

CSU44004-Formal Verification Assignment 2

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Question 1

- a) $\forall x(0 \leq x < |s| \Rightarrow \exists y \exists z(0 \leq y < z < |s'| \wedge s[x] = s'[y] \wedge s[x] = s'[z] \wedge y \neq z))$
- b) $\forall w(0 \leq w < |s| \Rightarrow \neg(\exists x \exists y \exists z(0 \leq x < y < z < |s'| \wedge s[w] = s'[x] \wedge s[w] = s'[y] \wedge s[w] = s'[z] \wedge x \neq y \wedge y \neq z \wedge x \neq z)))$
- c) $\forall x(0 \leq x < |s| \Rightarrow \exists y \exists z(0 \leq y < z < |s'| \wedge s[x] = s'[y] \wedge s[x] = s'[z] \wedge y \neq z \wedge \forall w(0 \leq w < |s'| \wedge y \neq w \wedge z \neq w \Rightarrow s[x] \neq s'[w])))$
- d) let a be the integer.
 $\forall x(\neg \exists y(0 \leq y < |s| \wedge s[y] = x) \Rightarrow \neg \exists z(0 \leq z < |s'| \wedge s'[z] = x))$

Question 2

Solution 2(a)

$$0 \leq lo < |s| \wedge \forall x (0 \leq x < |s| \Rightarrow s[lo] \geq s[x])$$

Solution 2(b)

Show $\vdash_{par} (\mid 0 < |s| \mid) \text{ findMax } (\mid T \mid)$.

Invariant = T

Variant = hi - lo

```
(\mid 0 < |s| \mid)
(\mid T \wedge 0 < |s| \mid) imp
lo := 0;
(\mid T \wedge lo < |s| \mid) asg
(\mid T \wedge lo \leq |s| - 1 \mid) imp
hi := |s| - 1 ;
(\mid T \wedge hi < |s| \mid) asg
(\mid T \wedge hi - lo \geq 0 \mid) imp
while ( lo < hi ) {
  (\mid T \wedge 0 \leq hi - lo = E_o \wedge lo < hi \mid) while
  (\mid 0 < hi - lo = E_o \mid) imp
  if (s[lo] \leq s[hi]) then {
    (\mid 0 < hi - lo = E_o \wedge s[lo] \leq s[hi] \mid) if-statement
    (\mid 0 \leq hi - (lo + 1) < E_o \mid) imp
    lo := lo + 1;
    (\mid 0 \leq hi - lo < E_o \mid) asg
  }
  else {
    (\mid 0 < hi - lo = E_o \wedge \neg(s[lo] \leq s[hi]) \mid) if-statement
    (\mid 0 \leq (hi - 1) - lo < E_o \mid) imp
    hi := hi - 1 ;
    (\mid 0 \leq hi - lo < E_o \mid) asg
  }
  (\mid 0 \leq hi - lo < E_o \mid) if-statement
  (\mid T \wedge 0 \leq hi - lo < E_o \mid) imp
}
(\mid T \wedge \neg(lo < hi) \mid) while
(\mid T \mid)
```

Solution 2(c)

Show $\vdash_{par} (\emptyset \ 0 < |s| \ \emptyset \ \text{findMax} \ (\emptyset \ \text{isMax}(s, lo))$.

Invariant = $\text{isMax}(s[..lo + 1] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi..], lo + 1)$

