## Properties of relations on a set determined by a matrix

Mr Ashlin Darius Govindasamy University of South Africa Department of Computer Science and Mathematics

November 20, 2022

#### Abstract

This paper introduces a technique of converting a relation set into a matrix of  $\mathbb{R}^n$  space using that matrix classifying the properties for the relational set. Using this technique, we can find the properties of a relation set by using some properties of the matrix. In programming, we can use this technique to determine our properties of a relation set by using the matrix. Code snippets are provided to show how this technique can be used in programming.

# Contents

1	Relations			2
	1.0.1	What is a re	elation?	. 2
		1.0.1.1 So	ome specific relations	
	1.0.2	Properties of	of relations	
		1.0.2.1 Re	eflexive	
		1.0.2.2 Irr	reflexive	. 4
		1.0.2.3 Sy	m rmmetric	. ;
		1.0.2.4 Ar	ntisymmetric	. 6
		1.0.2.5 Tr	ransitive	. 6
		1.0.2.6 Eq.	quivalence relations	. 7
			$\operatorname{rder}$ relations	

### Chapter 1

# Relations

### 1.0.1 What is a relation?

A (binary) relation R between sets A and B is a subset of  $A \times B$ . ( $A \times B$  is a Cartesian product.) Thus, a relation is a set of pairs.

The interpretation of this subset is that it contains all the pairs for which the relation is true. We write aRb if the relation is true for A and B (equivalently B, if  $(A, B) \in R$ ).

A and B can be the same set, in which case the relation is said to be "on" rather than "between": A binary relation R on a set A is a  $\subseteq A \times A$ .  $(A \times A$  is a Cartesian product.)

Example of a relation using  $A = \{0, 1, 2, 3\}$ 

$$R = \{(0,0), (1,1), (2,2), (3,3)\}$$

Relations may also be of other arities. An *n*-ary relation R between sets  $X_1, \ldots,$  and  $X_n \subseteq n$ -ary product  $X_1...X_n$ , in which case R is a set of n-tuples.

#### 1.0.1.1 Some specific relations

The empty relation between sets X and Y, or on E, is the empty set  $\emptyset$ .

The empty relation is false for all pairs.

The full relation (or universal relation) between sets X and Y is the set  $X \times Y$ .

The full relation on set E is the set  $E \times E$ .

The full relation is true for all pairs.

The identity relation on set E is the set  $(x, x)|x \in E$ .

The identity relation is true for all pairs whose first and second element are identical.

### 1.0.2 Properties of relations

A relation R is	if
reflexive	xRx
symmetric	xRy implies yRx
transitive	xRy and yRz implies xRz
irreflexive	xRy implies x\neq y
antisymmetric	xRy and yRx implies x=y
triohotomy	xRy or x=y or yRx

#### 1.0.2.1 Reflexive

Let Relation R on A

Where  $R \subseteq A \times A$ 

The reflexive property of a relation is that  $\forall a \in A, aRa$  is true.

Also could be written as  $R \subseteq A \times A$  and  $\forall a \in A$ , then  $(a, a) \in R$ 

Example:

Let  $A = \{1, 2, 3, 4\}$ 

Let  $R = \{(1,1), (2,2), (2,3), (3,4), (3,3), (4,4), (4,2)\}$ 

We could draw a matrix to show the relation R on A

To draw the matrix we plot the elements of A on the x-axis and y-axis.

Then we plot the pairs of R on the matrix as 1.

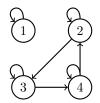
If the pair is not in R then we plot a 0.

	1	2	3	4
1	1	0	0	0
2	0	1	1	0
3	0	0	1	1
4	0	1	0	1

We can also draw a directional graph to show the relation R on A

We plot the elements of A as nodes.

We plot the pairs of R as directed edges.



Thus resulting in the relation R being reflexive.

This is because for all  $a \in A$ , aRa is true.

For example, 1R1 is true, 2R2 is true, 3R3 is true, 4R4 is true.

Therefore R is reflexive.

In the matrix, we can see that the diagonal is all 1's. If you notice, the diagonal is the pairs (a, a) for all  $a \in A$ . In programming we look at this problem as a 2D array.

Using Mathematics logic, we can write this as:  $R \subseteq A \times A$  and  $\forall a \in A$ , then  $(a,a) \in R$ 

Where R is a 2D array.

