## O. RECALL

We observe (Yi) =1,..., n such that

(1)  $y = \hat{f} + \xi$  ,  $\xi \sim W(0; \sigma^2 I_n)$ 

where  $f''=(f_1,...,f_n)\in\mathbb{R}^n$  is an unknown vector (called signal)  $\xi=(\xi_1,...,\xi_n)\in\mathbb{R}^n$  is a random noise

Dictionary:  $f' = \Phi \cdot \beta^*$  where  $\Phi \in \mathbb{R}^{n \times p}$  is a known design matrix

Sparsity: the dimension p of B\* is large, possibly larger than n, but many coordinates of B\* vanish:

5 = 11 p\*110 = \$ (1 p; 1 + 0) << p.

Lasso: Blasso E argmin { | |y- \P||\_2 + C ||B||\_1}

THEOREM: If  $\Phi$  is orthogonal, i.e.,  $\frac{1}{n}$   $\Phi\Phi = \mathbf{Ip}$ , then  $\hat{\beta}^{\text{Lasso}}$  coincides with the soft thresholding procedure. More precisely, for  $\lambda = \frac{C}{2n}$ , we have

 $\hat{\beta}_{j}^{\text{Lasso}} = \left( \left| \frac{1}{n} (\Phi^{\text{T}} y)_{j} \right| - \lambda \right)_{+} \operatorname{sign} \left( (\Phi^{\text{T}} y)_{j} \right), j = 1, \dots, p.$ 

## 1 Remarks

1) From now on, we will parameterize the Lasso by λ instead of C:
(2) β Lasso ε argmin { 1 | | y - Φβ || 2 + λ ||β ||<sub>1</sub> }.

2) The condition 1: of = Ip essentially means that the dictionary

3) Even if I is not ortho-normal (ON), the Lasso is efficiently computable even for very large values of p. Lemma: is a solution to (2) if and only if there exist (û, î, î) & RP × R" × R such that (\$, û, î, î) is a solution of the second-order cone program (SOCP): min  $\left\{\frac{1}{2n} \times \pm + \lambda \sum_{i=1}^{n} u_i\right\}$ subject to: -u; ≤ B; ≤ u; ¥ j=1,...,p (SOCP1) y- 48 = V  $\|(v_1 + \frac{1}{2})\|_2 \leq t + \frac{1}{2}$ Proof. 1) We prove first that if (\beta, \hat i, \beta, \hat i) is a solution of (50CP 1) then \$ is a solution of (2). Indeed, let \$ be any vector of RP. Define uj = 1p.1, j=1,..,p; v=y-\$B and t = ||y- pp 1/2. One early checks that all the constraints of (SOCPI) are fulfilled for (B, u, v, t) defined is such a way. Therefore, (p,u,v,t) is a feasible solution and  $\frac{1}{2n} \times \hat{t} + \lambda \sum \hat{u}_i \leq \frac{1}{2n} \times t + \lambda \sum u_i$ (3)  $= \frac{1}{2n} \|y - x \beta\|_2^2 + \lambda \|\beta\|_2$ On the other hand, since (\$, û, v, £) is a solution of (SOCPI) we have notice that the third constraint is equivalent to  $\|\mathbf{v}\|_{2}^{2} \leq t$  ;  $\|\mathbf{y} - \mathbf{v}\hat{\mathbf{p}}\|_{2}^{2} \leq t^{2}$  and  $\|\hat{\mathbf{p}}\|_{1} \leq \|\hat{\mathbf{u}}\|_{1}$ . Combined with (3), this implies that 1/2 | |y - \$\pi ||\_2 + \lambda || \pi ||\_1 < \frac{1}{2n} ||y - \pi (\pi) ||\_2 + \lambda || \pi ||\_1.

Since this holds true for every BERP, B is a solution of (2).

