

# HW1

● Graded

20 Hours, 34 Minutes Late

## Student

Adam Fenjiro

## Total Points

59 / 80 pts

### Question 1

Q1

8 / 10 pts

– 0 pts Correct

– 2 pts Incorrect rules and stack changes after step 3

### Question 2

Q2

21 / 30 pts

2.1

a

8 / 10 pts

– 0 pts Correct

– 2 pts Missing closure step in CFMS. The given grammar is SLR.

2.2

b

6 / 10 pts

– 0 pts Correct

– 4 pts Incomplete CFMS and missing first/follow sets

2.3

c

7 / 10 pts

– 0 pts Correct

– 3 pts Incomplete CFMS and parser table

### Question 3

Q3

30 / 40 pts

3.1

a

2 / 5 pts

– 0 pts Correct

– 3 pts Point adjustment

3.2

b

3 / 5 pts

– 0 pts Correct

– 2 pts Incomplete CFSM

3.3

c

4 / 5 pts

– 0 pts Correct

– 1 pt Point adjustment

3.4

d

2 / 5 pts

– 0 pts Correct

– 3 pts Point adjustment

3.5

e

10 / 10 pts

✓ – 0 pts Correct

3.6

f

8 / 8 pts

✓ – 0 pts Correct

3.7

g

1 / 2 pts

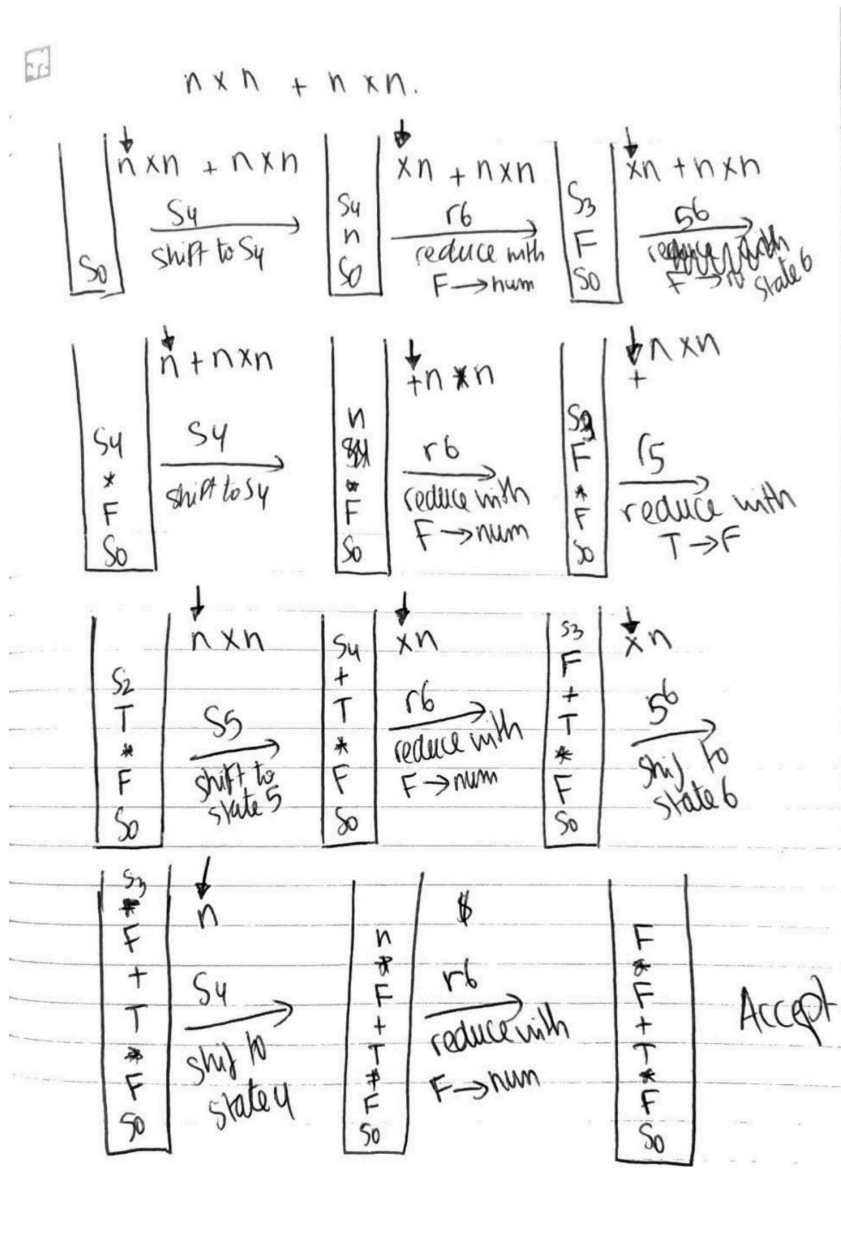
– 0 pts Correct

– 1 pt Point adjustment

Question assigned to the following page: [1](#)

