

CS 325 Week 3-1

More database fundamentals (DB Reading Packet 2)
(and a little beyond)

More database fundamentals

from Kroenke:

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 - that is, it also contains information about the structure of the data in that database;

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- self-describing: means the database contains METADATA as well as data –
 - that is, it also contains information about the structure of the data in that database;
 - this description of the data structure is called a data dictionary or a data directory

More database fundamentals

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- "a database is a self-describing collection of integrated records"
- self-describing: means the database contains METADATA as well as data –
 - why would we want it to be self-describing?
 - ...for one thing, it can be used by a program to determine what a database contains;

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- "a database is a self-describing collection of integrated records"
- self-describing: means the database contains METADATA as well as data –
 - why would we want it to be self-describing?
 - ...for one thing, it can be used by a program to determine what a database contains;
 - ...that is, the self-describing aspect can help promote application program/data INdependence

Last login: Fri Sep 6 01:46:07 on ttys001

dl313@CS-S34217 ~ % ssh dl313@nrs-projects-ssh.humboldt.edu

dl313@nrs-projects-ssh.humboldt.edu's password:

Activate the web console with: systemctl enable --now cockpit.socket

Last login: Fri Sep 6 09:10:35 2024 from 137.150.249.93

[dl313@nrs-projects ~]\$

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[dl313@nrs-projects ~]\$ mkdir f24-325lect03-1

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Last login: Fri Sep 6 09:10:35 2024 from 137.150.249.93

[dl313@nrs-projects ~]\$ mkdir f24-325lect03-1

[dl313@nrs-projects ~]\$ chmod 700 f24-325lect03-1

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[dl313@nrs-projects f24-325lect03-1]\$ vim 325lect03-1.sql

```
-- Oracle, like many relational DBMSs, provides some metadata
--   in the form of tables![]
```

2 2 2 2 2 2 2 2 2 2 2 2 2 2 2

```
-- INSERT --
```

```
-- Oracle, like many relational DBMSs, provides some metadata
--   in the form of tables!
--
-- for example: a user_tables table, and a users_objects table
```

```
describe user_tables
```

[illegible]

```
-- INSERT --
```

Last login: Sun Sep 8 21:10:42 on ttys000

dl313@CS-S34217 ~ % ssh dl313@nrs-projects-ssh.humboldt.edu

dl313@nrs-projects-ssh.humboldt.edu's password:

Activate the web console with: systemctl enable --now cockpit.socket

Last login: Sun Sep 8 21:11:11 2024 from 137.150.231.63

[dl313@nrs-projects ~]\$ cd f24-325lect03-1

[dl313@nrs-projects f24-325lect03-1]\$ sqlplus /

SQL*Plus: Release 19.0.0.0.0 – Production on Sun Sep 8 21:13:35 2024
Version 19.3.0.0.0

Copyright (c) 1982, 2019, Oracle. All rights reserved.

Last Successful login time: Fri Sep 06 2024 09:10:49 -07:00

Connected to:

Oracle Database 19c Enterprise Edition Release 19.0.0.0.0 – Production
Version 19.3.0.0.0

SQL> @ 325lect03-1.sql


```
SQL> @ 325lect03-1.sql
```

Name	Null?	Type
TABLE_NAME	NOT NULL	VARCHAR2(128)
TABLESPACE_NAME		VARCHAR2(30)
CLUSTER_NAME		VARCHAR2(128)
IOT_NAME		VARCHAR2(128)
STATUS		VARCHAR2(8)
PCT_FREE		NUMBER
PCT_USED		NUMBER
INI_TRANS		NUMBER
MAX_TRANS		NUMBER
INITIAL_EXTENT		NUMBER
NEXT_EXTENT		NUMBER
MIN_EXTENTS		NUMBER
MAX_EXTENTS		NUMBER
PCT_INCREASE		NUMBER
FREELISTS		NUMBER
FREELIST_GROUPS		NUMBER
LOGGING		VARCHAR2(3)
BACKED_UP		VARCHAR2(1)
NUM_ROWS		NUMBER
BLOCKS		NUMBER
EMPTY_BLOCKS		NUMBER


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```

```
describe user_objects
```

```
-- INSERT --
```

