Distributed Database Management Systems

Chapter 2

Distributed DBMS Architecture

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Distributed DBMS Architecture

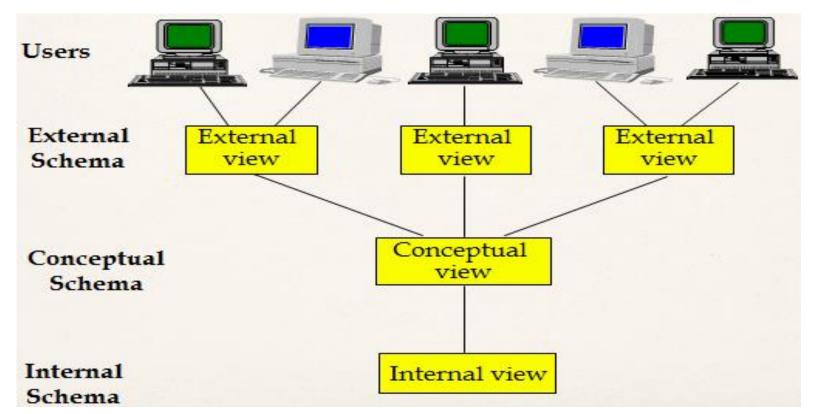
- Architecture of a system defines the structure of the system
 - → components identified
 - → functions of each component defined
 - → interrelationships and interactions between components defined
- "ANSI/SPARC architecture" is a *data logical* approach to defining a DBMS architecture.
 - → focuses on the different user classes and roles on data.
- A centralized DBMS architecture that extends to identify the set of alternative architectures for a distributed DBMS.
- Three "reference" architectures for a distributed DBMS:
 - → client/server systems
 - → peer-to-peer distributed DBMS, and
 - → multidatabase systems
- ANSI stands for *American National Standards Institute* and SPARC stand for *Standards Planning and Requirement Committee*.

ANSI/SPARC architecture:

- The interfaces of this architecture were proposed to be standardized.
- Defines a framework that contains 43 interfaces, 14 of which deal with physical storage subsystem
- There are three views of data:
 - → the external view, which is that of the end user, who might be a programmer;
 - → the internal view, that of the system or machine; and
 - → the conceptual view, that of the enterprise.
- For each of these views, an appropriate schema definition is required.
- At the lowest level of the architecture is the internal view
 - → deals with the physical definition and organization of data.

ANSI/SPARC architecture(cont....)

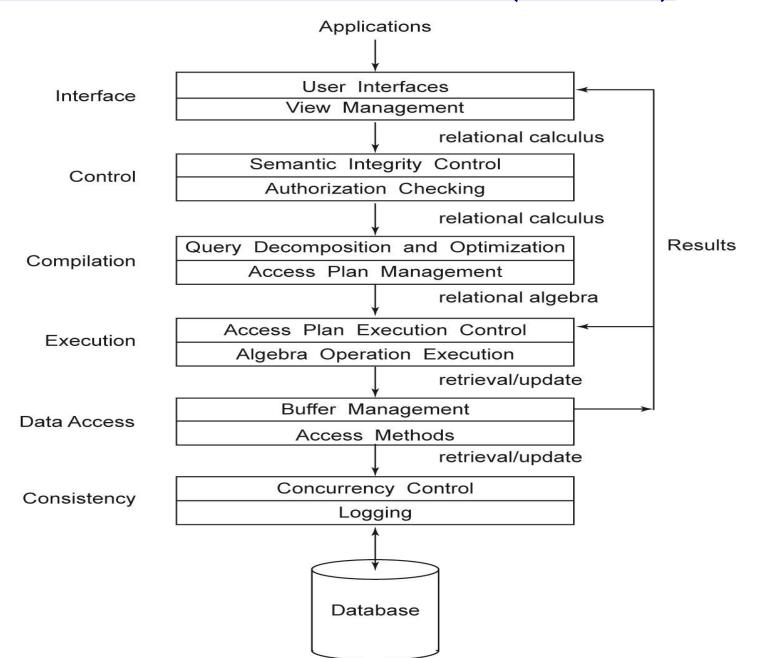
- At the other extreme is the external view, which is concerned with how users view the database.
- In between these two ends is the conceptual schema, which is an abstract definition of the database.
 - → It is the "real world" view of the enterprise being modeled in the database



Centralized DBMS Architecture

- A centralized DBMS is interfaced with two other components:
 - → the communication subsystem permits interfacing the DBMS with other subsystems in order to communicate with applications
 - → the operating system provides the interface between the DBMS and computer resources (processor, memory, disk drives, etc.).
- The function's performed by a DBMS can be layered: the layers are the
 - interface,
 - control,
 - compilation,
 - execution,
 - data access, and
 - consistency management

Centralized DBMS Architecture(cont...)



