

Question 1:

A* Algorithm

g=1 h=17									
g=2 h=16									
g=3 h=15									
g=4 h=14									
g=5 h=13									
g=6 h=12									
g=7 h=11									
g=8 h=10									
g=9 h=9	g=10 h=8	g=11 h=7	g=12 h=6	g=13 h=5	g=14 h=4	g=15 h=3	g=16 h=2	g=17 h=1	g=18 h=0

Explanation: In a A* search, because it takes into consideration both path cost and future cost, A* takes a more balanced approach. It expands nodes that balance shortest path so far $g(n)$ with closeness to goal $h(n)$.

Greedy Best-First

	g=0 h=17	g=0 h=16	g=0 h=15	g=0 h=14	g=0 h=13	g=0 h=12	g=0 h=11	g=0 h=10	g=0 h=9
									g=0 h=8
									g=0 h=7
									g=0 h=6
									g=0 h=5
									g=0 h=4
									g=0 h=3
g=0 h=11	g=0 h=10	g=0 h=9	g=0 h=8	g=0 h=7	g=0 h=6	g=0 h=5	g=0 h=4	g=0 h=3	g=0 h=2
g=0 h=10									
g=0 h=9	g=0 h=8	g=0 h=7	g=0 h=6	g=0 h=5	g=0 h=4	g=0 h=3	g=0 h=2	g=0 h=1	g=0 h=0

Explanation: In a Greedy Best-First search, it ignores path cost and only takes into consideration future cost. In doing so, the search attempts to “beeline” to the goal but is blocked by a wall. It is forced to travel along the wall until it isn’t blocked.

Removal of $g(n)$

```
#####
### Greedy Best-First Algorithm
#####
def find_path(self):
    open_set = PriorityQueue()

    ### Add the start state to the queue
    open_set.put((0, self.agent_pos))

    ### Continue exploring until the queue is exhausted
    while not open_set.empty():
        current_cost, current_pos = open_set.get()
        current_cell = self.cells[current_pos[0]][current_pos[1]]

        ### Stop if goal is reached
        if current_pos == self.goal_pos:
            self.reconstruct_path()
            break

        ### Agent goes E, W, N, and S, whenever possible
        for dx, dy in [(0, 1), (0, -1), (1, 0), (-1, 0)]:
            new_pos = (current_pos[0] + dx, current_pos[1] + dy)

            if 0 <= new_pos[0] < self.rows and 0 <= new_pos[1] < self.cols and not self.cells[new_pos[0]][new_pos[1]].is_wall:
                ### Eliminates cost of moving to a new position
                new_g = current_cell.g

                if new_g < self.cells[new_pos[0]][new_pos[1]].g:
                    ### Update the path cost g()
                    self.cells[new_pos[0]][new_pos[1]].g = new_g

                    ### Update the heuristic h()
                    self.cells[new_pos[0]][new_pos[1]].h = self.heuristic(new_pos)

                    ### Update the evaluation function for the cell n:  $f(n) = h(n)$ 
                    self.cells[new_pos[0]][new_pos[1]].f = new_g + self.cells[new_pos[0]][new_pos[1]].h ##  $g(n) = 0$  in search
                    self.cells[new_pos[0]][new_pos[1]].parent = current_cell

                    ### Add the new cell to the priority queue
                    open_set.put((self.cells[new_pos[0]][new_pos[1]].f, new_pos))
```

Explanation: To achieve a Greedy Best-First search, I altered the A* Algorithm in AStarMaze by changing this snippet of code:

$\text{new_g} = \text{current_cell.g} + 1 \rightarrow \text{new_g} = \text{current_cell.g}$

By doing so, $g(n)$ is constant and does not accumulate path cost and therefore is not calculated into the $f(n)$ function. The function now looks like this:

$f(n) = g(n) + h(n) \rightarrow f(n) = h(n)$

The algorithm no longer takes past cost $g(n)$ into consideration, only concerning itself with the estimate of future cost $h(n)$, making it greedy.

Conclusion: A* and Greedy algorithms differ in their best path to the goal. A* will take a balanced approach and consider both future and past costs. Greedy Best-First searches will result in a “beeline” to the goal, often resulting in inference with obstacles and resulting in a longer path. By making a small tweak to the code, by removing accumulation of $g(n)$, will result in a Greedy Best-First algorithm.