INMEMORY_DISTRIBUTE	VARCHAR2(15)
INMEMORY_COMPRESSION	VARCHAR2(17)
INMEMORY_DUPLICATE	VARCHAR2(13)
DEFAULT_COLLATION	VARCHAR2(100)
DUPLICATED	VARCHAR2(1)
SHARDED	VARCHAR2(1)
EXTERNAL	VARCHAR2(3)
HYBRID	VARCHAR2(3)
CELLMEMORY	VARCHAR2(24)
CONTAINERS_DEFAULT	VARCHAR2(3)
CONTAINER_MAP	VARCHAR2(3)
EXTENDED_DATA_LINK	VARCHAR2(3)
EXTENDED_DATA_LINK_MAP	VARCHAR2(3)
INMEMORY_SERVICE	VARCHAR2(12)
INMEMORY_SERVICE_NAME	VARCHAR2(1000)
CONTAINER_MAP_OBJECT	VARCHAR2(3)
MEMOPTIMIZE_READ	VARCHAR2(8)
MEMOPTIMIZE_WRITE	VARCHAR2(8)
HAS_SENSITIVE_COLUMN	VARCHAR2(3)
ADMIT_NULL	VARCHAR2(3)
DATA_LINK_DML_ENABLED	VARCHAR2(3)
LOGICAL_REPLICATION	VARCHAR2(8)

SQL> @ 325lect03-1.sql

Name	Null?	Type
OBJECT_NAME		VARCHAR2(128)
SUBOBJECT_NAME		VARCHAR2(128)
OBJECT_ID		NUMBER
DATA_OBJECT_ID		NUMBER
OBJECT_TYPE		VARCHAR2(23)
CREATED		DATE
LAST_DDL_TIME		DATE
TIMESTAMP		VARCHAR2(19)
STATUS		VARCHAR2(7)
TEMPORARY		VARCHAR2(1)
GENERATED		VARCHAR2(1)
SECONDARY		VARCHAR2(1)
NAMESPACE		NUMBER
EDITION_NAME		VARCHAR2(128)
SHARING		VARCHAR2(18)
EDITIONABLE		VARCHAR2(1)
ORACLE_MAINTAINED		VARCHAR2(1)
APPLICATION		VARCHAR2(1)
DEFAULT_COLLATION		VARCHAR2(100)
DUPLICATED		VARCHAR2(1)
SHARDED		VARCHAR2(1)
CREATED_APPID		NUMBER

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describe user_tables
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```

```
-- another form of the SQL select statement lets you specify
--     just what columns you want to see from a table;
-- for example, here's how I can see JUST the table_name
--     column from the user_tables metadata table:█
```

```
~
~
~
~
~
~
~
~
~
~
~
~
```

```
-- INSERT --
```

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-- for example, here's how I can see JUST the table_name
--     column from the user_tables metadata table:
```

```
select table_name
from   user_tables;
```

```
~
~
~
~
~
~
~
~
```

```
-- INSERT --
```

DATA_OBJECT_ID	NUMBER
OBJECT_TYPE	VARCHAR2(23)
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LAST_DDL_TIME	DATE
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ORACLE_MAINTAINED	VARCHAR2(1)
APPLICATION	VARCHAR2(1)
DEFAULT_COLLATION	VARCHAR2(100)
DUPLICATED	VARCHAR2(1)
SHARDED	VARCHAR2(1)
CREATED_APPID	NUMBER
CREATED_VSNID	NUMBER
MODIFIED_APPID	NUMBER
MODIFIED_VSNID	NUMBER

SQL> @ 325lect03-1.sql

TABLE_NAME

CHOCOLATE
CLIENT
COMIC
COMIC_LOAN
CUSTOMER
DEPT
EMPL
MOVIE
MOVIE_CATEGORY
PARTS
PART_ORDERS

TABLE_NAME

RENTAL
STUDENT
STUFF
VIDEO

15 rows selected.

SQL>

```
--  
-- for example: a user_tables table, and a users_objects table  
  
describe user_tables  
  
describe user_objects  
  
-- another form of the SQL select statement lets you specify  
--     just what columns you want to see from a table;  
-- for example, here's how I can see JUST the table_name  
--     column from the user_tables metadata table:  
  
select table_name  
from   user_tables;  
  
-- and here's how I can see just the object_name and object_type  
--     columns from the user_objects metadata table:  
  
