

CS5200 Homework 2 Dynamic Programming

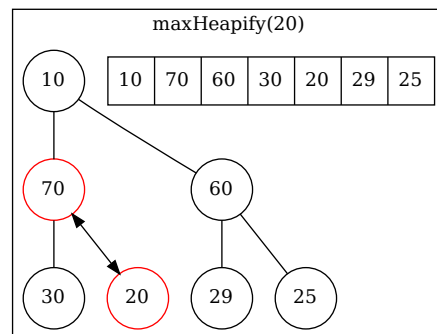
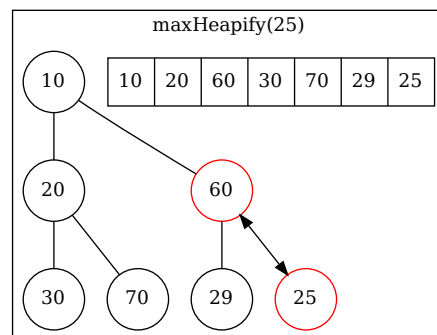
Adam McNeil

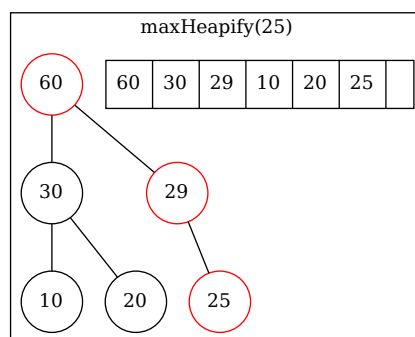
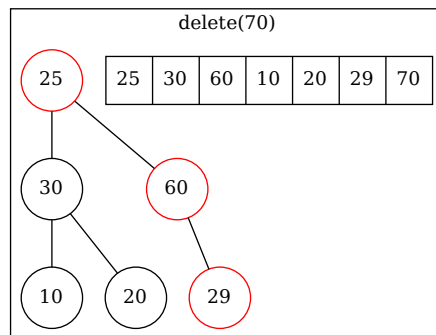
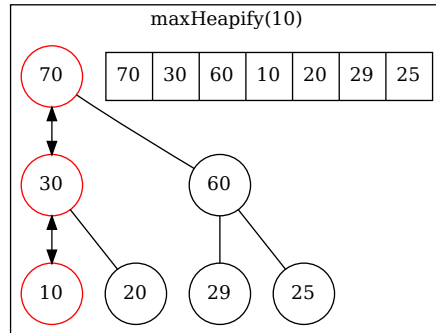
Question 1)

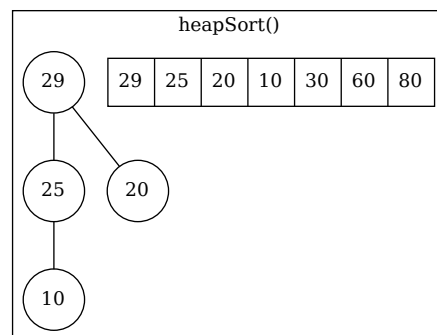
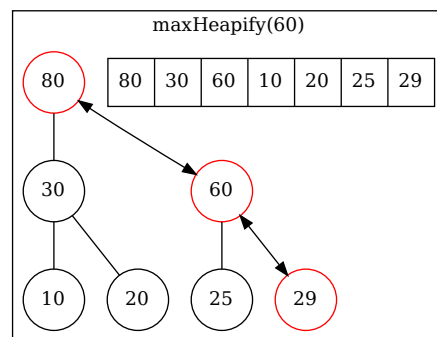
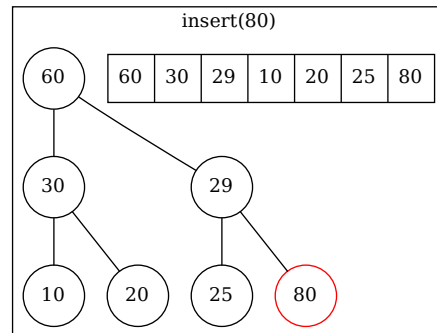
max heapify