Variant = $hi - lo$

```

( $\emptyset \ 0 < |s| \ \emptyset$ )
( $\emptyset \ 0 < |s| \wedge (\text{isMax}(s[..0 + 1] + s[|s| - 1..], 0) \vee \text{isMax}(s[..0 + 1] + s[|s| - 1..], |s| - 1)) \ \emptyset \ \text{imp}$ 
lo := 0;
( $\emptyset \ lo < |s| \wedge (\text{isMax}(s[..lo + 1] + s[|s| - 1..], lo) \vee \text{isMax}(s[..lo + 1] + s[|s| - 1..], |s| - 1)) \ \emptyset \ \text{asg}$ 
hi := |s| - 1 ;
( $\emptyset \ lo \leq hi < |s| \wedge (\text{isMax}(s[..lo + 1] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi..], lo + 1)) \ \emptyset \ \text{asg}$ 
while ( lo < hi ) {
  ( $\emptyset \ lo \leq hi \wedge (\text{isMax}(s[..lo + 1] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi..], lo + 1)) \ \emptyset \ \text{while}$ 
  ( $\emptyset \ (\text{isMax}(s[..lo + 1] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi..], lo + 1)) \ \emptyset \ \text{imp}$ 
  if (s[lo] ≤ s[hi]) then {
    ( $\emptyset \ (\text{isMax}(s[..lo + 1] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi..], lo + 1)) \wedge s[lo] \leq s[hi] \ \emptyset$ 
    if-statement
    ( $\emptyset \ \text{isMax}(s[..lo + 1] + s[hi..], lo + 1) \ \emptyset \ \text{imp}$ 
    lo := lo + 1;
    ( $\emptyset \ \text{isMax}(s[..lo] + s[hi..], lo) \ \emptyset \ \text{asg}$ 
    }
  else {
    ( $\emptyset \ (\text{isMax}(s[..lo + 1] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi], lo + 1)) \wedge \neg(s[lo] \leq s[hi]) \ \emptyset$ 
    if-statement
    ( $\emptyset \ \text{isMax}(s[..lo + 1] + s[hi..], lo) \ \emptyset \ \text{imp}$ 
    hi := hi - 1 ;
    ( $\emptyset \ \text{isMax}(s[..lo + 1] + s[hi + 1..], lo) \ \emptyset \ \text{asg}$ 
    }
     $\text{isMax}(s[..lo] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi + 1..], lo)$  if-statement
     $lo \leq hi \wedge \text{isMax}(s[..lo] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi + 1..], lo)$  imp
  }
  ( $\emptyset \ (\text{isMax}(s[..lo] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi + 1..], lo)) \wedge \neg(lo < hi) \ \emptyset \ \text{while}$ 
  ( $\emptyset \ (\text{isMax}(s[..lo] + s[hi..], lo) \vee \text{isMax}(s[..lo + 1] + s[hi + 1..], lo)) \wedge lo = hi \ \emptyset \ \text{imp}$ 
  ( $\emptyset \ \text{isMax}(s[..lo] + s[lo..], lo) \ \emptyset \ \text{imp}$ 
  ( $\emptyset \ \text{isMax}(s, lo) \ \emptyset$ 

```

Question 3

Solution 3(a)

$$\forall x(0 \leq x < \frac{|s|}{2} \Rightarrow s[x] = s[|s| - x - 1])$$

Solution 3(b)

Show $\vdash_{tot} \langle 0 \leq |s| \rangle \text{ checkPalindrome } \langle T \rangle$.

