
        if rnd_num < err_prob</pre>
          if msg_in(i) == 0
            msg(i) = 1;
9
          else
            msg(i) = 0;
11
          end
13
          msg(i) = msg_in(i);
        end
15
      end
16
    end
17
```

Listing 3: Per-packet Transmission Simulator

Listing 4: Per-packet Parity Bit Inserter

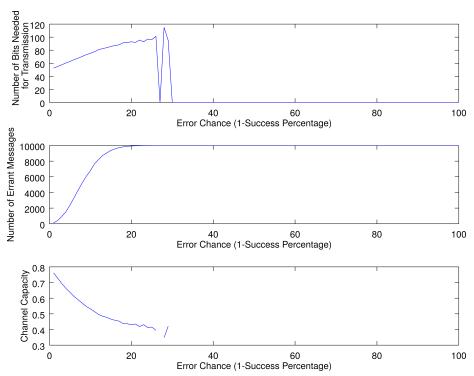
```
function par = parity(packet)
     num_ones = 0;
     for i = 1:size(packet)(2)
        if packet(i) == 1
          num_ones = num_ones + 1;
        end
      end
     if (mod(num\_ones, 2) == 0)
       par = 0;
9
      else
       par = 1;
11
      end
12
   end
13
```

Listing 5: Parity Bit Message Inserter

```
function msg = generate_random_msg(num_packets, packet_len)
msg = zeros(num_packets,packet_len);
for i = 1:num_packets
for j = 1:packet_len
msg(i,j) = floor(2*rand());
end
end
end
end
end
```

Listing 6: Generator of Random Messages with Packets

4 Results



Because the error correction was implemented using parity bits, the maximum channel capacity possible is $\frac{4}{5}$. This is since one new bit is added to each four bit transmission package thus lengthening the transmitted message at a $\frac{4}{5}$ rate.

5 Analysis