

Hashing

1) Count Frequency - Use dictionary

2) Find a pair with sum x in array

a) Sort and use 2 pointers ($O(n \log n)$)

b) Use dictionary ($O(n)$)

-> Take $d = \{\}$

-> loop through array. $l+j = x$; $j = x-i$

-> Chck if $x-i$ in d . If yes return the pair. Else add i to d .

$d = \{\}$

for i in arr :

 If $x-i$ in d :

 return $[i, x-i]$

$d[i] = 1$

3) Find whether an array is subset of another

Use dictionary

-> make frequency array of $arr1$ called map and then loop through $arr2$ and if element is in map then reduce $map[i] -= 1$ and delete element from $arr2$.

-> if in the end map elements have all 0 in frequency then $arr1$ is subset of $arr2$.

4) Maximum distance between 2 same elements in array

Input : $arr[] = \{3, 2, 1, 2, 1, 4, 5, 8, 6, 7, 4, 2\}$

Output: 10

// maximum distance for 2 is $11-1 = 10$

// maximum distance for 1 is $4-2 = 2$

// maximum distance for 4 is $10-5 = 5$

Use dictionary

-> if element in dict then $ans = \max(ans, i-d[li[i]])$

-> else add element to dict; $d[li[i]] = i$

```
Ans = -1
d = {}
for i in range(len(li)):
    if li[i] in d:
        ans = max(ans, i-d[li[i]])
    d[li[i]] = i
print(ans)
```

5) Min. operations to make all elements equal

-> result is $\text{len}(\text{arr}) - (\text{highest freq. element})$

```
def func(arr):  
    d = {}  
    for i in arr:  
        if i in d:  
            d[i] += 1  
        else:  
            d[i] = 1  
    return len(arr) - max(d.values())
```

6) Check if a given array contains duplicate elements within k distance from each other

Input: $k = 3$, $\text{arr}[] = \{1, 2, 3, 1, 4, 5\}$

Output: true

1 is repeated at distance 3.

```
def func(arr, k):  
    d = {}  
    for i in range(len(arr)):  
        if arr[i] in d:  
            if abs(i - d[arr[i]]) == k:  
                return True  
            d[arr[i]] = i  
        else:  
            d[arr[i]] = i  
    return False
```

7) Sum of elements in an array with frequencies greater than or equal to that element

```
def func(arr, k):  
    d = {}  
    for i in range(len(arr)):  
        if arr[i] in d:  
            d[arr[i]] += 1  
        else:  
            d[arr[i]] = 1
```

```
s = 0
for i in d.keys():
    if d[i] >= i:
        s += i
return s
```

8) First unique character of a string

-> Make a dict and store all indexes of elements and return the element with only single element and lowest index.

class Solution:

def firstUniqChar(self, s: str) -> int:

```
m = {}
for i in range(len(s)):
    if s[i] in m:
        m[s[i]].append(i)
    else:
        m[s[i]] = [i]
```

```
c = 2**31-1
print(m)
for i in m.keys():
    if len(m[i]) == 1:
        c = min(c, m[i][0])
```

```
if c == 2**31-1: return -1
return c
```

9) Count pairs with given sum

-> Create freq map and loop through array.

-> if k-i in d: count += d[k-i] as it can form as many as pair as count

```
def getPairsCount(self, arr, n, k):

    d = Counter()

    c = 0

    for i in arr:

        c += d.get(k-i, 0)

        d[i] += 1

    return c
```

10) Find Common Characters

-> Use dictionary

Input: words = ["bella", "label", "roller"]

Output: ["e", "l", "l"]

11) Count Number of Pairs With Absolute Difference K

-> Similar to question 9

class Solution:

```
def countKDifference(self, nums: List[int], k: int) -> int:
```

```
    d = {}
    c = 0
    for i in nums:
        c += d.get(i-k, 0) + d.get(i+k, 0)
        if i in d:
            d[i] += 1
        else:
            d[i] = 1
```

```
    return c
```

12) Number of Pairs of Strings With Concatenation Equal to Target

Brute -

class Solution:

```
def numOfPairs(self, nums: List[str], target: str) -> int:
```

```
    ans=0
    for i in range(len(nums)):
        for j in range(i+1, len(nums)):
            if nums[i]+nums[j]==target:
                ans+=1
            if nums[j]+nums[i]==target:
                ans+=1
    return ans
```

Optimal - -> Make a freq dict.

Loop through array and replace i from target

If i == target-i: cnt += d[target-i]-1 as we want to omit extra occurrence

else: cnt += d[target-i]

class Solution:

```
def numOfPairs(self, nums: List[str], target: str) -> int:
```

```
    ans=0
    cnt = Counter(nums)
    for i in nums:
        x = target.replace(i, "", 1)
        #print(x)
        if x in cnt:
            if x == i:
                ans += cnt[x] - 1
            else:
```

```

        if i+x == target:
            ans += cnt[x]
    return ans

```

13) Count maximum points on same line

-> calc. Slope of each point and store in map. Elements with same slope are in same line.

