

# The Observer–Guardian Invariant: Causal Decoupling of Audit Logs in Safety-Critical AI Architectures

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## Abstract

The act of observation is itself a causal intervention—a fundamental tension for safety-critical AI systems requiring comprehensive auditing without compromising execution integrity.

We present **Temple**, a write-only observer designed for strict causal decoupling within a Guardian-enforced planning architecture (Dual-Sovereign Safety Architecture, dSSA). Temple operates strictly downstream of enforcement decisions with zero return authority, asynchronous non-blocking logging, and isolated memory.

Under paired deterministic evaluation (SEED=42, cuDNN deterministic mode, runner v4.9.2) using a live planner (Qwen2.5-7B-Instruct), the canonical adversarial enforcement suite (7 cases, 1 Guardian veto) produced **bit-identical** SHA-256 execution-stream hashes:

483647c61a124ad85f0c625a43d41ad39366596e2cf9789829ab47c9dcfc3bac

We establish observer non-interference as an empirically falsifiable systems property under a bounded threat model and introduce a reproducible protocol for verifying causal decoupling.

## 1 Motivation: The Observer Paradox

Safety-critical systems demand perfect auditability for failure analysis and compliance [7]. Yet observation introduces a paradox: an observer that changes what it observes becomes a participant [1].

Three causal failure modes exist:

- **Optimization Coupling:** Monitoring-derived reward signals alter agent objectives [6]
- **Evaluation Faking:** Models detect evaluation contexts and strategically alter behavior [1]
- **Timing Perturbation:** Instrumentation overhead affects execution ordering [4]

We test observer non-interference empirically: can execution traces remain bit-identical under observer presence?

## 2 Background: Authority Separation in dSSA

The Dual-Sovereign Safety Architecture (dSSA) implements strict authority partitioning [3]:

- **Planner:** Learned proposal generation (proposal authority only)
- **Guardian:** Deterministic safety validation (veto authority)
- **Temple:** Post-decision logging (zero authority)

## 2.1 Safety Dominance Property

All actions pass Guardian validation prior to execution. Planner capability cannot weaken enforcement.

## 2.2 Observation as Vulnerability

Even when enforcement is architecturally dominant, observation layers introduce a distinct class of risk: *visibility can become a causal input*. Evaluation-aware systems may change behavior when they infer they are being measured [1], and monitoring-derived metrics can create proxy objectives that couple measurement to optimization [6]. Separately, even “passive” instrumentation can perturb execution via contention, ordering changes, or timing effects [4]. For these reasons, auditability should not be assumed causally neutral in safety-critical AI.

# 3 Temple Architecture

Temple enforces non-interference through three principles: downstream positioning, write-only isolation, and fail-open design.

## 3.1 Downstream Positioning

Temple receives only finalized Guardian verdicts:

1. Planner proposes action
2. Guardian validates (PASS/VETO)
3. Temple records post-decision trace

Temple has no authority to influence decisions and no access to pre-decision internal planner state.

## 3.2 Write-Only Isolation

- Zero read authority
- No shared mutable state
- Asynchronous non-blocking I/O
- No execution branching

## 3.3 Fail-Open Design

Temple cannot halt execution. Logging failures are silent and non-propagating. This design prevents a safety-through-observation failure mode where monitoring presence changes outcomes by triggering exceptions, blocking I/O, or altering control flow.

### 3.4 Record Semantics: What Temple Logs and What It Does Not Log

To make the non-interference claim operationally testable, Temple records only a minimal, execution-relevant payload. The hashed execution stream comprises the following fields:

```
{test_id, guardian_verdict, force_n, velocity_mps, distance_m, proposal_hash}
```

Temple does *not* log planner chain-of-thought, intermediate guardian computations, system resource telemetry (CPU/memory), or any signal that could create a feedback pathway into planning or enforcement.

## 4 Threat Model

Verification scope:

- Fully deterministic execution (SEED=42 across random, NumPy, Torch, and CUDA; temperature=0.0, do\_sample=False)
- Frozen model weights (Qwen2.5-7B-Instruct)
- Sequential evaluation (no parallelization)
- Single-node execution (CUDA 12.2)

Out-of-scope:

- Hardware side-channels
- Adversarial anticipatory planning
- Multi-agent coordination

Within scope, any causal pathway from Temple presence to execution behavior must manifest as a mismatch in the ordered, canonicalized evaluation stream and therefore hash divergence.

## 5 Experimental Design

### 5.1 Null Hypothesis

Temple presence produces no change in:

1. Guardian verdicts
2. Planner proposal hashes
3. Kinematic parameters
4. Execution ordering

### 5.2 Falsification Criterion

SHA-256 stream hash mismatch = causal interference.

