

# **Image Processing Assignment 8**

## **Image Compression and Decompression using Huffman and Arithmetic Coding**

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Date: October 18, 2025

## **1 Aim**

To write a Python program for Image Compression and Decompression using:

- Huffman Coding and Decoding
- Arithmetic Coding and Decoding

## **2 Theory**

Image compression reduces the amount of data required to represent an image while preserving visual quality. Lossless compression techniques such as Huffman and Arithmetic coding exploit statistical redundancy in the image data.

### **2.1 Huffman Coding**

Huffman coding is a variable-length lossless compression technique that assigns shorter binary codes to frequent symbols and longer codes to rare ones.

**Steps:**

1. Calculate frequency of each pixel intensity.
2. Construct a binary Huffman tree.
3. Assign binary codes to each pixel value.
4. Replace pixel values with their corresponding codes.

**Advantages:**

- Simple and effective for data with skewed probability distributions.
- Produces exact decompression without any loss of information.

**Disadvantages:**

- Requires storing the Huffman tree along with the compressed data.
- Inefficient if symbol probabilities are nearly uniform.

## 2.2 Arithmetic Coding

Arithmetic coding represents the entire message as a fractional number in the range [0, 1). It achieves higher compression efficiency than Huffman coding by allowing fractional bits per symbol.

### Steps:

1. Compute probabilities of all pixel values.
2. Define cumulative probability ranges for each symbol.
3. Narrow down the range for each successive symbol.
4. Output the final interval as the encoded value.

### Advantages:

- Provides near-optimal compression close to entropy limit.
- Handles fractional probabilities more accurately.

### Disadvantages:

- Computationally complex and slower than Huffman coding.
- Sensitive to rounding and floating-point precision.

## 3 Algorithm

### 3.1 Huffman Coding

1. Read image and flatten it into a 1D array.
2. Compute pixel frequency distribution.
3. Build a Huffman tree using frequencies.
4. Encode image pixels using generated codes.
5. Save encoded data and tree information to a file.

### 3.2 Arithmetic Coding

1. Compute frequency and cumulative tables.
2. Encode image data using scaled integer arithmetic.
3. Save encoded bitstream and header information.
4. Decode the bitstream using inverse arithmetic mapping.

## 4 Implementation

The implementation was done in Google Colab using Python and OpenCV.

## 4.1 Required Packages

```
1 !pip install opencv-python matplotlib numpy
```

## 4.2 Code Implementation

Google Colab Notebook: [Click here to view the Colab Notebook](#)

```
1 import numpy as np
2 import cv2
3 import heapq
4 import pickle
5 from google.colab.patches import cv2_imshow

1 img = cv2.imread("771476.jpg")
2 # Display
3 cv2_imshow(img)

1 class HuffNode:
2     def __init__(self, symbol, freq):
3         self.symbol = symbol
4         self.freq = freq
5         self.left = None
6         self.right = None
7     def __lt__(self, other):
8         return self.freq < other.freq
9
10 def build_huffman_tree(freqs):
11     heap = [HuffNode(sym, f) for sym, f in freqs.items()]
12     heapq.heapify(heap)
13     while len(heap) > 1:
14         a = heapq.heappop(heap)
15         b = heapq.heappop(heap)
16         parent = HuffNode(None, a.freq + b.freq)
17         parent.left = a
18         parent.right = b
19         heapq.heappush(heap, parent)
20     return heap[0]
21
22 def build_codes(node, code='', mapping=None):
23     if mapping is None:
24         mapping = {}
25     if node.symbol is not None:
26         mapping[node.symbol] = code or '0'
27     else:
28         build_codes(node.left, code + '0', mapping)
29         build_codes(node.right, code + '1', mapping)
30     return mapping

1 def huffman_compress(img, out_file="compressed_image.huff"):
```

