[Logic: Propositions and Operations Def: A proposition is a statement that can be unambiguously evaluated as true or false. Ex: p = "There is life on Mass" + unambiguously false q = "My name is Annalise" + unambiguously true Non-Ex: r= "What time is it now?" 5 = "More people have been in Morcow than I have" grammatically correct but non-sensical t="Hissipp (a+b) = a2+ Lab + b2 is a proposition in algorous p= "Pant do it = not a proposition be it can't be assigned a truth value P.= x+1 = 2" + not a proposition be me don't know what x is Ps= A square has 4 sides -> true proposition p4= x20" → not a proposition be x is not necessarily a real number 1. Negation: Det let p be a proposition. Then -p or p is the proposition, called regular of p, -p = "It is not the case that p' Ex let p = "Jane is healthy" -p= "It is not the case that Jane is healthy" or "Jane is sick" Ex let p="I got the highest score on the exam" -p= "I didn't the highest score on the exam" Truth Tables P -P 1.2 Legical AND, or Conjunction Def Let p and q be two propositions. Then the conjunction of p and q is the proposition p A q, given by PIGDAG Ex Let p="I am vegetorian" Let q = "Sarah likes apples" (false) (false) pAq="I am regetorion and Small like uple" I T

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	D Logic: Propositions and Operations	
	1.3 Logical OR, or disjunction	
	Def let p and q be two propositions. Then the disjunction of	19
No.	pand q is the proposition by a 13 given by	
	P= I was to visit fronce	_
en me	y I WALL TO ALL	The second second second
	pry= "I want to visit France or Spain" (t) T F T	
-	1.4 Exclusive OR, or XOR	
	Truth table p @ Q P Q p@Q fx p: "You can take a colfe F F F * specify by using "or	Lu
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	V A STATE OF THE S	
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-	a multin, but not leth. TITF	
1	1.5 Conditional Statement or implication	
	Def Let p, q be two propositions. Then the conditional statement	
	$p \rightarrow q$ is the proposition given by $p \mid q \mid p \rightarrow q$ Ex Let $p = ''I$ am the president' $F \mid F \mid T$	
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-	VIII (VIII)	
1	P Q p > Q T T T president fulfilled the promise	
1	TF F president lied	
-	FIT T president did not break promise	
1	IF F T president did not break pornise	
A Constitution of the last	Ex Let p = "We live on mass" (f)	
1	let q= "2+2=10119" (f)	
-	p=q="If we live on More, then "2+2=10119" = true	
1	* if you assume nonsense, any statement can be true *	
	paq: If pitte q . p is sufficient for q eqposited that p	
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(1) Logic: Propositions and Operations I In prog p is called condition or premise q to called conclusion or consequence Det Let p-) q be a conditional statement Then · q - p is called converse · - q -> - pis called contrapositive ·-p-1-q is called inverse 1.6 Bicanditional Statements Def let p an q be two propositions. Then the compound proposition p to q is called broaditional or equivalence and given by the table Ex p= "You are get a flight T q = "You have a ticket penq = "You can fly if and only if you have a ticket p is equivalent to q or p is exactly when "Sarah will go for a run unless weather is load" (T) · if weather is good, Sarah will run and if it is bad, she will not. " p= "Scrah rons" q= "Weather is good" p+q "You can't min the lettery unless you buy a firstet." if you don't buy you cont min and if you do you con p="you con't win a letting" q="you don't buy a ticket" q -> p "A lunka is healthy unless it has a lympoka" (T) · if lunka is healty lit doesn't have lympolea. What question if lunka is stell it has a lympolea · if lunka has lympoka it is sicle . if it doesn't have lympoke it is healthy ! Translation between logic and language · punless q means -q -> p · p is necessary for q means q > p · p only if q means p -> q · p is necessary and sufficient for q means p -> q · p is sufficient for q means p+q

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(1) Logic: Predicates and Quantifiers

