




Chapter 12

Distributed Database Management Systems



Learning Objectives

- After completing this chapter, you will be able to:
 - Explain the purpose and function of distributed database management systems (DDBMSs)
 - Summarize the advantages and disadvantages of DDBMSs
 - Describe the characteristics and components of DDBMSs
 - Explain how database implementation is affected by different levels of data and process distribution
 - Understand how transactions are managed in a distributed database environment
 - Describe how distributed database design balances performance, scalability, and availability
 - Explain the trade-offs of implementing a distributed data system


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The Evolution of Distributed Database Management Systems (1 of 5)

- A distributed database management system (DDBMS)
 - Governs storage and processing of logically related data over interconnected computer systems
 - Distributes data and processing functions among several sites
- Centralized database management system
 - Required corporate data be stored in a single central site
 - Provided data access through dumb terminals
 - Filled structured information needs of corporations; fell short when quickly moving events required faster response times and equally quick access to information


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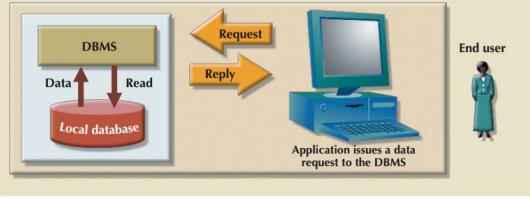

 The Evolution of Distributed Database Management Systems (2 of 5)

FIGURE 12.1 CENTRALIZED DATABASE MANAGEMENT SYSTEM



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 The Evolution of Distributed Database Management Systems (3 of 5)

- Changes that affected the nature of systems
 - Globalization of business operation
 - Increased market needs for an on-demand transaction style, based on web-based services
 - Rapid social and technological changes fueled by low-cost smart mobile devices
 - Converging data realms in the digital world
 - Advent of social media to reach new customers and markets

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 The Evolution of Distributed Database Management Systems (4 of 5)

- Database requirements in a dynamic business environment
 - Rapid ad hoc data access
 - Crucial in the quick-response decision-making environment
 - Distributed data access
 - Needed to support geographically dispersed business units

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Distributed Processing and Distributed Databases (2 of 3)

FIGURE 12.2 DISTRIBUTED PROCESSING ENVIRONMENT

Site 1
Miami user Joe
Computer A

Site 2
New York user Donna
Computer B
Update payroll data

Site 3
Atlanta user Victor
Computer C
Generate payroll report

DBMS
Employee database

Communications network

Database records are processed in different locations

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Distributed Processing and Distributed Databases (3 of 3)

FIGURE 12.3 DISTRIBUTED DATABASE ENVIRONMENT

Site 1
Miami user Alan
Computer A
DBMS
E1

Site 2
New York user Betty
Computer B
DBMS
E2

Site 3
Atlanta user Hernando
Computer C
DBMS
E3

Communications network

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Characteristics of Distributed Management Systems (1 of 3)

- A DBMS must have several functions to be classified as distributed
 - Application interface
 - Validation
 - Transformation
 - Query optimization
 - Mapping
 - I/O interface
 - Formatting
 - Security
 - Backup and recovery
 - DB administration
 - Concurrency control
 - Transaction management

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Characteristics of Distributed Management Systems (2 of 3)

- Functions of fully distributed DBMS
 - Receive the request of an application or end user
 - Validate, analyze, and decompose the request
 - Map request's logical-to-physical data components
 - Decompose request into several I/O operations
 - Search, locate, read and validate data
 - Ensure database consistency, security, and integrity
 - Validate data for conditions specified by request
 - Present data in required format
 - Handle all necessary functions transparently to user

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Characteristics of Distributed Management Systems (3 of 3)

FIGURE 12.4 A FULLY DISTRIBUTED DATABASE MANAGEMENT SYSTEM

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DDBMS Components

- The DDBMS must include at least the following components:
 - Computer workstations or remote devices
 - Network hardware and software components
 - Communications media
 - Transaction processor (TP)
 - Data processor (DP) or data manager (DM)

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The diagram illustrates a centralized architecture for single-site processing. On the left, a 'Host computer' contains a 'DBMS' (Database Management System) and a 'Database'. A 'Front-end processor' is connected to the Database. To the right, three 'Remote dumb terminal' units are shown, labeled T1, T2, and T3. Each terminal is connected to the Front-end processor. The terminals are collectively labeled as 'Dumb terminals'. The connection between the Front-end processor and the terminals is labeled 'Communication through DSL or T-1 line'.