# CS4121 - Spring25 - Adam Fenjiro - HW1

## Problem 1:



Question assigned to the following page: [2.1](#)

**Problem 2:****a)  $S \rightarrow abS$** **| ab****1. CFSM****State 0:** $S' \rightarrow \bullet S$  $S \rightarrow \bullet abS$  $S \rightarrow \bullet ab$ **State 1:** $S' \rightarrow S \bullet$ **State 2:** $S \rightarrow a \bullet bS$ **State 3:** $S \rightarrow ab \bullet S$ **State 4:** $S \rightarrow abS \bullet$ **State 5:** $S \rightarrow ab \bullet$ **2. SLR Parse Table**

State	a	b	\$	S
S0	s2	s5		1
S1			acc	
S2		s3		
S3	s2	s5		4
S4			r2	
S5			r3	

**3. SLR or not**

=> The grammar is not SLR because state 5 has a shift-reduce conflict. When the input is b the parser does not know whether to shift or reduce leading to ambiguity.

Question assigned to the following page: [2.2](#)

b)  $S \rightarrow AaAb$   
     $\quad \mid BbBa$

$A \rightarrow c$

$B \rightarrow c$

1. CFSM

**State 0:**

$S' \rightarrow \bullet S$

$S \rightarrow \bullet AaAb$

$S \rightarrow \bullet BbBa$

$A \rightarrow \bullet c$

$B \rightarrow \bullet c$

**State 1:**

$S' \rightarrow S \bullet$

**State 2:**

$S \rightarrow A \bullet aAb$

$A \rightarrow c \bullet$

**State 3:**

$S \rightarrow B \bullet bBa$

$B \rightarrow c \bullet$

**State 4:**

$S \rightarrow Aa \bullet Ab$

$S \rightarrow \bullet AaAb$

$S \rightarrow \bullet BbBa$

$A \rightarrow \bullet c$

$B \rightarrow \bullet c$

**State 5:**

$S \rightarrow Bb \bullet Ba$

$S \rightarrow \bullet AaAb$

$S \rightarrow \bullet BbBa$

$A \rightarrow \bullet c$

$B \rightarrow \bullet c$

**State 5:**

$S \rightarrow AaAb \bullet$

$S \rightarrow BbBa \bullet$



Question assigned to the following page: [2.2](#)

## 2. SLR Parse Table

State	a	b	c	\$	S	A	B
S0			S2		1	2	3
S1				acc			
S2	S4						
S3		S5					
S4			S2			6	7
S5			S3			8	9
S6				r1			

## 3. SLR or not

=> The grammar is not SLR. This is because there is a shift-reduce conflict in state 4 and 5.  
When the input is c the parser does not know whether to shift or reduce.

Question assigned to the following page: [2.3](#)

c)  $S \rightarrow ASB \mid ab$

$A \rightarrow a$

$B \rightarrow b$

1. CFSM

**State 0:**

$S' \rightarrow \cdot S$

$S \rightarrow \cdot ASB$

$S \rightarrow \cdot ab$

$A \rightarrow \cdot a$

$B \rightarrow \cdot b$

**State 1:**

$S' \rightarrow S \cdot$

**State 2:**

$S \rightarrow A \cdot SB$

$A \rightarrow a \cdot$

**State 3:**

$S \rightarrow B \cdot SB$

$B \rightarrow b \cdot$

**State 4:**

$S \rightarrow AS \cdot B$

$S \rightarrow \cdot ASB$

$S \rightarrow \cdot ab$

$A \rightarrow \cdot a$

$B \rightarrow \cdot b$

**State 5:**

$S \rightarrow ab \cdot$

2. SLR Parse Table

State	a	b	\$	S	A	B
S0	S2	S3		1	2	3
S1			acc			
S2		S4				
S3	S5					
S4			r2			
S5			r3			

3. SLR or not

=> The grammar is SLR because there are no conflicts in the parse table. The parsing table correctly differentiates between shift and reduce actions, ensuring smooth parsing.