- 2) Let  $\hat{\beta}$  be a solution of (2) and define  $\hat{u}_j = |\hat{\beta}_j|$ , j = 1, ..., p;  $\hat{v} = y X\hat{\beta}$  and  $\hat{t} = ||\hat{v}||_2^2$ . One easily checks that  $(\hat{\beta}, \hat{u}, \hat{v}, \hat{t})$  verifies the constraints of (SOCP1). Furthermore, if  $(\beta, u, v, t)$  is another vector satisfying the constraints of (SOCP1), we have
- (4)  $\frac{1}{2n} \|y \Phi \hat{p}\|_{2}^{2} + \lambda \|\hat{p}\|_{1} \le \frac{1}{2n} \|y \Phi p\|_{2}^{2} + \lambda \|p\|_{1}$ The left-hand side of (4) is equal to  $\frac{1}{2n} \hat{t} + \lambda \sum_{j=1}^{p} \hat{u}_{j}$

whereas the right-hand side of (4) is bounded as follows:  $\frac{1}{2n} \|y - \Phi p\|_2^2 + \lambda \|p\|_1 \le \frac{1}{2n} \|v\|_2^2 + \lambda \sum_{j \ge 1}^p u_j \le \frac{1}{2n} t + \lambda \sum_{j \ge 1}^p u_j.$ 

This implies that

 $\frac{1}{2n}\hat{t} + \lambda \sum_{j=1}^{p} \hat{u}_{j} \leq \frac{1}{2n} + \lambda \sum_{j=1}^{p} u_{j}$ 

for every (p,u,v,t) satisfying the constraints.

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3. Prediction loss of the Lasso: the "slow" rater.

 $y = \Phi \beta^* + \xi$ ;  $\hat{\beta} Lasso = argmin \left\{ \frac{1}{2n} \|y - \Phi \beta\|_2^2 + \lambda \|\beta\|_1 \right\}$ 

We wish to evaluate the prediction loss

$$\ln(\hat{x}, \hat{x}^*) = \frac{1}{n} \|\hat{x} - \hat{x}^*\|_2^2 = \frac{1}{n} \|\Phi(\hat{\beta} - \beta^*)\|_2^2 \triangleq \ln(\hat{\beta}, \beta^*)$$

Note: we divide by a since III- f\* 1/2 is a sum containing a terms.

Condition C1: For  $\forall j \in \{1,...,p\}$ , we have  $\|\vec{\Phi}^j\|_2^2 = \sum_{j=1}^n q_{ij}^2 \leq n$ .

Lemma For every  $S \in (0,1)$ , the event  $B = \{ \| \phi^T \xi \|_{\infty} \le \sigma \sqrt{2n \ln(p/\delta)} \}$  satisfies  $\mathbb{P}(B) \ge 1 - \delta$ .

Proof. We will prove that  $P(B^c) \leq S$ . Set  $z = \sqrt{2 \ln (P/S)}$  and  $S_j = (\Phi^j)^T \xi / (\sigma \cdot ||\Phi^j||_2)$ ; j = 1, ..., p. Since Gaussian distribution is stable by affine transformations, we have  $S_j \sim \mathcal{N}(0,1) \, \forall j$ . Therefore,

 $P(\mathcal{B}^{c}) = P\left(\max_{j} |(\phi^{j})^{T} \xi| > \sigma \sqrt{2n \ln(\rho/\delta)}\right)$   $= P\left(\bigcup_{j \ge 1}^{p} \{|(\phi^{j})^{T} \xi| > \sigma \sqrt{n} \cdot z\}\right)$   $\leq \sum_{j \ge 1}^{p} P\left(|(\phi^{j})^{T} \xi| > \sigma \sqrt{n} \cdot z\right)$ 

 $\leq \sum_{j=1}^{n} \mathbb{P}(|\phi^{ij}\xi| > \sigma|\phi^{i}|_{2} \cdot z)$   $= \sum_{j=1}^{n} \mathbb{P}(|\xi_{j}| > z) = 2p \mathbb{P}(\xi_{1} > z)$ 

cohere for the last equality we used the symmetry of the Gaussian distribution. To complete the proof we will use the Gaussian tail bound:  $P(\xi_1 > x) \leq \frac{1}{3} \exp(-x^2/2)$ ,  $\forall x > 0$ .