[  
~  
~  
~  
~  
~  
-- INSERT --
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select table_name  
from   user_tables;  
  
-- and here's how I can see just the object_name and object_type  
--     columns from the user_objects metadata table:  
  
select object_name, object_type  
from   user_objects;  
~  
~  
~  
-- INSERT --
```

```
[dl313@nrs-projects f24-325lect03-1]$ sqlplus /
```

```
SQL*Plus: Release 19.0.0.0.0 – Production on Sun Sep 8 21:25:35 2024  
Version 19.3.0.0.0
```

```
Copyright (c) 1982, 2019, Oracle. All rights reserved.
```

```
Last Successful login time: Sun Sep 08 2024 21:13:35 -07:00
```

```
Connected to:
```

```
Oracle Database 19c Enterprise Edition Release 19.0.0.0.0 – Production  
Version 19.3.0.0.0
```

```
SQL> @ 325lect03-1.sql
```

OBJECT_NAME

OBJECT_TYPE

CATEGORY_STATS
VIEW

CHOCOLATE
TABLE

CLIENT
TABLE

OBJECT_NAME

OBJECT_TYPE

COMIC
TABLE

COMIC_LOAN
TABLE

Question (3-1-Q1):

What does the term "**self-describing**" mean in the context of a database?

- 1) The database only stores data without any additional information.
- 2) The database includes metadata, which is information about the data and its structure.
- 3) The database cannot describe its own structure.
- 4) The database stores only user-entered data and not system information.

More database fundamentals

from Kroenke:

- "a database is a self-describing collection of integrated records"
- self-describing: means the database contains METADATA as well as data.
- Integrated: you not only have data that's self-describing, you also have a description of RELATIONSHIPS among the data records;

More database fundamentals

from Kroenke:

- "a database is a self-describing collection of integrated records"
- Integrated: you not only have data that's self-describing, you also have a description of RELATIONSHIPS among the data records;
 - foreign keys are one way to describe such relationships!

More database fundamentals

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More database fundamentals

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More database fundamentals

from Kroenke:

- "a database is a self-describing collection of integrated records"
- NOTE that a database is a model, not of reality, but of the user's model or view of what's important in some scenario;
 - model - "simplified abstraction of real-world events or conditions"
 - the degree of detail provided by a database should be based on the needs and desires of the users involved in the scenario that database is modeling or representing;

More database fundamentals

a LITTLE more database history:

More database fundamentals

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More database fundamentals

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More database fundamentals

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More database fundamentals

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 - and these often were VERY oriented towards TRANSACTIONS - it can be as simple as a representation of a class of significant events
 - these early organizational databases were VERY good at keeping track of regular (predicatable, regularly-occurring) transactions
 - also good at creating regularly scheduled reports;

More database fundamentals

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 - also good at creating regularly scheduled reports;
 - BUT! they were not very flexible,

More database fundamentals

- those early 1960's databases tended to be "organizational" in scope -- large companies, an important subset of such a large company;
 - also good at creating regularly scheduled reports;
 - BUT! they were not very flexible,
 - and application programs tended to need to be written in procedural languages such as COBOL and PL/I
 - ...users might think of a "new" question the data COULD conceivably answer -- but they might not be able or willing to wait for a programmer to get around to writing a procedural program to answer their questions;

Question (3-1-Q2):

Why were early databases not very flexible?

- 1) They were designed for small-scale, personal use
- 2) They could not handle scheduled tasks
- 3) Users needed custom programs to analyze data in new ways, which took time
- 4) They had built-in flexibility to handle any new queries instantly

More database fundamentals

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More database fundamentals

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- E.F. Codd developed this in 1970;
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 - instead of organizing a database using a hierarchical or network structure, this considers data to be stored in the form of RELATIONS, table-like structures that meet certain mathematical requirements;

More database fundamentals

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- E.F. Codd developed this in 1970;
 - It is based on relational algebra
 - instead of organizing a database using a hierarchical or network structure, this considers data to be stored in the form of RELATIONS, table-like structures that meet certain mathematical requirements;
 - relations have rows, and relations have columns, and relationship between rows of tables are visible in the data

More database fundamentals

- ...enter THE RELATIONAL MODEL!
- E.F. Codd developed this in 1970;
 - it turns out that relational algebra's relation operations provide a very useful way to express operations on relational tables that can map very nicely to QUESTIONS user might want to ask about the data;

More database fundamentals

- ...enter THE RELATIONAL MODEL!
- E.F. Codd developed this in 1970;
 - it turns out that relational algebra's relation operations provide a very useful way to express operations on relational tables that can map very nicely to QUESTIONS user might want to ask about the data;
 - and asking and answering questions over time remains very reasonable! And you can do so with relatively short SQL select statements (compared to procedural programs written in COBOL or PL/I or C++ or Java or JavaScript or Python, etc.

More database fundamentals

- Potential benefits of Codd's relational model:

More database fundamentals