Invariant = T

Variant = j - i

```
 $\langle 0 \leq |s| \rangle$   
 $\langle T \wedge 0 \leq |s| \rangle \text{ imp}$   
res := 1;  
var i := 0;  
 $\langle T \wedge i \leq |s| \rangle \text{ asg}$   
 $\langle T \wedge i - 1 \leq |s| - 1 \rangle \text{ imp}$   
var j := |s| - 1;  
 $\langle T \wedge -1 \leq j \rangle \text{ asg}$   
 $\langle T \wedge i - 1 \leq j \rangle \text{ imp}$   
 $\langle T \wedge -1 \leq j - i \rangle \text{ imp}$   
 $\langle T \wedge j - i + 1 \geq 0 \rangle \text{ imp}$   
while ( i < j & res = 1)  
{  
   $\langle T \wedge i < j \wedge \text{res} = 1 \wedge 0 \leq j - i + 1 = E_o \rangle \text{ while}$   
   $\langle T \wedge j - i > 0 \wedge \text{res} = 1 \wedge 0 \leq j - i + 1 = E_o \rangle \text{ imp}$   
   $\langle T \wedge \text{res} = 1 \wedge 0 < j - i < E_o \rangle \text{ imp}$   
  if ( s[i] != s[j] )  
  {  
     $\langle i < j \wedge \text{res} = 1 \wedge 0 \leq j - i + 1 = E_o \wedge s[i] \neq s[j] \rangle \text{ if-statement}$   
    res := 0  
     $\langle 0 \leq j - i + 1 < E_o \rangle \text{ imp}$   
  }  
  else  
  {  
     $\langle i < j \wedge \text{res} = 1 \wedge 0 \leq j - i + 1 = E_o \wedge \neg(s[i] \neq s[j]) \rangle \text{ if-statement}$   
    skip  
     $\langle 0 \leq j - i + 1 < E_o \rangle \text{ imp}$   
  }  
   $\langle T \wedge 0 \leq (j - 1) - (i + 1) + 1 < E_o \rangle \text{ imp}$   
  i := i + 1;  
   $\langle T \wedge 0 \leq (j - 1) - i + 1 < E_o \rangle \text{ asg}$   
  j := j - 1;  
   $\langle T \wedge 0 \leq j - i + 1 < E_o \rangle \text{ asg}$   
}  
 $\langle T \wedge (j \leq i \parallel \text{res} = 0) \rangle \text{ while}$   
 $\langle T \rangle$ 
```

Solution 3(c)

Show $\vdash_{par} \langle 0 \leq |s| \rangle \text{checkPalindrome} \langle (res == 1) \Leftrightarrow \text{isPal}(s) \rangle$.

Invariant = $res = 1 \Leftrightarrow \text{isPal}(s[..i] + s[j + 1..])$

Variant = $j - i + 1$

Let

$I_2 = 0 \leq i < j + 1 \leq |s|$

$\langle 0 \leq |s| \rangle$

$\langle I_2 \wedge 1 = 1 \Rightarrow \text{isPal}(s[..0] + s[0 - 1 + 1..]) \wedge \text{isPal}(s[..0] + s[0 - 1 + 1..]) \Rightarrow 1 = 1 \rangle \text{imp}$

res := 1;

var i := 0;

var j := |s| - 1;

$\langle I_2 \wedge res = 1 \Rightarrow \text{isPal}(s[..0] + s[0 - 1 + 1..]) \wedge \text{isPal}(s[..0] + s[0 - 1 + 1..]) \Rightarrow res = 1 \rangle \text{asg}$

while (i < j & res = 1)

{

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i] + s[j + 1..]) \wedge \text{isPal}(s[..i] + s[j + 1..]) \Rightarrow res == 1 \rangle \text{while}$

$\langle I_2 \wedge (s[i]! = s[j] \Rightarrow (0 \neq 1 \wedge \text{isPal}(s[..i] + s[j + 1..]))) \wedge \neg(s[i]! = s[j] \Rightarrow (1 \neq 1 \wedge \text{isPal}(s[..i] + s[j + 1..]))) \rangle \text{if-statement}$

if (s[i] != s[j])

{

$\langle 0 \neq 1 \wedge \text{isPal}(s[..i] + s[j + 1..]) \rangle \text{imp}$

res := 0

$\langle res == 0 \wedge \text{isPal}(s[..i + 1] + s[j..]) \rangle \text{imp}$

$\langle res == 1 \Rightarrow \text{isPal}(s[..i + 1] + s[j..]) \wedge \text{isPal}(s[..i + 1] + s[j..]) \Rightarrow res == 1 \rangle \text{imp}$

}

}

else

{

$\langle 1 == 1 \wedge \text{isPal}(s[..i] + s[j + 1..]) \rangle \text{imp}$

skip

$\langle res == 1 \wedge \text{isPal}(s[..i + 1] + s[j..]) \rangle \text{imp}$

$\langle res == 1 \Rightarrow \text{isPal}(s[..i + 1] + s[j..]) \wedge \text{isPal}(s[..i + 1] + s[j..]) \Rightarrow res == 1 \rangle \text{imp}$

}

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i + 1] + s[j..]) \wedge \text{isPal}(s[..i + 1] + s[j..]) \Rightarrow res == 1 \rangle \text{if-statement}$

