Unit 1 Introduction

Unit Outcomes. Here you will learn

- to define a distributed system (DS) and describe and justify the characteristics of all DSs
- to classify DSs by various important aspects, appreciate the challenges in accurately specifying the behaviour of a DS
- about some common ways to distribute responsibility among networked nodes
- about network abstractions avialable to DS developers
- about basic tools for deploying and managing DSs

Further Reading: CDK2005 1, 2.3, 3, 4

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Definition of Distributed System

A Distributed System (DS) is

- a software system
 - deployed on networked multiple computers (nodes)
 - making these nodes cooperate towards a shared goal
 - nodes communicate only by exchanging messages

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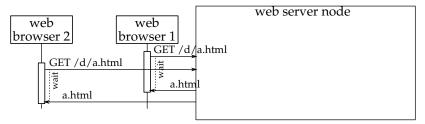
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Some common DSs

- phone networks (phones and exchanges)
- email (email clients, SMTP servers and mail-box servers)
- World-Wide Web (web servers and browsers)
- Internet search engines (high-performance clusters)
- enterprise systems (database servers and clients)

Characteristics of a DS Concurrency

 messages arrive concurrently, are processed concurrently example:



- multi-threading inside a node usually necessary
- access to internal data must be synchronised

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Independent failures

- failures occur in components of DS
- how to communicate failures?
 - when node dies completely
 - only by lack of response
 - when network or its part fails silently
 - slow or broken?
 - nodes may be stranded, DS fragmented

No global clock

- impossible to 100% synchronise clocks among nodes
- approximate synchronisation using Cristian protocol:
 - estimating message delays:

using radio or satellite signals:

• more generally: a node cannot fully know others' state

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Classification of DSs Purpose

- information transfer over distance
- resource sharing
- performance enhancement
 - processing of large data
 - complex scientific computation
- protection against node failure by duplicating resources

Scalability

- how does a DS cope with increasing load?
- load = eg count of nodes, requests, users, ...
- perfectly scalable DS:
 - can be adapted to take any load
 - without modifying protocols or sw
 - cost of hw grows proportionally: O(n)
 - decrease in performance logarithmic: $O(\log(n))$ time increase for basic ops

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Security — attack techniques

- typical techniques used in attacks:
 - deploy a rogue node to the DS
 - spy on messages in the DS
 - modify messages in the DS
 - flood the DS with legitimate but meaningless messages
 - send malicious mobile code

Security — attacks

- security = level of resistance to attacks
- attack = attempt to make a DS behave in unintended ways
- typical goals of attacks:
 - unauthorised access to information
 - unauthorised alteration of data
 - total system break-down

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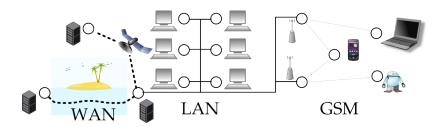
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Security — levels and solutions

- various levels of security in DS by combinations of:
 - relying on private network
 - authenticating nodes using central authority
 - secure channels no understanding without a key
 - combating isolated internal attacks (eg by voting) techniques)
 - · analysing any incoming mobile code
- beware: for some attacks no known cure

Heterogeneity support

- ? is a DS deployable through a variety of:
 - network types
 - hardware platforms + operating systems



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Transparency (1/2)

- transparency = certain aspect is not detectable (a tangible measure of abstraction)
- what to hide?
 - access tr.: local or remote?
 - location tr.: location in the net
 - mobility tr.: is it in a fixed location or moving?
 - failure tr.: failures of node hw/sw or communication hw/sw

Openness

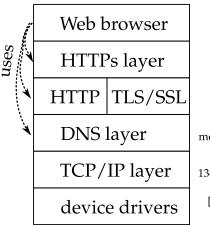
- how easy is it to extended and adapt a DS?
- a DS is open to the extreme if
 - all its protocols and APIs are public and well-documented
 - it provides well-documented access methods to all its resources
 - it is well-designed for extensibility, eg:
 - good decomposition (low coupling, cohesion, layers of abstraction)
 - no unnecessary complexity
 - it supports heterogeneity of networks, OSs and hardware
 - it is extensible in any common programming language

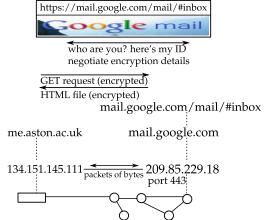
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Transparency (2/2)

- what to hide?
 - performance tr.: overall load of DS
 - scaling tr.: size and power of DS
 - replication tr.: using multiple copies of a single resource
 - concurrency tr.: doing something concurrently

