

① a) #pragma Parallel §
for (j=0; j < N; j++) §

Value + sum = 0

for (k=0; k < L; k++)

sum += A[get_xL + k] * B[k * N + j];

C[get_thread * N + j] = sum;

3

3

Parallelized the of each row in A, since there are no data dependencies in C[get_thread * N + j]

b) 1- Can't parallelize first loop since it is dependent on each loop & the result of the previous loop

2- Parallelizing second loop into one bigger loop

for (i=0; i < M; i++) §

#pragma Parallel §

if (get_thread < N-2) §

S[get_thread + 2] = 0

for (k = -2; k < 3; k++)

S[get_thread + 2] += 0.2 * V[get_thread + k];

3

V[get_thread] = S[get_thread]

3

2a) count the number of element that is less than the current element then put it in that element position

b) the outer most loop is the best one to parallelize

for($j=0$;) \Rightarrow Depends on $x[i]$

for($i=0$;) \Rightarrow Depends on nothing only temp array

There might be a race if the count is the same for two elements.