Multiple-Site Processing, Single-Site Data (MPSD) (1 of 2)

- Multiple processes run on different computers sharing a single data repository
 - Typically requires network file server running conventional applications
 - Accessed through LAN
- Client/server architecture
 - Reduces network traffic
 - Distributes processing
 - Supports data at multiple sites


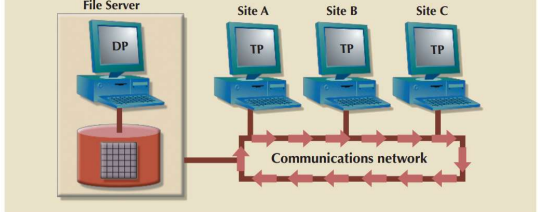
 Multiple-Site Processing, Single-Site Data (MPSD) (2 of 2)

FIGURE 12.7 MULTIPLE-SITE PROCESSING, SINGLE-SITE DATA




File Server
DP

Site A
TP


Site B
TP

Site C
TP


Communications network

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 Multiple-Site Processing, Multiple-Site Data (MPMD) (1 of 2)


- Fully distributed database management system
 - Support multiple data processors and transaction processors at multiple sites
- Classifications
 - Homogeneous: integrate multiple instances of same DBMS over a network
 - Heterogeneous: integrate different types of DBMSs over a network
 - Fully heterogeneous: support different DBMSs, each supporting different data model running under different computer systems

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 Multiple-Site Processing, Multiple-Site Data (MPMD) (2 of 2)


- Restrictions of DDBMS
 - Remote access is provided on a read-only basis
 - Restrictions on the number of remote tables that may be accessed in a single transaction
 - Restrictions on the number of distinct databases that may be accessed
 - Restrictions on the database model that may be accessed

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Distribution Transparency

- Allows management of physically dispersed database as if centralized
 - Levels: fragmentation, location, and local mapping
 - Unique fragment: each row is unique, regardless of the fragment in which it is located
 - Supported by distributed data dictionary (DDD) or distributed data catalog (DDC)
 - DDC contains the description of the entire database as seen by the database administrator
 - Distributed global schema: common database schema to translate user requests into subqueries

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23

Transaction Transparency

- Transaction transparency: DDBMS property that ensures database transactions will maintain the distributed database's integrity and consistency
 - Ensures transaction completed only when all database sites involved complete their part
- Distributed database systems require complex mechanisms to manage transactions
 - Ensure the database's consistency and integrity

Distributed Requests and Distributed Transactions

- Remote request
 - Single SQL statement accesses data processed by a single remote database processor
- Remote transaction
 - Accesses data at single remote site composed of several requests
- Distributed transaction
 - Requests data from several different remote sites on network
- Distributed request
 - Single SQL statement references data at several DP sites

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Distributed Concurrency Control (1 of 2)

- Concurrency control is especially important in distributed databases environment
 - Multi-site, multiple-process operations are more likely to create inconsistencies and deadlocked transactions
 - Solution to inconsistent database is a two-phase commit protocol

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Distributed Concurrency Control (2 of 2)

FIGURE 12.14 THE EFFECT OF A PREMATURE COMMIT

Data is committed


Can't roll back Sites A and B

Rollback at Site C

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Distributed Database Design


- Data fragmentation
 - How to partition database into fragments
 - Breaks a single object into two or more segments
- Data replication
 - Which fragments to replicate
 - Storage of data copies at multiple sites served by a computer network
 - Mutual consistency rule requires all copies of data fragments be identical
- Data allocation
 - Where to locate those fragments and replicas
 - Fully replicated database: stores multiple copies of each database fragment at multiple sites
 - Partially replicated database: stores multiple copies of some database fragments at multiple sites

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The CAP Theorem (1 of 2)

- CAP stands for:
 - Consistency
 - Availability
 - Partition tolerance
- Trade-off between consistency and availability generated in a new system in which data is basically available, soft state, eventually consistent (BASE)
 - Data changes are not immediate but propagate slowly through the system until all replicas are consistent


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The CAP Theorem (2 of 2)

**Table 12.8
Distributed Database Spectrum**

DBMS Type	Consistency	Availability	Partition Tolerance	Transaction Model	Trade-Off
Centralized DBMS	High	High	N/A	ACID	No distributed data processing
Relational DBMS	High	Relaxed	High	ACID (2PC)	Sacrifices availability to ensure consistency and isolation
NoSQL DDBMS	Relaxed	High	High	BASE	Sacrifices consistency to ensure availability
NewSQL DDBMS	High	High	Relaxed	ACID	Sacrifices partition tolerance to ensure transaction consistency and availability

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C. J. Date's 12 Commandments for Distributed Databases

Table 12.9	C.1 Data's 12 Commitments for Distributed Databases	
Rule Number	Rule Name	Rule Explanation
1	Local-site independence	Each local site can act as an independent, autonomous, centralized DBMS. Each site is responsible for security, concurrency control, backup, and recovery.
2	Central-site independence	Each local site can act as an independent, autonomous, centralized DBMS. Each site is responsible for security, concurrency control, backup, and recovery.
3	Failure independence	The system is not affected by node failures. The system is in continuous operation even in the case of a node failure or an expansion of the network.
4	Failure independence	The user does not need to know the location of data to retrieve that data.
5	Fragmentation transparency	Data fragmentation is transparent to the user, who sees only one logical database. The user does not need to know the name of the database fragment to which the data is assigned.
6	Replication transparency	The user sees only one logical database. The DBMS manager selectively selects the database fragment to access. To the user, the DBMS manages all fragments transparently.
7	Distributed query processing	A distributed query may be executed at several different DP sites. Query optimization is performed transparently by the DBMS.
8	Distributed transaction processing	A transaction may update data at several different sites, and the transaction is executed transparently.
9	Hardware independence	The system must run on any hardware platform.
10	Operating system independence	The system must run on any operating system platform.
11	Network independence	The system must run on any network platform.
12	Database independence	The system must support any vendor's database product.