

# 1 The Collision Probability of Y Set

**Glossary** The symbol, abbreviation and marks needed in expressing the shift stage and probability computation.

## 1.1 Preliminaries

**The Equality of Two Sets** Assume there are two sets, marked as  $X\_A$  and  $X\_B$ , having same NO. of elements and the number is  $M$ . Each element is given an index from 1 to  $M$ . The definition of Identical Index Equality is expressed here:

**Definition 1.1.** *Identical Index Equality(IIE): For any  $i \in [1, M]$ ,  $X\_A[i] = X\_B[i]$ .*

The definition of Cross Index Equality is expressed here:

**Definition 1.2.** *Cross Index Equality(CIE): For any  $i, j \in [1, M]$ ,  $X\_A[i] = X\_B[j]$ , where  $i \neq j$*

Any two sets that have same number of elements can be splited into several sub-sets and each sub-sets contains several  $(X\_A[i], X\_B[i])$  pairs. The number of sub-sets various from 1 to  $M$ . We give the definition of set level equality of two sets with same number of elements:

**Definition 1.3.** *Set Level Equality(SLE): Each sub-set have one of the two kinds of Equality.*

For two sets  $X\_A$  and  $X\_B$ , if there is no way of split that at least one sub-set is of CIE and there are at least one way of split that all sub-sets are of IIE, then such  $X\_A$  and  $X\_B$  are IIE only Sets. The definition of CIE only Sets is given in same way.

### 1.1.1 Replay Attack and Y Set Collision

**The Message Sets in CETD Under Replay Attack**

- $D\_A$  and  $D\_B$  are definitely IIE.
- $X\_A = D\_A$ ,  $X\_B = D\_B$
- $Y\_A$  and  $Y\_B$  are just SLE
- when  $Y\_A$  and  $Y\_B$  are SLE, tag collision

In replay attack, the adversary replace a data-tag pair on the memory with a pair copied from the same address at an old time point. That means for the two pairs at different time point, the two message block sets and related tags are identical respectively, while the nonce  $N_A$  and  $N_B$  are randomly generated and the equality is unpredictable if their generator is of high quality. That means the shifting bits parameter segment on the nonce  $R\_A$  and  $R\_B$ , are randomly generated.

In this scenario, the probability of a successful attack can be expressed as the equation 1.1:

**Definition 1.4.**  $Pr[Successful\ Replay\ Attack] = Pr[Tag_a = Tag_b \mid (D_A = D_B) \ \& \ (R_A \text{ and } R_B \text{ are random})] = Pr[Y_A = Y_B \mid (D_A = D_B) \ \& \ (R_A \text{ and } R_B \text{ are random})]$

$D_A$  and  $D_B$  sets are of Identical Index Equality.  $Y_A$  and  $Y_B$  is of Set Level Equality.

We can see that the set level equality of  $Y_A$  and  $Y_B$  will directly leads to the collision of tag and cause the succeed of replay attack. For easy understanding, we assume the shuffle stage does not work at first, which means  $D_A = X_A = D_B = X_B$ (IIE). Hence the  $Y$  set is the output of rotate shifting stage in CETD, we will analyze the properties of block rotate shifting and the cases of input sets that can result  $Y$  set collision(SLE).

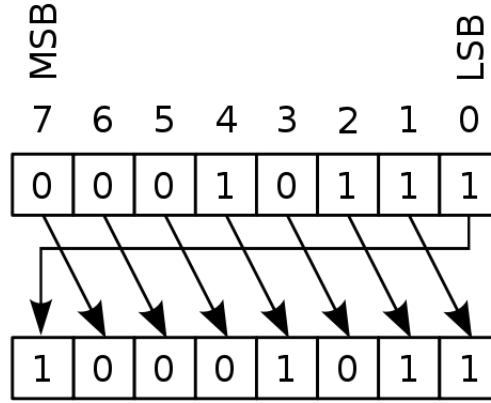


Figure 1: The Concept of Rotate Shifting(Right)

### 1.1.2 Rotate Shifting Introduction

- the concept of rotate shift
- The examples of X,R,Y

Unlike logical shifting and arithmetic shifting, rotate shifting behaves like whirling a wheel. The empty position in a message block shifted is filled by the bits shifted out. Figure 1 express the concept of rotate shift.