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 - in a well-deigned relational database, data are stored in a way that minimizes unnecessarily duplicated data and eliminates certain kinds of day-to-day processing errors that can occur compared to some other data-storage approaches;

More database fundamentals

- Potential benefits of Codd's relational model:
 - in a well-deigned relational database, data are stored in a way that minimizes unnecessarily duplicated data and eliminates certain kinds of day-to-day processing errors that can occur compared to some other data-storage approaches;
 - columns can be used to contain data that relate one row to another row
 - it is easier for people to think about relational/tabular data than that is earlier hierarchical/network database approaches;

More database fundamentals

- Potential benefits of Codd's relational model:
 - ...
 - It provide much better support for ad-hoc queries, spur-of-the-moment questions over time and can encourage more-creative use of one's data;

More database fundamentals

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More database fundamentals

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 - there IS overhead here! providing this relational abstraction and performing these relational operations;

More database fundamentals

- there was initial resistance to the relational model (and there is a little now, also)
 - there IS overhead here! providing this relational abstraction and performing these relational operations;
 - fortunately, computer hardware and memory speed and power increased, and prices for these decreased, to where relational implementations could be practical;

Question (3-1-Q3):

What is a key characteristic of data organization in the relational model?

- 1) Data is stored in a hierarchy, with parent-child relationships between records.
- 2) Data is stored in tables (relations), with rows and columns, and relationships between rows are managed through foreign keys.
- 3) Data is organized in a network, with multiple parent and child nodes.
- 4) Data is stored sequentially, with no relationships between different tables.

More database fundamentals

- 1970s: Relational Model (E.F. Codd)
- Microcomputers (e.g., Apple II, Commodore 64) made computers accessible to more people.

More database fundamentals

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- 1980s: Client-Server Architecture

More database fundamentals

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- 1990s: Internet and Distributed Databases

More database fundamentals

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- Object-Oriented Databases

More database fundamentals

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- Microcomputers (e.g., Apple II, Commodore 64) made computers accessible to more people.
- Early Personal Databases.
- 1980s: Client-Server Architecture
- 1990s: Internet and Distributed Databases
- Object-Oriented Databases
- Modern Era: NoSQL Databases

More database fundamentals

- Four main elements of a (relational) database:
 - user data
 - metadata
 - indexes of various kinds - especially foreign keys
 - application metadata

More database fundamentals

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 - user data - the most obvious element!
 - in a relational database, user data is represented as relational tables;

More database fundamentals

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 - user data - the most obvious element!
 - in a relational database, user data is represented as relational tables;
 - we can depict these tables and/or their structure in quite a few ways -- here are THREE such ways:

More database fundamentals

- Four main elements of a (relational) database:
 - user data - the most obvious element!
 - in a relational database, user data is represented as relational tables;
 - we can depict these tables and/or their structure in quite a few ways -- here are THREE such ways:
 - tabular: actually list a table's rows and columns!

More database fundamentals

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STUDENT:

Stu_id	Stu_lname	Stu_phone	Adviser_id
-----	-----	-----	-----
123456	Jones	123-4567	234567
123457	Nguyen	234-5678	234568
123458	Garza	345-6789	234569

More database fundamentals

- We can depict these tables and/or their structure in quite a few ways -- here are THREE such ways:
 - tabular: actually list a table's rows and columns!
 - straightforward!
 - can see contents!
 - you can see the basic table structure,

More database fundamentals

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 - But
 - you can infer domain information for column, but you don't know its physical domain

More database fundamentals

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 - tabular: actually list a table's rows and columns!
 - straightforward! can see contents! you can see the basic table structure,
 - But
 - you can infer domain information for column, but you don't know its physical domain
 - you might be able to infer how tables are related, but not explicit;
 - can be a pain to type;

Question (3-1-Q4):

Which of the following is **NOT** one of the limitations of the **tabular form** in relational databases?

- A) It does not explicitly show relationships between tables
- B) It does not display the data types of columns
- C) It does not allow for easy navigation of data (actual rows of data)
- D) It can be tedious to manually enter large amounts of data

More database fundamentals

- We can depict these tables and/or their structure in quite a few ways -- here are THREE such ways:
 - tabular: actually list a table's rows and columns!
 - there are times when you care more about the table structure and less about the current rows in a table - relation structure form can be useful!

More database fundamentals

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 - tabular: actually list a table's rows and columns!
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```
name_of_relation(PRIM_KEY_COL1, PRIM_KEY_COL2, another_col2,  
                 another_col2, ...)  
    foreign key (another_col2) references another_relation
```


