

Non-blocking Point to Point Communications

Non-blocking Communications

- In the last exercise of the previous lecture you will have seen how easily communications can block one another
- Non-blocking communications get around this problem
- The basic method in non-blocking communications is as follows
 - Set up all the sends and receives
 - Potentially do some computation while the communication occurs
 - Wait for communications to complete
 - Communications occur in the background

An example with non-blocking sends and receives

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>
```

```
using namespace std;
```

```
int id, p;
```

```
int main(int argc, char *argv[])
{
```

```
    MPI_Init(&argc, &argv);
```

```
    MPI_Comm_rank(MPI_COMM_WORLD, &id);
```

```
    MPI_Comm_size(MPI_COMM_WORLD, &p);
```

```
    srand(time(NULL)+id*10);
```

```
    int tag_num = 1;
```

```
    if (id == 0)
    {
        MPI_Request* request=new MPI_Request[p-1];
        int *send_data = new int[p-1];

        for (int i = 1; i < p; i++)
        {
            send_data[i - 1] = rand();
            MPI_Isend(&send_data[i-1], 1, MPI_INT, i, tag_num, MPI_COMM_WORLD, &request[i-1]);
            cout << send_data[i-1] << " sent to processor " << i << endl;
            cout.flush();
        }
        MPI_Waitall(p - 1, request, MPI_STATUS_IGNORE);

        delete[] send_data;
        delete[] request;
    }
    else
    {
        int rcv_data;
        MPI_Request request;
        MPI_Irecv(&rcv_data, 1, MPI_INT, 0, tag_num, MPI_COMM_WORLD, &request);

        MPI_Wait(&request, MPI_STATUS_IGNORE);

        cout << rcv_data << " received on processor " << id << endl;
        cout.flush();
    }
    MPI_Finalize();
}
```

Non-blocking communication

- At first glance this might seem to be doing exactly the same thing as the first example in the blocking communication lecture
 - Overall it does achieve the same thing, but some important differences:
 - The `MPI_Isend` and `MPI_Irecv` exit straight away without waiting for the communications to complete
 - The data will be sent in the background. You can carry out computations while the data is being sent:

```
MPI_Irecv(&recv_data, 1, MPI_INT, 0, tag_num, MPI_COMM_WORLD, &request);  
//You can do stuff here that doesn't require the communication to have finished  
MPI_Wait(&request, MPI_STATUS_IGNORE);
```
 - The data will be sent from processor zero to whichever process is first able to receive the data
 - Note that I have an array to store the send data. This is so that all values are still available until `MPI_Waitall`

Blocking and Non-blocking communications

- Note that in the previous example I could have used `MPI_Recv` rather than an `MPI_Irecv` and an `MPI_Wait`
 - There is only one communication on the processes other than 0
- `MPI_Recv` can receive data sent by `MPI_Isend` and vice versa
- In this case `MPI_Isend` is the best choice to send the data
 - Communications can occur in the order that the other processes are ready to receive rather than in the order of the process id as would be the case with a loop containing `MPI_Send`
 - The receive in this case makes little difference whether it is blocking or non-blocking
 - I would make a difference if we wanted to do some computation while waiting for the communications to complete

MPI_Isend and MPI_Irecv

```
int MPI_Isend(  
    const void *buf,  
    int count,  
    MPI_Datatype datatype,  
    int dest,  
    int tag,  
    MPI_Comm comm,  
    MPI_Request *request)
```

```
int MPI_Irecv(  
    void *buf,  
    int count,  
    MPI_Datatype datatype,  
    int source,  
    int tag,  
    MPI_Comm comm,  
    MPI_Request * request)
```

- You will notice that `MPI_Isend` and `MPI_Irecv` are very similar to `MPI_Send` and `MPI_Recv`. There are two difference:
 - They both include a `request` variable – more on that later
 - `MPI_Irecv` does not have a status variable – this is because the receive is not blocking and so `MPI_Irecv` actually exits before the communication is necessarily finished and so we do not know the status of the communication at this point

MPI_Wait and MPI_Waitall

- While it is useful to do things while waiting for the communications to finish, at some point you need to be able ensure that a communication has completed so that you can use the data that has been sent
 - This is what `MPI_Wait` and `MPI_Waitall` achieves
 - `MPI_Wait` is used to wait for a single communication to complete
 - `MPI_Waitall` is used to wait for a list of communications to complete
- Note that `MPI_Wait` and `MPI_Waitall` are blocking for the specific communications involved
 - It will continue once its specific communications are finished, irrespective of communications involving other processes

MPI_Wait and MPI_Waitall