We can refer that the result of rotate shifting a message block depends on the value of block and the bits shifted. If two message blocks whose value are identical are shifted with distinct  $R_i$ s, the result blocks may be identical. That means when  $R_i$  fixed, the mapping from message blocks to the shifted result blocks is not injection. This case is expressed in Figure 2(a).

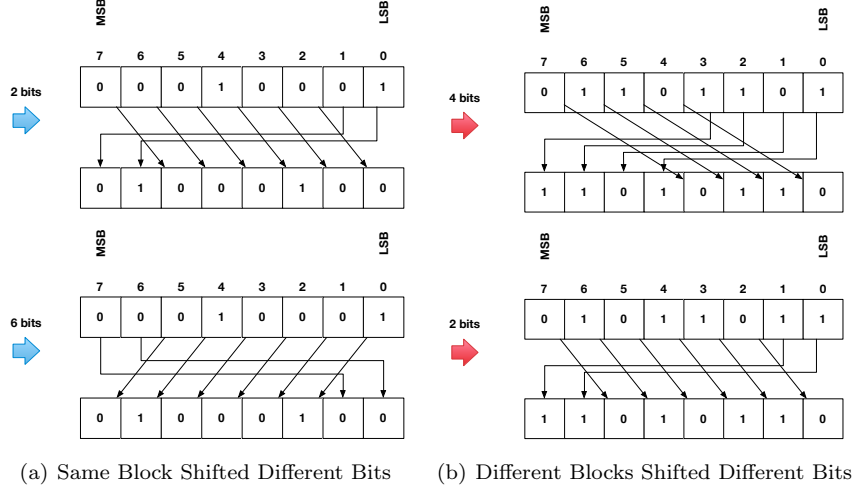


Figure 2: The Examples of Y Block Collision

For two distinct message blocks, however, their result blocks may be identical when their  $R_i$ s are distinct. This case is expressed in Figure 2(b). When the  $R_i$ s of two message blocks are identical, the equality of result blocks is same as the equality of their input blocks.

## 1.2 Rotate Shifting and Y Set Collision

- distinct pairs in R sets: IIE pairs and CIE pairs
- X block pairs causing IIE Y block pairs
- X block pairs causing CIE Y block pairs

Assume the shuffle stage does not work, then  $D\_A = X\_A = D\_B = X\_B$ , which means the following properties exist in  $X\_A$  and  $X\_B$ :

- $X\_A[i] = X\_B[i]$  for all  $i \in [1, M]$ ,  $M$  is the number of elements
- The equality of elements in a X set is uncertain.

That means  $X\_A$  and  $X\_B$  are of Identical Index Equality. All the analysis in this section is based on the assumption that shuffle stage does not work. If such  $X\_A$  and  $X\_B$  result two Y sets of Set Level Equality using two identical shifting bits parameter sets  $R\_A$  and  $R\_B$ , then the related tag  $T_a = T_b$ .

Figure 3 and Figure 4 express the examples of X sets that lead to SLE Y sets with distinct R sets. In this paper, we analyze the cases of the input set pair ( $X\_A$  and  $X\_B$ ) of shifting stage in CETD under replay attack that can lead to Y set pair collision (SLE).