Ex Domain: all people, P(x)="x is a chess player."
    IxP(x) = "there exists purson x such that x is a chess player"
             = "There exists a chass, player" (true)
Ex Domain : all real numbers, P(x) = "x > 3"
   FXP(x) = "there exists number x s.t. x>3
Ex Domain: all people on Mars, P(x) = "x has a leg"
    IxP(x)= "There exects a person on Mars with a leg." (false) (no ex.)
    YxP(x)= "Every person on Mars has a leg." (tre) (no counter-ex.
· to show that \exists x P(x) is true to usually easy ble it is enough
 to provide an example.
Quantifiers over Finite Domains
 Domain: Jessica, Mike, Bob, Alice P(x) - predicate
 · Yx P(x) is the same as P(Jesoica) A P(Mike) A P(Bob) A P(Alice)
 · 7 x P(x) is the same as P(Tessica) v P(Mike) V P(Bob) V P(Alice)
But if domain is infinite, you can't do this.
Common Mistake P(x)="x">0" * must always specify domain *
                                     Domain all pos. red num txP(x) - free
                \forall \times P(x)
               FXP(x)
                                                               FXP(x) - tive
                                     Domain: all reals YXP(x)- Falce FX (x) - tree
                                      Domain: all neg reals YXP(x) - Folse 3xP(x) - folse
                                     Domain empty VxP(x) - true ]xP(x) - fake
Logical Equivalence
Def Two expressions with predicates and quantifixes are logically equivalent
      if they are equivalent as propositions for all choices of domains,
      predicates, and specified variables.
E_{\times} (\forall \times P(\times)) \Lambda(\forall y Q(y)) \equiv \forall x \forall y (P(x) \Lambda P(y))
Notation is often simplified
Ex Domainiall reals Ex VX ((x <0) -> (x3 <0)), where x
     Yx(x2>,0)
                        a real number.
Ex ] x ((x2=2) 1 (x>0)), where x 13 a red runber.
Ex 4x40 (x3<0) (Domain: all 11115) Ex 7x>0(x2=2) (Domain: all 1116)
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0 1 Logic: Predicates and Quantifiers Logical Equivalence Ex Yx[P(x) AQ(x)] = [VxP(x)] A [VxQ(x)] De Morgen Laws for Quantifiers

[ExDomain: all students P(x) = "x took calculus" Vx P(x) = "All students took calculus" ~[VxP(x)] = "Not all students took calculus." = "There is a student who didn't take calculus =]x [¬P(x)]. ExDomain: all people P(x)="x is an honest politician." 7xP(x): "There exists an honest politician." -[]xP(x)]: "Every politician is distanst." = Yx[-P(x)] Def - [YxP(x)]=]x[-P(x)] 7 [3xP(x)] = \(\frac{1}{2}\) \(\frac{1}{2}\) Ex Domain: rational numbers (5/1, 1/2, -3/6) $\exists x(x^2 = 2)$ - [3x(x=2)] = \forall x [x \div \div 2] Ex Yx (x2>1x), where x is a geal number [7(p+g)=px(7g) $\neg \left[\forall x \left(x^2 > x \right) \right] \equiv \exists x \left[\left(x^2 < x \right) \right]$ $\mathbb{E}_{x} \cdot (\forall x [P(x) \rightarrow Q(x)]) = \exists x [\neg (P(x) \rightarrow Q(x))]$ $\equiv \exists x [P(x) \land (\neg Q(x))]$ Ex. Domein: all students ((x) = "x took colocles" VxS(x) -> t/xC(x) Yx C(x) = "All students took calculus." · Domain = all people S(x) = "x is a student" Yx [S(x) → C(x)]. correct. Yx [S(x) ∧ C(x)]. Wrong. Ex Domain: all students R(x) = "x visited Russia" FXR(x) = "There is a student who visited Russia · Domain : all pieple S(x) = "x is a student" Fx [S(x) A R(x)] correct. Fx[S(x) -> R(x)] wrong. * Remark about variable names Ex Domain: all cilius P(x)="x is in the U.S." $\exists s P(s)$ Ex P(x,y) = "x is a city in y" = 32 Jw P(z, w) (can call variables anything)