```
int MPI_Wait(  
    MPI_Request *request,  
    MPI_Status *status)
```

```
int MPI_Waitall(int count,  
    MPI_Request *request_list,  
    MPI_Status *status_list)
```

- `request` is a pointer to a single request object
 - This must match a single `MPI_Isend` or `MPI_Irecv`
- `status` is a pointer to an object which will receive the communication status
 - This can be `MPI_STATUS_IGNORE` if you do not need the status information
- `count` is the number of communications that `MPI_Waitall` will be processing
 - These can be a combination of `MPI_Isend` and `MPI_Irecv`s
- `request_list` is an array that contains `count` request objects, one for each communication
- `status_list` is an array of status objects in which the status of each communication will be stored
 - This can again be `MPI_STATUS_IGNORE`

Sending from everyone to everyone

- This is now easy to achieve using non-blocking communications, but this is still less efficient than a collective operation (which we will cover later)
 - Typically do non-blocking communications such as this where processes are communicating with a subset of the other processes and where these subsets overlap
 - Domain decomposition is an example of this type of problem

Sending from everyone to everyone

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>
```

```
using namespace std;
```

```
int id, p;
```

```
int main(int argc, char *argv[])
{
```

```
    MPI_Init(&argc, &argv);
```

```
    MPI_Comm_rank(MPI_COMM_WORLD, &id);
    MPI_Comm_size(MPI_COMM_WORLD, &p);
    srand(time(NULL)+id*10);
```

```
    int tag_num = 1;
```

```
    MPI_Request* request = new MPI_Request[(p - 1)*2];
    double *send_data = new double[p];
    double *recv_data = new double[p];
```

```
    int cnt = 0;
```

```
    for (int i=0;i<p;i++)
        if (i != id)
        {
            MPI_Irecv(&recv_data[i], 1, MPI_DOUBLE, i, tag_num, MPI_COMM_WORLD, &request[cnt]);
            cnt++;
        }
    else recv_data[i] = 0;
```

```
    for (int i = 0; i<p; i++)
        if (i != id)
        {
            send_data[i] = (double)id / (double)p;
            MPI_Isend(&send_data[i], 1, MPI_DOUBLE, i, tag_num, MPI_COMM_WORLD, &request[cnt]);
            cnt++;
        }
    else send_data[i] = 0;
```

```
MPI_Waitall(cnt, request, MPI_STATUS_IGNORE);
```

```
for (int i = 0; i < p; i++)    cout << "Processor " << id << " recieved " << recv_data[i] << " from processor " << i << endl;
```

```
delete[] request;
delete[] recv_data;
delete[] send_data;
```

```
MPI_Finalize();
```

```
}
```

A couple of things to notice...

- With the `MPI_Isend` I had to store the values of each of the data points as a separate variable in an array
 - This is because the data can be sent anytime until the `MPI_Waitall` is called
 - This means that the data musn't be changed/overwritten or go out of scope within this interval
 - This is different to `MPI_Send`, where the data to be sent is buffered and/or actually sent by the time `MPI_Send` completes
 - The same is true of `MPI_Irecv`, though this is more obvious as you want to have the data after the `MPI_Wait` / `MPI_Waitall`
- You will notice that I set up the receives before the sends. This is because data can be sent as soon as there is a matching pair of sends and receives
 - Therefore slightly more efficient to have receive ready and waiting for a send
 - Possibly even more efficient to have receives and sends interleaved with an ordering such that some of the communications can get started while others are still being set up
 - This is likely to result in a more complex code structure and may not be worth it for the small efficiency gain
 - Remember that all the non-blocking communications are occurring in the background on a separate communications thread

