

Homework 1
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Exercise 2. *Show that the number of monochromatic triangles in any 2-coloring of the edges of K_n is at least*

$$\frac{n(n-1)(n-5)}{24}$$

Proof. Let G be a K_n and Fix a vertex $v \in V(G)$. There are n such ways to fix a vertex v . Now color edges incident to v such that we maximize the number of triangles that are not monochromatic (i.e. contain edges of different colors). For each vertex $v \in V(G)$ color $\frac{n-1}{2}$ edges one color, and the remaining edges another color. This creates at most $\frac{n-1}{2}$ new triangles that are not monochromatic. For every vertex v we create at most $\frac{1}{2} \left(\frac{n-1}{2} \right)^2$ such triangles as we have double counted. The total number of triangles in G is $\binom{n}{3}$. Therefore, the total number of monochromatic triangles is

$$\begin{aligned} &\geq \binom{n}{3} - \frac{n}{2} \left(\frac{n-1}{2} \right)^2 = \frac{n(n-1)(n-2)}{6} - \frac{n(n-1)^2}{8} \\ &= \frac{n(n^2 - 6n + 5)}{24} = \frac{n(n-1)(n-5)}{24} \end{aligned}$$

□

Exercise 3. Show that the number of triangles in a graph with n vertices and m edges is at least

$$\frac{4m}{3n} \left(m - \frac{n^2}{4} \right)$$

Proof. Let G be a graph on n vertices and m edges, and consider $uv \in E(G)$. The $d(u) + d(v) > n$ if and only if \exists a triangle containing the edge uv . Therefore, it follows that the number of common neighbors of u and v is at least $d(u) + d(v) - n$. Each common neighbor of u and v form a triangle in G and there are greater than $d(u) + d(v) - n$ such neighbors for each $uv \in E(G)$. This implies the number of triangles

$$\geq \frac{1}{3} \sum_{uv \in E(G)} (d(u) + d(v) - n) = \frac{1}{3} \sum_{v \in V(G)} (d(v)^2) - \frac{1}{3} mn$$

By Cauchy-Schwarz we have

$$\geq \frac{1}{3n} \left(\sum_{v \in V(G)} d(v) \right)^2 - \frac{1}{3} mn = \frac{1}{3n} (2m)^2 - \frac{1}{3} mn = \frac{4m^2}{3n} - \frac{1}{3} mn$$

□

Exercise 4. *Confirm that $T_k(n)$ is the largest (most edges) n -vertex complete multipartite graph without a $(k + 1)$ -clique.*

Proof. Let G be a candidate graph for a complete multipartite graph without a $(k + 1)$ -clique. Observe that certainly an edge maximal graph without a $k + 1$ -clique requires the addition of only one edge to have a $k + 1$ -clique. This implies that such a graph must have a K_k embedded within it. Use a K_k as the starting point for G , certainly it has the most edges of any k -vertex multipartite graph. Now we will add vertices one by one, until we run out of vertices. We want to add a vertex in such a way that maximizes the number of edges we create. First add a vertex to any of the k -classes from the k -partite G for a new G . A new class wasn't created as it would have formed a $k + 1$ -clique with the other k classes. In general, each iteration we will add a vertex to a class of smallest size, which ensures that the number of edges that are created is maximal. If all classes at that current iteration have the same size, pick a class arbitrarily to add a vertex to. Repeat this procedure until we have used all of the vertices. Observe that each class differs by a size of at most 1, and more specifically the ending G will have class sizes $\lfloor \frac{n}{k} \rfloor$ and $\lceil \frac{n}{k} \rceil$. This is the Turán graph $T_k(n)$, so $G = T_k(n)$. \square

Exercise 6. *The maximum number of edges in a graph G on n vertices with no $(k+1)$ -clique subgraph is at most*

$$\left(1 - \frac{1}{k}\right) \frac{n^2}{2}$$

Proof. To each vertex $v \in V(G)$ assign a non-negative weight $w(x)$ such that

$$\sum_{x \in V(G)} w(x) = 1$$

We would like to determine the max value of the following

$$S = \sum_{xy \in E(G)} w(y)w(x)$$

Assigning $1/n$ to each vertex gives the maximum of S is

$$S \geq \frac{|E(G)|}{n^2}$$

Showing S cannot exceed $\frac{1}{2}\left(1 - \frac{1}{k}\right)$ will complete the proof. So we will employ the weight shifting technique to show this. Let x and y be non-adjacent vertices and let W_x and W_y be the sum of weights on vertices adjacent to x and y respectively. Assume $W_x \geq W_y$ and let $\epsilon > 0$, thus