Questions assigned to the following page: [3.1](#) and [3.2](#)

### Problem 3:

a.

$S' \rightarrow E$

$E \rightarrow +EE$

$E \rightarrow + (+EE) E$

$E \rightarrow + (+ (\sim E) E) E$

$E \rightarrow + (+ (\sim 5) 12) cy$

b.

State 0:

$S' \rightarrow \bullet E$

$E \rightarrow \bullet + EE$

$E \rightarrow \bullet \sim E$

$E \rightarrow \bullet F$

$F \rightarrow \bullet \text{num } 10$

$F \rightarrow \bullet \text{num } 26$

State 1:

$S' \rightarrow E \bullet$

$E \rightarrow E \bullet + E$

$E \rightarrow E \bullet \sim E$

$E \rightarrow E \bullet F$

State 2:

$E \rightarrow + \bullet EE$

$E \rightarrow + \bullet EE$

$E \rightarrow + \bullet \sim E$

$E \rightarrow + \bullet F$

$F \rightarrow \bullet \text{num } 10$

$F \rightarrow \bullet \text{num } 26$

State 3:

$E \rightarrow \sim \bullet E$

$E \rightarrow + \bullet EE$

$E \rightarrow + \bullet \sim E$

$E \rightarrow + \bullet F$

$F \rightarrow \bullet \text{num } 10$

$F \rightarrow \bullet \text{num } 26$

Questions assigned to the following page: [3.2](#), [3.3](#), and [3.4](#)

State 4:

$E \rightarrow F \bullet$

$E \rightarrow F \bullet + E$

$E \rightarrow F \bullet \sim E$

$E \rightarrow F \bullet F$

State 5:

$F \rightarrow \text{num } 10 \bullet$

State 6:

$F \rightarrow \text{num } 26 \bullet$

c.

**First Sets**

$\text{First}(E) = \text{First}(F) = \{ +, \sim, \text{num}_{10}, \text{num}_{26} \}$

**Follow Sets**

$\text{Follow}(S') = \{ \$ \}$

$\text{Follow}(E) = \text{Follow}(F) = \{ +, \sim, \text{num}_{10}, \text{num}_{26}, \$ \}$

d.

State	+	~	num10	num26	\$	E	F
S0	S2	S3	S4	S4		1	5
S1					acc		
S2			S6	S6		7	5
S3	S3	S4	S4			8	5
S4							
S5					r4		
S6			S6	S6			
S7					r2		
S8					r3		



Question assigned to the following page: [3.5](#)

e.

#### Flex Rules:

```
C/C++
%{
#include "y.tab.h"
%}

[0-9]+      { yylval.num = atoi(yytext); return NUM10; }
[a-z]+      { yylval.num = base26(yytext); return NUM26; }
"+"         { return '+'; }
"~"         { return '~'; }

%%
```

#### Bison Rules:

```
C/C++
%token NUM10 NUM26
%union { int num; }

%%
E : '+' E E { $$ = $2 + $3; }
  | '~' E   { $$ = -$2; }
  | F       { $$ = $1; }
F : NUM10   { $$ = $1; }
  | NUM26   { $$ = $1; }

%%
```

#### Base26 conversion in c

```
C/C++
int base26(char* str) {
    int result = 0;
    while (*str) {
        result = result * 26 + (*str - 'a');
        str++;
    }
    return result;
}
```

Questions assigned to the following page: [3.6](#) and [3.7](#)

(f) Here is the parse process of the input ++~512cy:

Start by pushing S' into the stack

Read '+' first symbol from the input

Shift '+' then go to state 2

Read '+' next symbol from the input

Shift '+' and go to state 2

Read '~' next symbol from the input

Shift '~' and go to state 3

Read '5' next symbol from the input

Shift '5' and go to state 5

Reduce by rule  $F \rightarrow \text{num } 10$  and go to state 4

Reduce by rule  $E \rightarrow F$  and go to state 1

Reduce by rule  $E \rightarrow \sim E$  and go to state 2

Read 'c' next symbol from the input

Shift 'c' and go to state 6

Reduce by rule  $F \rightarrow \text{num } 26$  and go to state 4

Reduce by rule  $E \rightarrow F$  go to state 1

Reduce by rule  $E \rightarrow +EE$  and go to state 2.

Reduce by rule  $E \rightarrow +EE$  and go to state 1.

Finally accept the input

g) They are both related since both of these find the same sequence of grammar rules to parse and derive the expression. The parsing process uses a stack and a parse table to systematically apply the grammar rules.