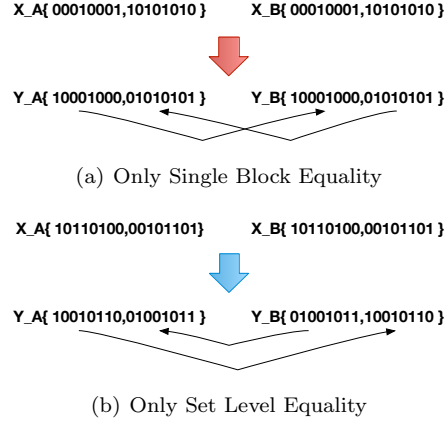


Figure 3: X Set Pairs with Only One Type of Y Equality

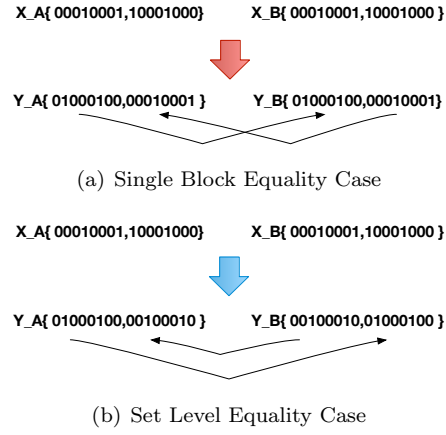


Figure 4: A X Set Pair with Two Types of Y Equality

### 1.2.1 Case of X Sets Resulting Y Set Collision

- If  $R_A \neq R_B$ , there is at least one pair distinct
- the related x blocks pairs should generate a SLE y pairs.
- if the Y pairs are iie, then each X block pair has pattern
- if the Y pairs are cie, then each X block in a sub-set is from a same base

Assume the number of distinct block pairs in  $R_A$  and  $R_B$  is  $R_d$ . For these  $R$  block pairs, their related  $X_A$  and  $X_B$  blocks formed a sub-group, marked as  $X_{A_s}$  and  $X_{B_s}$ , and this sub-group is IIE. The related sub-group in  $Y_A$  is marked as  $Y_{A_s}$  and the one in  $Y_B$  for  $Y_{B_s}$ .

In this part, we discuss the equality case of  $Y_{A_s}$  and  $Y_{B_s}$

**Identical Index Equality Case** If  $Y_{A_s}$  and  $Y_{B_s}$  are IIE, then for all  $i \in [1, R_d]$  the following two properties are met:

- $Y_{A_s}[i] = Y_{B_s}[i]$
- $R_{A_s}[i] \neq Y_{B_s}[i]$
- $X_{A_s}[i] = X_{B_s}[i]$

Figure 2(a) provides a instance of IIE Y sets. The identical block pairs ( $X_{A_s}[i], X_{B_s}[i]$ ) that lead to IIE  $Y_{A_s}$  and  $Y_{B_s}$  have the property in Theorem 1.1:

**Theorem 1.1.** Assume  $X_A[i]$  and  $X_B[i]$  are two identical block from two sub-group  $X_{A_s}$  and  $X_{B_s}$  with same index  $i$ .  $X_A[i]$  and  $X_B[i]$  have same number of bits  $N=2^n$ . The result of rotate shifting  $X_A[i]$  and  $X_B[i]$  with distinct shifting bit parameter  $R_A[i]$  and  $R_B[i]$  are marked as  $Y_A[i]$  and  $Y_B[i]$ .

Then  $Y_A[i]$  and  $Y_B[i]$  can be identical only when  $X_A[i]$  and  $X_B[i]$  are formed by repeating a binary pattern  $P$ , which is a binary segment. The length of this pattern  $P$  can be expressed as:

$$P_L = 2^p, p \in [0, n-1]$$

**Cross Index Equality Case** If  $Y_{A_s}$  and  $Y_{B_s}$  are CIE, then for all distinct  $(i, j)$  that  $i, j \in [1, R_d]$  the following two properties are met:

- $Y_{A_s}[i] = Y_{B_s}[j]$
- $R_{A_s}[i] \neq Y_{B_s}[i]$
- $X_{A_s}[i] = X_{B_s}[i]$

Figure 2(b) provides an instance of CIE Y sets. The blocks in  $X_{A_s}$  and  $X_{B_s}$  that lead to CIE  $Y_{A_s}$  and  $Y_{B_s}$  have the property in Theorem 1.2:

**Theorem 1.2.** If  $Y_{A_s}$  and  $Y_{B_s}$  are CIE, each element in  $X_{A_s}$  can be formed by rotate shifting another element in  $X_{A_s}$ .