More database fundamentals

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 - tabular: actually list a table's rows and columns!
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```
Student(STU_ID, Stu_lname, Stu_phone, Adviser_id)
```

More database fundamentals

- We can depict these tables and/or their structure in quite a few ways -- here are THREE such ways:
 - there are times when you care more about the table structure and less about the current rows in a table - relation structure form can be useful!
 - still pretty straightforward; nice and concise; pretty easy to type; CAN see the primary key; can often see the foreign key.

More database fundamentals

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 - there are times when you care more about the table structure and less about the current rows in a table - relation structure form can be useful!
 - still pretty straightforward; nice and concise; pretty easy to type; CAN see the primary key; can often see the foreign key
 - But
 - don't see the rows in this form
 - don't see information about the column domains!

More database fundamentals

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 - create-table-statement form:
 - represent a relation as the SQL create table statement you could use to create that table or relation!

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 - create-table-statement form:
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```
create table student
(stu_id      char(6),
 stu_lname   varchar2(30),
 stu_phone   char(8),
 adviser_id   char(6),
 primary key  (student_id),
 foreign key  (adviser_id) references adviser
);
```

More database fundamentals

- We can depict these tables and/or their structure in quite a few ways -- here are THREE such ways:
 - create-table-statement form: represent a relation as the SQL create table statement you could use to create that table or relation!
 - definitely shows the columns; gives SOME physical domain information about each column; OUGHT to give primary key and foreign keys explicitly; can even create actual tables from them!
 - But
 - don't see rows in this form; this is a LITTLE more a pain to type...
 - maybe a little less easy for people to read than relation structure form...

Question (3-1-Q5):

Consider the different ways of depicting a relation that we discussed. Which **most explicitly** shows the **data type/physical domain** of the elements within a particular column?

- 1) tabular form
- 2) relation structure form
- 3) create-table statement form
- 4) metadata form

More database fundamentals

- Relational Databases Have Potential: However, poor table structure can waste this potential.
- An example of poorly-structures vs. "better"-structured relations.

More database fundamentals

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More database fundamentals

- Relational Databases Have Potential: However, poor table structure can waste this potential.
- An example of poorly-structures vs. "better"-structured relations.
 - The quality of structure depends on what you want to achieve and how you use the data.
 - Operational vs. Read-Only Databases.
 - The best structure depends on the context.

More database fundamentals

- An example of poorly-structures vs. "better"-structured relations;
 - because some database structures will be harder to live with over time, or will be easier to maintain data integrity over time, or will make representing some information easier over time, etc.

More database fundamentals

- An example of poorly-structures vs. "better"-structured relations;

```
Student(STU_ID, Stu_lname, Stu_phone, Adviser_id,  
        Adviser_lname, Adviser_phone)
```

More database fundamentals

- An example of poorly-structures vs. "better"-structured relations;

```
Student(STU_ID, Stu_lname, Stu_phone, Adviser_id,  
        Adviser_lname, Adviser_phone)
```

VS.

```
Student(STU_ID, Stu_lname, Stu_phone, Adviser_id)  
    foreign key (Adviser_id) references Adviser
```

```
Adviser(ADVISER_ID, Adviser_lname, Adviser_phone)
```

More database fundamentals

- First Approach (Single Table)
 - Data duplication (e.g., repeating advisor details for each student).
 - Updates (e.g., phone number changes) must be made in multiple places, increasing risk of errors.