Centralized DBMS Architecture(cont...)

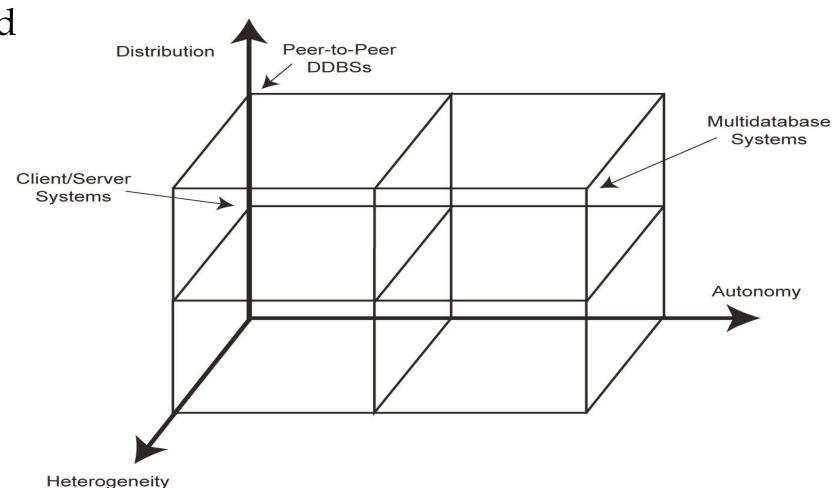
- The interface layer manages the interface to the applications.
- The control layer controls the query by adding semantic integrity predicates and authorization predicates.
- The query processing (or compilation) layer maps the query into an optimized sequence of lower-level operations.
- The execution layer directs the execution of the access plans, including transaction management and synchronization of algebra operations
- The data access layer manages the data structures that implement the files, indices, etc.
- The consistency layer manages concurrency control and logging for update requests.

Architectural Models for DDBMSs

- Architectural models for DDBMS can be characterized with respect to three different dimensions:
 - → the autonomy of local systems,

→ their distribution, and

→ their heterogeneity.



Architectural Models for DDBMSs(Autonomy)

- Autonomy refers to the distribution of control, not of data.
 - →It indicates the degree to which individual DBMSs can operate independently.
- Dimensions of autonomy can be specified as:
 - → Design autonomy: Ability of a component DBMS to decide on issues related to its own design.
 - → Communication autonomy: Ability of a component DBMS to decide whether and how to communicate with other DBMSs.
 - → Execution autonomy: Ability of a component DBMS to execute local operations in any manner it wants to.

Architectural Models for DDBMSs(Autonomy) cont...

- Dimensions of autonomy are covered in three different aspects:
 - Tight integration: a single-image of the entire database is available to any user
 - →Semiautonomous systems: Consist of DBMSs that can operate independently and determine what parts of their own database will be accessible to users of other DBMSs.
 - →Total isolation: Individual systems are stand-alone DBMSs that know neither of the existence of other DBMSs nor how to communicate with them.

Architectural Models for DDBMSs(Distribution)

- The distribution dimension of the taxonomy deals with the physical distribution of data over multiple sites.
- Two classes of distribution are:
 - →Client/server distribution: Concentrâtes data management duties at servers while the clients focus on providing the application environment including the user interface.
 - → Peer-to-peer distribution (or full distribution): There is no distinction of client machines versus servers.
 - ◆ Each machine has full DBMS functionality and can communicate with other machines to execute queries and transactions

