

Big-O Notation

Tuesday, August 25, 2020

8:00 AM

Definition **Case 1** $f(x) = O(g(x))$ as $x \rightarrow \infty$ if

$\exists x_0$ and $\exists M < \infty$ s.t. if $x > x_0$ then $|f(x)| \leq M \cdot g(x)$
↑ "exists"

Case 2 $f(x) = O(g(x))$ as $x \rightarrow a$ (usually $a=0$) if

$\exists \delta > 0$ and $\exists M < \infty$ s.t. if $|x-a| < \delta$ then $|f(x)| \leq M \cdot g(x)$

(Both cases, equivalent def'n)

$f = O(g)$ if $\limsup_{x \rightarrow \infty} \left| \frac{f(x)}{g(x)} \right| < \infty$ } valid

or, if $\lim_{x \rightarrow \infty} \left| \frac{f(x)}{g(x)} \right|$ exists and is $< \infty$ } sufficient

Interpretations and Examples

∞ "infinite case"

$f = O(g)$ means, eventually, f grows no faster than g (up to constants)

(typically, $g(x) \rightarrow \infty$ as $x \rightarrow \infty$)

Use case: n = size of input, $f(n)$ = how long it takes to run

$$\begin{aligned} f(n) &= n^3 + 3n^2 - 4n + 7, & f(n) &= O(n^3) \\ f(n) &= 10n^3 + 3n^2 \dots & f(n) &= O(n^3), \quad f(n) \neq O(n^2) \end{aligned}$$

$f(n) = n^3, g(n) = n^2$, is $f = O(g)$? No

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{n^3}{n^2} = \lim_{n \rightarrow \infty} n = \infty. \quad \underline{\text{No}}$$

$$\begin{aligned} n^2 &= O(n^3) \quad \checkmark \\ n^2 &= O(n^2) \end{aligned}$$

$$x^2 = O(e^x), \quad e^x \neq O(x^2)$$

$a=0$ "infinitesimal case"

typically, $a=0$, write " h " instead of " x ", $g(h) \rightarrow 0$ as $h \rightarrow 0$
discretization

$f=O(g)$ means, eventually, f decays to 0 at least as fast as g
(up to a constant)

Ex: $\left. \begin{array}{l} 3h^2 = O(h^2), \quad h^2 = O(h) \\ h^2 \neq O(h^3) \end{array} \right\}$ check: $\lim_{h \rightarrow 0} \frac{h^2}{h} = \lim_{h \rightarrow 0} h = 0 < \infty$



WARNING

$x = O(x^2)$ as $x \rightarrow \infty$, $x^2 \neq O(x)$ as $x \rightarrow \infty$ ∞ case

$h \neq O(h^2)$ as $h \rightarrow 0$, $h^2 = O(h)$ as $h \rightarrow 0$ $a=0$ case

Smaller exponent is "better"

larger exponent is "better" (usually $f(h)$ is error term in Taylor series)

Ex: $f(x_0+h) = f(x_0) + \underbrace{f'(x_0)h}_P + \frac{f''(x_0)h^2}{2!} + \frac{f'''(\xi)h^3}{3!}$

then we can say
 $f(x_0+h) = P(h) + O(h^3)$

if $|f'''(\xi)|$ is bounded
for all $\xi \in [x_0, x_0+h]$

Variants

little-o notation:

$f=o(g)$ as $x \rightarrow \infty$ means (asymptotically) f grows slower than
 $c \cdot g(x)$ for all constants c

i.e., $\forall c > 0, \exists x_0$ s.t. $x > x_0, |f(x)| \leq c \cdot g(x)$

more 'precise' than big-O notation

$x^2 = O(x^2)$, $x^2 \neq o(x^2)$, $x^2 = o(x^3)$

big-theta θ

$f = \theta(g)$ means $f = O(g)$ and $g = O(f)$

ex: $5x^3 = \theta(10x^3)$ $x^3 \neq \theta(x^2)$, $x^2 \neq \theta(x^3)$

$f \sim g$ even stronger: means $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 1$

$5x^3 \not\sim 10x^3$, $5x^3 + 3x^2 \sim 5x^3 - x$

\tilde{O} ← ignore log factors

Ex: is $e^x = O(x!)$ or is $x! = O(e^x)$? ($x \rightarrow \infty$)

Use Stirling's formula: $x! \sim \sqrt{2\pi x} \left(\frac{x}{e}\right)^x$,

ie, $x! = \theta\left(\sqrt{x} \left(\frac{x}{e}\right)^x\right)$

x^x grows faster than e^x

so $x!$ grows faster than e^x , $e^x = O(x!)$, $x! \neq O(e^x)$
 $x! \neq \theta(e^x)$