

## Theoretical Questions

① 1- First we fill each cell with the byte obtained by joining together its row index & column index

2- Replace the value in each cell by its multiplicative inverse in  $GF(2^8)$  based on the irreducible polynomial  $x^8 + x^4 + x^3 + x + 1$ .  
 $0 \times 00 \rightarrow$  itself since no inverse

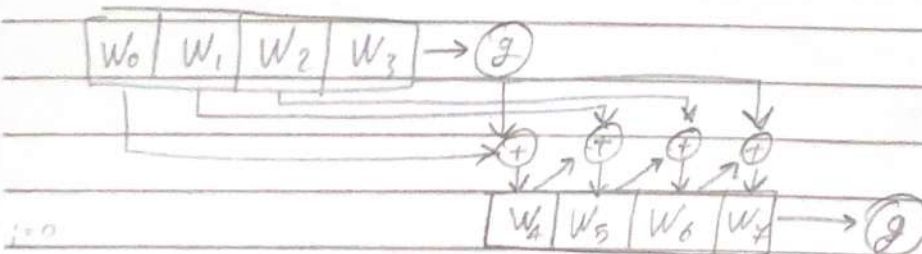
3- Represent each byte as  $b_7 b_6 b_5 b_4 b_3 b_2 b_1 b_0$ ,  $b_7 = \text{MSB}$  &  $b_0 = \text{LSB}$

4- Bit scramble using the following transformation for each bit  $b_i$ :

$$b_i' = b_i \otimes b_{(i+4) \bmod 8} \otimes b_{(i+5) \bmod 8} \otimes b_{(i+6) \bmod 8} \otimes b_{(i+7) \bmod 8} \otimes C_i$$

$C_i$  is the  $i$ th bit of sparsity designated byte  $C$  of value  $0 \times 63$

② We obtain them like so



$$\rightarrow W_{i+5} = W_{i+4} \otimes W_{i+1}$$

$$\rightarrow W_{i+6} = W_{i+5} \otimes W_{i+2}$$

$$\rightarrow W_{i+7} = W_{i+6} \otimes W_{i+3}$$

$$\rightarrow W_{i+4} = W_i \otimes g(W_{i+3})$$

③ The function  $g()$  consists of the following steps:

1- Perform a one byte left circular rotation on the argument 4-byte word

2- Perform a byte substitution for each byte of the word returned by the previous step by using the same  $16 \times 16$  table as used in SubBytes step in encryption round

3- XOR the bytes obtained from the previous step with what is known as a round constant. The constant is a word of whose 3 rightmost bytes are always zero.