

CS-E5740 Complex Networks, Answers to exercise set 1

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October 1, 2018

Acknowledgements: Lukas Haug, Bianca Lachennmann, Christoph Berger,
some

1 Basic network properties

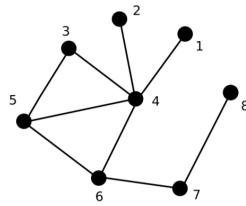


Figure 1: The graph for exercise 1.

a. Adjacency matrix

A network data structure in the form of a matrix used to represent a graph

$$a_{ij} = \begin{cases} 1 & \text{if } (j, i) \in E \\ 0 & \text{if } (j, i) \notin E \end{cases}$$

Hence, the adjacency matrix for the graph in figure 1 is as follows:asd

$$\begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

b. Edge density

The edge density of a network is the fraction of edges out of possible edges:

$$\rho = \frac{m}{\binom{N}{2}} = \frac{2m}{N(N-1)}$$

where m is the number of edges and N is the number of vertices.

Total possible edges is $\frac{N(N-1)}{2}$ which is $1 + 2 + 3 \dots + (N-2) + (N-1)$

The edge density ρ of graph 1 is $\frac{18}{56} = 0.321$

c. Degree and Degree Distribution

The degree k_i of vertex v_i is the number of edges it is incident to.

Vertices	Degree
1	1
2	1
3	2
4	5
5	3
6	3
7	2
8	1

The degree distribution $P(k)$ is the probability that the degree of a randomly chosen node is $2k$.

$$P(k) = \frac{N_k}{N}$$

where N_k is the number of nodes of degree k .

The degree distribution $P(k)$ of graph 1:

k	P(k)
1	0.375
2	0.25
3	0.25
4	0
5	0.125

d. The mean degree $\langle k \rangle$ of the graph

The average degree $\langle k \rangle$ of a network is:

$$\langle k \rangle = \sum_i \frac{k_i}{N} = \frac{2m}{N}$$

where m is the number of edges, N is the number of vertices, k_i is the degree for vertice i

The mean degree $\langle k \rangle$ for graph 1 is 2.25

e. The diameter d of the graph.

Diameter d is the largest distance in the network (It is the shortest distance between the two most distant nodes in the network.): $\max_{i,j \in V} d_{ij}$

Diameter of graph 1: 4

f. The clustering coefficient C_i

The clustering coefficient C_i for each node $i \in V$ that has degree $k_i > 1$. For nodes with $k_i = 0,1$, we define $C_i = 0$.

Clustering coefficient defined for node i as the fraction of edges between its neighbours out of possible edges between its neighbours:

$$C_i = \frac{E_i}{\binom{k_i}{2}} = \frac{2E_i}{k_i(k_i - 1)}$$

Average clustering coefficient (averaged over all nodes) is $c = \frac{1}{N} \sum_i c_i = \frac{2.2}{8} = 0.275$

Vertices	C_i
1	0
2	0
3	1
4	0.2
5	0.66
6	0.33
7	0
8	0

2 Computing network properties programmatically

a. Load the edge list and visualize the network

```
# python3
import networkx as nx

fn = 'karate_club_network_edge_file.edg'
graph = nx.read_weighted_edgelist(fn)

fig, ax = plt.subplots(figsize=(12, 8))
nx.draw(graph, with_labels=True)
ax.set_title('The Karate Club network')
```

Network:

b. Calculate the edge density of the karate club network.

First, write your own code without using `networkx.density` and then compare your result to the output of `networkx.density` (the corresponding NetworkX function).

```
m = graph.number_of_edges()
n = graph.number_of_nodes()
density = 2*m/(n*(n-1))
```