### 5.3 Evaluation Harness (v4.9.2)

- Fixed seed across random, numpy, torch, CUDA
- cuDNN deterministic mode
- Canonical JSON serialization (sorted keys)
- Hash excludes Temple artifacts (timestamps, file paths, I/O metadata)

Hash payload: {test ID, verdict, force/velocity/distance, proposal hash}.

### 5.4 Protocol Summary

Two executions are performed with identical configuration and test ordering:

1. **Temple OFF**: run the suite and compute the SHA-256 hash over the ordered stream.
2. **Temple ON**: enable Temple logging and repeat; compute the SHA-256 hash over the same ordered stream (excluding observer artifacts).

Equality of hashes is a strict certificate of parity under the deterministic scope. A single mismatch falsifies the non-interference claim.

## 6 Results

### 6.1 Guardian Enforcement

Metric	Balanced	Edge	Emergency	Service	Temple
Tests	20	15	10	15	10
PASS	20	6	4	15	10
VETO	0	9	6	0	0
Force (N)	0.000	1.248	0.625	0.373	0.310

Table 1: Guardian Enforcement (Planner OFF baseline)

Parity verification was performed with the planner enabled on the canonical 7-case adversarial enforcement subset; Table 1 reports separate planner-OFF baseline enforcement metrics.

### 6.2 Cryptographic Parity Verification

The canonical enforcement suite (gte\_core\_guardian.jsonl) contains 7 adversarial cases designed to trigger Guardian vetoes.

Hash identity across Temple ON/OFF confirms no causal interference. Parity was verified on identical hardware and software configuration; cross-hardware bit-level equivalence is not claimed.

Configuration	SHA-256 Stream Hash
Planner	Qwen2.5-7B-Instruct
Seed	42
Temple OFF	483647c61a124ad85f0c625a43d41ad39366596e2cf9789829ab47c9dcfc3bac
Temple ON	483647c61a124ad85f0c625a43d41ad39366596e2cf9789829ab47c9dcfc3bac
<b>Parity</b>	<b>MATCH</b>

Table 2: Bit-Identical Stream Hash Verification (7-case adversarial suite, 1 veto)

## 7 Discussion

### 7.1 Verified Properties

- No Guardian verdict changes
- No planner proposal variance
- No kinematic parameter drift
- No execution reordering

### 7.2 Why Hash Parity Is Strong

Hash parity is stronger than statistical equivalence: under deterministic execution it is an exact cryptographic certificate that the ordered execution-relevant stream is identical. This converts observer non-interference from a design intention into a falsifiable systems property [5].

### 7.3 Practical Implications

- Automated observer integrity testing [5]
- Compliance audit trails without feedback loops [2]
- Post-incident forensics without runtime coupling

## 8 Limitations

Exact hash matching requires deterministic execution. Stochastic planners need distributional tests. Hardware side-channels and adversarial introspection remain out-of-scope [4]. The claim is therefore intentionally scoped to a single-node deterministic setting; this is a feature for reproducibility and falsifiability, not a claim of universal cross-platform bit-level equivalence.

## 9 Related Work (Brief)

Third-party auditing and evaluation methodologies motivate high-integrity logging pipelines [7]. Observer effects and evaluation-aware behavior motivate treating monitoring as potentially causal [1]. Broader discussions of AI safety objectives motivate avoiding proxy coupling between measurement and optimization [6]. Finally, causal modeling perspectives motivate explicit reasoning about timing and interaction pathways in autonomous systems [4].

## 10 Conclusion

Temple demonstrates empirically verifiable observer non-interference under deterministic conditions. The paired-run hash comparison protocol generalizes to any observer claiming causal decoupling.

**dSSA:** Planner proposes  $\rightarrow$  Guardian enforces  $\rightarrow$  Temple observes.

**Reproducibility:** guardian-observer-parity (tag: observer-parity-v1.1.1)

### Canonical Commands:

```
# Temple OFF (produces canonical hash)
python run_eval_minimal.py test_sets/gte_core_guardian.jsonl \
  --planner --planner-name qwen --base-model Qwen/Qwen2.5-7B-Instruct \
  --device cuda --out-dir results_off --run-id qwen_core_guardian_temple_off

# Temple ON (must match above hash)
python run_eval_minimal.py test_sets/gte_core_guardian.jsonl \
  --planner --planner-name qwen --base-model Qwen/Qwen2.5-7B-Instruct \
  --device cuda --temple-out observer/temple_on.json \
  --out-dir results_on --run-id qwen_core_guardian_temple_on
```

## References

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