Network abstraction Stacking abstraction layers





network = nodes + links + communication facilities

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Internet protocols for DS developers

- TCP provides logical channels between pairs of remote nodes
 - as if they were directly linked
 - can stream sequences of bytes of any length either way
 - hiding IP packets from users
- establishing a TCP channel:
 - one computer listens, the other initiates
 - multiple listeners → listening on ports 1-
- DNS lesser-known facts:
 - several DNS names can share one IP
 - one DNS name can refer to multiple computers
 - DNS-IP mapping sensitive to location

Role of protocols

- protocol is a specification of
 - message formats
 - valid sequences of messages
- that allow multiple nodes cooperate to achieve a goal
- protocol usually defined in terms of lower-level protocol(s) layers of protocols
- abstract networks defined by a set of protocols
- eg Internet =
 - TCP/IP: reliable messaging + numerical addressing of nodes
 - DNS: hierarchical and human-friendly naming of nodes

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Web layer — resources and HTTP

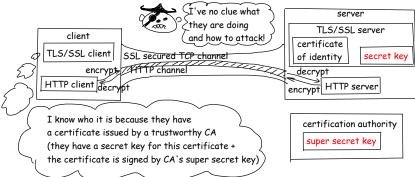
- resources = passive or active entities accessible over Internet
- can be viewed as logical nodes/sub-nodes for a DS
- URL resource locator http://a.com/d.pdf?page=1#fig1

scheme, authority (DNS name + TCP port), path, query, fragment

- URI = URL or URN, URN has no DNS or TCP information, only a name, eg urn:isbn:0-486-27557-4
- HTTP: a request-response protocol for manipulating resources
- most common requests and their purpose:
 - GET: obtain a remote resource
 - POST: initiate operation on a remote resource
 - PUT: store a local resource at a remote location
 - DELETE: permanently remove a remote resource

Web layer — HTTPS

HTTP over a secure channel:



protocol based on HTTP, TLS/SSL, TCP/IP, DNS

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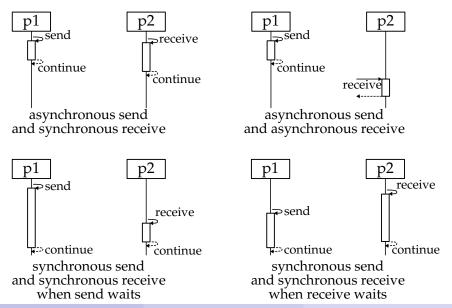
Network abstraction Middleware — network abstraction

- characterised by
 - addressing logical nodes, not physical nodes
 - eg processes, resources, objects
 - results in location & access transparency
 - convenient messaging, transporting typed structured data
 - types of data declared
 - → greater openness and PL heterogeneity
 - synchronisation of communicating nodes

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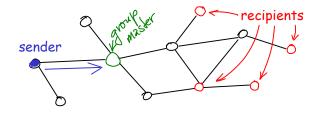
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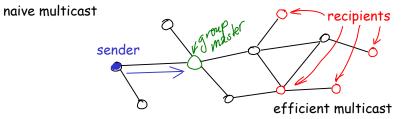
Synchronisation of communicating nodes



Multicast

- most efficient when implemented at TCP/IP level
- sometimes available at middleware level





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Remote administration Why remote administration?

- install/update/configure/deploy a node's software
- read logs local to a node to diagnose a problem
- tweak some data files to fix a problem or simulate a problem
- monitor a node's load to tune a DS algorithm

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Remote execution of GUI applications

- X Windows
 - thin-client: very efficient remote graphical clients
 - can have secure connection via SSH.
 - not for MS Windows applications, only X-Windows applications
- Remote desktop
 - desktop completely taken over by remote computer
 - powerful but slow unless on very fast network
 - works on all OS, common in MS Windows administration

Remote shell

- command line is simple but powerful easy remote operation
- SSH: security achieved via secured TCP channels
- demo: install prg, view load, search files, rename files
- advantages/disadvantages?

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Learning Outcomes

Learning Outcomes. You should now be able to

- define a DS
- name, describe and justify DS characteristics
- name and define important DS classifying aspects
- evaluate a given well-known DS against a set of classifying aspects
- list several levels of network abstraction behind a typical DS; for each level identify the key concepts that are relevant to a DS developer
- define and correctly use terms: TCP port, TCP channel, URL, secure channel, multicast
- explain difference between synchronous and asynchronous send
- give the main idea of an efficient multicast implementation
- start remote programs, copy files using SSH