-> return the slope with maximum points.

```

class Solution:
    def slopeMaxElements(self, points: List[int], k: int) -> int:
        d = {}
        for i in points:
            x = i[0]
            y = i[1]
            slope = (y-0)/(x-0)
            if slope in d:
                d[slope] += 1
            else:
                d[slope] = 1

        return max(d.values())

```

14) Smallest Subarray with given sum

```

def subArraylen(arr, n, K):
    mp = defaultdict()
    currPrefixSum = 0
    result = sys.maxsize
    for i in range(n):
        currPrefixSum += arr[i]
        if(currPrefixSum == K):
            currLen = i + 1
            result = min(result, currLen)
            requirePrefixSum = currPrefixSum - K
        if(requirePrefixSum in mp.keys()):
            foundIdx = mp[requirePrefixSum]
            currIdx = i
            result = min(result, currIdx - foundIdx)
        mp[currPrefixSum] = i
    return result

```

15) Triplet Sum in Array (Better with 2 pointers)

(2 Pointer Approach TC - $O(n^2)$)

```

#User function Template for python3
class Solution:

```

```

#Function to find if there exists a triplet in the
#array A[] which sums up to X.
def find3Numbers(self,A, n, X):
    # Your Code Here
    A.sort()
    for i in range(n):
        l = i+1
        r = n-1
        while l<r:
            s = A[i]+A[l]+A[r]
            if s == X:
                return 1
            elif s > X:
                r -= 1
            else:
                l += 1
    return 0

```

(Hashmap TC - $O(n^2)$)

->Run loop in from i -> 0-n

->take a set and make curr_sum = sum-A[i] so we just need to find a pair with sum as curr_sum and again run a loop from i+1-n.

->see if curr_sum-A[j] in set. If yes return true

->add A[j] to set.

```

def find3Numbers(A, arr_size, sum):
    for i in range(0, arr_size-1):
        # Find pair in subarray A[i + 1..n-1]
        # with sum equal to sum - A[i]
        s = set()
        curr_sum = sum - A[i]
        for j in range(i + 1, arr_size):
            if (curr_sum - A[j]) in s:
                print("Triplet is", A[i],
                    ", ", A[j], ", ", curr_sum-A[j])
                return True
            s.add(A[j])
    return False

```

16) Subarray sum equals k

-> take sums = 0 and loop through array

-> at every iteration add element to sums

If $k \neq 0$ then make $\text{sums} = \text{sums} \% k$

If sums is there in dict return True if $i - d[\text{sums}] > 1$ as we want a subarray with length 2 or more

else make $d[\text{sums}] = i$

Example:

```
nums = [23, 2, 4], k = 6
```

Lets walk through the code with the example.

```
(i=0) : sums = 23 => 23%6 => (sums = 5)
```

```
(i=1) : sums = 5+2=7 => 7%6 => (sums = 1)
```

```
(i=2) : sums = 1+4=5 => 5%6 => (sums = 5)
```

We have encountered the same sums(remainder) again which means we have the subarray of $\text{sums} \% k = 0$.

But, there's another aspect to this problem. The subarray must have a minimum size of 2.

That is why we check if $(i - d[\text{sums}]) > 1$.

In the above example, this if loop is executed when $(i=2)$ and $(d[\text{sums}]=1)$.

In other words, the same remainder($\text{sums}=5$) has been encountered twice and then we check for the respective difference in indices.

Counter example to understand this. Lets take $\text{nums} = [23, 6], k = 6$

```
(i=0) : sums = 23 => 23%6 => (sums = 5)
```

```
(i=1) : sums = 5+6=11 => 11%6 => (sums = 5)
```

```
class Solution:

    def checkSubarraySum(self, nums: List[int], k: int) -> bool:

        curr = 0

        d = {0: -1}

        for i in range(len(nums)):

            curr += nums[i]

            if k != 0:

                curr = curr % k

            if curr in d:

                if i - d[curr] > 1:

                    return True

            else:

                d[curr] = i
```

```
return False
```

17) Number of subarray with sum 0

First of all, the basic idea behind this code is that, whenever sums has increased by a value of k, we've found a subarray of sums=k.

I'll also explain why we need to initialise 0 in the hashmap.

Example: Let's say our elements are [1,2,1,3] and k = 3.

and our corresponding running sums = [1,3,4,7]

Now, if you notice the running sums array, from 1->4, there is increase of k and from 4->7, there is an increase of k. So, we've found 2 subarrays of sums=k.

But, if you look at the original array, there are 3 subarrays of sums==k. Now, you'll understand why 0 comes in the picture.

In the above example, $4-1=k$ and $7-4=k$. Hence, we concluded that there are 2 subarrays.

However, if sums==k, it should've been $3-0=k$. But 0 is not present in the array. To account for this case, we include the 0.

Now the modified sums array will look like [0,1,3,4,7]. Now, try to see for the increase of k.

1. 0->3
2. 1->4
3. 4->7

Hence, 3 sub arrays of sums=k

This clarified some confusions I had while doing this problem.

```
class Solution(object):
    def subarraySum(self, nums, k):
        """
        :type nums: List[int]
        :type k: int
        :rtype: int
        """
        count = 0
        sums = 0
        d = dict()
        d[0] = 1

        for i in range(len(nums)):
            sums += nums[i]
            count += d.get(sums-k, 0)
            d[sums] = d.get(sums, 0) + 1

        return(count)
```