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i + 1] + s[j + 1 - 1..]) \wedge \text{isPal}(s[..i + 1] + s[j + 1 - 1..]) \Rightarrow res == 1 \rangle \text{imp}$

i := i + 1;

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i] + s[j + 1 - 1..]) \wedge \text{isPal}(s[..i + 1] + s[j + 1 - 1..]) \Rightarrow res == 1 \rangle \text{asg}$

j := j - 1;

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i] + s[j + 1..]) \wedge \text{isPal}(s[..i] + s[j + 1..]) \Rightarrow res == 1 \rangle \text{asg}$

}

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i] + s[j + 1..]) \wedge \text{isPal}(s[..i] + s[j + 1..]) \Rightarrow res == 1 \wedge \neg(i < j \wedge res = 1) \rangle \text{while}$

$\langle I_2 \wedge res == 1 \Rightarrow \text{isPal}(s[..i] + s[j + 1..]) \wedge \text{isPal}(s[..i] + s[j + 1..]) \Rightarrow res == 1 \wedge i = j \rangle \text{imp}$

$\langle res == 1 \Rightarrow \text{isPal}(s) \wedge \text{isPal}(s) \Rightarrow res == 1 \rangle \text{imp}$

(| `res` == 1 \Leftrightarrow `isPal(s)` |)

Question 4

Solution 4(a)

Show $\vdash_{tot} (\mid 0 < |s| \mid) \text{ findMax } (\mid \text{isMax}(s, lo) \mid)$.

```
method findMax(s: seq<int>) returns(lo: int)
requires |s| > 0
ensures isMax(s, lo)
{
    assert(|s| > 0);
    assert(0 < 1 <= |s| && isMax(s[..1], 0));
    lo := 0;
    assert(0 <= lo < |s| && isMax(s[..1], lo));
    var hi : int := |s| - 1 ;
    assert(0 <= hi < |s| && isMax(s[hi..], 0));
    assert(hi - lo >= 0);
    while (lo < hi)
        decreases hi - lo
        invariant 0 <= lo <= hi < |s|
        invariant (isMax2(s, 0, lo, lo) && isMax2(s, hi, |s|-1, lo)) || (
            isMax2(s, 0, lo, hi) && isMax2(s, hi, |s|-1, hi))
        {
            if(s[lo] <= s[hi])
            {
                lo := lo + 1;
            }
            else
            {
                hi := hi - 1;
            }
        }
    }

predicate isMax2(s: seq<int>, lo: int, hi: int, max: int){
    0 <= max < |s| && 0 <= lo <= hi < |s| && forall x : int :: lo <= x <= hi
    ==> s[max] >= s[x]
}

predicate isMax(s : seq<int>, lo: int){
    0 <= lo < |s| && forall x: int :: 0 <= x < |s| ==> s[lo] >= s[x]
}
```

Solution 4(b)

Show $\vdash_{par} \langle T \rangle \text{checkPalindrome} \langle (res == 1) \Leftrightarrow \text{isPal}(s) \rangle$.

```
method checkPalindrome(s: seq<int>) returns (res: bool)
requires |s| >= 0
ensures isPalindrome(s) <==> res == true
{
    res := true;
    var i := 0 ;
    var j := |s| - 1 ;
    while ( i < j && res == true)
    invariant i == |s| - 1 - j && i <= |s| && ((forall k :: 0 <= k < i ==> s[k]
        == s[|s|-k-1]) <==> res == true)
    decreases |s| - i
    {
        if (s[i] != s[j])
        {
            res := false;
        }
        else {
            i := i + 1 ;
            j := j - 1;
        }
    }
}

predicate palindrome(s1: seq<int>, s2: seq<int>)
{
    |s1| == |s2| && forall x : int :: 0 <= x < |s1| ==> s1[x] == s2[|s2|-x-1]
}

predicate isPalindrome(s: seq<int>)
{
    forall x : int :: 0 <= x < |s| ==> s[x] == s[|s|-x-1]
}
```