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2     flat = img.flatten().tolist()
3     freqs = {}
4     for val in flat:
5         freqs[val] = freqs.get(val, 0) + 1
6
7     root = build_huffman_tree(freqs)
8     codes = build_codes(root)
9
10    bitstream = ''.join(codes[v] for v in flat)
11    padding = 8 - len(bitstream) % 8
12    bitstream += '0' * padding
13
14    data = bytearray()
15    for i in range(0, len(bitstream), 8):
16        data.append(int(bitstream[i:i+8], 2))
17
18    header = {'shape': img.shape, 'codes': codes, 'padding': padding}
19
20    with open(out_file, 'wb') as f:
21        pickle.dump((header, data), f)
22
23    orig_size = img.size
24    comp_size = len(data)
25    ratio = orig_size / comp_size
26    print(f'Huffman compressed: {orig_size} bytes      {comp_size} bytes (ratio: {ratio:.2f}:1)')
27    return out_file

```

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1 def huffman_decompress(huff_file, out_image="reconstructed_from_huffman.png"):
2     with open(huff_file, 'rb') as f:
3         header, data = pickle.load(f)
4
5     shape = header['shape']
6     codes = header['codes']
7     padding = header['padding']
8
9     reverse_codes = {v: int(k) for k, v in codes.items()}
10    bitstring = ''.join(format(byte, '08b') for byte in data)
11    bitstring = bitstring[:-padding] if padding > 0 else
12        bitstring
13
14    current = ''
15    decoded = []
16    for bit in bitstring:
17        current += bit
18        if current in reverse_codes:
19            decoded.append(reverse_codes[current])
            current = ''

```

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21     arr = np.array(decoded, dtype=np.uint8).reshape(shape)
22     cv2.imwrite(out_image, arr)
23     print(f"Huffman decompressed and saved as: {out_image}")
24     return arr
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45
1 def arithmetic_compress(img, out_file="compressed_image.arith"):
2     flat = img.flatten().tolist()
3     freqs = {}
4     for val in flat:
5         freqs[val] = freqs.get(val, 0) + 1
6     total = len(flat)
7     symbols = sorted(freqs.keys())
8
9     cumul = {}
10    cum = 0
11    for s in symbols:
12        cumul[s] = cum
13        cum += freqs[s]
14
15    SCALE = 1 << 16
16    low = 0
17    high = SCALE - 1
18    output_bits = []
19    pending_bits = 0
20
21    for pixel in flat:
22        range_size = high - low + 1
23        high = low + (range_size * (cumul[pixel] + freqs[pixel]))
24            // total - 1
25        low = low + (range_size * cumul[pixel]) // total
26
27        while True:
28            if high < SCALE // 2:
29                output_bits.append(0)
30                output_bits.extend([1] * pending_bits)
31                pending_bits = 0
32                low = low * 2
33                high = high * 2 + 1
34            elif low >= SCALE // 2:
35                output_bits.append(1)
36                output_bits.extend([0] * pending_bits)
37                pending_bits = 0
38                low = (low - SCALE // 2) * 2
39                high = (high - SCALE // 2) * 2 + 1
40            elif low >= SCALE // 4 and high < 3 * SCALE // 4:
41                pending_bits += 1
42                low = (low - SCALE // 4) * 2
43                high = (high - SCALE // 4) * 2 + 1
44            else:
45                break

```

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46     pending_bits += 1
47     if low < SCALE // 4:
48         output_bits.append(0)
49         output_bits.extend([1] * pending_bits)
50     else:
51         output_bits.append(1)
52         output_bits.extend([0] * pending_bits)
53
54     bitstring = ''.join(str(b) for b in output_bits)
55     padding = (8 - len(bitstring) % 8) % 8
56     bitstring += '0' * padding
57
58     encoded_bytes = bytearray()
59     for i in range(0, len(bitstring), 8):
60         encoded_bytes.append(int(bitstring[i:i+8], 2))
61
62     compressed_data = {
63         'shape': img.shape,
64         'freqs': freqs,
65         'cumul': cumul,
66         'total': total,
67         'symbols': symbols,
68         'encoded': bytes(encoded_bytes),
69         'padding': padding
70     }
71
72     with open(out_file, 'wb') as f:
73         pickle.dump(compressed_data, f)
74
75     orig_size = img.size
76     comp_size = len(encoded_bytes)
77     ratio = orig_size / comp_size
78     print(f"Arithmetic compressed: {orig_size} bytes      {"
79           f"comp_size} bytes (ratio: {ratio:.2f}:1)")
80     return out_file