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 - May have quick access to advisor information from the student table

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```


More database fundamentals

- Second Approach (Two Related Tables)
 - Advisor information stored once, avoiding duplication and ensuring data integrity.

```
Student(STU_ID, Stu_lname, Stu_phone, Adviser_id)  
    foreign key (Adviser_id) references Adviser
```

```
Adviser(ADVISER_ID, Adviser_lname, Adviser_phone)
```

More database fundamentals

- First Approach (Single Table)
 - Data duplication (e.g., repeating advisor details for each student).
 - Updates (e.g., phone number changes) must be made in multiple places, increasing risk of errors.
 - May have quick access to advisor information from the student table
 - Advisor information is lost if all students are deleted.
- Second Approach (Two Related Tables)
 - Advisor information stored once, avoiding duplication and ensuring data integrity.
 - Requires an extra step to link student and advisor data.

More database fundamentals

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 - Advisor information is lost if all students are deleted.
- Second Approach (Two Related Tables)
 - Advisor information stored once, avoiding duplication and ensuring data integrity.
 - Requires an extra step to link student and advisor data.
 - Better for long-term data management in day-to-day operations.

Question (3-1-Q6)

What is a major **disadvantage** of storing both student and advisor information in a **single table**?

- A) It requires more tables to be created.
- B) It can result in unnecessary duplication of advisor information, leading to potential data inconsistency.
- C) It makes retrieving advisor information impossible.

More database fundamentals

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More database fundamentals

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 - -- these can vary a lot, BUT:
 - if it is going to provide an interface between users/applications and the data,
 - abstracting out physical storage details,
 - then at the very least, it needs to provide:

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 - some kind of data DEFINITION language (DDL)

More database fundamentals

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 - some kind of data MANIPULATION language (DML)

More database fundamentals

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 - some kind of data CONTROL language (DCL)

More database fundamentals

- going BEYOND the DB Reading Packet 2 a little bit:
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 - some kind of data DEFINITION language (DDL)
 - some kind of data MANIPULATION language (DML)
 - some kind of data CONTROL language (DCL)
 - many DBMSs implement these by providing an implementation of SQL! (some also provide various GRAPHICAL DDL, DML and/or DCL)

More database fundamentals

- many DBMSs implement these by providing an implementation of SQL! (some also provide various GRAPHICAL DDL, DML and/or DCL)
 - DDL - let users define the database!
 - for a relational database, this lets you define tables!
 - e.g. SQL's create table statement

More database fundamentals

- DDL - let users define the database!
 - for a relational database, this lets you define tables!
 - e.g. SQL's create table statement
- DML - lets users insert, update, delete, and retrieve data
 - e.g. SQL's insert statement
 - and its update statement
 - and its delete statement
 - and its select statement

More database fundamentals

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- DML - lets users insert, update, delete, and retrieve data
 - e.g. SQL's insert statement
- DCL - lets users control access to the database

More database fundamentals

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 - security features, data integrity features, concurrency control, recovery system, and more

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More database fundamentals

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More database fundamentals

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```
GRANT select, insert, update  
ON abc123.stuff  
TO fg3, gh5, st10;
```

- -- says users fg3, gh5, st10 can select, insert, and update, -- (but NOT delete) rows from user abc123's stuff table

More database fundamentals

- DCL - lets users control access to the database
 - security features, data integrity features, concurrency control, recovery system, and more
 - SQL provides *some* DCL support - creation of views, ability to GRANT or REVOKE privileges on tables
 - GRANT lets you say particular users can be GRANTED select, insert, update, and/or delete access to a particular table.
 - REVOKE lets you remove previously-GRANTED access

```
REVOKE select
ON abc124.stuff
FROM gh5;
```
 - -- says user gh5 can NO LONGER select rows from user abc123's -- stuff table (but they can still insert rows and update -- them...!)

Question (3-1-Q7):

Consider the SQL create table statement. Which DBMS capability is this most reasonably considered to be part of?

- 1) its DDL, data definition language
- 2) its DML, data manipulation language
- 3) its DCL, data control language
- 4) its MDL, metadata language