Edge Density: 0.13903743315508021

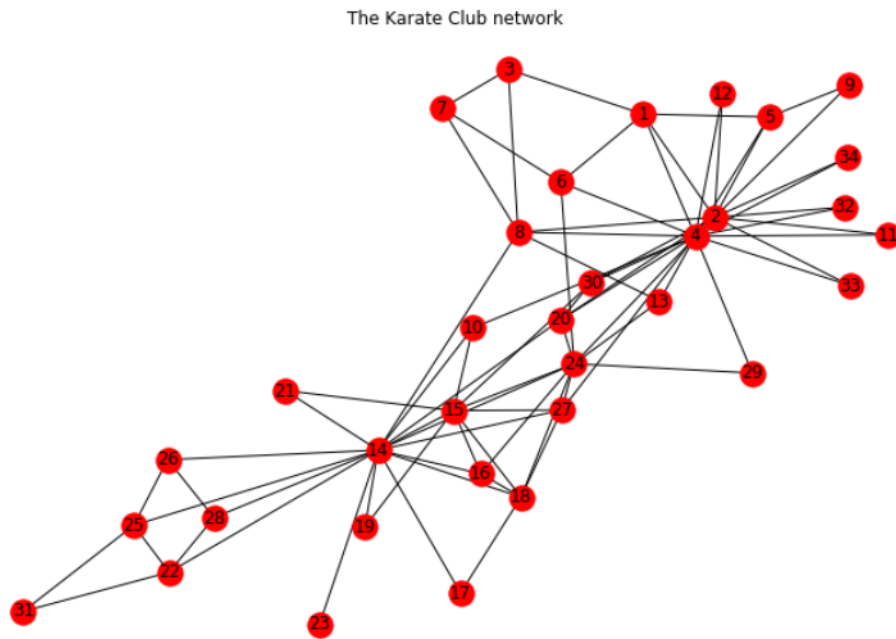


Figure 2: Karate Club Network graph

c. Calculate the average clustering coefficient

Calculate the average clustering coefficient with your own algorithm and compare it to the output of the corresponding NetworkX function.

```
def clustering_coefficient(node):
    n_links = 0
    neighbours = list(graph.neighbors(node))
    k = len(neighbours)
    if k < 2:
        return 0
    for neighbour_node in neighbours:
        for other_node in neighbours:
            if graph.has_edge(neighbour_node, other_node):
                n_links += 1

    n_links = n_links / 2 # double counting

    return 2 * n_links / (k * (k-1))

# apply function to every single node
```

```

all_nodes = list(graph.nodes)
node_ls = {}
for node in all_nodes:
    cc = clustering_coefficient(node)
    node_ls[node] = {"coefficient": cc}

sum([node_ls[node]["coefficient"] for node in node_ls])/ len(node_ls)
Average Clustering Coefficient: 0.5706384782076824

```

d. Calculate the Degree Distribution

Calculate the degree distribution $P(k)$ and the complementary cumulative degree distribution-CDF(k) of the network. Visualize the distributions using matplotlib.pyplot.

Degree	Degree Distribution (%)
1	0.029
2	0.324
3	0.176
4	0.176
5	0.088
6	0.059
9	0.029
10	0.029
12	0.029
16	0.029
17	0.029

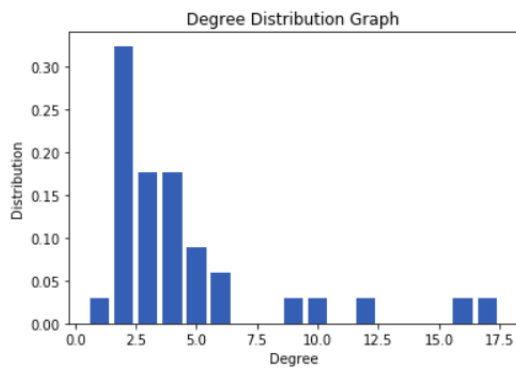


Figure 3: Degree Distribution graph

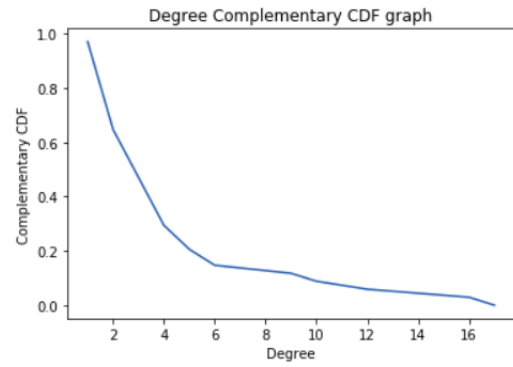


Figure 4: Degree Complementary CDF graph

e. Calculate the average shortest path length

```
length = nx.average_shortest_path_length(graph)
```

Average shortest path length: 2.408199643493761

f. Create a scatter plot of C_i as a function of k_i .