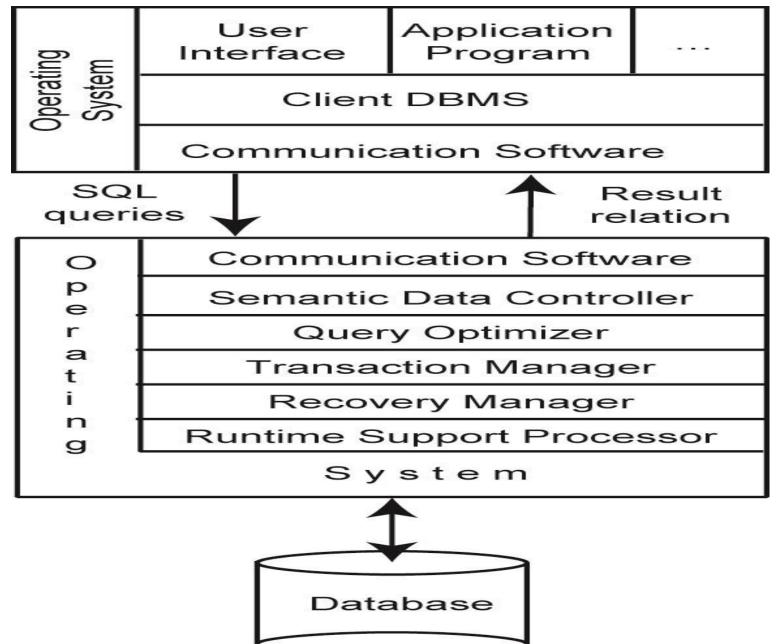
Architectural Models for DDBMSs(Heterogeneity)

- Heterogeneity may occur in various forms in DDBS:
 - →Heterogeneity in data models Created by the data representation with different modeling tools.
 - →Heterogeneity in query languages involves differences in languages.
 - →Heterogeneity in transaction management protocols the use of different data access paradigms in different data models.

Architectural Alternatives

- Three alternative architectures are:
 - →Client/server distributed DBMS provides a two-level architecture based on server functions and client functions
 - →Peer-to-peer distributed DBMS no differentiation between the functionality of each site in the system.
 - →Heterogeneous multidatabase system individual DBMSs (whether distributed or not) are fully autonomous and have no concept of cooperation.

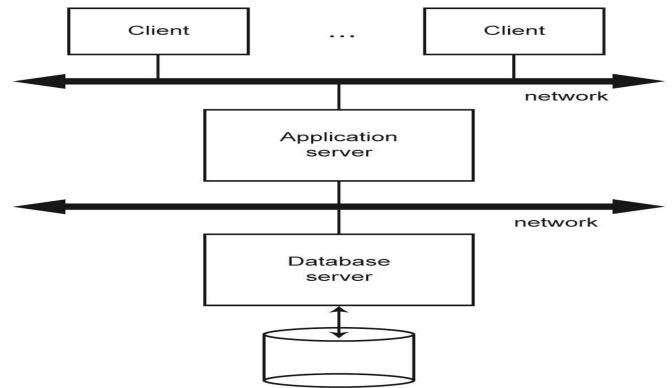
- Provides a two-level architecture based on server functions and client functions.
- Server stores and administers the data.
- All of query processing and optimization, transaction management and storage management is done at the server.
- The client provides:
 - → the application and the user interface
 - → a DBMS client module responsible for managing the data that is cached to the client
 - → The management of transaction lock that is cached to the client
 - consistency checking of user queries at the client side
- Communication between the clients and the server(s) is at the level of SQL statements.
 - → Client passes SQL queries to the server, the server returns the result



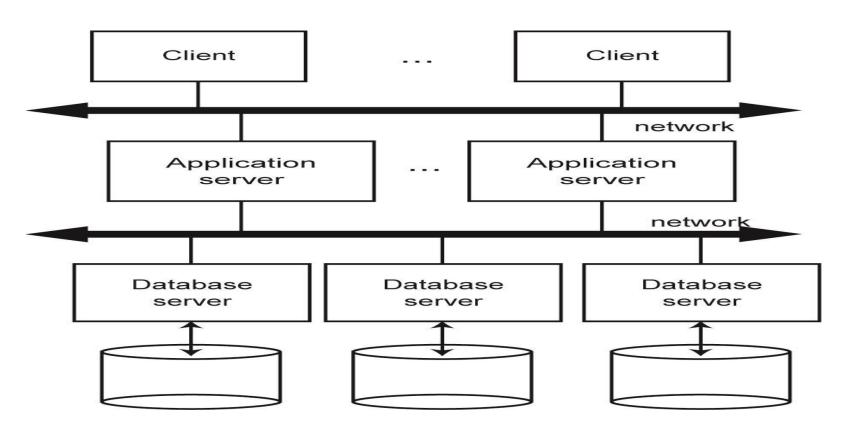
- Two simple types of client/server architecture:
 - →multiple client/single server: the database is stored only on the server that also hosts the software to manage it.
 - →multiple client/multiple server :
 - two alternative management strategies:
 - ✓ either each client manages its own connection to the appropriate server or
 - ✓ each client knows of only its "home server" which then communicates with other servers as required