$$(w(x) + \epsilon)W_x + (w(y) - \epsilon)W_y \geq w(x)W_x + w(y)W_y$$

This implies we can shift all of the weight from one vertex y to some non-adjacent vertex x and not decrease S (assuming $W_y \leq W_x$). The graph has no K_{k+1} so we can shift all of the weight to the vertices of a K_k and not decrease S . Certainly S is maximized when each vertex has equal weight. A K_k has k vertices and the weights sum to 1 so each vertex will have weight equal to $(1 - \frac{1}{k})$. Therefore, S is maximized at $\frac{1}{2}\left(1 - \frac{1}{k}\right)$. Substitute this value for S and we get the following

$$\frac{1}{2}\left(1 - \frac{1}{k}\right) \geq \frac{|E(G)|}{n^2}$$

Which is the theorem. □

Exercise 7. Let $n \geq k + 1$ and let G be an n -vertex graph. If $e(G) = ex(n, K_k) + 1$, then G contains the graph H formed by $K_k + 1$ minus an edge.

Proof. Proceed by induction on n , the number of vertices. The base case is when $n = k + 1$. Observe that G is simply the Turán graph with $k - 1$ classes, and $k + 1$ vertices. In other words

$$e(G) = ex(n, K_k) + 1 = e(T_{k-1}(n)) + 1$$

Therefore, it follows that $k - 3$ of the $k - 1$ classes contain only a single vertex. The remaining 2 classes have just 2 vertices each, for a total of $k + 1 = n$. Let V_1 be the class with the added edge. Trivially V_1 must be a class of size two. It follows by the definition of the Turán graph that this edge must also have both its endpoints in V_1 , as all other edges between vertices of different classes are already present. Let V_2 be the other class with 2 vertices and observe that all the vertices of G are connected to one another except for the vertices in V_2 . Therefore, G contains H a K_{k+1} minus the edge between vertices in V_2 .

Let $n > k + 1$ and assume the statement above holds for graphs on fewer vertices. Consider a vertex v of minimum degree in G .

Consider the following cases:

1. $k - 1$ divides n .

Once again from above we know that $e(G) = ex(n, k_k) + 1 = e(T_{k-1}(n)) + 1$. Therefore G is simply the Turán graph with $k - 1$ classes each of size $\frac{n}{k-1}$, one with an edge between two of its vertices. Consider removing a vertex v of minimum degree to get G' . Observe that removing v doesn't decrease the number of classes, as each class has ≥ 2 vertices (assuming $k - 1$ divides n). However G' is simply $T_{k-1}(n - 1)$ plus an edge. By induction G' contains H , formed by K_{k+1} minus an edge, and G' is contained in G . Therefore, G also contains H !

2. $k - 1$ doesn't divide n

Once more from above we know that $e(G) = ex(n, k_k) + 1 = e(T_{k-1}(n)) + 1$. Therefore G is simply the Turán graph with $k - 1$ classes each of size $\lfloor \frac{n}{k-1} \rfloor$ or $\lceil \frac{n}{k-1} \rceil$. Observe that there must be ≥ 3 classes of size at least 2. Let V_1 be a larger class that we have added an edge between two of its vertices. Let V_2 be a larger class that we remove a vertex v of minimum degree from, and let V_3 be one of the remaining largest classes. Observe that removing v doesn't decrease the number of classes, so we are left with $G' = T_{k-1}(n - 1)$ plus an edge. This is a smaller graph so we apply induction to find H which is in G' . However, G contains G' , so H is also in G !

□

Exercise 11. Use Turán's theorem to show that among $3n$ points on the unit disk at least $3\binom{n}{2}$ of them are distance at most $\sqrt{2}$

Proof. Build a graph G with $V(G)$ = points on the unit disk. Connect two points iff the sum of their euclidean distances is greater than $\sqrt{2}$.

Claim: G contains no K_4 .

Proof. Pick a subset $S \subset V(G)$ of 4 distinct points. It will be sufficient to show that at least 1 edge between any of these points cannot be greater than $\sqrt{2}$. Thus this edge is not present in G and these 4 vertices cannot form a K_4 .

Consider the following cases:

1. S contains a vertex at the center of the unit disk

This case is easily resolved as this vertex is distance at most 1 from any other vertex, and cannot be contained in any K_4 (as its degree is 0 in G).

2. No vertex in S is the center of the unit disk

Draw a line from the center of the disk through each individual point to the edge of the disk. This partitions the disk into potentially four parts P_1, P_2, P_3 , and P_4 . If there are less than 4 parts then ≥ 2 vertices lie on the same line and cannot possibly be distance greater than $\sqrt{2}$ from one another.