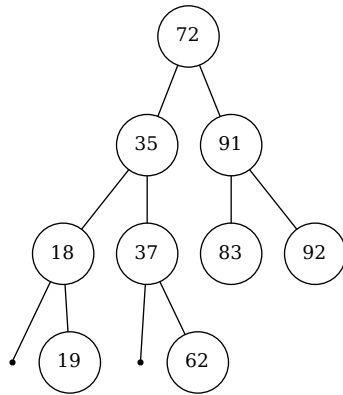
Call max heapify on all the internal nodes starting at the bottom

maxHeapify(25)







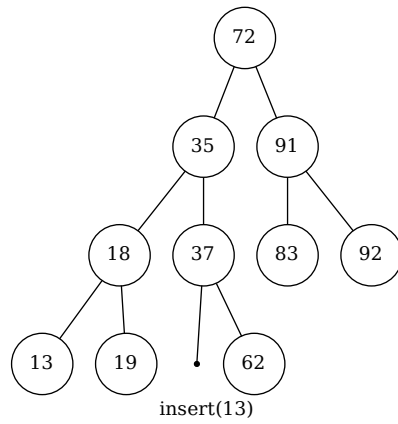


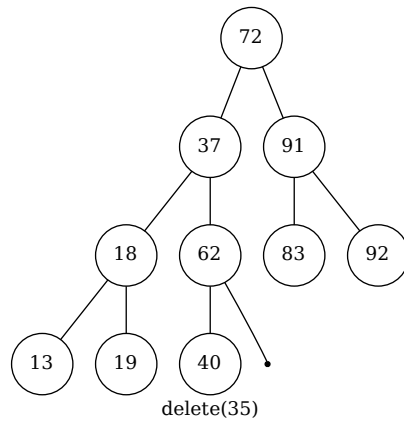
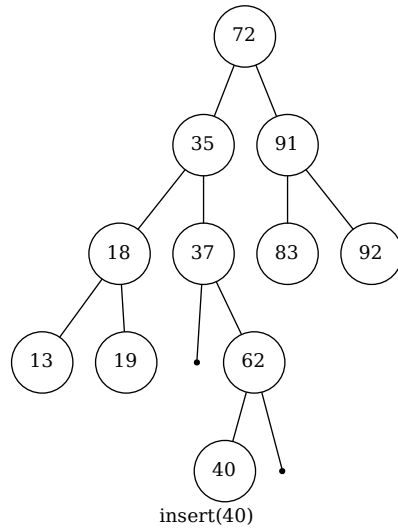
Question 2)

Pre-order: 72 35 18 19 37 62 91 83 92

In-order: 18 19 35 37 62 72 83 91 92

Post-order: 19 18 62 37 35 83 92 91 72





Question 3)

$p_0 = 4$ $p_1 = 10$ $p_2 = 3$ $p_3 = 12$ $p_4 = 7$

4	1	2	3	4	4	1	2	3	4
3	0	0	252	0	3	2	3	0	0
2	0	360	0		2	1	0		
1	120	0			1	0			

m(1, 3) i=1 j=3

$$\begin{aligned}
& k=1 \\
& m(1, 1) + m(2, 3) + p_0 p_1 p_3 \\
& 0 + 360 + 4*10*12 = 840
\end{aligned}$$

$$\begin{aligned}
& k=2 \\
& m(1, 2) + m(3, 3) + p_0 p_2 p_3 \\
& 120 + 0 + 4*3*12 = 264
\end{aligned}$$

$$\begin{aligned}
& \mathbf{m(2, 4) \ i=2 \ j=4} \\
& k=2 \\
& m(2, 2) + m(3, 4) + p_1 p_2 p_4 \\
& 0 + 252 + 10*3*7 = 462
\end{aligned}$$

$$\begin{aligned}
& k=3 \\
& m(2, 3) + m(4, 4) + p_1 p_3 p_4 \\
& 120 + 0 + 10*12*7 = 462
\end{aligned}$$

$$\begin{array}{c|c|c|c|c}
& 1 & 2 & 3 & 4 \\
4 & 0 & 462 & 252 & 0 \\
3 & 264 & 360 & 0 & \\
2 & 120 & 0 & & \\
1 & 0 & & &
\end{array}
\begin{array}{c|c|c|c|c}
& 1 & 2 & 3 & 4 \\
4 & & 2 & 3 & 0 \\
3 & 2 & 2 & 0 & \\
2 & 1 & 0 & & \\
1 & 0 & & &
\end{array}$$