- Three types of servers in client/server architecture:
 - → client servers run the user interface (e.g., web servers),
 - → application servers run application programs, and
 - → database servers run database management functions.
- Data Base Server Approach: An extension of the client/server architecture, with application servers connected to one database server via

a communication network.



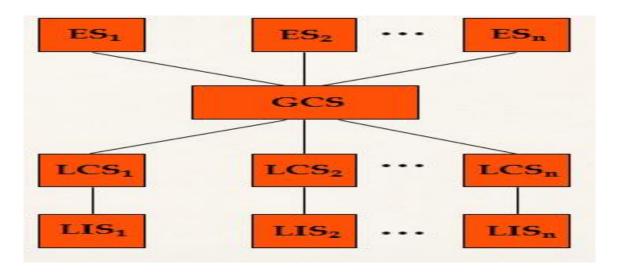
- Distributed Data Base Servers Approach:
 - → Introduces multiple database servers and multiple application servers
 - → Each application server is dedicated to one or a few applications, while database servers operate in the multiple server fashion



Architectural Alternatives (Peer-to-peer distributed DBMS)

- In peer-to-peer architecture:
 - → there is no differentiation between the functionality of each site in the system.
 - there is inherent heterogeneity of every aspect of the sites and their autonomy
 - → the physical data organization on each machine may be different:
 - → local internal schema (LIS) an individual internal schema definition at each site.
 - → global conceptual schema (GCS) The enterprise view of the data, which is global and describes the logical structure of the data at all the sites.
 - ◆ local conceptual schema (LCS) handles data fragmentation and replication at each site
 - external schemas (ESs) supports user applications and user access to the database

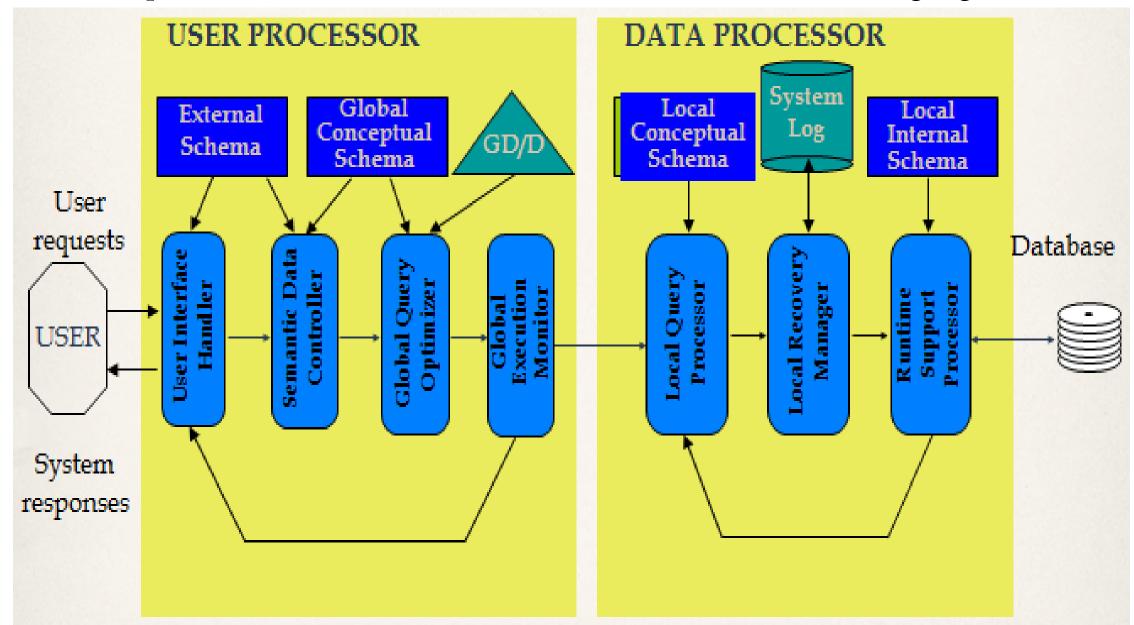
<u>Architectural Alternatives (Peer-to-peer distributed DBMS)cont....</u> Peer-to-peer Datalogical DDBS Architecture:



- Data independence is supported as the model is an extension of ANSI/SPARC
- Location and replication transparencies are supported by the definition of LCSs and GCSs schemas and the mapping in between.
- Network transparency is supported by the definition of the GCS
- The DDBS translates global queries into a group of local queries, which are executed by DDBMS components at different sites that communicate with one another.

Architectural Alternatives (Peer-to-peer distributed DBMS)cont....

The detailed components of a distributed DDBMS are shown in the following Figure:



Architectural Alternatives (Peer-to-peer distributed DBMS)cont....

- Two first major components of peer-to-peer DDBMS are:
 - → User processor
 - → Data processor
- User processor consists of four elements:
 - → user interface handler interprets user commands and formats the result data as it is sent to the user.
 - → semantic data controller uses the integrity constraints and authorizations (defined as part of GCS) to check if the user query can be processed.
 - → global query optimizer and decomposer determines an execution strategy to minimize a cost function, and translates the global queries into local ones using GCS, LCS and global directory
 - → distributed execution monitor (or distributed transaction manager)-coordinates distributed execution of the user request.

Architectural Alternatives (Peer-to-peer distributed DBMS)cont....

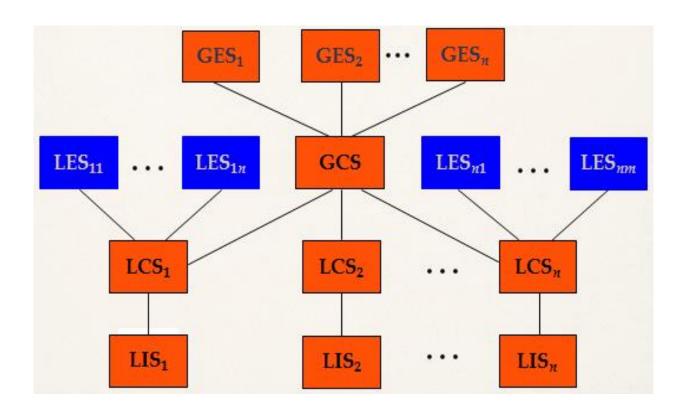
- Data processor consists of three elements:
 - →local query optimizer acts as the access path selector responsible for choosing the best access path to access any data item
 - →local recovery manager responsible for making sure that the local database remains consistent even when failures occur
 - →run-time support processor physically accesses the database according to the physical commands in the schedule generated by the query optimizer

Architectural Alternatives (Multidatabase System Architecture)

- In Multidatabase systems (MDBS), the individual DBMSs (whether distributed or not) are fully autonomous and have no concept of cooperation.
- Distributed DBMS vs. Distributed Multi-DBMS:
 - → In distributed DBMS, the GCS defines the conceptual view of the entire database;
 - → in distributed MDBMS, it represents only the collection of some of the local databases that each local DBMS wants to share.
 - →Global database in the latter is equal to the union of local databases,
 - ◆ whereas in the former it is only a (possibly proper) subset of the same union.

Architectural Alternatives (Multidatabase System Architecture) cont...

- Distributed DBMS vs. Distributed Multi-DBMS:
 - → Difference between the design of the GCS in multi-DBMSs and distributed DBMSs:
 - → in the former the mapping is from LCS to a GCS. In the latter, however, mapping is in the reverse direction.
- Datalogical Multi-DBMS Architecture :



Architectural Alternatives (Multidatabase System Architecture) cont....

- If heterogeneity exists in the system, then two implementation alternatives exist:
 - -unilingual MDBMS and
 - -Multilingual MDBMS