```

```

1 def arithmetic_decompress(arith_file, out_image="reconstructed_from_arithmetic.png"):
2     with open(arith_file, 'rb') as f:
3         compressed_data = pickle.load(f)
4
5     shape = compressed_data['shape']
6     freqs = compressed_data['freqs']
7     cumul = compressed_data['cumul']
8     total = compressed_data['total']
9     symbols = compressed_data['symbols']
10    encoded = compressed_data['encoded']
11    padding = compressed_data['padding']
12
13    bitstring = ''.join(format(byte, '08b') for byte in encoded)
14    if padding > 0:

```

```

15     bitstring = bitstring[:-padding]
16
17     SCALE = 1 << 16
18     low = 0
19     high = SCALE - 1
20     code = 0
21
22     for i in range(min(16, len(bitstring))):
23         code = (code << 1) | int(bitstring[i])
24
25     bit_pos = 16
26     decoded = []
27     pixels_needed = shape[0] * shape[1]
28
29     for _ in range(pixels_needed):
30         range_size = high - low + 1
31         scaled_value = ((code - low + 1) * total - 1) // range_size
32         symbol = None
33         for s in symbols:
34             if cumul[s] <= scaled_value < cumul[s] + freqs[s]:
35                 symbol = s
36                 break
37         if symbol is None:
38             symbol = symbols[-1]
39
40         decoded.append(symbol)
41         high = low + (range_size * (cumul[symbol] + freqs[symbol])) // total - 1
42         low = low + (range_size * cumul[symbol]) // total
43
44     while True:
45         if high < SCALE // 2:
46             low = low * 2
47             high = high * 2 + 1
48             code = (code * 2) % SCALE
49             if bit_pos < len(bitstring):
50                 code |= int(bitstring[bit_pos])
51                 bit_pos += 1
52         elif low >= SCALE // 2:
53             low = (low - SCALE // 2) * 2
54             high = (high - SCALE // 2) * 2 + 1
55             code = ((code - SCALE // 2) * 2) % SCALE
56             if bit_pos < len(bitstring):
57                 code |= int(bitstring[bit_pos])
58                 bit_pos += 1
59         elif low >= SCALE // 4 and high < 3 * SCALE // 4:
60             low = (low - SCALE // 4) * 2
61             high = (high - SCALE // 4) * 2 + 1
62             code = ((code - SCALE // 4) * 2) % SCALE
63             if bit_pos < len(bitstring):

```

```

64         code |= int(bitstring[bit_pos])
65         bit_pos += 1
66     else:
67         break
68
69     arr = np.array(decoded, dtype=np.uint8).reshape(shape)
70     cv2.imwrite(out_image, arr)
71     print(f"Arithmetic decompressed and saved as: {out_image}")
72     return arr

```

```

1 print("*"*50)
2 print("HUFFMAN CODING")
3 print("*"*50)
4 gray = cv2.cvtColor(img, cv2.COLOR_BGR2GRAY)
5 huff_file = huffman_compress(gray)
6 recon_huff = huffman_decompress(huff_file)
7
8 print("*"*50)
9 print("ARITHMETIC CODING")
10 print("*"*50)
11 arith_file = arithmetic_compress(gray)
12 recon_arith = arithmetic_decompress(arith_file)

```

## 5 Results and Analysis

### 5.1 Observations

- Huffman and Arithmetic Coding both achieved lossless image reconstruction.
- Arithmetic coding provided a slightly better compression ratio.
- Huffman coding was simpler and faster in execution.

## 6 Conclusion

This experiment successfully implemented Huffman and Arithmetic Coding algorithms for image compression and decompression. Both methods preserved image quality while reducing storage requirements. Huffman coding is efficient for simple probability distributions, whereas Arithmetic coding achieves higher compression for continuous-valued data.