$$\begin{aligned}
& \mathbf{m(1, 4) \ i=1 \ j=4} \\
& k=1 \\
& m(1, 1) + m(2, 4) + p_0 p_1 p_4 \\
& 0 + 462 + 4*10*7 = 742
\end{aligned}$$

$$\begin{aligned}
& k=2 \\
& m(1, 2) + m(3, 4) + p_0 p_2 p_4 \\
& 120 + 252 + 4*3*7 = 456
\end{aligned}$$

$$\begin{aligned}
& k=3 \\
& m(1, 3) + m(4, 4) + p_1 p_3 p_4 \\
& 264 + 0 + 4*12*7 = 600
\end{aligned}$$

(A₁ A₂) (A₃ A₄)
Question 4)

Question 5)

$$\begin{aligned}w[1, 1] &= w[1, 0] + p_1 + q_1 = 0.07 + 0.05 + 0.07 = 0.19 \\w[3, 2] &= w[3, 1] + p_2 + q_2 = 0.26 + 0.30 + 0.06 = 0.62 \\w[4, 4] &= w[4, 3] + p_4 + q_4 = 0.06 + 0.20 + 0.06 = 0.32\end{aligned}$$

w	1	2	3	4	5
4	1.00	0.88	0.69	0.32	0.06
3	0.74	0.62	0.43	0.06	
2	0.52	0.26	0.07		
1	0.19	0.07			
0	0.07				

r is from i to j

save the lowest r to root table and record the lowest value in the c table

$$c[i, j] = c[i, r-1] + c[r+1, j] + w[i, j]$$

r = 1

$$c[1, 1] = c[1, 0] + c[2, 1] + w[1, 1] = 0.07 + 0.07 + 0.19 = 0.33$$

r = 1

$$c[1, 2] = c[1, 0] + c[2, 2] + w[1, 2] = 0.07 + 0.40 + 0.52 = 0.99$$

r = 2

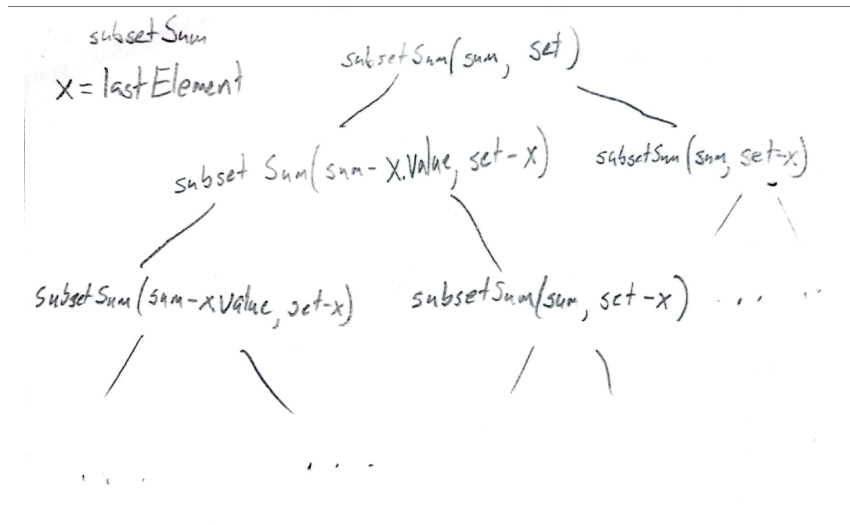
$$c[1, 2] = c[1, 1] + c[3, 2] + w[1, 2] = 0.33 + 0.07 + 0.52 = 0.92$$

c	1	2	3	4	5
4	2.36	1.76	1.20	0.44	0.06
3	1.63	1.08	0.56	0.06	
2	0.92	0.40	0.07		
1	0.33	0.07			
0	0.07				

root	1	2	3	4
4	3	3	3	4
3	2	3	3	
2	2	2		
1	1			

Bonus:

optimal substructure the solution either contains the number or does not contain the number



Pseudo code:

```

sumOfSubset(set, n)
listOfSums = [0]
for i in set:
    for j in listOfSums:
        if i + j is not in listOfSums:
            listOfSums.addFront(i+j)
return is n in listOfSums

```