

**Topics of the week**

1. Pseudo-inverse, definition and properties;
2. Pseudo-inverse and minimum norm solution;
3. Pseudo-inverse and projection;
4. Polyhedron, projections of sets, Farkas lemma.

**Left inverse** of  $A : U \rightarrow V$  is the matrix  $A^\dagger$  s.t.  $A^\dagger A = I$ .

**Interpretation:** Given  $y \in \text{im } A$ , find the **unique** pre-image  $x = A^\dagger y$ .

**Criterion:** There is  $A^\dagger \iff A$  is injective  $\iff$  columns of  $A$  are linearly independent.

**Reason:** If it's not injective, the pre-image of  $y = Ax$  is non-unique.

**Description:**  $(AA^\dagger)^2 = AA^\dagger$  is a projection on  $\text{im } A$ .

**Left pseudoinverse** is  $A^\dagger = (A^\top A)^{-1} A^\top$ . For each  $y$ , finds  $x$  that minimizes  $\|Ax - y\|$ .

**Extra:**  $AA^\dagger$  is the orthogonal projection on  $\text{im } A$ .

**Right inverse** of  $A : U \rightarrow V$  is the matrix  $A^\dagger$  s.t.  $AA^\dagger = I$ .

**Interpretation:** Given  $y \in V$ , find **any** pre-image  $x = A^\dagger y$ .

**Criterion:** There is  $A^\dagger \iff A$  is surjective  $\iff$  rows of  $A$  are linearly independent.

**Reason:** If it's not surjective, there are  $y$  that without pre-image.

**Description:**  $(A^\dagger A)^2 = A^\dagger A$  is a projection on a "representative" subspace  $\text{im } A^\dagger \subset U$ .

**Right pseudoinverse** is  $A^\dagger = A^\top (AA^\top)^{-1}$ . For each  $y$ , finds  $x$  s.t.  $Ax = y$  and  $\|x\| \rightarrow \min$ .

**Extra:**  $A^\dagger A$  is the orthogonal projection on  $(\ker A)^\perp = \text{im } A^\top$ .

**Pseudoinverse** of a matrix  $A = CR$  is  $A^\dagger = R^\dagger C^\dagger = R^\top (C^\top AR^\top)^{-1} C^\top$ . Solves the problem:

$$\begin{aligned} \|x\| &\rightarrow \min, \\ \text{s.t. } \|Ax - y\| &\rightarrow \min. \end{aligned}$$

**Implied equations** Let  $Ax = b$ , then  $(y^\top A)x = y^\top b$  for any  $y$ .

**Fredholm theorem** For a matrix  $A$  and a vector  $b$ , exactly one of the following holds:

1. There is  $x$  s.t.  $Ax = b$ ;
2. There is  $y$  s.t.  $y^\top A = 0$  and  $y^\top b \neq 0$ .

**Interpretation:**  $Ax = b$  is infeasible  $\iff$  there is an infeasible implied equation.

**Proof:** If there is such  $x$ , it means  $y^\top b = y^\top (b - Ax) = 0$  for any  $y$  s.t.  $y^\top A = 0$ . If there is no such  $x$ , use Gaussian elimination to find  $y$ .

**Variant:** if there is no such  $x$ , use  $y = Ax - b$ , where  $\|Ax - b\| \rightarrow \min$ .

**Conic combination** of  $v_1, \dots, v_n$  is  $x_1 v_1 + \dots + x_n v_n$ , where  $x_1, \dots, x_n \geq 0$ .

**Implied inequalities** Let  $Ax \leq b$  and  $y \geq 0$ , then  $y^\top Ax \leq y^\top b$ .

**Conic combinations:** When  $x \geq 0$ , the linear combination  $x_1 v_1 + \dots + x_n v_n$  is called **conic**.

**Fourier-Motzkin elimination** Consider a set of inequalities  $Ax \leq b$ . We can eliminate  $x_n$ :

$$\begin{cases} x_n \geq b_i - a_i^\top x', \\ x_n \leq b_j - a_j^\top x' \end{cases} \implies b_j - a_j^\top x' \leq x_n \leq b_i - a_i^\top x'$$

Add an inequality without  $x_n$  for each such  $(i, j)$ . If there is a solution  $x = (x', x_n)$ , then it also satisfies the new system. Correspondingly, if there is a solution  $x'$  of the new system, it's possible to find  $x_n$  s.t.  $x = (x', x_n)$  is the solution to the old system.

**Note:** Each new inequality is an implied inequality of the previous ones.

**Note:** Very inefficient in practice, compared to standard linear programming solvers.

**Farkas lemma** For a matrix  $A$  and a vector  $b$ , exactly one of the following holds:

1. There is  $x$  s.t.  $Ax \leq b$ ;
2. There is  $y \geq 0$  s.t.  $y^\top A = 0$  and  $y^\top b < 0$ .

**Interpretation:**  $Ax \leq b$  is infeasible  $\iff$  there is an infeasible implied inequality.

**Proof:** If there is such  $x$ , it means  $y^\top b = y^\top (b - Ax) \geq 0$  for any  $y \geq 0$  s.t.  $y^\top A = 0$ . If there is no such  $x$ , use Fourier-Motzkin elimination to find  $y$ .

**Variant:** If there is no such  $x$ , use  $y = Ax - b'$ , where  $\|Ax - b'\| \rightarrow \min$  s.t.  $b' \leq b$ .

**In-class exercises** Let  $A \in \mathbb{R}^{m \times n}$  and  $B \in \mathbb{R}^{n \times p}$  be arbitrary matrices.

1. Prove that if  $\text{rank } A = \text{rank } B = n$ , then  $(AB)^\dagger = B^\dagger A^\dagger$ .
2. Prove that  $A^\dagger A A^\dagger = A^\dagger$ .
3. Prove that  $(A^\top)^\dagger = (A^\dagger)^\top$ .
4. Prove that  $A^\dagger A$  is symmetric and is the projection matrix for  $\text{im } A^\top$ .