

Homework 1

CMPS130 Computational Models, Spring 2015

0.1

- a. The set of odd natural numbers greater than zero.
- b. The set of even integers.
- c. The set of even natural numbers great than one.
- d. The set of natural numbers divisible by six great than one.
- e. The set of palindromes over the alphabet $\{0,1\}$.
- f. The empty set.

0.2

- a. $\{1, 10, 100\}$
- b. $\{n \mid n \in \mathcal{Z} \wedge n > 5\}$
- c. $\{n \mid n \in \mathcal{N} \wedge n < 5\}$
- d. $\{\text{aba}\}$
- e. $\{\epsilon\}$
- f. $\{\}$

0.3

- a. No
- b. Yes
- c. $\{x, y, z\}$
- d. $\{x, y\}$
- e. $\{(x, x), (x, y), (y, x), (y, y), (z, x), (z, y)\}$

- f. $\{\{\}, \{x\}, \{y\}, \{x, y\}\}$

0.4

$$|A \times B| = |\{(a, b) \mid a \in A, b \in B\}| = \sum_{a \in A} \sum_{b \in B} 1 = |A| * |B|$$

0.5

$$|\mathcal{P}(C)| = |\{D \cup \{c \cup d \mid d \in D\} \mid c \in C \wedge D = \mathcal{P}(C \setminus c)\}| = 2 * |\mathcal{P}(C \setminus c)| = 2^{|C|}$$

0.6

a. $f(2) = 7$

b. $\text{dom}(f) = \{1, 2, 3, 4, 5\} \quad \text{range}(f) = \{6, 7\}$

c. $g(2, 10) = 6$

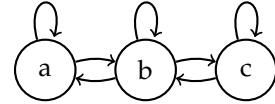
d. $\text{dom}(g) = \left\{ \begin{array}{l} (1, 6), (1, 7), (1, 8), (1, 9), (1, 10), \\ (2, 6), (2, 7), (2, 8), (2, 9), (2, 10), \\ (3, 6), (3, 7), (3, 8), (3, 9), (3, 10), \\ (4, 6), (4, 7), (4, 8), (4, 9), (4, 10), \\ (5, 6), (5, 7), (5, 8), (5, 9), (5, 10) \end{array} \right\} \quad \text{range}(g) = \{6, 7, 8, 9, 10\}$

e. $g(4, f(4)) = g(4, 7) = 8$

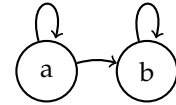
0.7

Using a finite set S and a relation $R \subseteq S \times S$.

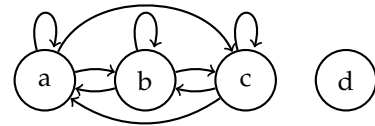
a. $S = \{a, b, c\}$
 $R = \{(a, a), (a, b), (b, a), (b, b), (b, c), (c, b), (c, c)\}$

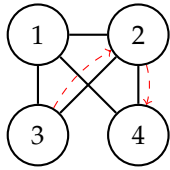


b. $S = \{a, b\}$
 $R = \{(a, a), (a, b), (b, b)\}$



c. $S = \{a, b, c, d\}$
 $R = \{(a, a), (a, b), (a, c), (b, a), (b, b), (b, c), (c, a), (c, b), (c, c)\}$



0.8

$$\deg(1) = 3 \quad \deg(3) = 2$$

0.9

$$G = (V, E) = (\{1, 2, 3, 4, 5, 6\}, \{(x, y) \mid x \in \{1, 2, 3\} \wedge y \in \{4, 5, 6\}\})$$

0.10

You cannot divide by $(a - b)$ unless you know that $a \neq b$. However, this contradicts with the assumption that $a = b$.

0.11

In the induction step, the proof assumes that H_1 and H_2 must have the same color if because all the horses in the set H_1 and H_2 have the same color and there is only one horse difference between H_1 and H_2 . However, this does not work for the base case $k = 2$ in which H_1 and H_2 will both only include a single horse. These sets are therefore trivially one-colored but might have different colors, therefore the set H might have more than one color.

0.12

Here, we only consider graphs that do *not* have self-loops: $\forall v \in V : (v, v) \notin E$. Otherwise, there is a simple counter-example with two nodes and exactly one edge: $G = (\{a, b\}, \{(a, a)\})$.

Let $G = (V, E)$ be a graph with $|V| \geq 2$ nodes, such that v_{\min} has the minimum degree of all nodes and v_{\max} has the maximum degree of all nodes.

Case 1: If $v_{\min} = 0$, then the node v_{\max} can at most be connected to all nodes other than v_{\min} in the graph, so $v_{\max} \leq |V| - 2$ and $v_{\max} - v_{\min} \leq |V| - 2$.

Case 2: If $v_{\min} \geq 1$, then the node v_{\max} can at most be connected to all nodes in the graph, so $v_{\max} \leq |V| - 1$ and $v_{\max} - v_{\min} \leq |V| - 1 - 1$.

In both cases, there are $|V|$ nodes but only $v_{\max} - v_{\min} + 1 = |V| - 1$ possible degrees, so by the pigeonhole principle, at least two nodes have the same degree.

Set intersection distribution $\overline{(A \cap B)} = (\overline{A} \cup \overline{B})$

If an element a is in $a \in \overline{(A \cap B)}$ then it is not in both A and B , so either it is not in A ($a \in \overline{A}$) or it is not in B ($a \in \overline{B}$) – in any case, it is in the union of both complements $(\overline{A} \cup \overline{B})$.

If an element a is in $\overline{A} \cup \overline{B}$ then it is either not in A ($a \in \overline{A}$) or not in B ($a \in \overline{B}$) – in any case, it is not in the intersection of both $((A \cap B))$.

Countable Sets

According to the definition, a set is countable if it is either finite or has the same size as \mathcal{N} .

Taking the set \mathcal{N} , the resulting set of multiplying all elements by two ($\{2n \mid n \in \mathcal{N}\}$) has the same size as \mathcal{N} and it is therefore countable. Subtracting 1 from each element of this set does not change its size, so the resulting set $\{2n - 1 \mid n \in \mathcal{N}\}$ includes all odd numbers but is countable.

Induction Example

Basis: For $n = 1$,

$$\sum_{i=1}^n i^2 = \sum_{i=1}^1 i^2 = 1^2 = \frac{6}{6} = \frac{1 \cdot 2 \cdot 3}{6} = \frac{1 \cdot (1+1)(2 \cdot 1 + 1)}{6} = \frac{n(n+1)(2n+1)}{6}$$

Induction Step: For $n = k + 1$, assuming that $\sum_{i=1}^k i^2 = \frac{k(k+1)(2k+1)}{6}$.

$$\begin{aligned} \sum_{i=1}^n i^2 &= \sum_{i=1}^{k+1} i^2 = (k+1)^2 + \sum_{i=1}^k i^2 \\ &= (k+1)^2 + \frac{k(k+1)(2k+1)}{6} && \text{(I.H.)} \\ &= \frac{6(k^2 + 2k + 1) + (k^2 + k)(2k + 1)}{6} \\ &= \frac{6k^2 + 12k + 6 + 2k^3 + k^2 + 2k^2 + k}{6} \\ &= \frac{2k^3 + 9k^2 + 13k + 6}{6} \\ &= \frac{(k+1)(2k^2 + 7k + 6)}{6} \\ &= \frac{(k+1)(k+2)(2k+3)}{6} \\ &= \frac{(k+1)((k+1)+1)(2(k+1)+1)}{6} \\ &= \frac{n(n+1)(2n+1)}{6} \end{aligned}$$