Or in psuedocode:

```
\begin{array}{l} function \ is Reflexive \left( R \right) \\ bValid \ = \ True \\ for \ i \ = \ 0 \ to \ R. \, length \\ for \ j \ = \ 0 \ to \ R. \, length \\ if \ R[\,i\,][\,j\,] \ = \ 0 \\ bValid \ = \ False \\ end \ if \\ end \ for \\ end \ for \\ return \ bValid \end{array}
```

#### 1.0.2.2 Irreflexive

Let Relation R on A

Where  $R \subseteq A \times A$ 

The irreflexive property of a relation is that  $\forall a \in A, aRa$  is false. Also could be written as  $R \subseteq A \times A$  and  $\forall a \in A$ , then  $(a, a) \notin R$ 

Example:

Let 
$$A = \{1, 2, 3, 4\}$$
  
Let  $R = \{(1, 2), (2, 3), (3, 4), (4, 1)\}$ 

We could draw a matrix to show the relation R on A

To draw the matrix we plot the elements of A on the x-axis and y-axis.

Then we plot the pairs of R on the matrix as 1.

If the pair is not in R then we plot a 0.

	1	2	3	4
1	0	1	0	0
2	0	0	1	0
3	0	0	0	1
4	1	0	0	0

We can also draw a directional graph to show the relation R on A We plot the elements of A as nodes.

We plot the pairs of R as directed edges.



Thus resulting in the relation R being irreflexive.

This is because for all  $a \in A$ , aRa is false.

For example, 1R1 is false, 2R2 is false, 3R3 is false, 4R4 is false.

Therefore R is irreflexive.

In the matrix, we can see that the diagonal is all 0's.

If you notice, the diagonal is the pairs (a, a) for all  $a \in A$ .

In programming we look at this problem as a 2D array.

 $\begin{array}{c} bValid = False\\ end \ if\\ end \ for \end{array}$ 

end for return bValid

If you want to be real and save time coding.

```
\begin{array}{c} function \quad is Irreflexive \, (R) \\ \quad return \quad is Reflexive \, (R) \\ end \quad function \end{array}
```

#### 1.0.2.3 Symmetric

Let Relation R on A be symmetric if  $\forall a, b \in A$  then  $(a, b) \in R$   $(b, a) \in R$ 

**Example:** 
$$R = (1, 1), (1, 3), (2, 3), (2, 4), (3, 1), (3, 2), (4, 2)$$

We can draw a matrix to show the relation R on A

To draw the matrix we plot the elements of A on the x-axis and y-axis.

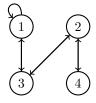
Then we plot the pairs of R on the matrix as 1.

If the pair is not in R then we plot a 0.

	1	2	3	4
1	1	0	1	0
2	0	0	1	1
3	1	1	0	0
4	0	1	0	0

We can also draw a directional graph to show the relation R on A We plot the elements of A as nodes.

We plot the pairs of R as directed edges.



Notice that the relation R is symmetric.

This is because for all  $a, b \in A$ , aRb is true.

aRb is true if and only if bRa is true.

For example, 1R1 is true, 1R3 is true, 2R3 is true, 2R4 is true, 3R1 is true, 3R2 is true, 4R2 is true. Therefore R is symmetric.

In the matrix, we can see that the matrix is symmetric.

Using our knowledge of matrices, we can write this as:  $R \subseteq A \times A$  and  $R = R^T$ 

Where R is a 2D array.

Remember that  $R^T$  is the transpose of R.

The transpose of a matrix is the matrix found by interchanging its rows into columns or columns into rows.

In programming we look at this problem as a 2D array.

```
\begin{array}{ccc} function & isSymmetric\left(R\right) \\ & RT = transpose\left(R\right) \\ & return & R \Longrightarrow RT \\ end & function \end{array}
```

#### 1.0.2.4 Antisymmetric

Let Relation R on A be antisymmetric if  $\forall a,b \in A$  then  $(a,b) \in R$   $(b,a) \in R$  Or in psuedocode:

```
function isAntisymmetric(R)
  bSymmetric = isSymmetric(R)
  if bSymmetric == True
      return False
  else
      return True
  end if
end function
```

#### 1.0.2.5 Transitive

A relation R on A is transitive if  $\forall a, b, c \in A$  then  $(a, b) \in R$  and  $(b, c) \in R$  then  $(a, c) \in R$ 

```
Example: R = (1, 1), (1, 3), (2, 3), (2, 4), (3, 1), (3, 2), (4, 2)
```

We can draw a matrix to show the relation R on A

To draw the matrix we plot the elements of A on the x-axis and y-axis.

Then we plot the pairs of R on the matrix as 1.

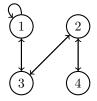
If the pair is not in R then we plot a 0.