Observe that the interior angle of one of these partitions must be $\leq 90^\circ$ as they all cannot have interior angle $\geq 90^\circ$. Let x and y be vertices on the boundary of such a partition with interior angle $\leq 90^\circ$. The distance between x and y is maximized when the angle of the partition is greatest (90°) and the points are pushed to the perimeter of the unit disk. Solving for the distance between them gives $d(x, y) = \sqrt{1^2 + 1^2} = \sqrt{2}$. Therefore, the edge xy is not present in G . Repeating this argument for all subsets S completes the claim.

□

Therefore G is K_4 free, and by Turán's theorem we get the following

$$e(G) \leq \left(1 - \frac{1}{4}\right) \frac{n^2}{2} = \frac{3}{8}n^2$$

The compliment of G has edges between points who's distance is at most $\sqrt{2}$. As we can have at most $\binom{n}{2}$ edges in any graph, the number of edges in \overline{G} is

$$\geq \binom{n}{2} - \frac{3}{8}n^2 = \frac{4n^2 - 3n}{8}$$

Therefore, selecting $3n$ points guarantees the number of edges

$$\geq \frac{12n^2 - 9n}{8} \geq 3\binom{n}{2}$$

As edges are pairs with distance at most $\sqrt{2}$ this completes the proof. □

Exercise 20. For a fixed graph F , show that the following function is decreasing as n increases

$$\frac{ex(n, F)}{\binom{n}{2}}$$

Proof. Let G be an extremal graph for F and count the number of pairs (e, v) where e is an edge of G and v is a vertex not incident to e . First pick an edge e , that edge is not incident to at most $n - 2$ vertices. Summing over all edges of G gives the following lower bound on the number of pairs

$$\geq \sum_{xy \in E(G)} (n - 2) = ex(n, F)(n - 2)$$

Now instead fix a vertex $v \in V(G)$, there are n such ways to do this. Observe that once we have fixed a vertex the maximum number of edges that v is not incident to is the extremal number for a graph on $n - 1$ vertices. This gives an upper bound on the number of pairs which is

$$\leq ex(n - 1, F)n$$

Combining these bounds we get

$$ex(n, F)(n - 2) \leq ex(n - 1, F)n$$

WLOG multiply each side by $\frac{(n-1)}{2}$

$$\frac{1}{2}(ex(n, F)(n - 1)(n - 2)) \leq \frac{1}{2}(ex(n - 1, F)n(n - 1))$$

Rearranging terms gives

$$\frac{ex(n, F)}{\frac{n(n-1)}{2}} \leq \frac{ex(n - 1, F)}{\frac{(n-1)(n-2)}{2}}$$

\implies

$$\frac{ex(n, F)}{\binom{n}{2}} \leq \frac{ex(n - 1, F)}{\binom{n-1}{2}}$$

Therefore, the function is decreasing as n grows larger. □

Exercise 24. Prove Lemma 1.29. Let G be a C_4 -free bipartite graph with class sizes a and b , then

$$e(G) \leq a\sqrt{b} + b$$

Proof. Let A be the class of size a and B the class of size b . Consider counting V structures centered on a vertex $v \in A$ in two different ways. First, fix a vertex $v \in A$ and count the number of V 's centered at that vertex. There are at least $\binom{d(v)}{2}$ such V 's. Summing over all vertices in A , we get the number of V 's

$$\begin{aligned} &\geq \sum_{v \in V(A)} \binom{d(v)}{2} \geq a \binom{\frac{1}{a} \sum d(v)}{2} = a \binom{e(G)/a}{2} \\ &\geq a \binom{(e(G)/(a-1))^2}{2} \geq a \binom{(\frac{e(G)}{a} - 1)^2}{2} \end{aligned}$$

Now count V 's coming from B to get a related upper bound. Observe that for every two vertices x and y in B they have at most 1 common neighbor (else we have a C_4). Therefore the number of V 's is

$$\leq \binom{b}{2} \leq \frac{b^2}{2}$$

Combining these bounds we get

$$a \binom{(\frac{e(G)}{a} - 1)^2}{2} \leq \frac{b^2}{2}$$

Solving for $e(G)$ gives the statement of lemma 1.29.

$$a \left(\frac{e(G)}{a} - 1 \right)^2 \leq b^2$$

$$\frac{e(G)}{\sqrt{a}} - \sqrt{a} \leq b$$

$$E(G) \leq b\sqrt{a} + a$$

Swapping the classes that we counted gives the other bound

$$E(G) \leq a\sqrt{b} + b$$

□