This leads to

P(B°) < 2p × f × exp(- 1/2 (√2ln(p/5))2) = 8. 1

For completeness, we provide the proof of the Gaussian tail bound.

 $\mathbb{P}(S_1 > \infty) = \frac{1}{\sqrt{2\pi}} \int_{\infty}^{\infty} e^{-u^2/2} du$ 

Set  $G(x) = \mathbb{P}(5, > x) - \frac{1}{2} e^{-x^2/2}$ . One early checks that  $G'(x) = \frac{-1}{\sqrt{2\pi}} e^{-x^2/2} + \frac{x}{2} e^{-x^2/2} = \frac{e^{-x^2/2}}{2} \left(x - \sqrt{\frac{2}{\pi}}\right)$ .

This implies that G is decreasing on  $JO, \sqrt{2/\pi} \left[ \text{ and increasing on } JO, \sqrt{2/\pi$ 

From now on, we will always work on the event B.

Theorem: If  $\Phi$  satisfies condition C1, then on  $\mathcal{B}$ , it holds that (5)  $\ln(\hat{\beta}^L, \beta^*) \leqslant \inf_{\beta \in \mathbb{R}^p} \left\{ \ln(\beta, \beta^*) + 4\lambda \|\beta\|_1 \right\}$ provided that  $\lambda$  is chosen  $\geq \sigma \sqrt{\frac{2}{n}} \ln(\rho/\delta)$ .

Proof. Let  $\beta \in \mathbb{R}^{p}$  be any vector. Then  $\frac{1}{2n} \| y - \varphi \beta \|_{z}^{2} + \lambda \| \beta \|_{1} \leq \frac{1}{2n} \| y - \varphi \beta \|_{z}^{2} + \lambda \| \beta \|_{1}$ .  $\frac{1}{2n} \| \varphi (\beta^{*} - \hat{\beta}) + \xi \|_{z}^{2} + \lambda \| \beta \|_{1} \leq \frac{1}{2n} \| \varphi (\beta^{*} - \beta) + \xi \|_{2}^{2} + \lambda \| \beta \|_{1}$   $\frac{1}{n} \| \varphi (\beta^{*} - \hat{\beta}) \|_{z}^{2} + 2\lambda \| \beta \|_{1} \leq \frac{2}{n} \xi^{T} \varphi (\hat{\beta} - \beta) + 2\lambda \| \beta \|_{1} + \frac{1}{n} \| \varphi (\beta^{*} - \beta) \|_{2}^{2}$   $C_{n}(\hat{\beta}, \beta^{*}) \leq C_{n}(\beta, \beta^{*}) + 4\lambda \| \beta \|_{1} + \frac{2}{n} \xi^{T} \varphi (\hat{\beta} - \beta) - 2\lambda (\| \beta \|_{1} + \| \hat{\beta} \|_{1})$   $\leq n \text{ on } 35$   $\leq n \text{ effet}$   $\leq n \text{ of } \xi$   $\leq n \text{ on } 35$   $\leq n \text{ on } 35$   $\leq n \text{ on } 35$   $\leq n \text{ on } 35$ 

Commentaires

① On dit que (5) est une Oracle Inequality with "slow" rate. If an oracle tells us what is the best sparse approximation  $\overline{\beta}$  of  $\beta^*$ , then can use it as estimator of  $\beta^*$  and get the loss  $l_n(\overline{\beta}, \beta^*)$ . Since we do not know  $\overline{\beta}$ , we use  $\overline{\beta}$  lasso and get the loss bounded by  $l_n(\overline{\rho}, \beta^*) + Const \times \sqrt{\frac{l_n(\overline{\rho})}{n}} \times ||\overline{\rho}||_{\underline{\delta}}$  this is called "slow" rate.