Using matplotlib.pyplot library

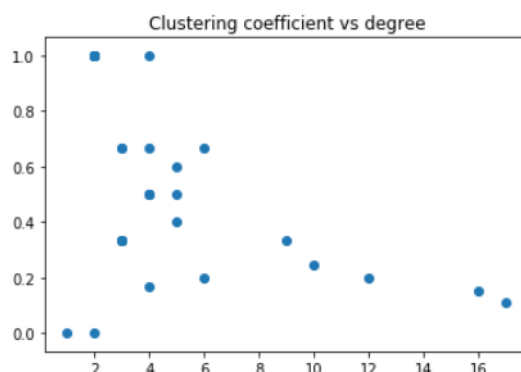


Figure 5: C_i as a function of k_i

3 Number of walks

a. Draw the induced subgraph G^*

where G^* is induced by the the vertices $V^* = \{14\}$ of network visualized in Figure 1. Calculate by hand the number of walks of length two between all node pairs $(i, j), i, j \in \{1...4\}$; a link can be travelled in both directions and the walk can visit a node multiple times. Remember to consider also walks, where $i = j$. Then, compute the matrix A^2 (you may do this also using a computer), where A is the adjacency matrix of the network G^* . Compare your results; what do you notice?

Node 1	Node 2	Node 3	Node 4
1-4-1	2-4-2	3-4-3	4-1-4
1-4-2	2-4-1	3-4-1	4-2-4
1-4-3	2-4-3	3-4-2	4-3-4

The adjacency matrix for the sub graph in figure 2 is as follows:

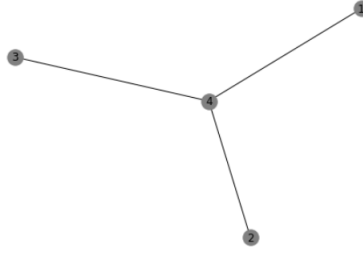


Figure 6: Induced subgraph G^*

$$A = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} \quad A^2 = \begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 3 \end{bmatrix}$$

A_{ij}^2 refers to the number of walks from i to j where the length = 2. E.g. $A_{11}^2 = 1$ where there is only 1 possible walk with length 2, from node 1 to 1 (1-4-1). For $A_{44}^2 = 3$ there are 3 such walks, as shown in the table above

b. Compute the number of walks of length 3 from node 3 to node 4

Compute the number of walks of length three from node 3 to node 4 in G^* . Then, starting from matrices A^2 and A , compute by hand the value of $(A^3)_{3,4}$ showing also the intermediate steps for computing the matrix element.

Walks
3-4-3-4
3-4-1-4
3-4-2-4

$$A^3 = \begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 3 \end{bmatrix} \cdot \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 & 3 \\ 0 & 0 & 0 & 3 \\ 0 & 0 & 0 & 3 \\ 3 & 3 & 3 & 0 \end{bmatrix}$$

From the matrix above, $(A^3)_{3,4} = 3$

c. Proving

Now, let's consider a general network with adjacency matrix A . **Show** that the element $(A^m)_{ij}$, $m \in \mathbb{N}$ corresponds to the number of walks of length m between nodes i and j .

Hint: Make use of mathematical induction: show that the statement holds for general m and prove that it holds for $m + 1$.

Base case

For $m = 1$ the statement holds, for A^1 is the adjacency matrix.

Inductive step

For a general m , first we assume $a_{ij}^{(m)}$ gives the number of walks of length m

For $m + 1$,

$$\begin{aligned}(A^{m+1})_{ij} &= (A^m \cdot A)_{ij} \\ &= \sum_{l=1}^N (a_{il}^m * a_{lj}) \\ &= a_{ij}^{m+1}\end{aligned}\tag{1}$$

Conclusion

Hence, for the matrix A^m , each element $(A^m)_{ij}$ corresponds to number of walks of length m between the nodes i and j .