	1	2	3	4
1	1	0	1	0
2	0	0	1	1
3	1	1	0	0
4	0	1	0	0

We can also draw a directional graph to show the relation R on A

We plot the elements of A as nodes.

We plot the pairs of R as directed edges.



For this problem we calculate the  $M^k$  where each element identifies the number of paths of length. So for example  $M^2$  is the number of paths of length 2.

$$M^2=M\cdot M$$

$$M = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

$$M^{2} = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

$$M^{2} = \begin{bmatrix} 1 \cdot 1 + 0 \cdot 0 + 1 \cdot 1 + 0 \cdot 0 & 1 \cdot 0 + 0 \cdot 0 + 1 \cdot 1 + 0 \cdot 1 & 1 \cdot 1 + 0 \cdot 1 + 1 \cdot 0 + 0 \cdot 0 & 1 \cdot 0 + 0 \cdot 1 + 1 \cdot 0 + 0 \cdot 0 \\ 0 \cdot 1 + 0 \cdot 0 + 1 \cdot 1 + 1 \cdot 1 & 0 \cdot 0 + 0 \cdot 0 + 1 \cdot 1 + 1 \cdot 1 & 0 \cdot 1 + 0 \cdot 1 + 1 \cdot 0 + 1 \cdot 0 & 0 \cdot 0 + 0 \cdot 1 + 1 \cdot 0 + 1 \cdot 0 \\ 1 \cdot 1 + 1 \cdot 0 + 0 \cdot 1 + 0 \cdot 0 & 1 \cdot 0 + 1 \cdot 0 + 0 \cdot 1 + 0 \cdot 1 & 1 \cdot 1 + 1 \cdot 1 + 0 \cdot 0 + 0 \cdot 0 & 1 \cdot 0 + 1 \cdot 1 + 0 \cdot 0 + 0 \cdot 0 \\ 0 \cdot 1 + 1 \cdot 0 + 0 \cdot 1 + 0 \cdot 0 & 1 \cdot 0 + 1 \cdot 0 + 0 \cdot 1 + 0 \cdot 1 & 1 \cdot 1 + 1 \cdot 1 + 0 \cdot 0 + 0 \cdot 0 & 1 \cdot 0 + 1 \cdot 1 + 0 \cdot 0 + 0 \cdot 0 \end{bmatrix}$$

$$M^{2} = \begin{bmatrix} 2 & 1 & 2 & 0 \\ 0 & 0 & 2 & 2 \\ 2 & 2 & 0 & 0 \\ 0 & 2 & 0 & 0 \end{bmatrix}$$

Looking at the matrix there must be at least one path of length 2 and length 1. Thus this is a transitive relation.

Looking at our psuedocode code.

```
def transitive (R):
    Rs = R**2
    bOnePass = False
    for i in range (len(Rs)):
        for j in range (len (Rs)):
            if Rs[i][j] = 1:
                 bOnePass = True
    if bOnePass:
        bTwoPass = False
        for i in range (len(R)):
            for j in range (len(R)):
                 if R[i][j] == 2:
                     bTwoPass = True
                     break
    if bOnePass and bTwoPass:
        return True
    else:
        return False
```

#### 1.0.2.6 Equivalence relations

An equivalence relation is a relation that is reflexive, symmetric, and transitive.

An equivalence relation partitions its domain E into disjoint equivalence classes. Each equivalence class contains a set of elements of E that are equivalent to each other, and all elements of E equivalent to any element of the equivalence class are members of the equivalence class. The equivalence classes are disjoint: there is no  $x \in E$  such that x is in more than one equivalence class. The equivalence classes exhaust E: there is no  $x \in E$  such that x is in no equivalence class. Any element of an equivalence class may be its representative; the representative stands for all the members of its equivalence class.

#### 1.0.2.7 Order relations

An order (or partial order) is a relation that is antisymmetric and transitive.

A strict order is one that is **irreflexive** and **transitive**; such an order is also trivially **antisymmetric** because there is no x and y such that xRy and yRx.

A non-strict (weak) order is one that is reflexive, antisymmetric, and transitive.

An order relation R on E is a **total order** if either xRy or yRx  $\forall$ x,y $\in$  E.

An order relation R on E is a **partial order** if there is a  $x,y \in E$  for which neither xRy nor yRx.

And order relation R on E is a **Triohotomy** if either xRy, yRx, or  $x=y \ \forall x,y \in E$ .

A weak total order is reflexive, antisymmetric, transitive and trichotomy.

A strict total order is irreflexive, transitive, and trichotomy.