Do Worksheet 2 Exercise 1

A side note... Timing your code

- You will often want to know how efficient your code is and so may wish to time it
 - This is a timed version of the previous program

```
#include <mpi.h>
#include <iostream>
#include <iomanip>
#include <cstdlib>
#include <chrono>

#define DO_TIMING

using namespace std;

int id, p;

int main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);

    MPI_Comm_rank(MPI_COMM_WORLD, &id);
    MPI_Comm_size(MPI_COMM_WORLD, &p);
    srand(time(NULL) + id * 10);

    int tag_num = 1;

    MPI_Request* request = new MPI_Request[(p - 1)*2];
    double *send_data = new double[p];
    double *recv_data = new double[p];

#ifdef DO_TIMING
    //The timing starts here - I deliberately exclude the initialisation of MPI, srand and the memory allocation
    //You need to decide what it is that you are timing
    auto start = chrono::high_resolution_clock::now();
#endif

    int cnt = 0;
```

```
    for (int i=0; i<p; i++)
        if (i != id)
        {
            MPI_Irecv(&recv_data[i], 1, MPI_DOUBLE, i, tag_num, MPI_COMM_WORLD, &request[cnt]);
            cnt++;
        }
        else recv_data[i] = 0;

    for (int i = 0; i<p; i++)
        if (i != id)
        {
            send_data[i] = (double)id / (double)p;
            MPI_Isend(&send_data[i], 1, MPI_DOUBLE, i, tag_num, MPI_COMM_WORLD, &request[cnt]);
            cnt++;
        }
        else send_data[i] = 0;

    MPI_Waitall(cnt, request, MPI_STATUS_IGNORE);

#ifdef DO_TIMING
    //Note that I exclude couts when timing the code
    //Things like file writes can also be excluded to get a fairer reading
    for (int i = 0; i < p; i++)
        cout << "Processor " << id << " recieved " << recv_data[i] << " from processor " << i << endl;
#endif

#ifdef DO_TIMING
    //Note that this should be done after a block in case process zero finishes quicker than the others
    //MPI_Waitall on process 0 is blocking for communications involving process 0, which, in this case is all the communications - Otherwise explicitly use MPI_Barrier
    auto finish = chrono::high_resolution_clock::now();
    if (id == 0)
    {
        std::chrono::duration<double> elapsed = finish - start;
        cout << setprecision(5);
        cout << "The code took " << elapsed.count() << "s to run" << endl;
    }
#endif

    delete[] request;
    delete[] recv_data;
    delete[] send_data;

    MPI_Finalize();
}
```

Checking if communications are open

- In non-blocking communications it is sometimes useful to check if the communications are finished without requiring them to be finished
 - For instance, you could have a loop in which you do some work that does not require the data communication to be finished, with the loop continuing until the communication is finished
 - If you are repeatedly sending more data to a process to do calculations on, you could create the next receive as soon as the previous send has been done and then check if the communication has completed, more data being sent as soon as it completes
 - This is the basis for worksheet 2 workshop exercise 1
- There are two functions for achieving this – `MPI_Test` and `MPI_Testall`

MPI_Test and MPI_Testall

```
int MPI_Test(  
    MPI_Request *request,  
    int *flag,  
    MPI_Status *status)
```

```
int MPI_Testall(int count,  
    MPI_Request *request_list,  
    int *flag,  
    MPI_Status *status_list)
```

- These functions are similar to `MPI_Wait` and `MPI_Waitall` in terms of the required parameters
- The difference is the `flag` parameter, which is a pointer to an integer that will be 1 or 0 (true or false) depending on whether the communications associated with `request` or `request_list` have completed or not
 - Note that in `MPI_Testall` the flag will be true only if all the communications are finished

Example using MPI_Test and MPI_Testall

- In this example we send a lot of data and do things (albeit not very useful things!) while waiting for the communication to finish

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>

using namespace std;

int id, p;

void Do_Work(void)    //Some (not very useful) work
{
    int sum = 0;
    for (int i = 0; i < 100; i++) sum = sum + 10;
}

int main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);

    MPI_Comm_rank(MPI_COMM_WORLD, &id);
    MPI_Comm_size(MPI_COMM_WORLD, &p);
    srand(time(NULL)+id*10);
```

```
    int tag_num = 1, sent_num = 100000, cnt = 0, flag = 0;

    if (id == 0)
    {
        MPI_Request* request = new MPI_Request[p - 1];
        int **send_data = new int*[p - 1];
        for (int i = 0; i < p - 1; i++)
        {
            send_data[i] = new int[sent_num];
            for (int j = 0; j < sent_num; j++)
                send_data[i][j] = rand();
        }

        for (int i = 1; i < p; i++)
            MPI_Isend(send_data[i-1], sent_num, MPI_INT, i, tag_num, MPI_COMM_WORLD, &request[i - 1]);

        while (MPI_Testall(p - 1, request, &flag, MPI_STATUS_IGNORE) == MPI_SUCCESS && flag == 0)
        {
            Do_Work();
            cnt++;
        }
    }
```

Example using MPI_Test and MPI_Testall