If each element in  $X_{A_s}$  is formed by a pattern can be formed by rotate shifting another element in  $X_{B_s}$  then the related  $Y_{A_s}$  and  $Y_{B_s}$  can be either IIE or CIE or SLE. The proof of Theorem 1.1 and Theorem 1.2 can be refered in appendix.

### 1.3 The Probability of Y Set Collision

- Assume the sub-set contains M blocks, the probability that they are both pattern
- if they are both pattern, the probability that Y set are iie
- the probability that they are from same base
- if they are same base, the pro that Y set is cie
- the probability that bothe pattern and same base
- if both, the pro of Y SLE

Hence the shift bit parameter set  $R_A$  and  $R_B$  are randomly generated, the probability of Y set collision for each input case is determined by the number of specific combinations of  $R_A$  and  $R_B$ . The computation of  $\Pr[Y_A = Y_B]$  is based on this idea.

#### 1.3.1 Identical Index Equality Case

For ease analysis, we assume  $R_A[i] \neq R_B[i]$  for all  $i \in [1, M]$ , where M is the No. of elements in  $R_A$ . From Theorem 1.1 we know that  $Y_A$  and  $Y_B$  are IIE if each element in  $X_A$  is formed by a pattern P whose length  $P_L$  is  $2^p$ . Assume two identical sets  $X_A$  and  $X_B$  satisfy the following properties:

- Each element is formed by a pattern with same length  $P_L = 2^p$
- Each element cannot be formed by rotate shifting any other element.

We call this kind of  $X_A$  and  $X_B$  Identical Pattern Length(IPL) X Sets. IPL X sets can form IIE Y sets. Assume  $D_A$  set is randomly generated, then the probability of a IPL  $D_A$  is expressed as:

**Definition 1.5.**  $\Pr[IPL D(X)] =$

The probability of IIE of  $Y_A$  and  $Y_B$  is marked as  $\Pr[IIE Y sets]$  and expressed as the following way:

**Definition 1.6.**  $\Pr[IIE Y sets] =$

$$\prod_{i=1}^M \Pr[Y_A[i] = Y_B[i]]$$

*M is the No. of elements in a set*

As  $D\_A[i]=X\_A[i]=X\_B[i]=D\_B[i]$  for all  $i \in [1, M]$ ,  $\Pr[Y\_A[i] = Y\_B[i]]$  can be expressed as  $\Pr[Y\_A[i] = Y\_B[i] \mid X\_A[i] = X\_B[i] \ \& \ R\_A \text{ and } R\_B \text{ are randomly generated}]$ . We use  $\Pr[Y \text{ block collision}]$  to express this probability.

$\Pr[Y \text{ block collision}]$  is expressed in Theorem 1.3:

**Theorem 1.3.** *If the No. of bits of each block in  $X\_A[i]$ - $X\_B[i]$  pair is  $N=2^n$ , the pattern length  $P_l=2^p$  where  $p \in [0, n-1]$ . The pattern contains no internal sub-pattern, then  $\Pr[Y \text{ block collision}] = 1/2^p$*

If each X set contains M elements,  $\Pr[SBE] = (1/2^p)^M$ . The proof of theorem can be referred in appendix.

### 1.3.2 Cross Index Equality Case

#### What to say in this section

- In general condition,  $Y\_A$  and  $Y\_B$  are multisets. If  $Y\_A$  and  $Y\_B$  are identical in set level, then  $Y\_B$  is a permutation of  $Y\_A$ .
- how to make two identical Y block with two distinct X block and R block: if there is no pattern, for each value of block, the responding Y block is definite. 1 r map to 1 Y value
- How to represent the expression of  $\Pr[\text{set level collision} \mid X \text{ collision} \ \& \ R \text{ distinct}]$
- If  $Y\_A$  is identical to  $Y\_B$  in set level, element in  $Y\_B$  can be regarded as a multiset permutation of set  $Y\_A$ .

