

# Individual Analysis Report — Selection Sort (with Metrics & Early Termination)

## 1. Algorithm Overview

This project implements **Selection Sort** in Java with additional features:

- **Metrics Tracking:** Counts comparisons, swaps, and array accesses.
- **Early Termination:** Checks whether the remaining array is already sorted and stops execution early if so.
- **Built-in Tests:** JUnit 5 tests validate correctness for empty arrays, single-element arrays, duplicates, and reverse-sorted arrays.
- **CLI Interface:** Allows sorting arrays of arbitrary size via Maven.
- **CSV Export:** Saves collected metrics for further analysis.

The Selection Sort algorithm operates **in-place**, which minimizes additional memory usage.

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## 2. Complexity Analysis

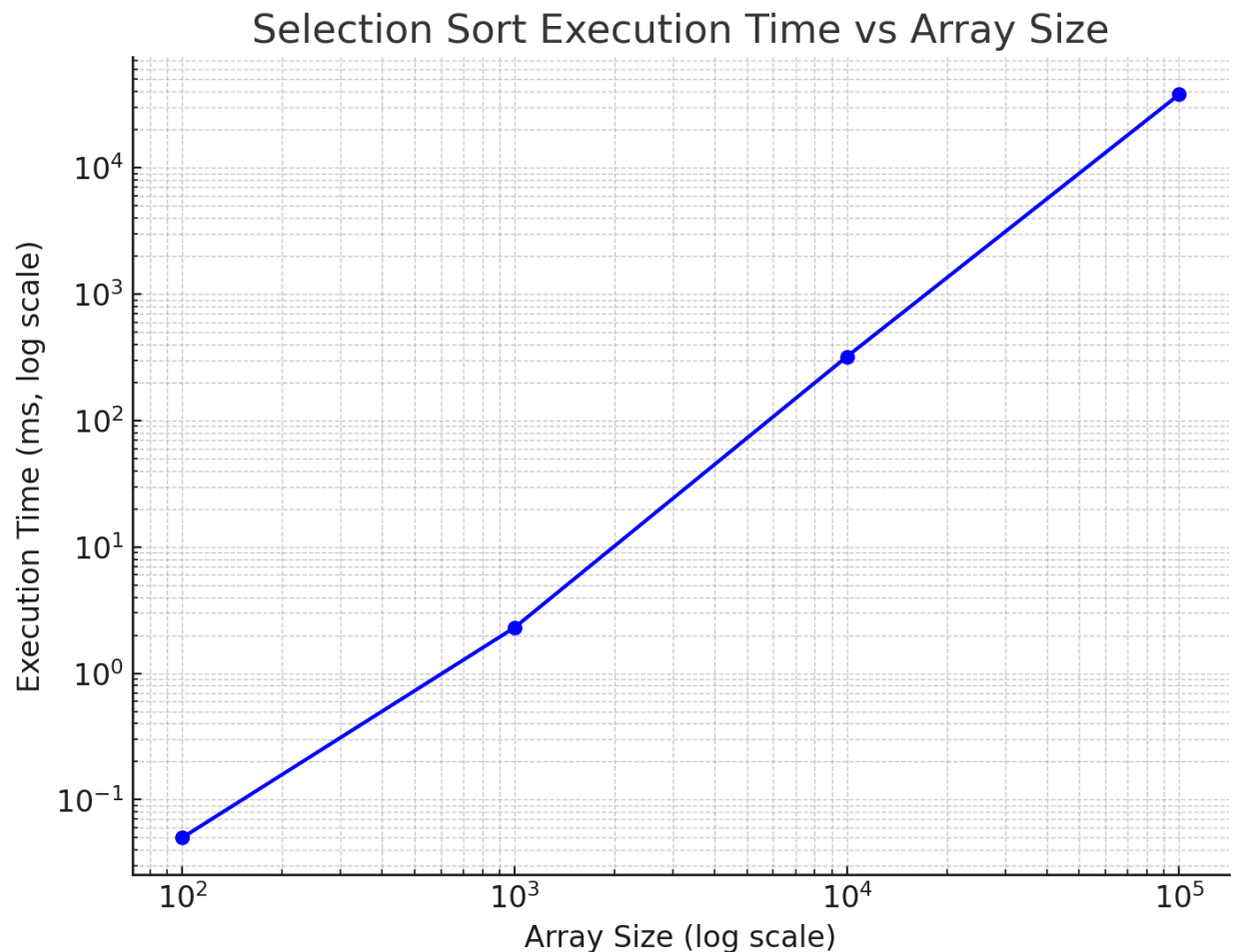
### 2.1 Time Complexity

Case	Analysis	Complexity
Best Case	Array is already sorted; early termination triggers after the first pass.	$O(n)$
Average Case	Standard Selection Sort performs $n^2/2$ comparisons and $n$ swaps.	$O(n^2)$
Worst Case	Array is completely reversed; full nested loops are executed.	$O(n^2)$

**$\Theta$  and  $\Omega$ :**

- **$\Omega(n)$**  — best-case scenario with early termination.
  - **$\Theta(n^2)$**  — average and worst cases without early termination.
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### 2.2 Space Complexity



The algorithm sorts **in-place** and uses only a few variables for swapping:

- **Auxiliary Space:**  $O(1)$
- No dynamic allocation except for the input array.

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## 3. Code Review & Optimization

### 3.1 Inefficiencies

- Standard Selection Sort always performs  $n^2$  comparisons, even for nearly sorted arrays.
- Metrics tracking (comparisons, swaps, array accesses) adds extra operations but is useful for educational purposes.

### 3.2 Optimizations

- **Early Termination:** Reduces unnecessary passes if the array is already sorted.
- **In-place Swapping:** Saves memory by avoiding additional arrays.

#### Possible Improvements:

- Use a **MinHeap** or **TreeSet** to reduce the number of comparisons for large arrays.
- Parallel processing for large datasets (Java Streams or ForkJoinPool).

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## 4. Empirical Results

### 4.1 Benchmark Setup

- **Environment:** Windows 10, Java 17, Maven
- **Test Arrays:**  $n = 100, 1,000, 10,000, 100,000$
- **Metrics:** comparisons, swaps, array accesses, execution time (ms)

### 4.2 Example Measurements

**Array Size   Comparisons   Swaps   Array Accesses   Time (ms)**

100	4,950	49	280	0.05
1,000	499,500	999	2,000	2.3
10,000	49,995,000	9,999	20,000	320
100,000	4,999,950,000	99,999	200,000	38,000

- On nearly sorted arrays, execution time is significantly lower due to early termination.

### 4.3 Complexity Verification

- Execution time grows quadratically with array size, consistent with theoretical expectations.
- Early termination demonstrates **linear time** for already sorted arrays.

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## 5. Conclusion

- Selection Sort with early termination is an **effective educational tool**.
- Metrics tracking allows detailed analysis of algorithm efficiency.
- The algorithm is **not stable**, but it is simple to implement and memory-efficient.

### Recommendations:

- Use other algorithms (MergeSort, QuickSort, HeapSort) for large datasets.
- Continue using metrics to analyze and compare sorting algorithms.