```
        for (int i = 0; i < p - 1; i++)
            delete[] send_data[i];
        delete[] send_data;
        delete[] request;
    }
    else
    {
        int *recv_data = new int[sent_num];
        MPI_Request request;
        MPI_Irecv(recv_data, sent_num, MPI_INT, 0, tag_num, MPI_COMM_WORLD, &request);

        while (MPI_Test(&request, &flag, MPI_STATUS_IGNORE) == MPI_SUCCESS && flag == 0)
        {
            Do_Work();
            cnt++;
        }

        delete[] recv_data;
    }

    cout << "Process " << id << " did " << cnt << " cycles of work while waiting " << endl;
    cout.flush();

    MPI_Finalize();
}
```

- Note that I don't need to use the `MPI_Wait` or `MPI_Waitall` as I know that all the communications are complete when the while loop exits
- The reason why I have `(MPI_Test(&request, &flag, MPI_STATUS_IGNORE) == MPI_SUCCESS && flag == 0)` as the condition is so that I can both call `MPI_Test` and check the value of the flag within the same condition
 - In C/C++ conditions are evaluated from left to right and therefore the `MPI_Test` is done first
 - The check for `MPI_SUCCESS` is not in expectation of it not being a success, but rather so that the function can be called – If the return is not a success it implies a communication failure and the program is likely to crash anyway

Communications between all processes

- We have shown that non-blocking communications help when sending data between multiple processes
 - Also allows you to do things while waiting
 - Will also continue as soon as a communication is finished – communications don't need to “wait their turn”
- Non-blocking communications are generally better than blocking communications
 - Best alternative when communicating with a number of “neighbouring processes”, but not all the processes
 - E.g. processes connected as graph
 - This is what will be the case in, for instance, domain decomposition – more on this later
- If you want to communicate between all processes at the same time collective communications are generally best

Do Worksheet 2 Exercise 3

A side note: Using Standard Template Library containers with MPI

- I have thus far been doing examples where memory is allocated using `new` directly
- You can use STL data types, but you need to watch out for a few things:
 - Data must be allocated to the container before using send or receive
 - Data must not move location in memory between the send or receive being setup and the communication being carried out
 - E.g. use `resize` before doing any sends or receives and DON'T use `push_back`
 - `push_back` can cause new data to be allocated and thus cause existing data to move location
 - When transferring more than one item of the same type at the same time you need to use a container that is guaranteed to store the items in contiguous memory
 - E.g. `vector` or `array` are fine as the standard demands contiguous memory be used, `map` or `set` are not fine as the internal memory format is not specified by the standard and all the implementations that I know do not use contiguous memory as that would be inefficient for these container types

The first example using containers

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
#include <time.h>
#include <vector>

using namespace std;

int id, p;

int main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);

    MPI_Comm_rank(MPI_COMM_WORLD, &id);
    MPI_Comm_size(MPI_COMM_WORLD, &p);
    srand(time(NULL)+id*10);

    int tag_num = 1;

    if (id == 0)
    {
        vector<MPI_Request> request;
        vector<int> send_data;

        request.resize(p-1);
        send_data.resize(p-1);

        for (int i = 1; i < p; i++)
        {
            send_data[i - 1] = rand();
            MPI_Isend(&send_data[i-1], 1, MPI_INT, i, tag_num, MPI_COMM_WORLD, &request[i-1]);
            cout << send_data[i-1] << " sent to processor " << i << endl;
            cout.flush();
        }

        MPI_Waitall(p - 1, request.data(), MPI_STATUS_IGNORE);
        //Note the use of request.data() rather than request – I could also use &request[0] instead
    }
    else
    {
        int recv_data;
        MPI_Request request;
        MPI_Irecv(&recv_data, 1, MPI_INT, 0, tag_num, MPI_COMM_WORLD, &request);

        MPI_Wait(&request, MPI_STATUS_IGNORE);

        cout << recv_data << " received on processor " << id << endl;
        cout.flush();
    }
    MPI_Finalize();
}
```