Assume two identical sets  $X\_A$  and  $X\_B$  satisfy the following properties:

- Each element can be formed by rotate shifting any other element in the set
- None of the elements is formed by pattern.

Such two sets  $X\_A$  and  $X\_B$  can only result set level identical  $Y\_A$  and  $Y\_B$ . The probability of set level equality of  $Y\_A$  and  $Y\_B$  is marked as  $\Pr[SLE]$  and expressed as the following way:

**Definition 1.7.**  $\Pr[SLE] =$

$$\prod_{i=1, j=1}^M \Pr[Y\_A[i] = Y\_B[j]]$$

where  $M$  is the No. of elements in a set,  $i, j \in [1, M]$  and  $i \neq j$

#### what to say next

- for any  $Y\_A$ , the set level identical  $Y\_B$  is a permutation of  $Y\_A$
- The No. of  $Y\_B$  is effected by the value distribution of elements in  $Y\_A$

### 1.3.3 The Intersection Case

#### What to say in this section

- If there is pattern in the block, for each element, the mapping between  $r$  and  $Y$  is not 1-to-1.
- Compute set level equality first, then solve the several  $r$  to 1  $Y$  problem
- express the Probability with combinatorics

### 1.4 A General Condition of X Set and Related Y Set Collision

#### what to say in this section

- If shuffle stage works, the properties in some block pairs in  $D_A$  and  $D_B$  can be eliminated, while the remaining blocks maintain the pattern
- the no-pattern block pairs follow the analysis in Main Case 3.
- these patterned block pairs can be split into several, each group is a kind of pattern.
- the pattern in any two group is distinct
- assume the No. of element in each group is expressed as  $M_i$ , then the theorem of each case of X sets can be applied in each group

## A Proof of Pattern Introduction

### A.1 Proof of Theorem 1

This part proves that if two identical block  $X_A[i]$  and  $X_B[i]$  are shifted different bits and the result blocks remain identical,  $X_A[i]$  is formed by pattern. The pattern length and  $\delta = |R_A[i] - R_B[i]|$  has such correlation:

- If  $\delta = P_L$  then  $Y_A[i] = Y_B[i]$  where  $P_L$  the length of pattern

**Proof of Corollary 1.2** According to the definition of set block equality,  $Y_A[i] = Y_B[i]$  for all  $i \in [1, M]$ . As  $R_A \neq R_B$ , there is at least one pair  $(R_A[i], R_B[i])$  is distinct. As the related  $X_A[i]$  is identical to  $X_B[i]$ ,  $X_A[i]$  is formed by pattern based on Theorem 1.1.



## A.2 Proof of Theorem 2

Assume  $Y\_A[i]$  and  $Y\_B[j]$  are identical. That means the following equations are met:

- $Y\_A[i][b] = X\_A[i][(b+R\_A[i]) \bmod N]$
- $Y\_B[j][b] = X\_B[j][(b+R\_A[j]) \bmod N]$
- $Y\_A[i]_b = Y\_B[j]_b$
- $X\_A[i][(b+R\_A[i]) \bmod N] = X\_B[j][(b+R\_A[j]) \bmod N]$

**Proof of Corollary 1.3** According to the definition of set level equality,  $Y\_A[i]=Y\_B[j]$  for  $i,j \in [1,M]$ ,  $i \neq j$  and  $X\_A[i] \neq .$  As  $R\_A \neq R\_B$ , there is at least one pair  $(R\_A[i], R\_B[i])$  is distinct. As the related  $X\_A[i]$  is identical to  $X\_B[i]$ ,  $X\_A[i]$  is formed by pattern based on Theorem 1.1.

## B Proof of Probability Computation

### B.1 Proof of Theorem 1.4

### B.2 Proof of Theorem 1.5